Batch Name : PreCAT OM22 - Fast Track Batch

Module Name: Data Structures

Q. What is a data structure?

We want to store marks of 100 students int m1, m2, m3, m4,, m100;//400 bytes - we want to sort marks of 100 students in a descenfding order sorting

int marks[100];//400 bytes

we want to store info of 100 students:

rollno: int name: char[] marks: float

struct employee emp[100];

- to learn data structure is not to learn any programming language, it is nothing but to learn **algorithms**.

Q. What is a Program?

- a program is a finite set of instructions written in any programming language given to the machine to do specific task.

Q. What is an algorithm?

- an algorithm is a finite set of instructions written in human understandable language like english, if followed, acomplishesh given task.
- A Program is an impementation of an Algorithm
- An Algorithm is like a blueprint of a Program.

Algorithm ==> User / Pseudocode ==> Programmer User Program ==> Machine

Q. What is a Pseudocode?

- an algorithm is a finite set of instructions written in human understandable language like english with some **programming constraints**, if followed, acomplishesh given task, such algo is also called as pseudocode.
- pseudocode is a special form of an algo

traversal on an array/to scan array => to visit each array element sequentially from first element max till last element.

```
Algorithm: to do sum of array elements:
step-1: initially take sum as 0.
step-2: start traversal of an array and add each array element into the sum
sequentially.
Step-3: return final sum
Pseudocode/Special form of an algo: ==> Programmer User
Algorithm ArraySum(A, n)//A is an array having size n
{
     sum = 0:
     for( index = 1; index \leq n; index++){
           sum += A[index];
     }
     return sum;
}
Program: ==> Machine
int array_sum( int arr[ ], int size ){
     int sum = 0;
     for(int index = 0; index < size; index++){
           sum += arr[index];
     return sum:
}
- An algorithm is a solution of a given problem
- algorithm = solution
- Problem: can we have multiple solutions for single problem
Pune => Mumbai
multiple paths exists between Pune & Mumbai
efficient/optmized path ==>
     distance in km
     cost required for
     medium
     traffic conditions
```

- One problem may has multiple solutions

Searching: to search given key element into a collection/list of elements

- 1. lienar search
- 2. binary search

Sorting: to arrange data elements in a collection/list of elements either in an ascending order (or in a desceding order).

- 1. bubble sort
- 2. selection sort
- 3. insertion sort
- 4. merge sort
- 5. quick sort
- 6. heap sort
- 7. radix sort
- 8. shell sort

etc....

- When we have multiple solustions/algo's for a single problem, we need to select an efficient solution/algo out of them, and to decide efficiency of an algo's we need to do their analysis.
- analysis of an algo is a work of calculating how much time i.e. computer time and space i.e. computer memory it needs to run to completion.
- there are two measures of analysis of an algo:
- **1. time complexity** of an algo is the amount of time i.e. computer time it needs to run to completion.
- **2. space complexity** of an algo is the amount of space i.e. computer memory it needs to run to completion.

Linear Search:

```
Best case: occurs if key is found at first position in only 1 comparison: O(1) for size of an array = 10 \Rightarrow no. of comparisons = 1 for size of an array = 20 \Rightarrow no. of comparisons = 1 for size of an array = 30 \Rightarrow no. of comparisons = 1.

.

for size of an array = 50 \Rightarrow no. of comparisons = 1 for size of an array = 100 \Rightarrow no. of comparisons = 1
```

Worst case : occurs if either key is found at last position or key is not found O(n).

```
for size of an array = 10 \Rightarrow no. of comparisons = 10 for size of an array = 20 \Rightarrow no. of comparisons = 20 for size of an array = 30 \Rightarrow no. of comparisons = 30.

for size of an array = 50 \Rightarrow no. of comparisons = 50 for size of an array = 100 \Rightarrow no. of comparisons = 100 for size of an array = n \Rightarrow no. of comparisons = no
```

best case: if an algo takes min amount of time to run to completion **worst case:** if an algo takes max amount of time to run to completion **average case:** if an algo takes neither min nor max amount of time to run to completion

for size of an array =
$$10 \Rightarrow 20$$
 bytes/40 bytes ==> 10 units for size of an array = $20 \Rightarrow 40$ bytes/80 bytes ==> 20 units

for size of an array = n ==> n units Space Complexity = O(n).

+ Rule: if running time of an algo is having any additive/substractive/multiplicative/divisive constant then it can be neglected. e.g.

$$O(n+3) => O(n)$$

 $O(n-4) => O(n)$
 $O(n/3) => O(n)$
 $O(2*n) => O(n)$

+ Binary Search:

by means of calculating mid pos big size array gets divided logically into two subarrays: left subarray & right subarray

for left subarray => value of left remains same, right = mid-1 for right subarray => value of right remains same, left = mid+1

```
if size of an array = 1000
iteration-1: [ 0 1 2 3 ..... 1000 ] => mid -> 1 comparison => 500
[ 0.. 499 ] 500 [ 501 .... 1000 ]
iteration-2: [ 501 .... 1000 ] => mid = 750 => 1 comparison => 250
[ 501...... 749 ] 750 [ 751 .... 1000 ]
iteration-3: [ 501 .... 750] 1 comparison => 125
```

```
after iteration-1: no. of cmp = 1, n/2 => T(n/2^1) + 1 after iteration-2: no. of cmp = 2, n/4 => T(n/2^2) + 2 after iteration-3: no. of cmp = 3, n/8 => T(n/2^3) + 3
```

let us assume k no. of iterations takes after iteration-k: no. of cmp = k, $n/2^K$ => $T(n/2^K) + K$

