

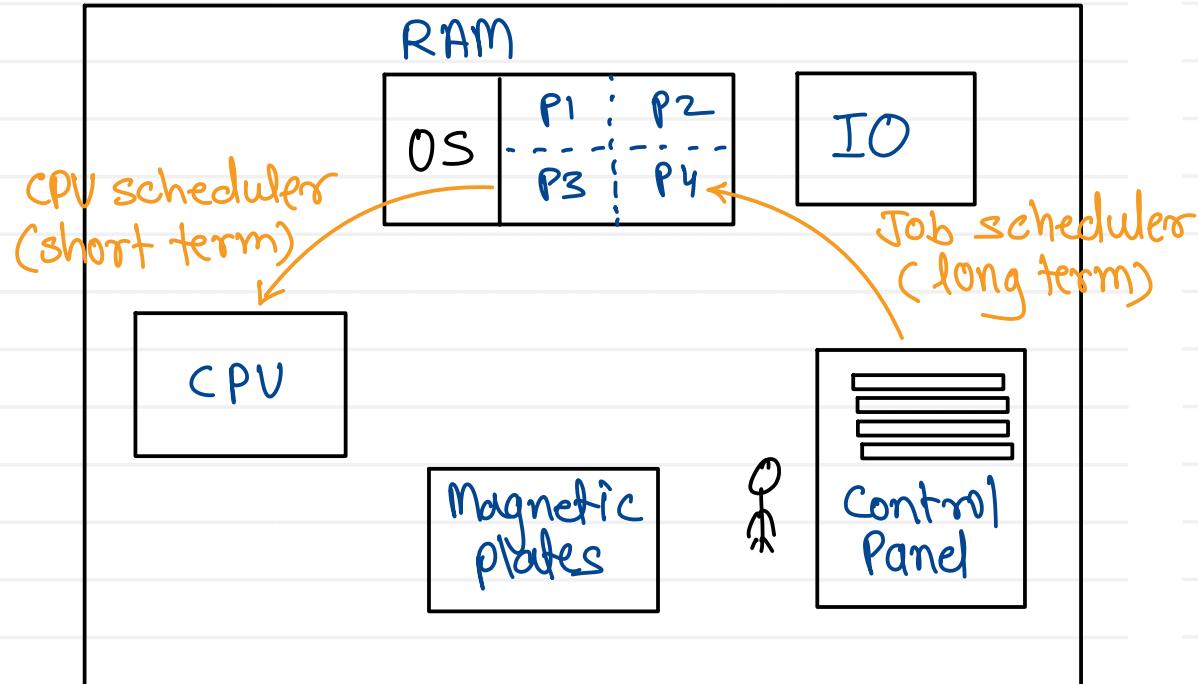


**Sunbeam Institute of Information Technology  
Pune and Karad**

## **Module - Concepts of Operating System**

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# Types of Operating system



1. Resident monitor
2. Batch system
3. Multiprogramming system:
  - multiple programs are loaded inside RAM (memory) at a time.

CPV burst: time spent by process on CPU

IO burst : time spent by process to perform IC

CPV burst > IO burst : CPV bound process

IO burst > CPV burst : IO bound process

Degree of multiprogramming :  
No. of programs loaded into RAM

- mixture of CPU & IO bound processes was getting loaded inside RAM to balance CPU & IO load.



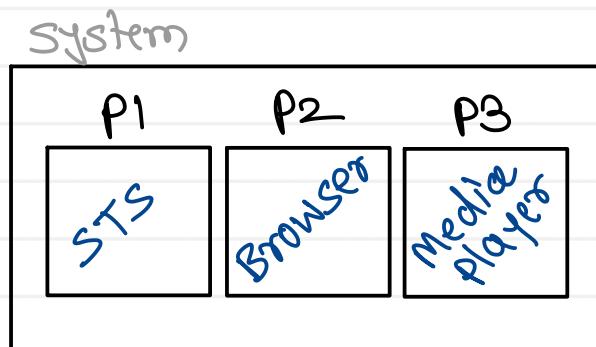
# Types of Operating system

4. Time Sharing system ( Multitasking system) :

- CPU time is shared in all the processes of RAM
- Response time < 1 sec

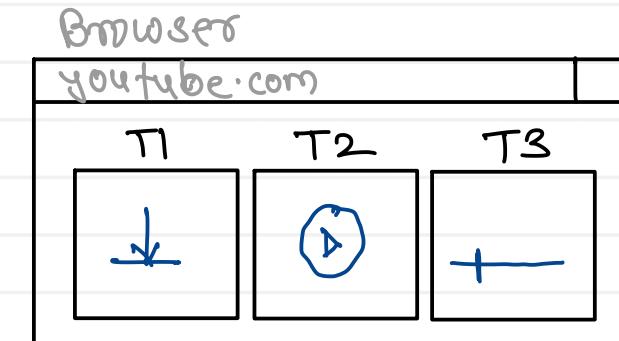
There are two types of Multitasking.

1. Process based multitasking



- system wide multitasking

2. Thread based Multitasking  
( Multithreading)

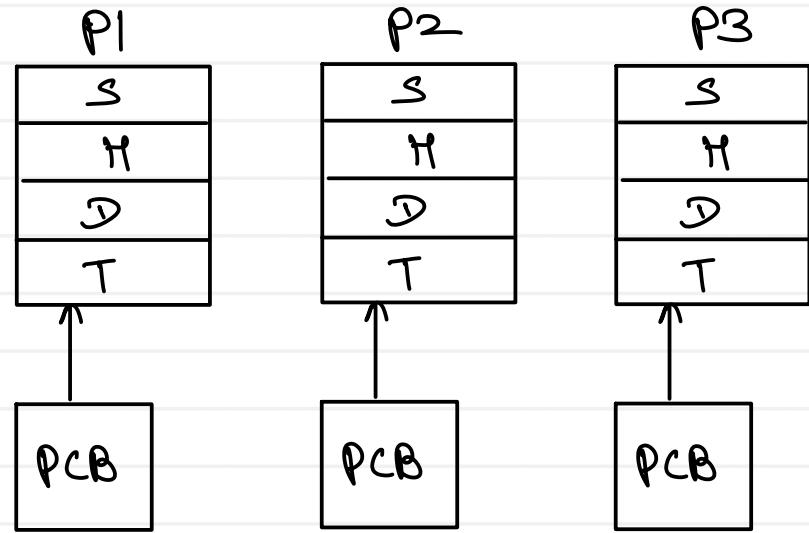


- multitasking within process

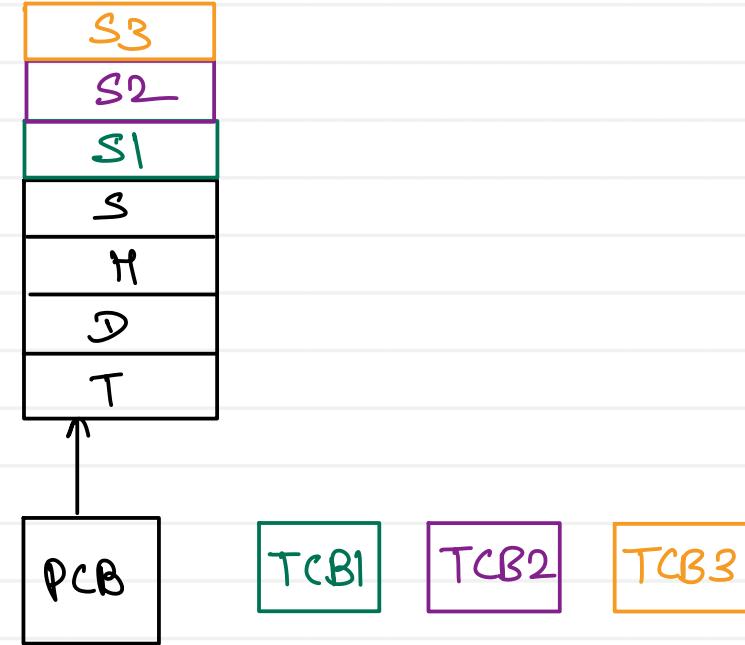


# Process Vs Thread

- processes are created inside system



- Thread are always created inside process



- Process is container of resources
- thread is alive entity in that container.
- process never runs on CPU, it is a thread which runs on CPU.

- Thread is light weight process

# Types of Operating system

## 5. Multiprocessing system :

- multiple processors (CPUs) are fitted on single chip. such chips are called as "multiprocessor" / "multicore".
- OS starts scheduling processes for multiple cores.
- multiple instructions / processes can be processed parallelly.

### i. Symmetric multiprocessing :

- OS treats every core equally & schedules separate process for every core.

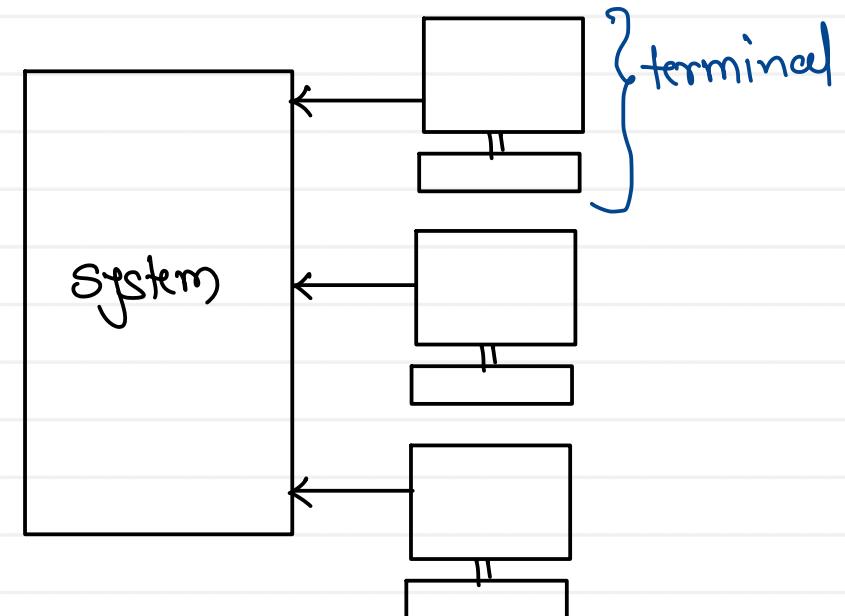
### ii Asymmetric multiprocessing :

- OS schedules process for only one core & internally that core divides job into remaining cores

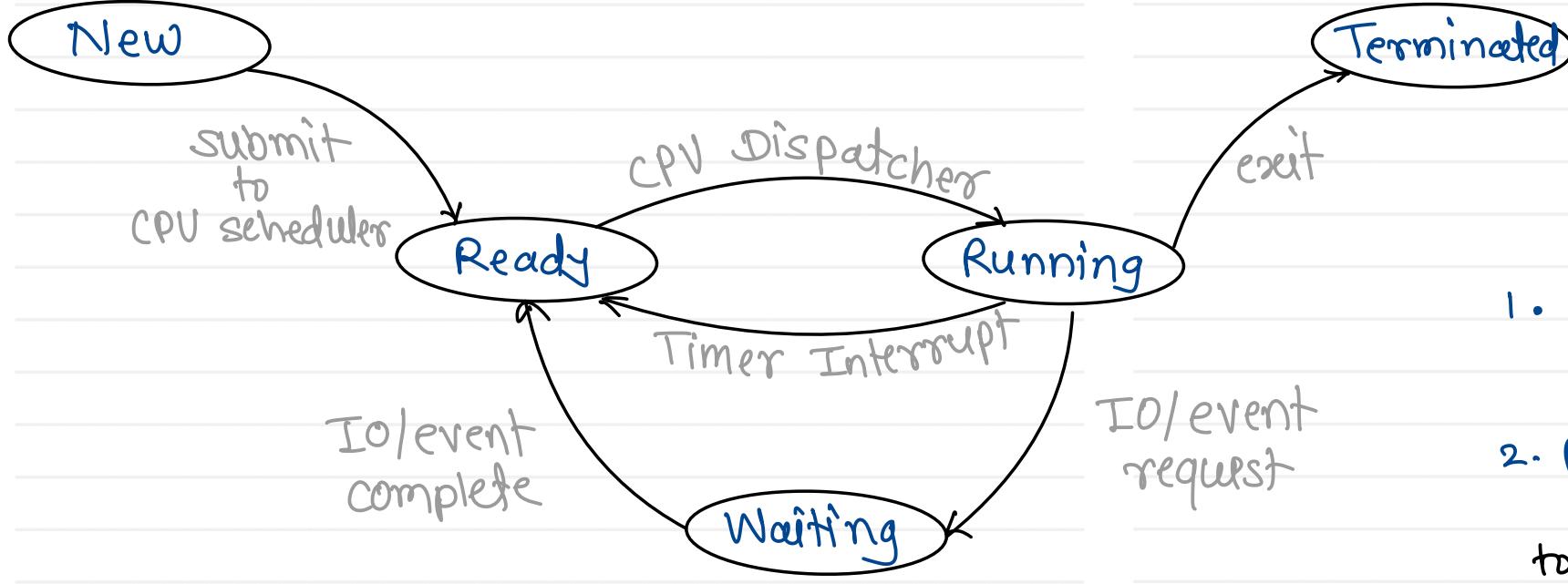
Windows Vista , Linux 2.6

## 6. Multiuser system :

multiple users can operate single system at a time.



# Process life cycle



OS data structures:

1. Job queue / process list
  - all processes of system
2. Ready queue
  - processes which are ready to execute on CPU
3. Waiting queues
  - processes which are waiting for IO / event.
  - waiting queues are created for every IO device / event separately.

# Context switching

Execution context :

- values of CPU registers

↑  
small memory elements  
available inside CPU

1. General purpose registers : (operands / results)

AH, AL, BH, BL, CH, CL, DH, DL

2. Segment registers : (base addresses of sections)

CS, DS, SS, ES

3. Offset registers : (displacement of variables)

BP, SP, SI, DI

4. Special registers : (address of next instruction)

PC / IR

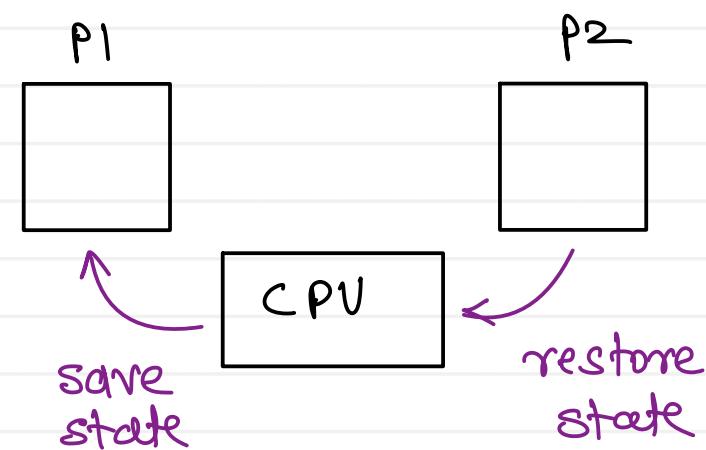
5. Flag register : (status of current operation)

Z, S, P, O

- CPU dispatcher loads , execution context from PCB of process on CPU

Context switching :

- changing the process of CPU
- execution context of running process is stored into its PCB & execution context of selected process is restored on CPU





# Types of scheduling

CPU scheduler is called in below 4 conditions:

1. Running → Terminated } voluntarily
2. Running → Waiting
3. Running → Ready } forcefully
4. Waiting → Ready

CPU scheduling algorithms:

1. FCFS
2. SJF (SRTF)
3. Priority
4. RR

Types of scheduling :

1. Non pre emptive scheduling
  - CPU access is given to next process voluntarily.
  - cooperative scheduling.

(case 1 & 2)

2. Pre emptive scheduling

- CPU access is given to next process forcefully.

(all cases)



# CPU scheduling criteria's

## 1. CPU Utilization : (Max)

- it determines how much time CPU is busy or idle.  
e.g. CPU Utilization = 70 %

70 % CPU busy & 30 % CPU is idle  
server OS - CPU Utilization = 90 %  
desktop OS - CPU Utilization = 70 %

## 2. Throughput : (Max)

- amount of work done in unit time  
- it measured as no. of processes completed in unit time

## 3. Waiting time : (Min)

- time spent by process in ready queue.

$$WT = TAT - CPU \text{ burst}$$

## 4. Response time : (Min)

- time from arrival in ready queue to first time executed on CPU

$$RT = ST - AT$$

## 5. Turn Around Time (TAT) : (Min)

- time spent by process inside RAM

$$TAT = CT - AT$$

$$TAT = \frac{CPU}{Waiting \text{ time}} + \frac{CPU}{burst} + \frac{IO}{Waiting \text{ time}} + \frac{IO}{burst}$$

# FCFS (First Come First Serve) (Non preemptive)

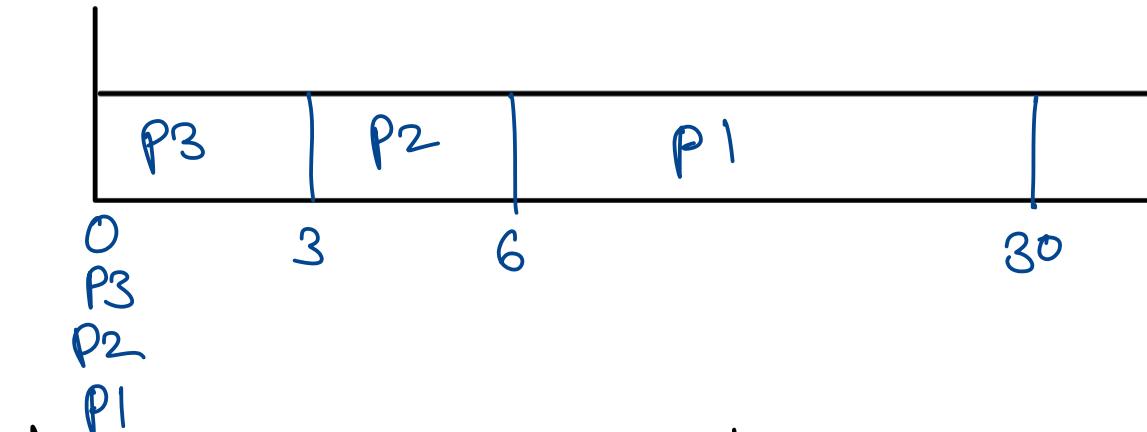
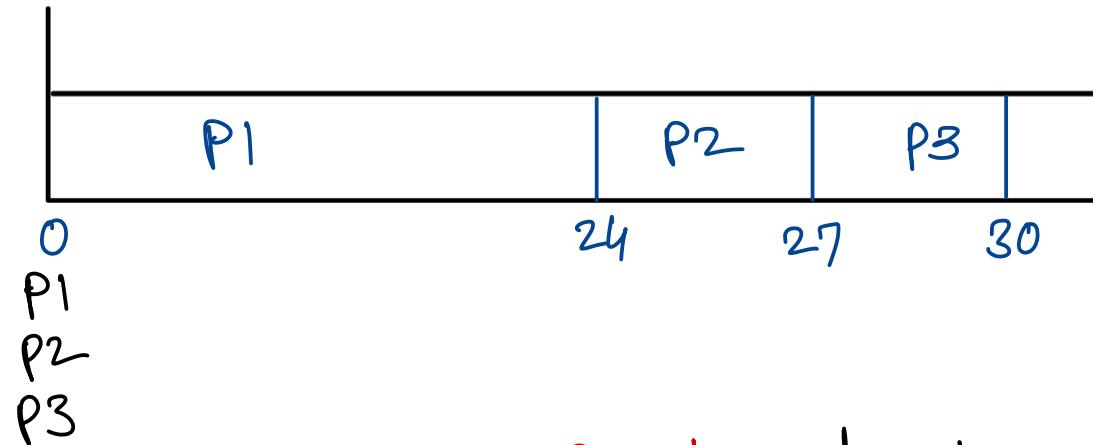
Process	Arrival	CPU Burst
P1	0	24
P2	0	3
P3	0	3

RT WT TAT

0	0	24
24	24	27
27	27	30

Process	Arrival	CPU Burst	RT	WT	TAT
P3	0	3	0	0	3
P2	0	3	3	3	6
P1	0	24	6	6	30

Gantt's chart

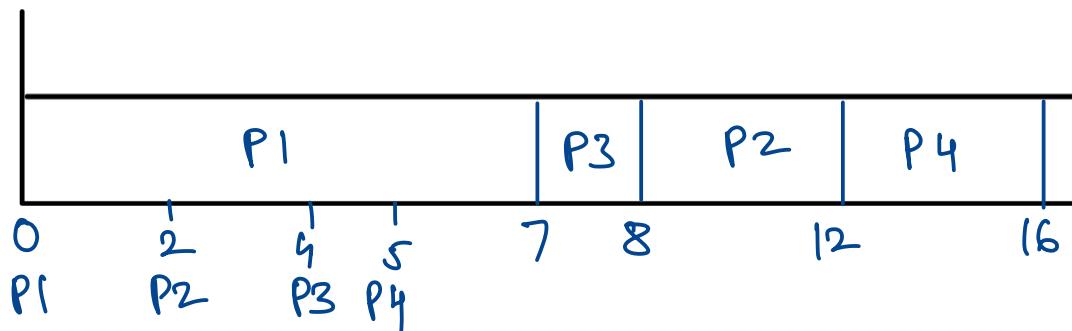


**Convoys effect:** due to arrival of longer process early, all other processes has to wait for longer time

# SJF (Shortest Job First)

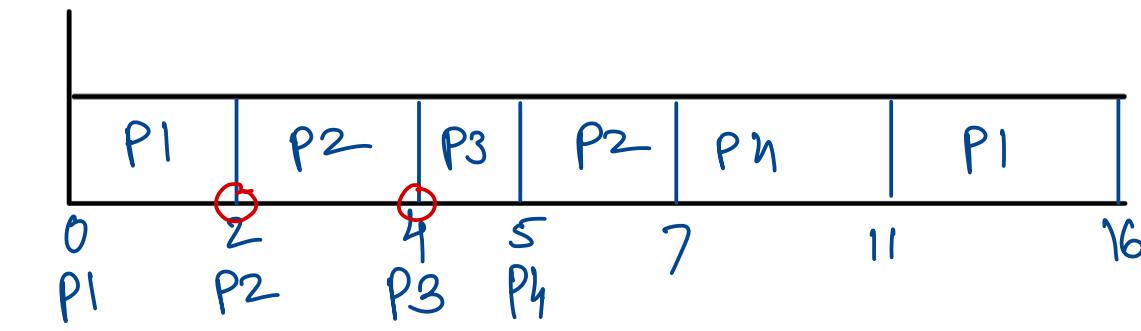
(Non pre emptive)

Process	Arrival	CPU Burst	WT	RT	TAT
P1	0	7	0	0	7
P2	2	4	6	6	10
P3	4	1	3	3	4
P4	5	4	7	7	11



Shortest Remaining Time First (SRTF)  
(preemptive)

Process	Arrival	CPU Burst	Remain time	WT	RT	TAT
P1	0	7	5	9	0	16
P2	2	4	2	1	0	5
P3	4	1	0	0	0	1
P4	5	4	2	2	2	6



Starvation: due to longer CPU burst process has to wait for longer time to get CPU access

# Priority

(Non pre emptive)

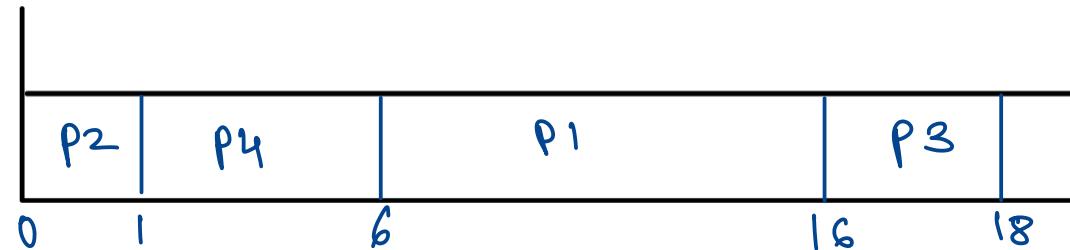
Process	Arrival	CPU Burst	Priority
P1	0	10	3
P2	0	1	1 (H)
P3	0	2	4 (L)
P4	0	5	2

	WT	RT	TAT
P1	6	6	16
P2	0	0	1
P3	16	16	18
P4	1	1	6

(Pre emptive)

Process	Arrival	CPU Burst	Priority
P1	0	10	3
P2	1	1	1
P3	3	2	4
P4	0	5	2

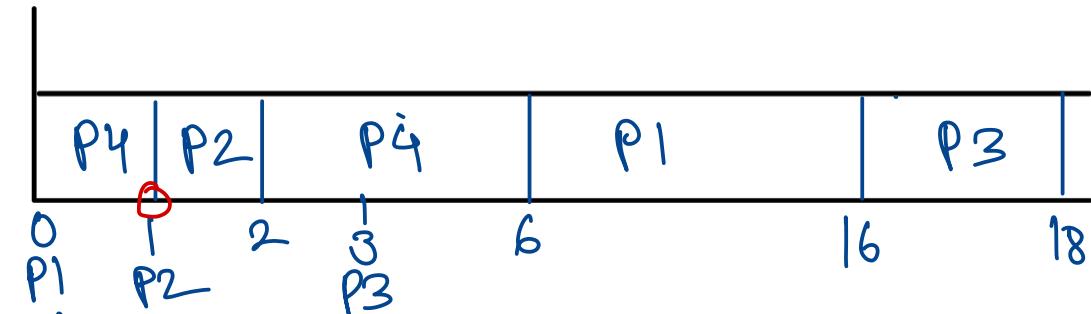
	WT	RT	TAT
P1	6	6	16
P2	0	0	1
P3	13	13	15
P4	1	0	6



P1(6)  
 P2(8)  
 P3(7)  
 P4(6)  
 P5(5)  
 P6(5)  
 P7(6)

P1  
 P5  
 P4  
 P3  
 P6  
 P7  
 P2

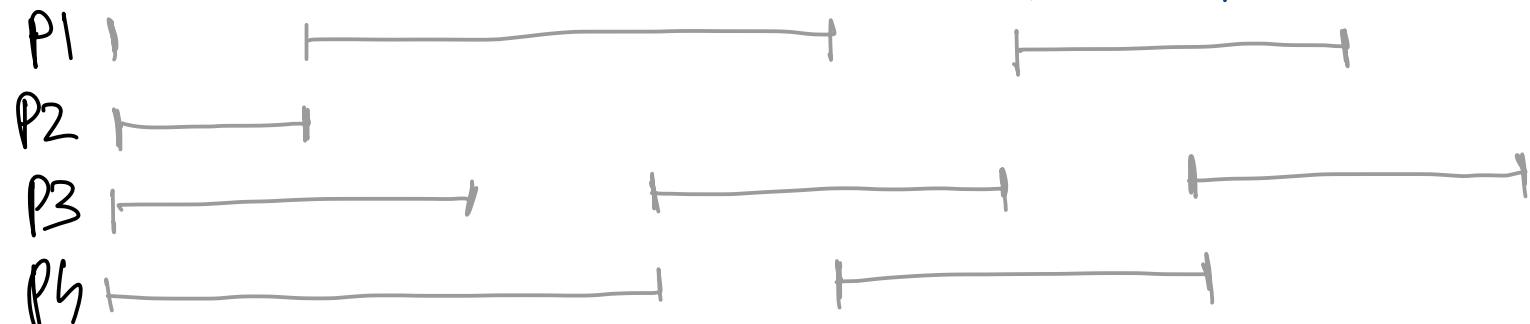
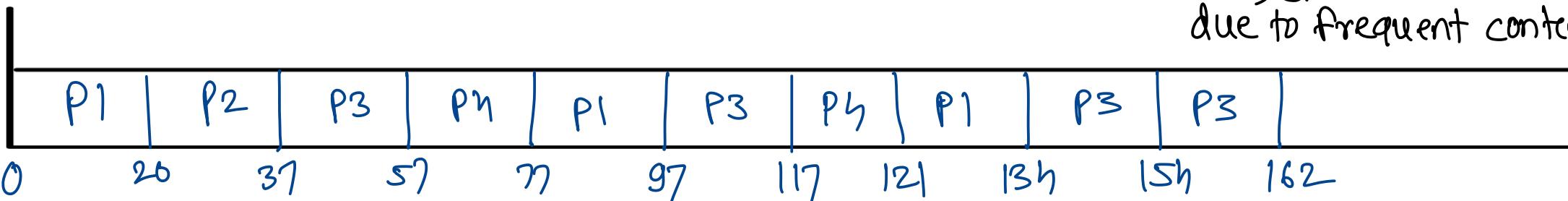
**starvation:**  
 due to low priority process  
 do not get enough time to execute  
 on CPU  
**aging:**  
 priority of process is increased  
 gradually.



# RR (Round Robin)

Process	CPU Burst
P1	53
P2	17
P3	68
P4	24

33, 13, 0      0 + 57 + 24  
 0                  20  
 48, 28, 8      37 + 40 + 17  
 4, 0              57 + 40



Time quantum : CPU time slice

$$TQ = 20$$

- next process will be scheduled in each time quantum.

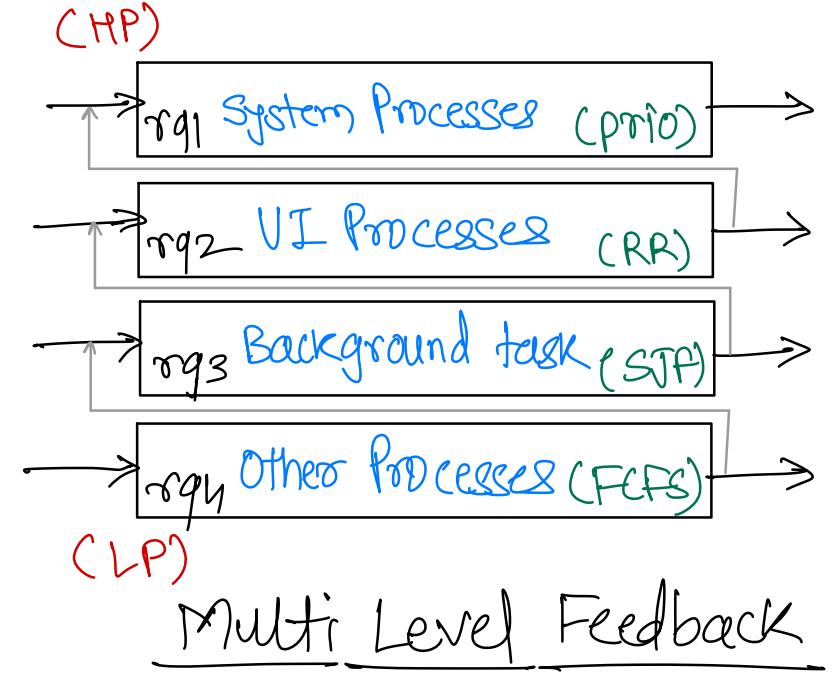
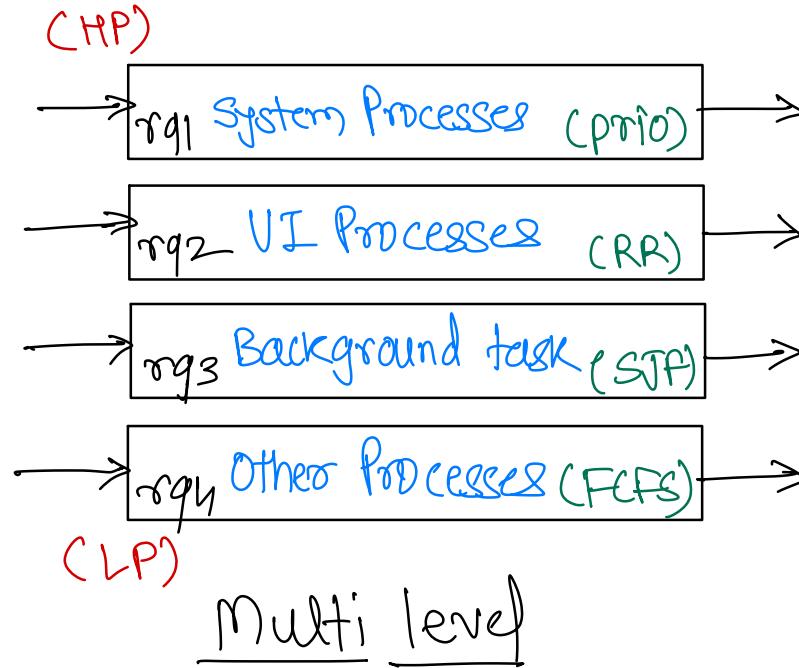
$$TQ = 100$$

↳ behave like FCFS

$$TQ = 4$$

↳ CPU overhead will increase due to frequent context switch

# Multi Level Ready Queue





Thank you!!!

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