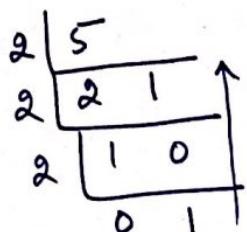


Module 4 - Digital Electronics

Number System & Representation

Decimal to binary

$$\bullet (5)_{10} \longrightarrow (?)_2.$$



$$\bullet (5)_{10} \longrightarrow (101)_2.$$

$$\bullet (0.61)_{10} \longrightarrow (?)_2$$

$$0.61 \times 2 = 1.22$$

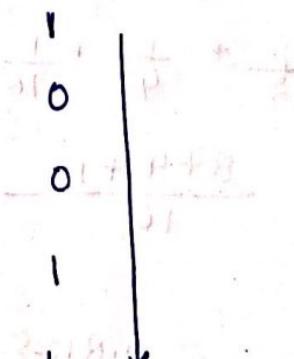
$$0.22 \times 2 = 0.44$$

$$0.44 \times 2 = 0.88$$

$$0.88 \times 2 = 1.76$$

$$0.76 \times 2 = 1.52$$

$$0.52 \times 2 = 1.04$$



$$(0.61)_{10} \longrightarrow (0.100111)_2.$$

• Binary to Decimal

• $(11100)_2 \rightarrow (?)_{10}$

$$\Rightarrow \begin{array}{r} 11100 \\ \hline 0 \times 2^0 + 0 \times 2^1 + 1 \times 2^2 + 1 \times 2^3 + 1 \times 2^4 \\ = 28. \end{array}$$

$(11100)_2 \rightarrow (28)_{10}.$

• $(0.1101)_2 \rightarrow (?)_{10}.$

$$\begin{array}{r} 0.1101 \\ \hline 1 \times 2^{-1} + 1 \times 2^{-2} + 0 \times 2^{-3} + 1 \times 2^{-4} \\ = \frac{1}{2} + \frac{1}{4} + \frac{1}{16} \\ = \frac{8+4+1}{16} = \frac{13}{16} \\ = 0.8125. \end{array}$$

$(0.1101)_2 \rightarrow (0.8125)_{10}$

(2)

- Decimal to octal

- $(246)_{10} \rightarrow (?)_8$

$$\begin{array}{r} 246 \\ 8 \overline{)30 \quad 6} \\ 8 \overline{)3 \quad 6} \\ 0 \quad 3 \end{array}$$

$$(246)_{10} \rightarrow (366)_8$$

- $(0.32)_{10} \rightarrow (?)_8$

$$0.32 \times 8 = 2.56 \quad 2$$

$$0.56 \times 8 = 4.48 \quad 4$$

$$0.48 \times 4 = 1.92 \quad 1$$

$$0.92 \times 4 = 3.68 \quad 3.$$

$$(0.32)_{10} \rightarrow (0.2413)_8$$

- Octal to Decimal

- $(746)_8 \rightarrow (?)_{10}$

$$\begin{aligned} \underline{746} &= 6 \times 8^0 + 4 \times 8^1 + 7 \times 8^2 \\ &= 486. \end{aligned} \Rightarrow (746)_8 \rightarrow (486)_{10}$$

$$\bullet (0.123)_8 \rightarrow (?)_{10}$$

$$1 \times 8^{-1} + 2 \times 8^{-2} + 3 \times 8^{-3}$$

$$= \frac{1}{8} + \frac{2}{64} + \frac{3}{512}$$

$$= 0.1621$$

$$(0.123)_8 \rightarrow (0.1621)_{10}.$$

Decimal to Hexadecimal

$$\bullet (2338)_{10} \rightarrow (?)_{16}.$$

$$\begin{array}{r} 2338 \\ 16 \overline{)146} \quad 2. \\ 16 \overline{)9} \quad 2. \\ 16 \overline{)0} \quad 9 \end{array}$$

$$(2338)_{10} \rightarrow (922)_{16}.$$

$$\bullet (0.122)_{10} \rightarrow (?)_{16}.$$

$$0.122 \times 16 = 1.952$$

1

$$0.952 \times 16 = 15.232$$

15 → F.

$$0.232 \times 16 = 3.712$$

3

$$0.712 \times 16 = 11.392$$

11 → B.

$$(0.122)_{10} \rightarrow (0.1F3B)_{16}$$

Hexadecimal to Decimal

$$\cdot (2AF)_{16} \rightarrow (?)_{10}$$

←

$$F \times 16^0 + A \times 16^1 + 2 \times 16^2$$

$$15 \times 1 + 10 \times 16 + 2 \times 256$$

$$(2AF)_{16} = (687)_{10}$$

$$\cdot (0.F27)_{16} \rightarrow (?)_{10}.$$

→

$$F \times 16^{-1} + 2 \times 16^{-2} + 7 \times 16^{-3}$$

$$\frac{15}{16} + \frac{2}{256} + \frac{7}{4096}$$

$$= 0.94702$$

$$(0.F27)_{16} \rightarrow (0.94702)_{10}.$$

Binary to Hex

$$\cdot (1100110 \cdot 011010)_2 \rightarrow (?)_{16}$$

$$\begin{array}{c} \uparrow \quad \uparrow \\ 1 \ 1 \ 0 \ 0 \ 1 \ 1 \ 0 \end{array} \cdot \begin{array}{c} \uparrow \quad \uparrow \\ 0 \ 1 \ 1 \ 0 \ 1 \ 0 \ 0 \ 0 \end{array}$$

6 6 6 8

$$\Rightarrow (66.68)_{10}.$$

Hex to Binary

• $(FEA)_{16} \rightarrow (?)_2$.

F	E	A
↓	↓	↓
1111	1110	1010

$$(FEA)_{16} \rightarrow (1111 \ 1110 \ 1010)_2$$

• $(0.A25)_{16}$

A25		
↓	↓	
1010	0010	0101

$$\Rightarrow (0.1010 \ 0010 \ 0101)_2$$

Binary to octal

$(1001010.01100101)_2 \rightarrow (?)_8$.

001	001	010	.	011	001	010
↑	↑	↑		↑	↑	↑
112	.	312				

$$\Rightarrow (112.312)_8$$

• Octal to Binary

(4)

$$(123.567)_8 \rightarrow (?)_2$$

$$\begin{array}{ccccccc} 1 & 2 & 3 & . & 5 & 6 & 7 \\ \downarrow & \downarrow & \downarrow & & \downarrow & \downarrow & \downarrow \\ 0 & 0 & 1 & 0 & 1 & 0 & 0 \end{array} \rightarrow (001\ 010\ 011\ .\ 101\ 110\ 111)_2$$



$$\boxed{(1\ 010\ 011\ .\ 101\ 110\ 111)_2}$$

• Hex to octal

(i) Hex \rightarrow Binary

(ii) Binary \rightarrow octal.

$$(BB.0A)_{16} \rightarrow (?)_8$$

(i) $(1011\ 1011\ .\ 0000\ 1010)_2$

(ii) $\underbrace{010}_{2}\ \underbrace{111}_{7}\ \underbrace{011}_{3}\ .\ \underbrace{000}_{4}\ \underbrace{010}_{2}\ \underbrace{100}_{4}$

273. 024

$$\boxed{(BB.0A)_{16} \rightarrow (273.024)_8}$$

• Octal to Hex

(i) Octal \rightarrow Binary

(ii) Binary \rightarrow Hex

$$(436.71)_8 \rightarrow (?)_{16}$$

$$\begin{array}{ccccccc} \text{(i)} & 4 & 3 & 6 & . & 7 & 1 \\ & \downarrow & \downarrow & \downarrow & & \downarrow & \downarrow \\ 100 & 011 & 10110 & 111 & 001 & & \end{array}$$

$$\begin{array}{ccccccc} \text{(ii)} & 0001 & 0001 & 1110 & . & 1110 & 0100 \\ & \overbrace{\quad\quad\quad}^1 & \overbrace{\quad\quad\quad}^1 & \overbrace{\quad\quad\quad}^1 & & \overbrace{\quad\quad\quad}^1 & \overbrace{\quad\quad\quad}^0 \\ 1 & 1 & E & . & E & 4 & \end{array}$$

$$\boxed{(11E \cdot E4)_{16}}$$

is complement.

- Replace '0' by '1' and '1' by '0's.

2's complement

(i) Find 1st complement

(ii) Add '1' to it.

Boolean Algebra.

Duality

$$\overline{A \cdot B} \neq \overline{A} \cdot \overline{B} \quad \text{AND laws.}$$

$$A \cdot 0 = 0$$

$$A \cdot 1 = A$$

$$A \cdot A = A$$

$$A \cdot \overline{A} = 0$$

OR laws.

$$A + 0 = A$$

$$A + 1 = 1$$

$$A + A = A$$

$$A + \overline{A} = 1$$

De Morgan's Theorem.

$$\overline{A+B} = \overline{A} \cdot \overline{B}$$

$$\overline{AB} = \overline{A} + \overline{B}$$

$$\Rightarrow (C+A)(C+D) \\ (C+A)(C+D)$$

Commutative law.

$$A+B = B+A$$

$$A \cdot B = B \cdot A$$

Associative law.

$$(A+B) + C = A + (B+C)$$

$$A(BCD) = (ABC)D.$$

Distributive law.

$$A(B+C) = AB + AC.$$

$$A + BC = (A+B)(A+C)$$

$$A + \bar{A}B = (A+\bar{A})(A+B) = 1 \cdot (A+B) = A+B.$$

Q. Minimize the following expressions using boolean algebra:

$$\begin{aligned} (i) Y &= \bar{A}B + A\bar{B} + AB \\ &= \bar{A}B + A(\bar{B} + B) \\ &= \bar{A}B + A \cdot 1 \\ &= A + \bar{A}B \\ &= (A+\bar{A})(A+B) \\ &= A+B. \end{aligned}$$

$$\begin{aligned} (ii) Y &= \bar{A}\bar{B}C + \bar{A}B\bar{C} + A\bar{B}\bar{C} + A\bar{B}C \\ &= \bar{B}C(\bar{A}+A) + \bar{A}B\bar{C} + A\bar{B}C \\ &= \bar{B}C + \bar{C}(\bar{A}B + A\bar{B}) \\ &= \bar{B}\bar{C} + \bar{C}\bar{A}B + A\bar{B}\bar{C} \end{aligned}$$

$$= \bar{B} (C + A\bar{C}) + A\bar{B}\bar{C}$$

$$= \bar{B} [(C + \bar{C})(C + A)] + \bar{A}\bar{B}\bar{C}$$

$$= \bar{B} (C + A) + \bar{A}\bar{B}\bar{C}$$

$$= \boxed{A\bar{B} + A\bar{B}C + \bar{A}\bar{B}\bar{C}}$$

$$(iii) Y = \bar{A}\bar{B}\bar{C} + \bar{A}\underline{B}\bar{C} + A\bar{B}\bar{C} + \underline{A}\bar{B}\bar{C}$$

$$= \bar{B}\bar{C}(\bar{A} + A) + B\bar{C}(\bar{A} + A)$$

$$= \bar{B}\bar{C} + B\bar{C}$$

$$= \bar{C}(B + \bar{B})$$

$$= \bar{C}$$

$$(iv) F = \bar{Y}\bar{Z} + \bar{W}\bar{X}\bar{Z} + \bar{W}\bar{X}\bar{Y}\bar{Z} + Y\bar{Z}$$

$$= \cancel{\bar{Y}\bar{Z}} / (1 + \bar{W}\bar{X}) + \bar{W}\bar{X}\bar{Z}(1 + \bar{Y})$$

$$= \cancel{\bar{Y}\bar{Z}} + \bar{W}\bar{X}\bar{Z}$$

$$= \bar{Z}(Y + \bar{W}\bar{X})$$

$$= \bar{Y}\bar{Z}(1 + \bar{W}\bar{X}) + \bar{W}\bar{X}\bar{Z} + Y\bar{Z}$$

$$= \bar{Y}\bar{Z} + \bar{W}\bar{X}\bar{Z} + Y\bar{Z}$$

$$= \bar{Z}(Y + \bar{Y}) + \bar{W}\bar{X}\bar{Z}$$

$$= \bar{Z} + \bar{W}\bar{X}\bar{Z}$$

$$= \bar{z}(1 + \bar{w}\bar{x})$$

$$\boxed{F = \bar{z}}$$

b (v) $f = A + \bar{A}B\bar{C} + \overline{B+C}$

$$= A + \bar{A}B\bar{C} + \bar{B}\bar{C}$$

$$= (A + \bar{A})(A + B\bar{C}) + \bar{B}\bar{C}$$

$$= A + B\bar{C} + \bar{B}\bar{C}$$

$$= A + \bar{C}(B + \bar{B})$$

$$= A + \bar{C}$$

(vi) $y = AB + AC + BC + B.$

$$= A(B + C) + B(1 + C).$$

$$= A(B + C) + B.$$

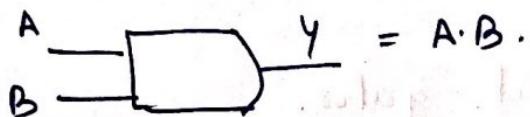
$$= AB + AC + B$$

$$= B(1 + A) + AC.$$

$$= \underline{\underline{B + AC.}}$$

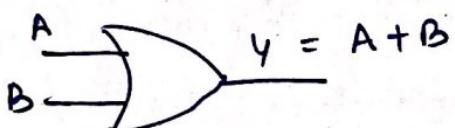
logic gates

AND



$$Y = A \cdot B$$

OR



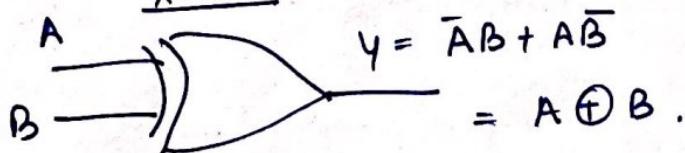
$$Y = A + B$$

NOT



$$Y = \bar{A}$$

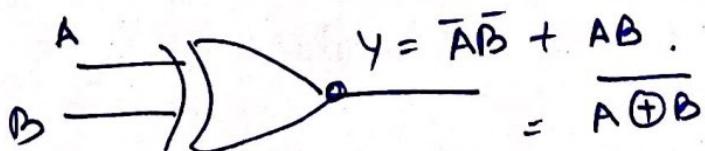
XOR.



$$Y = \bar{A}B + A\bar{B}$$

$$= A \oplus B$$

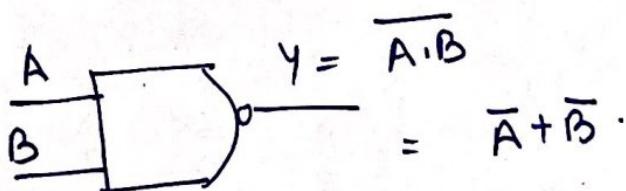
XNOR.



$$Y = \bar{A}\bar{B} + A\bar{B}$$

$$= \overline{A \oplus B}$$

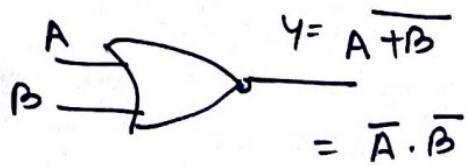
NAND



$$Y = \overline{A \cdot B}$$

$$= \bar{A} + \bar{B}$$

NOR.



NAND & NOR \rightarrow universal gates.

2 variable K Map.

$$(i) Y = A\bar{B} + A\bar{B} \quad (\text{Ans: } A)$$

$$(ii) Y = \bar{A}\bar{B} + A\bar{B} + A\bar{B} \quad (\text{Ans: } A + B)$$

$$(iii) Y = \sum m(0,3) \quad (\text{Ans: } A\bar{B} + AB)$$

$$(iv) Y(A,B) = \sum m(0,1,2,3) \quad (\text{Ans: } 1)$$

3 Variable, 4 Variable.

Q. i) $f(A,B,C) = \sum m(0,2,4,6)$

		BC	00	01	11	10			
		A	0	1	0	1	3	1	2
		0	1	1	1	5	7	1	6
		1	1	1	1	5	7	1	6

$$= \bar{C}$$

(ii) $f(A,B,C) = \sum m(1,2,4,5)$

		BC	00	01	11	10			
		A	0	0	1	1	3	1	2
		0	1	1	1	5	7	1	6
		1	1	1	1	5	7	1	6

$$= \bar{B}C + A\bar{B} + \bar{A}BC$$

(iii) $f(A,B,C,D) = \sum m(0,1,2,5,13,15)$

		CD	00	01	11	10		
		AB	00	1	1	3	1	2
		0	1	1	1	5	7	6
		1	1	1	1	5	7	6
		10	8	9	11	14		

$$Y = \bar{A}CD + ABD + \bar{A}\bar{B}\bar{D}$$

$$\text{Ans: } B\bar{D} + \bar{B}\bar{C}D$$

$$(iv) \quad f(A, B, C, D) = \sum m(1, 3, 7, 8, 10, 12, 13, 15)$$

		CD	00	01	11	10
		AB	00	01	11	10
	00	0	1	1		2
	01	4	5	1	6	7
	11	12	13	1	15	14
	10	8	9	11	1	10

Q: $f(A, B, C, D)$

$$= \sum m(0, 8, 2, 10)$$

(Ans: $\overline{B} \overline{D}$)

$$= \overline{A} \overline{B} D + B C D + A B \overline{C} + A \overline{B} \overline{D}$$

$$(v). \quad f(A, B, C, D) = \sum m(1, 3, 5, 9, 11, 13)$$

		CD	00	01	11	10
		AB	00	01	11	10
	00	0	1		1	3
	01	4	5		7	6
	11	12	13	1	15	14
	10	8	9	11	1	10

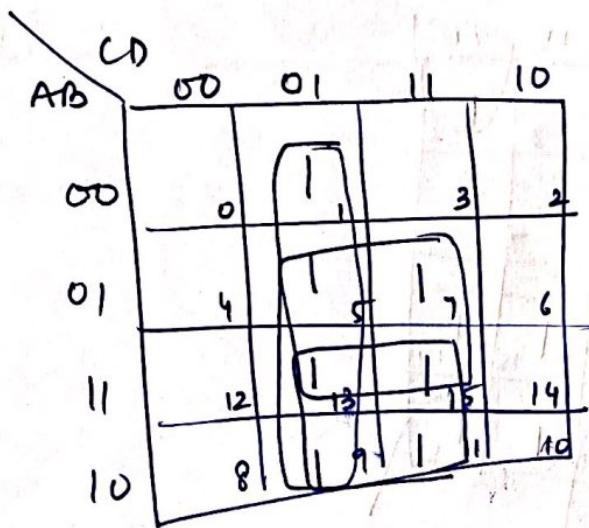
$$= \overline{C} D + \cancel{\overline{B} B} \overline{B} D.$$

Q. (i) $f(A, B, C, D) = \sum m(1, 4, 6, 12, 14, 9)$ (Ans: $B \overline{D} + \overline{B} \overline{C} D$)

Q. (ii) $f(A, B, C, D) = \sum m(0, 8, 2, 10)$ (Ans: $B \overline{D}$)

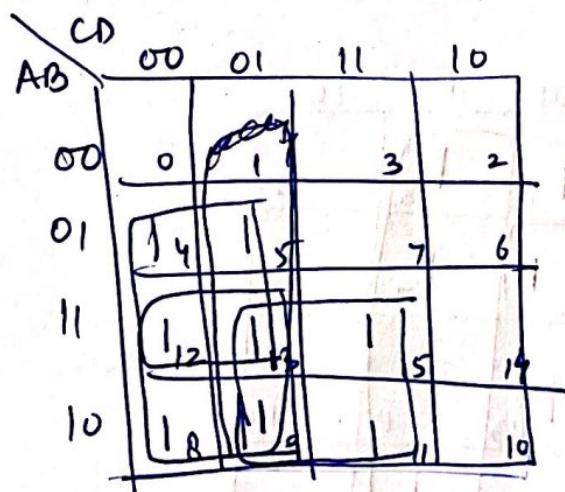
(21)

$$(vi) f(A, B, C, D) = \sum m (1, 5, 7, 9, 11, 13, 15)$$



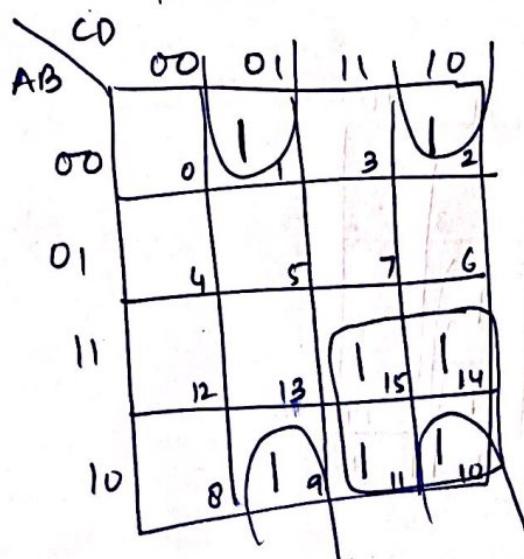
$$\bar{C}D + BD + AD.$$

$$(vii) f(A, B, C, D) = \sum m (4, 5, 8, 9, 11, 12, 13, 15)$$



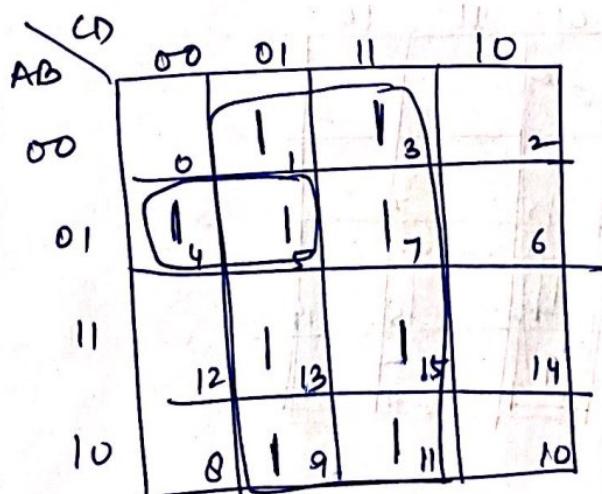
~~$$\bar{C}\bar{D} + BC\bar{C} + A\bar{C} + AD.$$~~

$$(viii) f(A, B, C, D) = \sum m(1, 2, 9, 10, 11, 14, 15)$$



$$= AC + \bar{B}C\bar{D} + \bar{B}\bar{C}D.$$

$$(ix) f(A, B, C, D) = \sum m(1, 3, 4, 5, 7, 9, 11, 13, 15)$$



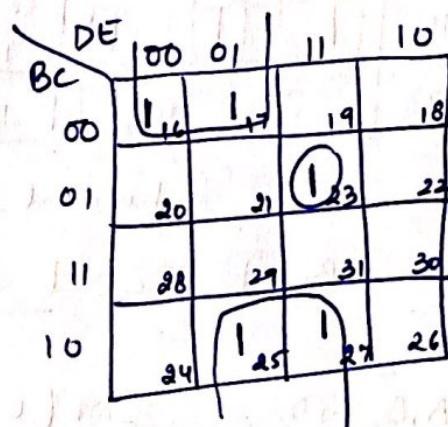
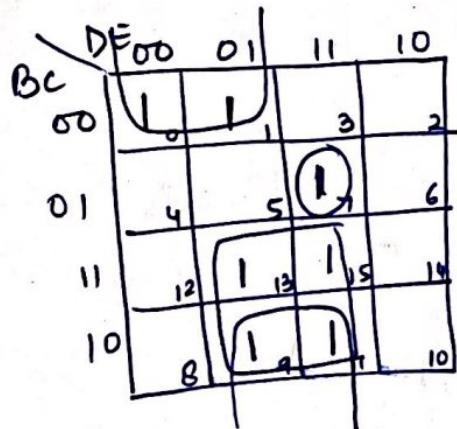
$$= \boxed{D + \bar{A}B\bar{C}}$$

$$\text{Q. } f(A, B, C, D) = \sum m(1, 3, 4, 5, 7, 9, 11, 13, 15) + \sum d(12, 14) \\ (\text{Ans: } D + \bar{B}\bar{C})$$

$$\text{Q. } f(A, B, C, D) = \sum m(1, 5, 7, 9, 11, 13, 15) + \sum d(3, 2, 6, 14, 10) \\ (\text{Ans: } f = D).$$

K Map (5 Variable).

$$(i) f(A, B, C, D, E) = \sum m(0, 1, 7, 9, 11, 13, 15, 16, 17, 23, 25, 27)$$

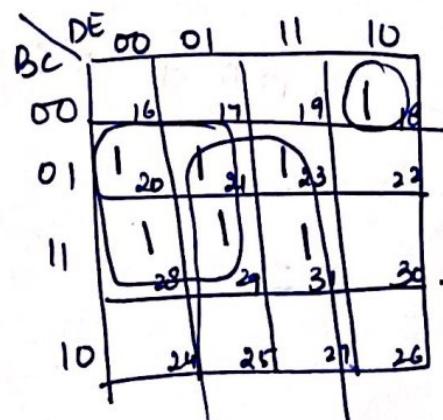
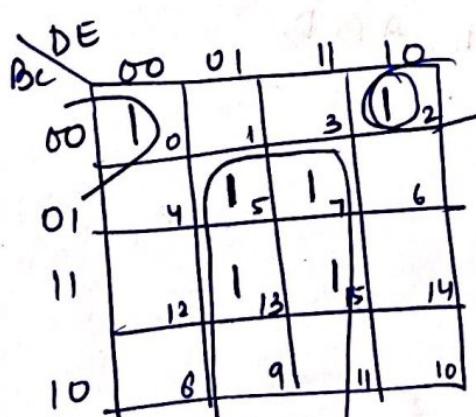


$$A=1$$

$$\bar{A}=0$$

$$\bar{A}BE + B\bar{C}E + \bar{B}\bar{C}\bar{D} + \bar{B}CDE$$

$$(ii) f(A, B, C, D, E) = \sum m(0, 2, 5, 7, 13, 15, 18, 20, 21, 23, 28, 29, 31).$$



$$= [CE + ACD + \frac{A}{ABCE} + \bar{B}\bar{C}D\bar{E}]$$

$$(iii). f = \sum m(1, 5, 6, 7, 11, 12, 13, 15) \\ (A, B, C, D)$$

AB\CD	00	01	11	10
00	0	1	3	2
01	4	1, 5	1, 6	
11	1, 12	1, 13	1, 15	14
10	8	9	11	10

Redundant.

$$= \bar{A}BC + \bar{A}\bar{C}D + AB\bar{C} + ACD.$$

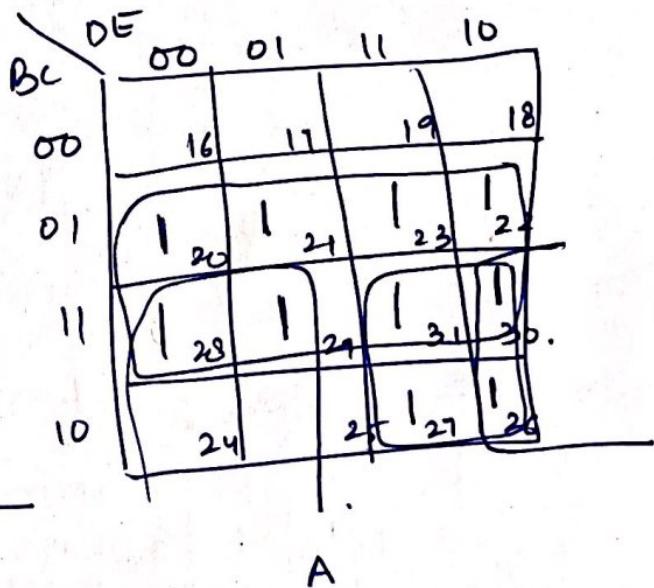
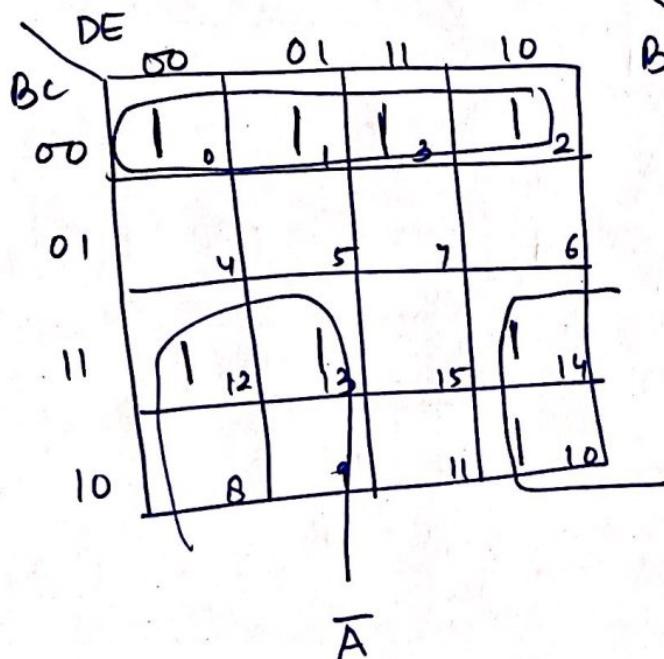
$$(iv). f(A, B, C, D) = \sum m(0, 1, 2, 5, 13, 15)$$

AB\CD	00	01	11	10
00	0	1	3	2
01	4	1, 5	7	6
11	12	1, 13	1, 15	14
10	8	9	11	10

$$= \bar{A}\bar{B}\bar{D} + \bar{A}\bar{C}D + AB\bar{D}.$$

(2)

$$(V) F(A, B, C, D, E) = \sum m (0, 1, 2, 3, 10, 12, 13, 14, \\ 20, 21, 22, 23, 26, 27, 28, 29, \\ 30, 31)$$



$$A' C + B' C' \bar{D} + B' D \bar{E} + \bar{A} \bar{B} \bar{C} \\ - \\ A C + B C \bar{D} + A B D.$$

K Map (6 variable).

		B'					
		CD	EF	00	01	11	10
A'	00	0	1	3	2		
	01	4	5	7	6		
11	12	13	15	14			
10	8	9	11	10			

		B					
		CD	EF	00	01	11	10
A'	00	16	17	19	18		
	01	20	21	23	22		
11	28	29	31	30			
10	24	25	27	26			

		A					
		CD	EF	00	01	11	10
A'	00	32	33	35	34		
	01	36	37	39	38		
11	44	45	47	46			
10	40	41	43	42			

		A					
		CD	EF	00	01	11	10
A'	00	48	49	51	50		
	01	52	53	55	54		
11	60	61	63	62			
10	56	57	59	58			

• first check the corners of all K Maps.

$$(i) Y = \sum m (0, 1, 2, 3, 4, 5, 8, 9, 12, 13, 16, 17, 18, 19, 24, 25, 36, 37, 38, 39, 52, 53, 60, 61)$$

	\bar{B}	D		
CD	00	01	11	10
00	1	1	1	1
01	14	15	7	6
11	12	13	15	14
10	18	19	11	10

	B	D		
CD	00	01	11	10
00	16	17	19	18
01	20	21	23	22
11	28	29	31	30
10	24	25	27	26

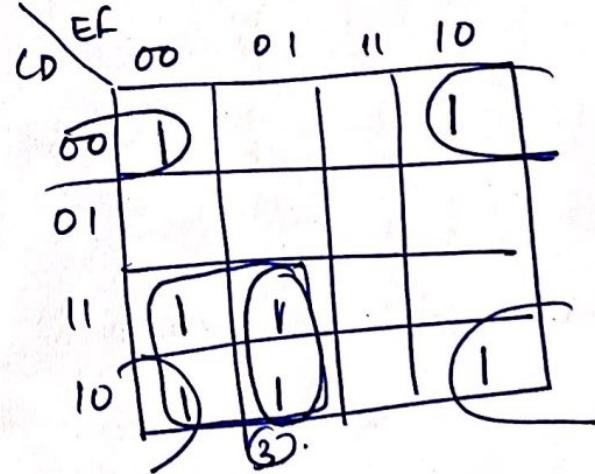
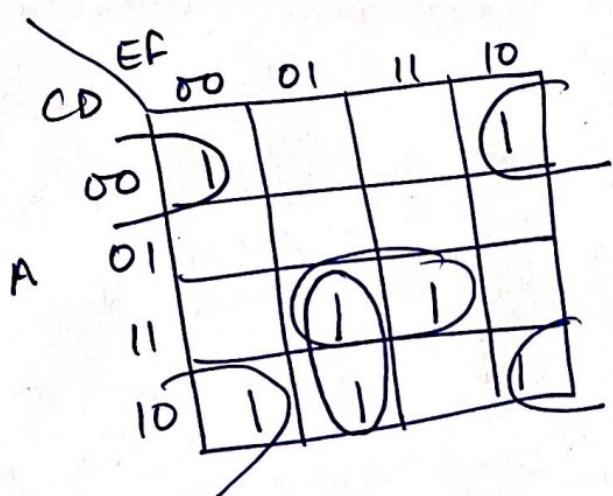
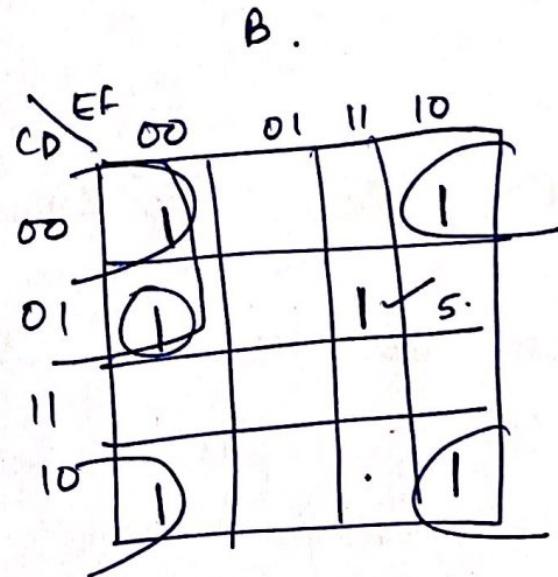
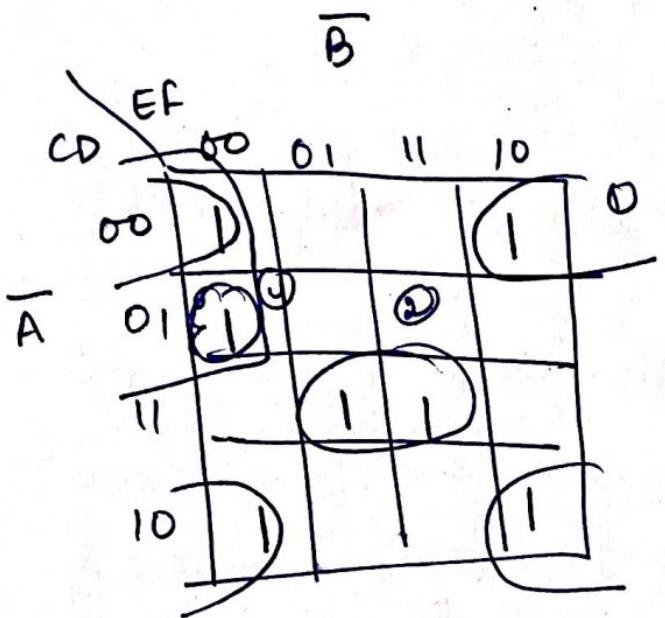
	E	F	CD	00	01	11	10
A	00	01	00	32	33	35	34
01	136	137	139	138			
11	44	45	47	46			
10	40	41	43	42			

	E	F	CD	00	01	11	10
00	48	49	51	50			
01	52	53	55	54			
11	60	61	63	62			
10	56	57	59	58			

$$= \overline{A}\overline{B}E + \overline{A}\overline{C}\overline{D} + A\overline{B}\overline{C}D + ABDE \\ + \overline{A}\overline{C}\overline{D} \cdot \overline{A}\overline{D}E$$

(ii) $Y = \{0, 2, 4, 8, 10, 13, 15, 16, 18, 20, 23, 24, 26, 32, 34, 40, 41, 42, 45, 47, 48, 50, 56, 57, 58, 60, 61\}$

(2)



$$\begin{aligned}
 &= \overline{D}\overline{F} + \overline{B}\overline{C}\overline{D}F + \overline{A}\overline{B}\overline{C}\overline{E}\overline{F} + \overline{A}\overline{C}\overline{D}\overline{E}\overline{F} \\
 &\quad + \overline{A}\overline{B}\overline{C}DEF. \\
 &\quad + A\overline{C}\overline{E}F.
 \end{aligned}$$

Introduction to IC technology : SSI, MSI, LSI



The first integrated circuits contained only a few transistors and so were called "Small scale integration" (SSI). They used circuits containing transistors numbering in the tens. They were very crucial in development of early computers. SSI was followed by introduction of the devices which contained hundreds of transistors on each chip and so were called "Medium Scale Integration (MSI)".

MSI were attractive economically. Next development was of Large Scale Integration (LSI). The development of LSI was driven by economic factors and each chip comprised of tens of thousands of transistors. It was in 1970s when LSI started getting manufactured in huge quantities.

⇒ An integrated circuit or monolithic integrated circuit (IC, chip or microchip) is a set of electronic circuits on one small flat piece of semiconductor material that is normally Silicon.

The integration of large no. of tiny MOS transistors into a small chip results in units that are orders of magnitude smaller, faster and less expensive than those constructed of discrete electronic components. The IC's mass production capability, reliability and building-block approach to integrated circuit design has ensured the rapid adoption of standardized IC's in place of designs using discrete transistors. IC's are now used in virtually all electronic equipments.

Advantages of IC's over discrete circuits.

→ cost and performance.

⇒ cost is low because the chips, with all their components are printed as a unit by photolithography rather than being constructed one transistor at a time.

Performance is high because the IC's components switch quickly and consume comparatively little power because of their small size and proximity.

The main disadvantage of IC's is the high cost to design them and fabricate the required photomasks.

		Year	Transistor count	Logic gates no.
SSI	small scale Integration	1964	1 to 10	- 12
MSI	Medium scale Integration	1968	10-500	13 - 99
LSI	Large scale Integration	1971	500 - 20,000	100 - 9999
VLSI	very large scale Integration	1980	20,000 - 1,00,0000	10,000 - 99,999

VLSI (very large scale Integration)

The final step in the development process, starting in the 1980s and continuing through the present is, "very large scale integration" (VLSI). The development started with hundreds of thousands of transistors in the early 1980s. As of 2016, transistor counts continue to grow beyond ten billion transistors per chip.

Multiple developments were required to achieve this increased density. Manufacturers moved to smaller MOSFET design rules and cleaner fabrication facilities so that they could make chips with more transistors and maintain adequate yield. Electronic design tools improved enough to make it practical to finish these designs in a reasonable time.

The more energy efficient CMOS replaced NMOS and PMOS, avoiding a prohibitive increase in power consumption.

Modern VLSI devices contain so many transistors, layers, interconnections and other features that it is no longer feasible to check the masks or do the original design by hand. Instead, engineers use EDA tools to perform most functional verification work.