# Modeling Electronic Systems on Quantum Computers

University of Florida, October

### Recap of last time

▶ We discussed solving for the ground state  $|\Psi_0\rangle$  and ground state energy of a molecule with n electrons,

$$\hat{H} = \sum_{ij} h_1^{ij} a_i^{\dagger} a_j + \frac{1}{2} \sum_{ijkl} h_2^{ijkl} a_i^{\dagger} a_j^{\dagger} a_k a_l$$

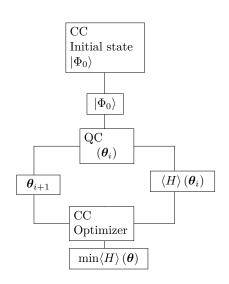
- $ightharpoonup a_i^{\dagger}$  creates a particle in the MO  $|\chi_i\rangle$ .
- ▶ For a calculations on a computer we need to to expand these orbital in some finite set of basis (AO). Let the number of basis be *m*.
- ▶ We use the occupation number representation

$$|n_1, n_2, \ldots, n_{m-1}, n_m\rangle$$

- ▶ The Hilbert space dimensions is  $\mathcal{O}(e^n)$ .
- ► The Hartree-Fock approximation give us the state with the lowest energy that can also be written as a Slater determinant.

# Where quantum computer can help

- Classical computer can do HF efficiently.
- In going beyond HF the bottle necks are:
  - 1. Sampling the Hilbert space.
  - Calculating the expectation value of the Hamiltonian.
- A variational quantum eigensolver algorithm (VQE), put the exponentially hard part of the calculation on a quantum computer.



#### Outline

1. Introduction to circuit model for quantum computing

Converting fermionic Hamiltonians to qubit Hamiltonians Jordan-Wigner

3. VQE algorithms circuit designs for UCCSD

#### Outline

1. Introduction to circuit model for quantum computing

Converting fermionic Hamiltonians to qubit Hamiltonians Jordan-Wigner

3. VQE algorithms circuit designs for UCCSE

### Qubits

A qubit lives in a 2 dimensional Hilbert space,  $\mathcal{H}^{\mathbf{1}}.$  An arbitrary qubit state can be described as,

$$|\psi\rangle = c_0 |0\rangle + c_1 |1\rangle.$$

### Qubits

A qubit lives in a 2 dimensional Hilbert space,  $\mathcal{H}^1$ . An arbitrary qubit state can be described as,

$$|\psi\rangle = c_0 |0\rangle + c_1 |1\rangle.$$

A system of N qubits lives in a  $2^{N}$  dimensional Hilbert space,  $\mathcal{H}^{N}$  such that,

$$\mathcal{H}^N = \bigotimes_i^N \mathcal{H}_i^1$$

An arbitrary state describing the system of N qubits can be written as,

$$|\Psi\rangle = \sum_{\{s_i\}} c_{1...N} |s_1...s_N\rangle, \ s_1,...s_N = \{0, 1\}$$

#### Gates

Operators can act on one or two (or more) qubits independently.

► These operators acting on the qubits are called quantum gates.

#### Gates

Operators can act on one or two (or more) qubits independently.

► These operators acting on the qubits are called quantum gates.

#### Examples of 1-qubit gates:

- ▶ The X-gate:  $X|0\rangle = |1\rangle$ ,  $X|1\rangle = |0\rangle$ .
- ▶ The Z-gate:  $Z |0\rangle = |0\rangle$ ,  $Z |1\rangle = -|1\rangle$ .
- The *H*-gate:  $H|0\rangle = \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle), H|1\rangle = \frac{1}{\sqrt{2}} (|0\rangle |1\rangle).$









#### Gates

Operators can act on one or two (or more) qubits independently.

► These operators acting on the qubits are called quantum gates.

#### Examples of 1-qubit gates:

▶ The 
$$X$$
-gate:  $X|0\rangle = |1\rangle$ ,  $X|1\rangle = |0\rangle$ .

▶ The 
$$Z$$
-gate:  $Z \left| 0 \right\rangle = \left| 0 \right\rangle, \ Z \left| 1 \right\rangle = - \left| 1 \right\rangle.$ 

The *H*-gate: 
$$H |0\rangle = \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle), \ H |1\rangle = \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle).$$

#### Example of 2-qubit gates:

▶ Controlled-X gate:  $cX \mid 00 \rangle = \mid 00 \rangle$ ,  $cX \mid 01 \rangle = \mid 01 \rangle$ ,  $cX \mid 10 \rangle = \mid 11 \rangle$ ,  $cX \mid 11 \rangle = \mid 10 \rangle$ . It adds the two qubits and puts the result in the second qubit



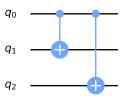






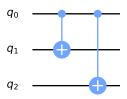
### Circuits

Example 1:  $|0\rangle \to |000\rangle$ ,  $|1\rangle \to |111\rangle$ . (Quantum repetition code.)

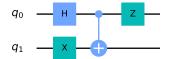


#### Circuits

Example 1:  $|0\rangle \rightarrow |000\rangle$ ,  $|1\rangle \rightarrow |111\rangle$ . (Quantum repetition code.)



Example 2:  $|00\rangle \rightarrow \frac{1}{\sqrt{2}} (|01\rangle - |10\rangle)$ , (Adds entanglement.)



- ▶ The most prevalent model for quantum computations is the circuit model.
- ► A circuit model is similar to how classical computers operate.

### Outline

1. Introduction to circuit model for quantum computing

2. Converting fermionic Hamiltonians to qubit Hamiltonians Jordan-Wigner

3. VQE algorithms circuit designs for UCCSD

### Fermionic and qubit degrees of freedom

Fermionic system with m orbitals:  $(2^m \text{ dimensional})$ 

$$\begin{split} \{a_i,a_j^\dagger\} &= \delta_{ij}, \ \{a_i,a_j\} = 0. \\ a_i^\dagger \left| n_1 \dots n_i \dots n_m \right\rangle &= \begin{cases} \zeta_i \left| n_1 \dots n_i = 1 \dots n_m \right\rangle, \ n_i = 0 \\ 0, \qquad n_i = 1 \end{cases} \\ \zeta_i &= (-1)^{\sum_{j=1}^{i-1} n_j} \ \text{Parity counting} \end{split}$$

## Fermionic and qubit degrees of freedom

Fermionic system with m orbitals:  $(2^m \text{ dimensional})$ 

$$\begin{split} \{a_i,a_j^\dagger\} &= \delta_{ij}, \ \{a_i,a_j\} = 0. \\ a_i^\dagger \left| n_1 \dots n_i \dots n_m \right\rangle &= \begin{cases} \zeta_i \left| n_1 \dots n_i = 1 \dots n_m \right\rangle, \ n_i = 0 \\ 0, \qquad n_i = 1 \end{cases} \\ \zeta_i &= (-1)^{\sum_{j=1}^{i-1} n_j} \ \text{Parity counting} \end{split}$$

Qubits system with m qubits:

$$\sigma^{+} = \frac{1}{2}(\sigma_{x} + i\sigma_{y}), \ \sigma^{-} = \frac{1}{2}(\sigma_{x} - i\sigma_{y}),$$

$$\{\sigma_{i}^{-}, \sigma_{j}^{+}\} = 1, \ \{\sigma_{i}^{-}, \sigma_{j}^{-}\} = 0, \ i = j,$$

$$[\sigma_{i}^{-}, \sigma_{j}^{+}] = 0, \ [\sigma_{i}^{-}, \sigma_{j}^{-}] = 0, \ i \neq j$$

$$\sigma_{i}^{+} |s_{1} \dots s_{i} \dots s_{m}\rangle = \begin{cases} |n_{1} \dots s_{i} = 1 \dots n_{m}\rangle, & s_{i} = 0\\ 0, & s_{i} = 1 \end{cases}$$

### Jordan-Wigner transformation

We need a map from the fermionic operators the qubits operators.

It seems reasonable to try,

$$a_i^{\dagger} = \sigma_i^+, \quad a_i = \sigma_i^-.$$

This doesn't produce the parity counting (the phase  $\zeta_i$ ) that the fermionic operators generate when acting on basis states.

## Jordan-Wigner transformation

We need a map from the fermionic operators the qubits operators.

It seems reasonable to try,

$$a_i^{\dagger} = \sigma_i^+, \quad a_i = \sigma_i^-.$$

This doesn't produce the parity counting (the phase  $\zeta_i$ ) that the fermionic operators generate when acting on basis states. Jordan-Wigner transformation:

$$a_{i}^{\dagger} = \prod_{j=1}^{i-1} \sigma_{j}^{z} \ \sigma_{i}^{+} \ a_{i} = \prod_{j=1}^{i-1} \sigma_{j}^{z} \ \sigma_{i}^{-}.$$

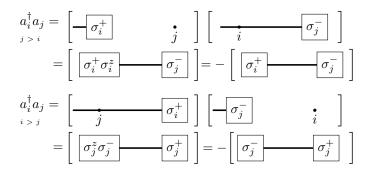
$$\prod_{j=1}^{i-1} \sigma_{j}^{z} \ \sigma_{i}^{+} | s_{1} \dots s_{i} \dots s_{m} \rangle = \begin{cases} \zeta_{i} | n_{1} \dots s_{i} = 1 \dots n_{m} \rangle, & s_{i} = 0 \\ 0, & s_{i} = 1 \end{cases}$$

# Opertators transformation

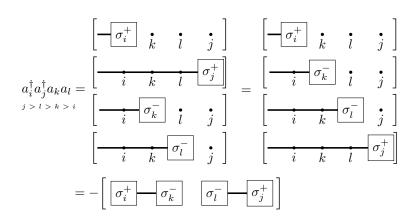
$$a_i^{\dagger} = \overline{\qquad} \sigma_i^{+} \qquad a_i = \overline{\qquad} \sigma_i^{-}$$

## Opertators transformation

$$a_i^{\dagger} = \overline{\hspace{1cm}} \sigma_i^+$$
  $a_i = \overline{\hspace{1cm}} \sigma_i^-$ 



## Opertators transformation



### The qubit Hamiltonian

Remember the fermionic Hamiltonian is,

$$\hat{H} = \sum_{ij} h_1^{ij} a_i^{\dagger} a_j + \frac{1}{2} \sum_{ijkl} h_2^{ijkl} a_i^{\dagger} a_j^{\dagger} a_k a_j$$

- ► For our minimal model for the Hydrogen molecule we have 4 basis, and hence the system will be represented by 4 qubits.
- After Jordan-Wigner transformation the uibit Hamiltonian will look something like this,

$$\hat{H} = \sum_{i,j,k,l=0}^{3} j^{ijkl} \sigma_1^i \sigma_2^j \sigma_3^k \sigma_4^l$$

To measure the  $\langle \hat{H} \rangle$ , our quantum computer will need to measure  $\langle \sigma_1^i \sigma_2^j \sigma_3^k \sigma_4^l \rangle$ .

There are other proposed maps.

# Setting up the quantum circuit

For the Hydrogen example.

$$|\Phi_0\rangle = |1100\rangle$$

$$a_i^{\dagger} a_i = \frac{1}{2} (1 + \sigma_i^z).$$

#### Outline

1. Introduction to circuit model for quantum computing

Converting fermionic Hamiltonians to qubit Hamiltonians Jordan-Wigner

3. VQE algorithms circuit designs for UCCSD

# Unitary coupled-cluser (UCC) expansion

The UCC method uses the following ansatz,

$$\begin{split} |\Psi_0\rangle &= e^{T(\boldsymbol{\theta}) - T^\dagger(\boldsymbol{\theta})} \, |\Phi_0\rangle \\ T(\boldsymbol{\theta}) &= T_1(\boldsymbol{\theta}) + T_2(\boldsymbol{\theta}) + \dots \\ T_1(\boldsymbol{\theta}) &= \sum_{\substack{m \in \mathsf{emp} \\ i \in \mathsf{occ}}} \theta^{mi} a_m^\dagger a_i \\ T_2(\boldsymbol{\theta}) &= \frac{1}{2} \sum_{\substack{m,n \in \mathsf{emp} \\ i,j \in \mathsf{occ}}} \theta^{mnij} a_m^\dagger a_n^\dagger a_i a_j \\ &\vdots \end{split}$$

$$e^{T(\boldsymbol{\theta})-T^{\dagger}(\boldsymbol{\theta})} = e^{\sum_{i}\theta_{i}(\tau_{i}-\tau_{i}^{\dagger})},$$

where  $\tau_i$  include all kinds of excitations.

# Unitary coupled-cluser (UCC) expansion

The UCC method uses the following ansatz,

$$\begin{split} |\Psi_0\rangle &= e^{T(\theta)-T^\dagger(\theta)}\,|\Phi_0\rangle \\ T(\theta) &= T_1(\theta) + T_2(\theta) + \dots \\ T_1(\theta) &= \sum_{\substack{m \in \mathsf{emp} \\ i \in \mathsf{occ}}} \theta^{mi} a_m^\dagger a_i & \mathsf{where} \; \tau_i \; \mathsf{include} \; \mathsf{all} \; \mathsf{kinds} \; \mathsf{of} \\ T_2(\theta) &= \frac{1}{2} \sum_{\substack{m,n \in \mathsf{emp} \\ i \in \mathsf{occ}}} \theta^{mnij} a_n^\dagger a_n^\dagger a_i a_j \end{split}$$

:

- ► For this ansatz to be computationally efficient we need to include only a few excitation operators.
- UCCSD include single and double excitations.
- Next step would be to design a circuit that implement this unitary operation.

#### **Trotterization**

Different terms of  $\tau_i$  don't necessarily commute which make simulating the exponential with quantum circuits not easy. We make use of the following identity,

$$e^{\sum_{i} \theta_{i}(\tau_{i} - \tau_{i}^{\dagger})} = \lim_{N \to \infty} \left( \prod_{i} e^{\frac{\theta_{i}(\tau_{i} - \tau_{i}^{\dagger})}{N}} \right)^{N}$$

#### **Trotterization**

Different terms of  $\tau_i$  don't necessarily commute which make simulating the exponential with quantum circuits not easy. We make use of the following identity,

$$e^{\sum_{i} \theta_{i}(\tau_{i} - \tau_{i}^{\dagger})} = \lim_{N \to \infty} \left( \prod_{i} e^{\frac{\theta_{i}(\tau_{i} - \tau_{i}^{\dagger})}{N}} \right)^{N}$$

lacktriangle Taking N to infinity is not possible for computations. Good results can be achieved by taking just a few trotter steps.

$$e^{\sum_{i} \theta_{i}(\tau_{i} - \tau_{i}^{\dagger})} = \left(\prod_{i} e^{\frac{\theta_{i}(\tau_{i} - \tau_{i}^{\dagger})}{\rho}}\right)^{\rho}$$

- lacktriangle The smaller the  $heta_i$ 's are the better the approximation.
- It's better to start with a good initial guess.

#### one-body exponentials

One body exponential  $e^{\theta(a_m^{\dagger}a_i-a_i^{\dagger}a_m)}$ . Note: m>i.

$$a_m^{\dagger} a_i - a_i^{\dagger} a_m = -\sigma_i^{-} \left( \prod_{s=i+1}^{m-1} \sigma_s^z \right) \sigma_m^{+} + \sigma_i^{+} \left( \prod_{s=i+1}^{m-1} \sigma_s^z \right) \sigma_m^{-}$$

$$a_m^{\dagger} a_i - a_i^{\dagger} a_m = \frac{i}{2} \left[ \sigma_i^y \left( \prod_{s=i+1}^{m-1} \sigma_s^z \right) \sigma_m^x - \sigma_i^x \left( \prod_{s=i+1}^{m-1} \sigma_s^z \right) \sigma_m^y \right]$$

Note that the two terms commute.

#### one-body exponentials

One body exponential  $e^{\theta(a_m^\dagger a_i - a_i^\dagger a_m)}$ . Note: m > i.

$$a_m^{\dagger} a_i - a_i^{\dagger} a_m = -\sigma_i^{-} \left( \prod_{s=i+1}^{m-1} \sigma_s^z \right) \sigma_m^{+} + \sigma_i^{+} \left( \prod_{s=i+1}^{m-1} \sigma_s^z \right) \sigma_m^{-}$$

$$a_m^{\dagger} a_i - a_i^{\dagger} a_m = \frac{i}{2} \left[ \sigma_i^y \left( \prod_{s=i+1}^{m-1} \sigma_s^z \right) \sigma_m^x - \sigma_i^x \left( \prod_{s=i+1}^{m-1} \sigma_s^z \right) \sigma_m^y \right]$$

Note that the two terms commute.

We rotate  $\sigma^x$  and  $\sigma^y$  to  $\sigma^z$ ,

$$e^{i\theta\sigma_i^y \left(\prod_{s=i+1}^{m-1} \sigma_s^z\right)\sigma_m^x} = H_m R_i^x(\pi/2) \left(e^{i\theta\sigma_i^z \left(\prod_{s=i+1}^m \sigma_s^z\right)}\right) H_m^{\dagger} R_i^{x\dagger}(\pi/2),$$

using,

$$H\sigma_z H^{\dagger} = \sigma_x$$
$$R^x (\pi/2) \sigma_z R^{x\dagger} (\pi/2) = \sigma_y$$

#### one-body exponentials

Need to implement  $\exp \left[i\theta\sigma_i^z\left(\prod_{s=i+1}^m\sigma_s^z\right)\right]$ .

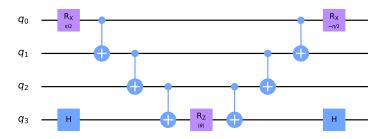
- $ightharpoonup \prod_{s=i+1}^m \sigma_s^z$  counts the parity between qubits i and m. The answer is plus or minus.
- $ightharpoonup e^{\pm i heta \sigma_i^z}$  makes a rotation around the z-axis on the i-th qubit accordingly.

#### one-body exponentials

Need to implement  $\exp \left[i\theta\sigma_i^z\left(\prod_{s=i+1}^m\sigma_s^z\right)\right]$ .

- $ightharpoonup \prod_{s=i+1}^m \sigma_s^z$  counts the parity between qubits i and m. The answer is plus or minus.
- $ightharpoonup e^{\pm i heta\sigma_i^z}$  makes a rotation around the z-axis on the i-th qubit accordingly.

Example:  $\exp\left[i\theta\sigma_1^y(\sigma_2^z\sigma_3^z)\sigma_4^x\right]$  for Hydrogen molecule.



two-body exponentials

Two-body exponential  $e^{\theta(a_m^\dagger a_n^\dagger a_j a_i - a_i^\dagger a_j^\dagger a_n a_m)}$ , and assume j>i, and m>n

$$\begin{split} a_m^\dagger a_n^\dagger a_i a_j - a_i^\dagger a_j^\dagger a_n a_m = & \sigma_i^- \left(\prod \sigma^z\right) \sigma_j^- \ \sigma_n^+ \left(\prod \sigma^z\right) \sigma_m^+ \\ - & \sigma_i^+ \left(\prod \sigma^z\right) \sigma_j^+ \ \sigma_n^- \left(\prod \sigma^z\right) \sigma_m^- \\ = & \frac{i}{4} \sigma_i^x \left(\prod \sigma^z\right) \sigma_j^x \ \sigma_n^x \left(\prod \sigma^z\right) \sigma_m^y \\ + \text{terms with odd } \sigma^y \end{split}$$

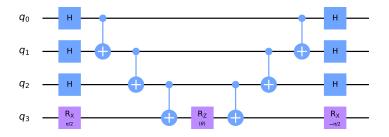
This generate 8 terms, all of which will commute. As before we rotate  $\sigma^x$  and  $\sigma^y$  to  $\sigma^z$ ,

$$H_i H_i H_n Y_m e^{\frac{i\theta}{4} \left(\sigma_i^z \left(\prod_{i=1}^j \sigma^z\right) \left(\prod_n^m \sigma^z\right)\right)} H_i^{\dagger} H_i^{\dagger} H_n^{\dagger} Y_m^{\dagger}$$

two-body exponentials

Note we only add parity from m to n and from j to i+1 then perform the rotation on the i-th qubit.

Example:  $\exp\{[i\theta\sigma_1^x\sigma_2^x\sigma_3^x\sigma_4^y\ ]\}$  for the Hydrogen molecule,



#### Measuerment

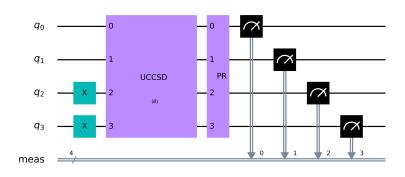
The Hamiltonian expectation value is written as,

$$\langle \hat{H} \rangle = \sum_{i,j,k,l=0}^{3} j^{ijkl} \langle \sigma_1^i \sigma_2^j \sigma_3^k \sigma_4^l \rangle$$

The computational basis is the z-axis. To measure the expectation value a general string of Pauli operators, we need a post-rotation circuit.

- ▶ Measuring the qubits gives a string of 0's and 1's,  $\{s_i\}$ .
- $P(\lbrace s_i \rbrace) = |\langle \lbrace s_i \rbrace | \Psi_0 \rangle|^2$
- ▶ To measure the i-th qubit with respect to the x or y-axes we need to rotate our basis for the i-th qubit using  $H_i$  or  $R_i^x(\pi/2)$  respectively.
- ▶ Any Pauli string has eigenvalues of either  $\pm 1$ .

#### The end result look something like this:



- 1. Initialize the qubits
- 2. Apply the UCCSD gate with parameters  $oldsymbol{ heta}$
- 3. Apply post-rotations .
- 4. Measure the qubits.

#### Comparison:

$$E_{\rm HF} = -1.116$$

$$E_{\mathsf{UCCSD}} = -1.137$$

$$E_{exact} = -1.166$$

### Summary

The VQE method can be summarized as follows:

- A classical computer calculate the Hamiltonian and come up with an initial state.
- 2. A quantum computer samples the Hilber space and measure the energy.
- A classical computer will run an optimization routine to minimize the energy expectation value.

