Modeling Electronic Systems on Quantum Computers

University of Florida, October

Outline

1. Hartree-Fock and beyond on classical computers
   Where quantum computers can speed things up

2. Introduction to circuit model for quantum computing

3. Converting fermionic Hamiltonians to qubit Hamiltonians
   Jordan-Wigner

4. VQE algorithms circuit designs for UCCSD

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The ground state $|\Psi_0\rangle$ is the state with the lowest $\langle H \rangle$.

Second-quantized form of the Hamiltonian

- For some molecular orbitals (MO) $\{\chi_i(\boldsymbol{r})\}$. $\langle\chi_i|\chi_j\rangle = \delta_{ij}$
- Define $a_i^\dagger |0\rangle = |\chi_i\rangle$
- $a_i |0\rangle = 0$
- $\{a_i, a_j^\dagger\} = \delta_{ij},\ \{a_i, a_j\} = \{a_i^\dagger, a_j^\dagger\} = 0$

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The Second-quantized Hamiltonian (dropping V_{nn}),

$$\hat{H} = \sum_{ij} h_1^{ij} a_i^\dagger a_j + \frac{1}{2}\sum_{ijkl} h_2^{ijkl} a_i^\dagger a_j^\dagger a_k a_l$$

$$h_1^{ij} = \langle h_1\rangle = \int d\boldsymbol{r}\ \chi^{\dagger i}(\boldsymbol{r}) \left[\frac{\boldsymbol{P}^2}{2m} - \sum_\alpha \frac{e^2 Z_\alpha}{|\boldsymbol{r} - \boldsymbol{R}_\alpha|}\right] \chi^j(\boldsymbol{r})$$

$$h_2^{ijkl} = \langle h_2\rangle = \int d\boldsymbol{r} d\boldsymbol{r}' \chi^{\dagger i}(\boldsymbol{r}) \otimes \chi^{\dagger j}(\boldsymbol{r}') \frac{e^2}{|\boldsymbol{r} - \boldsymbol{r}'|} \chi^k(\boldsymbol{r}') \otimes \chi^l(\boldsymbol{r})$$

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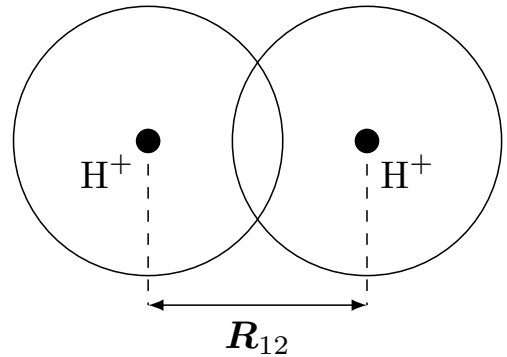
H_2 molecule, a minimal model.

- A complete set of molecular orbits are infinite, and hence computationally not efficient.



- We expand the (relevant) molecular orbits in some finite basis set $\{\phi_i(r)\}$,

$$|\chi_i\rangle = \sum_j C_{ji} |\phi_j\rangle \,, \;\; \langle \phi_i | \phi_j \rangle = S_{ij}$$

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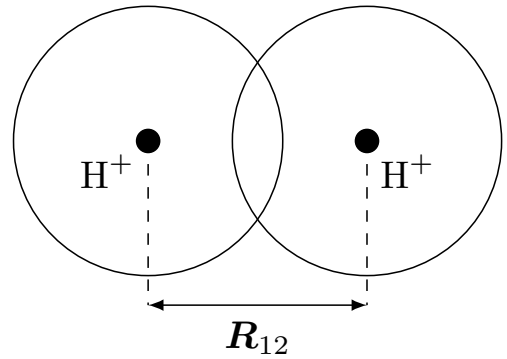

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A convenient choice of basis (atomic orbitals AO) for H_2 is,

$$|\phi_1\rangle = |\phi_L^{1s} \uparrow\rangle \,, \;\; |\phi_2\rangle = |\phi_L^{1s} \downarrow\rangle \,, \;\; |\phi_3\rangle = |\phi_R^{1s} \uparrow\rangle \,, \;\; |\phi_4\rangle = |\phi_R^{1s} \downarrow\rangle$$

In general, the more AO we include, the more accurate the calculation is.

Solving the problem on a classical computer

1. Choose an initial guess for $|\Psi_0\rangle$.
2. Calculate $\langle\Psi_0|H|\Psi_0\rangle$.
3. Have a way of varying $|\Psi_0\rangle$.
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How hard is the problem? The size of the Hilbert space can give us an idea.

- For our H_2 example with 4 basis functions and 2 electrons dimension of relevant Hilbert subspace is $\binom{4}{2} = 6$.
- In general, for n electrons and m basis functions, the Hilbert space dimension is $\binom{m}{n}$ which is $\mathcal{O}(e^n)$.

This make steps no. 2 and 3 computationally very costly for systems with more than a few electrons.

Hartree-Fock approximation

The Hartree-Fock approximation is very popular and offers an exponential reduction in the complexity of the problem.

> The Hartree-Fock approximation is the assumption that $|\Psi_0\rangle$ can be written as a Slater determinant.

Note: A Slater determinant state means no entanglement between the particles.

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$$|\Psi_0\rangle \approx |\Phi_0\rangle = \prod_{i=1}^{n} a_i^\dagger |0\rangle ,$$

- $|\Phi_0\rangle$ is the Slater determinant state with lowest $\langle H \rangle$.
- $\langle \Phi_0 | H | \Phi_0 \rangle = E_{\mathsf{HF}}$
- The problem is to find the optimal $\{\chi_i(\boldsymbol{r})\}$ HF orbitals.
- HF offers an exponential reduction in size of Hilbert space.

Dimension of product states space

A general Slater determinant state can be written as,

$$|\Phi_0\rangle = \prod_{i=1}^{n} a_i^{'\dagger} |0\rangle = \prod_{i=0}^{n} \left(\hat{U} a_i^\dagger \hat{U}^{-1}\right) |0\rangle = \hat{U} \prod_{i=1}^{n} a_i^\dagger |0\rangle$$

A general change of basis operation can be written as,

$$a_i^{'\dagger} = \hat{U} a_i^\dagger \hat{U}^{-1} = \sum_j U_{ji}\ a_j^\dagger,$$

$$\hat{U} = e^{i \sum_{i,j} \kappa^{ij} a_i^\dagger a_j}, \quad e^{i\kappa} = U$$

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- The state $\prod_{i=1}^{n} a_i^{\dagger} |0\rangle$ is invariant under transformations that mix the occupied states together or mix the empty states together.
- We are interested in the coset space $U(m)/U(n) \times U(m-n)$ which is $n(m-n)$ dimensional.

Hydrogen molecule Hartree-Fock

The Hartree-Fock orbitals for the Hydrogen atom:

$$|\Phi_0\rangle = a_1^\dagger a_2^\dagger |0\rangle = |1100\rangle$$

$$|\chi_1\rangle = N\left(|\phi_L^{1s} \uparrow\rangle + |\phi_R^{1s} \uparrow\rangle\right)$$
$$|\chi_2\rangle = N\left(|\phi_L^{1s} \downarrow\rangle + |\phi_R^{1s} \downarrow\rangle\right)$$
$$|\chi_3\rangle = N\left(|\phi_L^{1s} \uparrow\rangle - |\phi_R^{1s} \uparrow\rangle\right)$$
$$|\chi_4\rangle = N\left(|\phi_L^{1s} \downarrow\rangle - |\phi_R^{1s} \downarrow\rangle\right)$$

Some of the one-body, and two-body integrals evaluated at $\boldsymbol{R}_{12} = 0.74$A, and all energies measured in Hartrees. (Using a python package called PySCF):

$h_1^{ij} = \mathsf{diag}(-1.253, -1.253, -0.475, -0.475).$

$h_2^{1221} = 0.675, \quad h_2^{2121} = 0, h_2^{1331} = 0.664, \quad h_2^{3131} = 0.181, \ldots$

$E_{\mathsf{HF}} = -1.117$

Where quantum computer can help

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  1. Sampling the Hilbert space.
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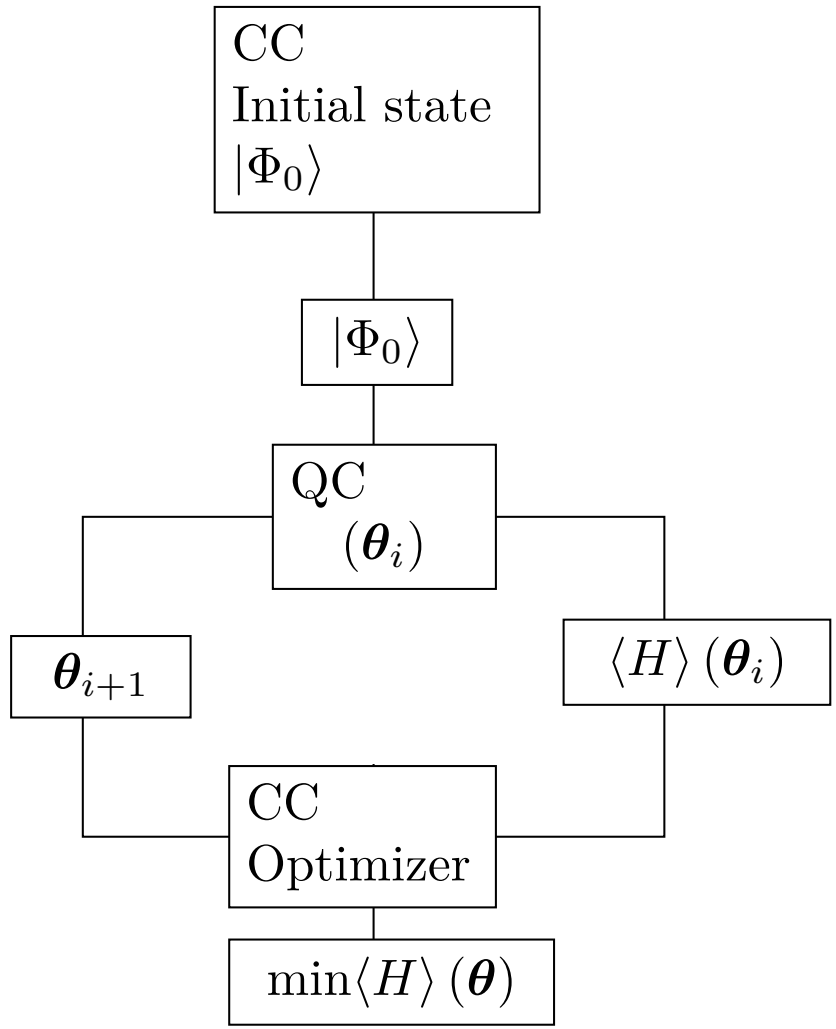
- Classical computer can do HF efficiently.
- In going beyond HF the bottle necks are:
  1. Sampling the Hilbert space.
  2. Calculating the expectation value of the Hamiltonian.
- A variational quantum eigensolver algorithm (VQE), put the exponentially hard part of the calculation on a quantum computer.



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A qubit lives in a 2 dimensional Hilbert space, $\mathcal{H}^1$. An arbitrary qubit state can be described as,

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A system of N qubits lives in a 2^N dimensional Hilbert space, $\mathcal{H}^N$ such that,

$$\mathcal{H}^N = \bigotimes_i^N \mathcal{H}_i^1$$

An arbitrary state describing the system of N qubits can be written as,

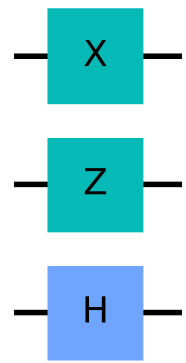
$$|\Psi\rangle = \sum_{\{s_i\}} c_{1\ldots N} |s_1 \ldots s_N\rangle , \;\; s_1, \ldots s_N = \{0, \; 1\}$$

Gates

Operators can act on states in $\mathcal{H}^N$ can act on one or two (or more) qubits independently. These operators acting on the qubits are called quantum gates.

Examples of 1-qubit gates:

- The X-gate: $X|0\rangle = |1\rangle\,,\ X|1\rangle = |0\rangle.$
- The Z-gate: $Z|0\rangle = |0\rangle\,,\ Z|1\rangle = -|1\rangle.$
- The H-gate:
  $H|0\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)\,,\ H|1\rangle = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle).$



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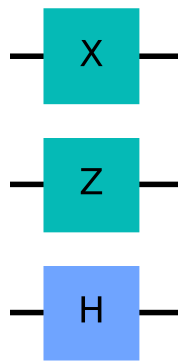
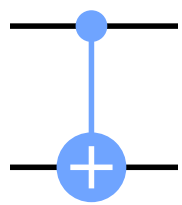


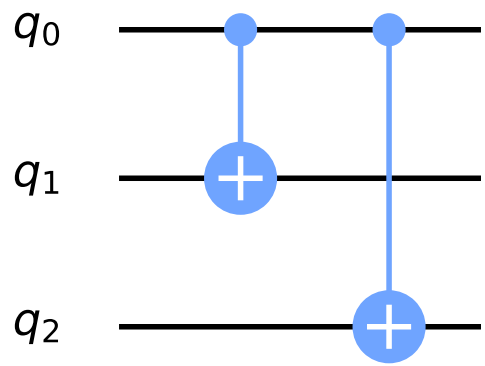
Example of 2-qubit gates:

- Controlled-X gate: $cX|00\rangle = |00\rangle\,,\ cX|01\rangle = |01\rangle\,, cX|10\rangle = |11\rangle\,,\ cX|11\rangle = |10\rangle$. It *adds* the two qubits and puts the result in the second qubit



Circuits

Example 1: $|0\rangle \to |000\rangle\,,\ |1\rangle \to |111\rangle$. (Quantum repetition code.)



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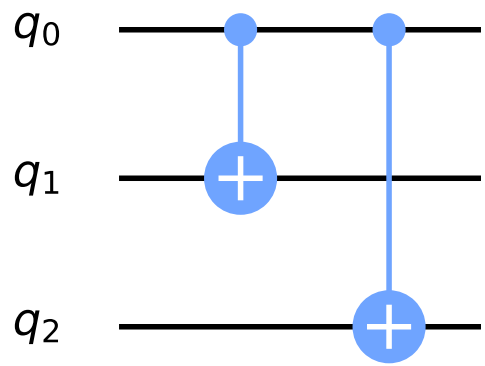
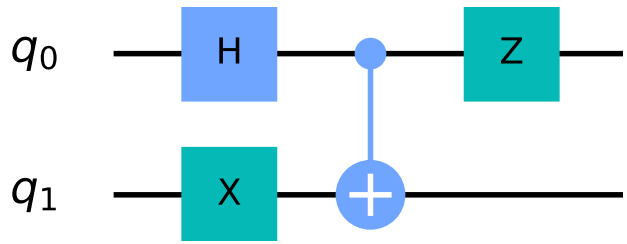


Example 2: $|00\rangle \to \frac{1}{\sqrt{2}}\left(|01\rangle - |10\rangle\right)$, (Adds entanglement.)



- The most prevalent model for quantum computations is the circuit model.
- A circuit model is similar to how classical computers operate.

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Fermionic and qubit degrees of freedom

Fermionic system with m orbitals: (2^m dimensional)

$$\{a_i, a_j^\dagger\} = \delta_{ij}, \ \{a_i, a_j\} = 0.$$

$$a_i^\dagger \left|n_1 \dots n_i \dots n_m\right\rangle = \begin{cases} \zeta_i \left|n_1 \dots n_i = 1 \dots n_m\right\rangle, & n_i = 0 \\ 0, & n_i = 1 \end{cases}$$

$$\zeta_i = (-1)^{\sum_{j=1}^{i-1} n_j} \quad \text{Parity counting}$$

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Qubits system with m qubits:

$$\sigma^+ = \frac{1}{2}(\sigma_x + i\sigma_y), \ \sigma^- = \frac{1}{2}(\sigma_x - i\sigma_y),$$

$$\{\sigma_i^-, \sigma_j^+\} = 1, \ \{\sigma_i^-, \sigma_j^-\} = 0, \quad i = j,$$

$$[\sigma_i^-, \sigma_j^+] = 0, \quad [\sigma_i^-, \sigma_j^-] = 0, \quad i \neq j$$

$$\sigma_i^+ \left|s_1 \dots s_i \dots s_m\right\rangle = \begin{cases} \left|n_1 \dots s_i = 1 \dots n_m\right\rangle, & s_i = 0 \\ 0, & s_i = 1 \end{cases}$$

Jordan-Wigner transformation

We need a map from the fermionic degrees of freedom to the qubits degrees of freedom.

It seems reasonable to try,

$$a_i^\dagger = \sigma_i^+, \quad a_i = \sigma_i^-.$$

This doesn't produce the parity counting (the phase ζ_i) that the fermionic operators generate when acting on basis states.

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Jordan-Wigner transformation:

$$a_i^\dagger = \prod_{j=1}^{i-1} \sigma_j^z \; \sigma_i^+ \quad a_i = \prod_{j=1}^{i-1} \sigma_j^z \; \sigma_i^- .$$

$$\prod_{j=1}^{i-1} \sigma_j^z \; \sigma_i^+ \, |s_1 \dots s_i \dots s_m\rangle = \begin{cases} \zeta_i \, |n_1 \dots s_i = 1 \dots n_m\rangle \, , & s_i = 0 \\ 0, & s_i = 1 \end{cases}$$

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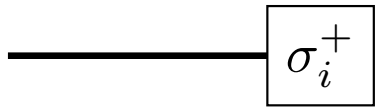
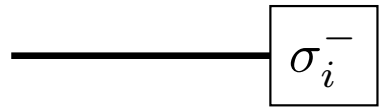
$a_i^\dagger =$ 

$a_i =$ 

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$$a_i^\dagger = \sigma_i^+ \qquad a_i = \sigma_i^-$$

$$a_i^\dagger a_j = \left[\sigma_i^+ \quad \overset{\bullet}{j} \right] \left[\overset{\bullet}{i} \quad \sigma_j^- \right]$$
$$j > i$$

$$= \left[\sigma_i^+ \sigma_i^z \quad \sigma_j^- \right] = - \left[\sigma_i^+ \quad \sigma_j^- \right]$$

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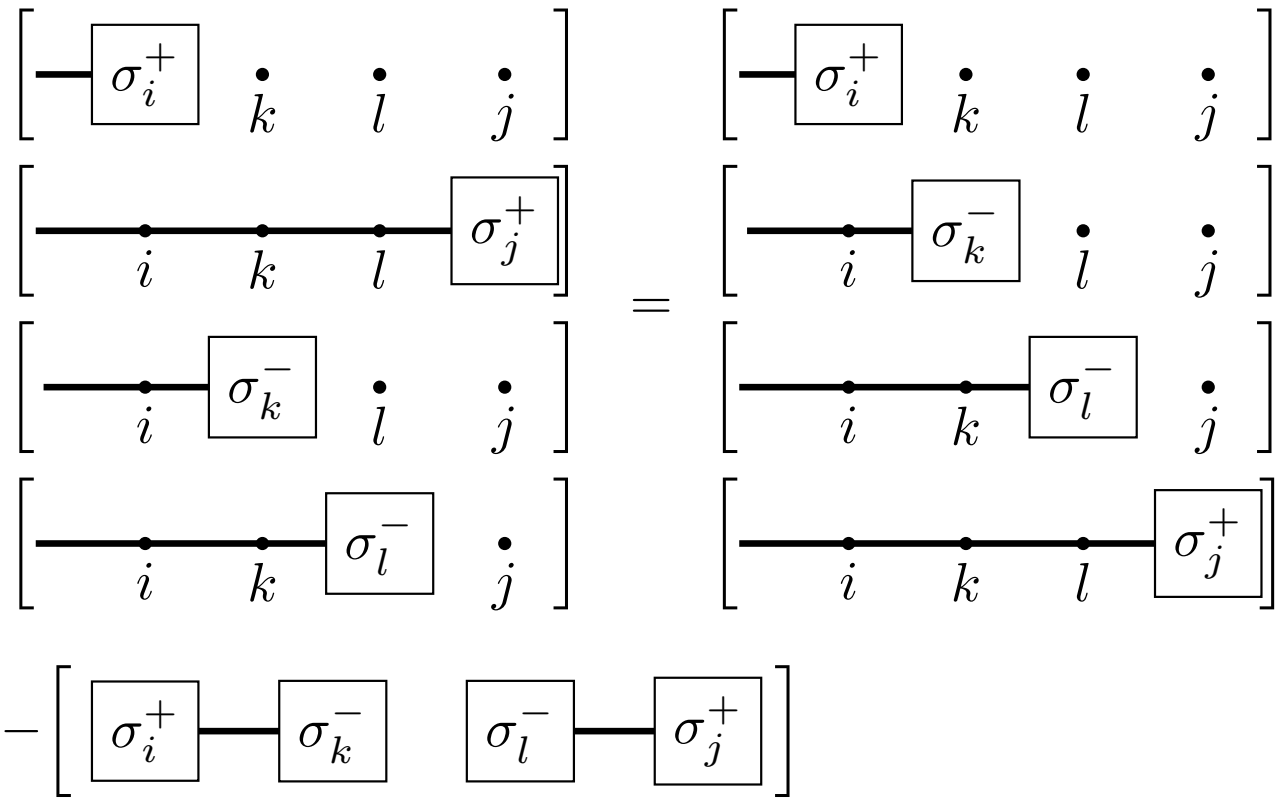
Opertators transformation

$$a_i^\dagger a_j^\dagger a_k a_l = \underset{j>l>k>i}{}$$



$$= -\Big[\ \sigma_i^+ \text{—} \sigma_k^- \quad \sigma_l^- \text{—} \sigma_j^+\ \Big]$$

The qubit Hamiltonian

Remember the fermionic Hamiltonian is,

$$\hat{H} = \sum_{ij} h_1^{ij} a_i^\dagger a_j + \frac{1}{2} \sum_{ijkl} h_2^{ijkl} a_i^\dagger a_j^\dagger a_k a_j$$

- For our minimal model for the Hydrogen molecule we have 4 basis, and hence the system will be represented by 4 qubits.
- After Jordan-Wigner transformation the uibit Hamiltonian will look something like this,

$$\hat{H} = \sum_{i,j,k,l=0}^{3} j^{ijkl} \sigma_1^i \sigma_2^j \sigma_3^k \sigma_4^l$$

To measure the $\langle \hat{H} \rangle$, our quantum computer will need to measure $\langle \sigma_1^i \sigma_2^j \sigma_3^k \sigma_4^l \rangle$.

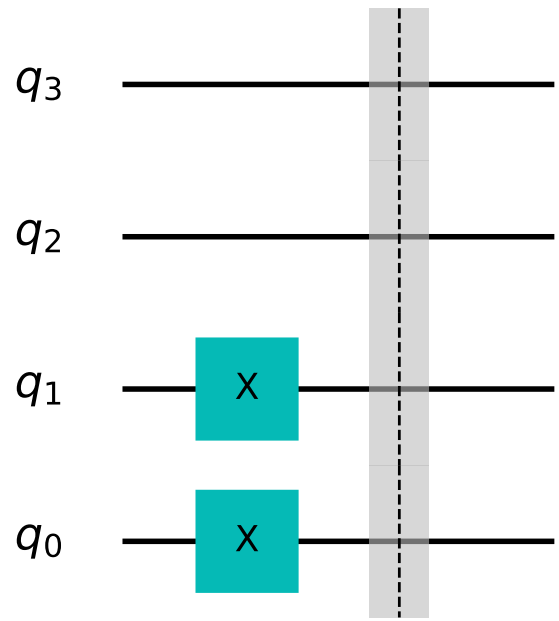
- There are other proposed maps.

Setting up the quantum circuit

For the Hydrogen example.

$$|\Phi_0\rangle = |1100\rangle$$

$$a_i^\dagger a_i = \frac{1}{2}(1 + \sigma_i^z).$$



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Unitary coupled-cluser (UCC) expansion

The UCC method uses the following ansatz,

$$|\Psi_0\rangle = e^{T(\boldsymbol{\theta})-T^\dagger(\boldsymbol{\theta})}\,|\Phi_0\rangle$$

$$T(\boldsymbol{\theta}) = T_1(\boldsymbol{\theta}) + T_2(\boldsymbol{\theta}) + \ldots$$

$$T_1(\boldsymbol{\theta}) = \sum_{\substack{m\in\text{emp}\\ i\in\text{occ}}} \theta^{mi} a_m^\dagger a_i$$

$$T_2(\boldsymbol{\theta}) = \frac{1}{2}\sum_{\substack{m,n\in\text{emp}\\ i,j\in\text{occ}}} \theta^{mnij} a_m^\dagger a_n^\dagger a_i a_j$$

$$\vdots$$

$$e^{T(\boldsymbol{\theta})-T^\dagger(\boldsymbol{\theta})} = e^{\sum_i \theta_i(\tau_i-\tau_i^\dagger)},$$

where τ_i include all kinds of excitations.

Unitary coupled-cluser (UCC) expansion

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where τ_i include all kinds of excitations.

- For this ansatz to be computationally efficient we need to include only a few excitation operators.
- UCCSD include single and double excitations.
- Next step would be to design a circuit that implement this unitary operation.

Trotterization

Different terms of τ_i don't necessarily commute which make simulating the exponential with quantum circuits not easy. We make use of the following identity,

$$e^{\sum_i \theta_i(\tau_i - \tau_i^\dagger)} = \lim_{N\to\infty} \left(\prod_i e^{\frac{\theta_i(\tau_i - \tau_i^\dagger)}{N}} \right)^N$$

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- Taking N to infinity is not possible for computations. Good results can be achieved by taking just a few trotter steps.

$$e^{\sum_i \theta_i(\tau_i - \tau_i^\dagger)} = \left(\prod_i e^{\frac{\theta_i(\tau_i - \tau_i^\dagger)}{\rho}} \right)^\rho$$

- The smaller the θ_i's are the better the approximation.
- It's better to start with a good initial guess.

Converting to qubits operations

one-body exponentials

One body exponential $e^{\theta(a_m^\dagger a_i - a_i^\dagger a_m)}$. Note: $m > i$.

$$a_m^\dagger a_i - a_i^\dagger a_m = -\sigma_i^- \left(\prod_{s=i+1}^{m-1} \sigma_s^z \right) \sigma_m^+ + \sigma_i^+ \left(\prod_{s=i+1}^{m-1} \sigma_s^z \right) \sigma_m^-$$

$$a_m^\dagger a_i - a_i^\dagger a_m = \frac{i}{2} \left[\sigma_i^y \left(\prod_{s=i+1}^{m-1} \sigma_s^z \right) \sigma_m^x - \sigma_i^x \left(\prod_{s=i+1}^{m-1} \sigma_s^z \right) \sigma_m^y \right]$$

Note that the two terms commute.

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Note that the two terms commute.

We rotate σ^x and σ^y to σ^z,

$$e^{i\theta\sigma_i^y \left(\prod_{s=i+1}^{m-1} \sigma_s^z\right)\sigma_m^x} = H_m R_i^x(\pi/2) \left(e^{i\theta\sigma_i^z \left(\prod_{s=i+1}^{m} \sigma_s^z\right)} \right) H_m^\dagger R_i^{x\dagger}(\pi/2),$$

using,

$$H\sigma_z H^\dagger = \sigma_x$$

$$R^x(\pi/2)\,\sigma_z R^{x\dagger}(\pi/2) = \sigma_y$$

Converting to qubit operations

one-body exponentials

Need to implement $\exp\left[i\theta\sigma_i^z\left(\prod_{s=i+1}^{m}\sigma_s^z\right)\right]$.

- $\prod_{s=i+1}^{m}\sigma_s^z$ counts the parity between qubits i and m. The answer is plus or minus.
- $e^{\pm i\theta\sigma_i^z}$ makes a rotation around the z-axis on the i-th qubit accordingly.

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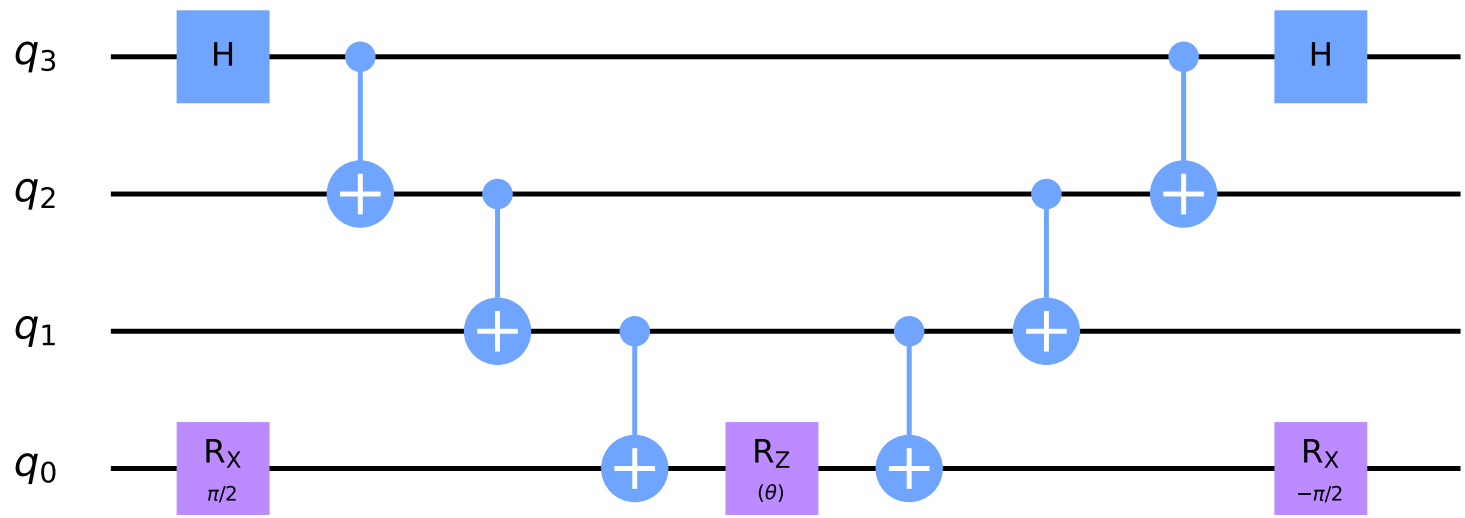
Example: $\exp\left[i\theta\sigma_1^y(\sigma_2^z\sigma_3^z)\sigma_4^x\right]$ for Hydrogen molecule.



Converting to qubit operations

two-body exponentials

Two-body exponential $e^{\theta(a_m^\dagger a_n^\dagger a_j a_i - a_i^\dagger a_j^\dagger a_n a_m)}$, and assume $j > i$, and $m > n$

$$\begin{aligned} a_m^\dagger a_n^\dagger a_i a_j - a_i^\dagger a_j^\dagger a_n a_m =& \sigma_i^- \left(\prod \sigma^z\right) \sigma_j^- \; \sigma_n^+ \left(\prod \sigma^z\right) \sigma_m^+ \\ & - \sigma_i^+ \left(\prod \sigma^z\right) \sigma_j^+ \; \sigma_n^- \left(\prod \sigma^z\right) \sigma_m^- \\ =& \frac{i}{4} \sigma_i^x \left(\prod \sigma^z\right) \sigma_j^x \; \sigma_n^x \left(\prod \sigma^z\right) \sigma_m^y \\ & + \text{terms with odd } \sigma^y \end{aligned}$$

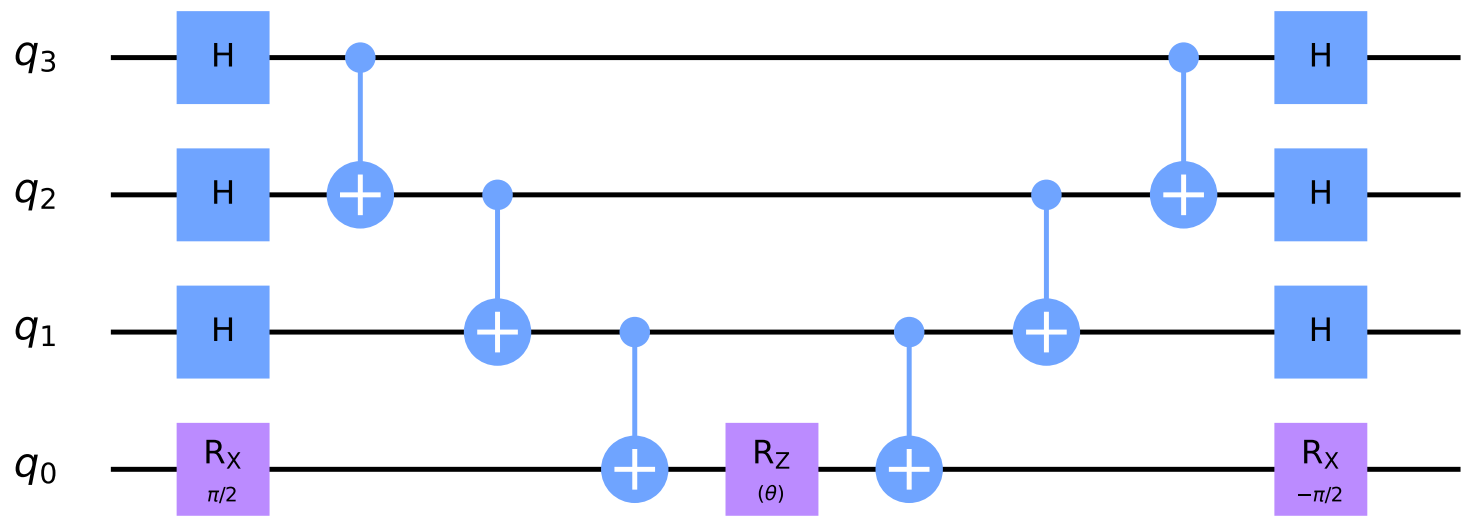
This generate 8 terms, all of which will commute. As before we rotate σ^x and σ^y to σ^z,

$$H_i H_i H_n Y_m e^{\frac{i\theta}{4}\left(\sigma_i^z \left(\prod_{i+1}^{j} \sigma^z\right) \left(\prod_n^m \sigma^z\right)\right)} H_i^\dagger H_i^\dagger H_n^\dagger Y_m^\dagger$$

Converting to qubit operations

two-body exponentials

Note we only add parity from m to n and from j to $i+1$ then perform the rotation on the i-th qubit.
Example: $\exp\left[i\theta\sigma_1^x\sigma_2^x\sigma_3^x\sigma_4^y\right]$ for the Hydrogen molecule,



Measuerment

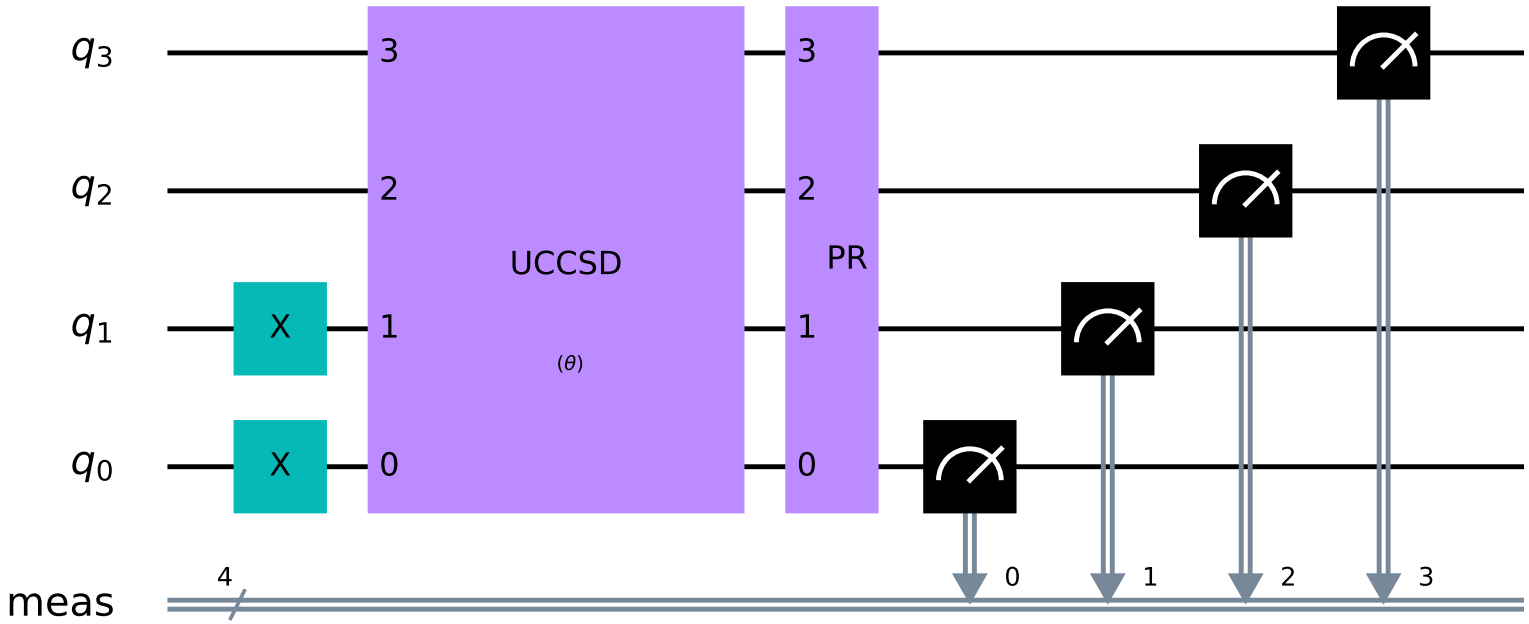
The Hamiltonian expectation value is written as,

$$\langle \hat{H} \rangle = \sum_{i,j,k,l=0}^{3} j^{ijkl} \langle \sigma_1^i \sigma_2^j \sigma_3^k \sigma_4^l \rangle$$

The computational basis is the z-axis. To measure the expectation value a general string of Pauli operators, we need a post-rotation circuit.

- Measuring the qubits gives a string of 0's and 1's, $\{s_i\}$.
- $P(\{s_i\}) = |\langle \{s_i\} | \Psi_0 \rangle|^2$
- To measure the i-th qubit with respect to the x or y-axes we need to rotate our basis for the i-th qubit using H_i or $R_i^x(\pi/2)$ respectively.
- Any Pauli string has eigenvalues of either ± 1.
- $\langle \sigma_1^i \sigma_2^j \sigma_3^k \sigma_4^l \rangle = P(1) - P(-1)$

The end result look something like this:



1. Initialize the qubits
2. Apply the UCCSD gate with parameters $\boldsymbol{\theta}$
3. Apply post-rotations .
4. Measure the qubits.

Comparison:

$E_{\text{HF}} = -1.116$

$E_{\text{UCCSD}} = -1.137$

$E_{exact} = -1.166$

Summary

The VQE method can be summarized as follows:

1. A classical computer calculate the Hamiltonian and come up with an initial state.
2. A quantum computer samples the Hilber space and measure the energy.
3. A classical computer will run an optimization routine to minimize the energy expectation value.