

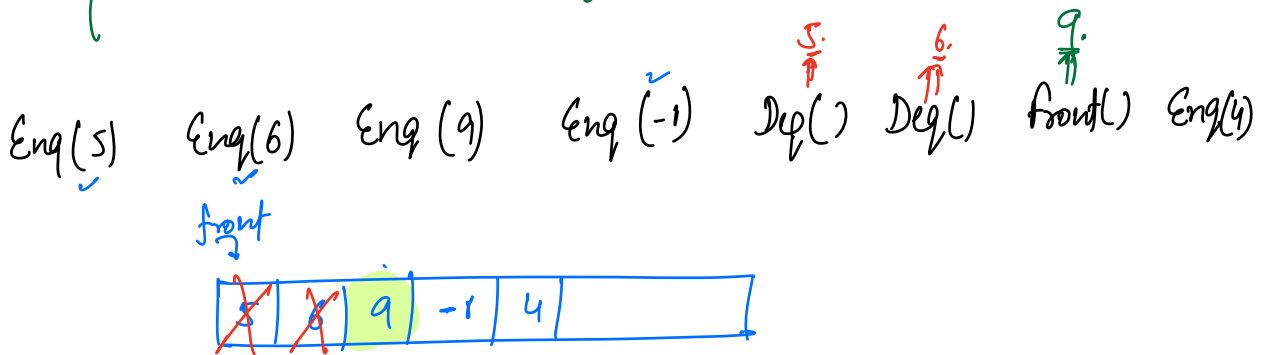
FIFO - [First In First Out]

- Movie Ticket / Car Wash
- Call centre
- Printer.
- Job scheduling.
- Message Queue.

Operations by Queue.

o(i) {

- Enqueue(x) → insert the data at rear end.
- Dequeue() → remove data from front.
- front() → data present at front end.
- size() → no. of elements in the queue.

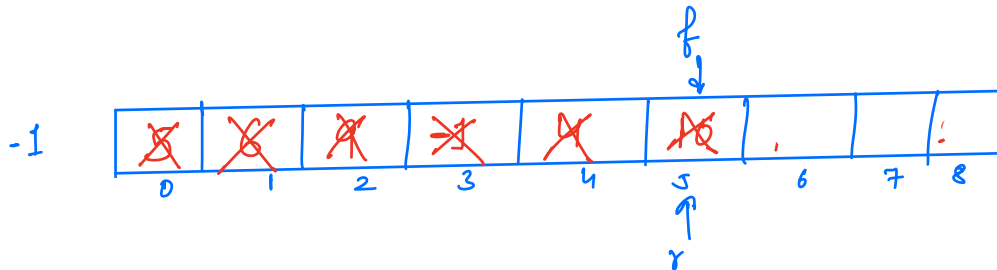


Implementation of Queue

Arrays:

$f = -1 \rightarrow$ idx of last element which was just removed

$r = -1 \rightarrow$ idx of last element which was just inserted



Enq(5) Enq(6) Enq(9) Enq(-1) Deq() Deq() Enq(4)

Enqueue(x) {

if ($r == N-1$) { // Queue is full }

$r++$

$arr[r] = x;$

}

Dequeue() {

if ($f == r$) { // Queue is Empty }

$f = f + 1;$

}

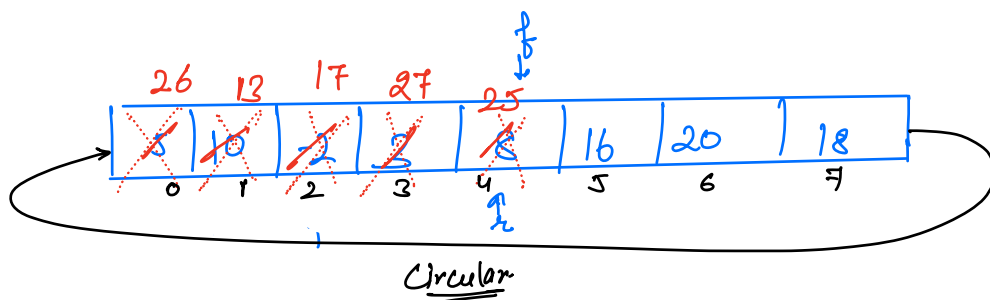
front() {

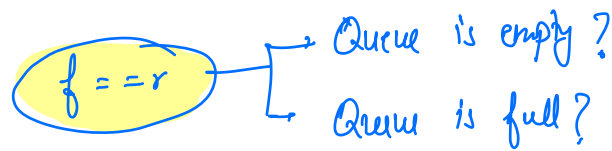
if ($f == r$) { // Queue is Empty }

return $arr[f+1];$

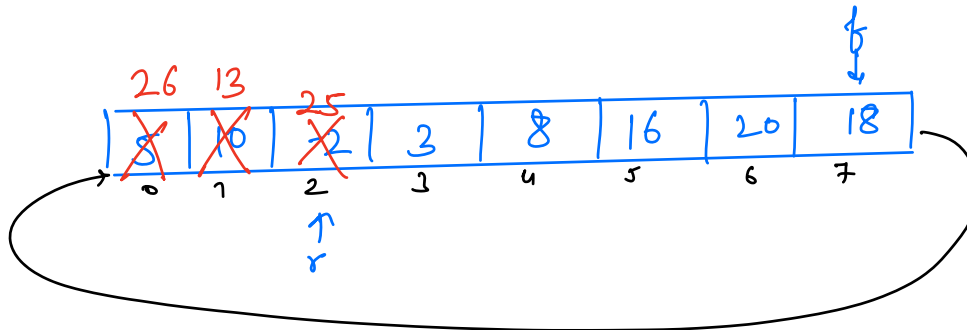
}

Eg:-





size \rightarrow ~~8~~ ~~7~~ ~~6~~ ~~5~~
 Capacity \rightarrow 8



enq(26)
 enq(13)
 enq(25)
 enq(-5);

Deq()
 Deq()
 Deq()
 Deq()
 Deq()
 Deq()

```

Enqueue(x) {
    if (size == capacity) { // Queue is full }
        r = (r + 1) % N;
        arr[r] = x;
        size++;
}
  
```

```

Dequeue() {
    if (size == 0) { // Queue is Empty }
        f = (f + 1) % N;
        size--;
}
  
```

```

front() {
    if (size == 0) { // Queue is Empty }
        return arr[(f + 1) % N];
}
  
```

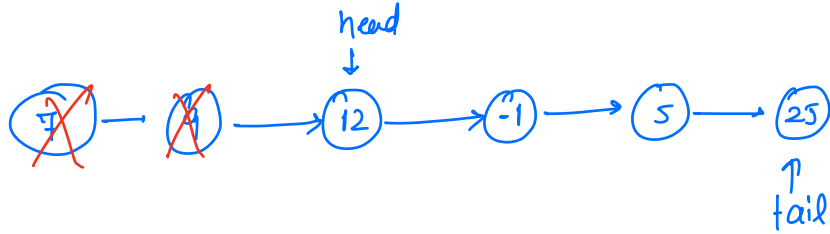
→ Implementation Using Linked-list

→ Enqueue insertion → tail / head

→ Dequeue deletion → tail / head

↓

$O(N)$



insertion

Node n = new Node(s)?

$$\text{tail.next} = n;$$

tail = tail.next

Engl(s)

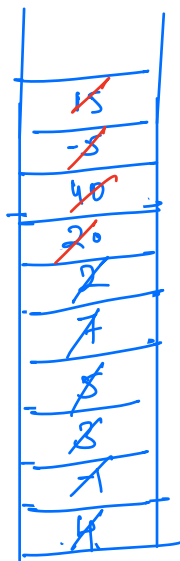
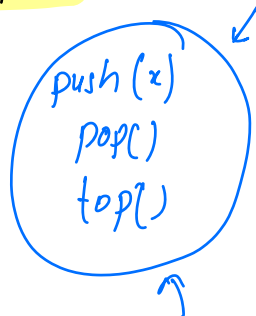
Exq(25)

deletion

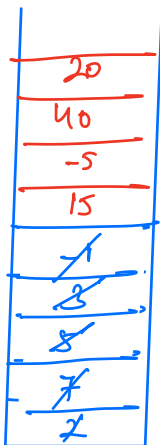
head = head.next

Q. Implement queue using stacks.

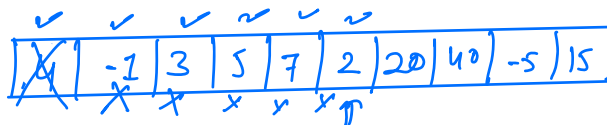
$O(1)$ { enqueue()
dequeue()
front()



st1



st2



deq() → 4.

eng(20);

eng(40)

eng(-5);

eng(15)

deq() → -1

deq() → 3

deq() → 5

deq() → 7

deq() → 2

deq() → 20.

N dequ → $2N + 1$.

1 deq → $\frac{2N+1}{N} \Rightarrow \underline{2} \therefore O(1)$

$O(1)$ { Enqueue(x) → push x in st1

$O(1)$ { Dequeue() → if (st2.size() == 0)

in ans
avg. case.

yes
transfer everything from st1
to st2 and then
↳ st2.pop()

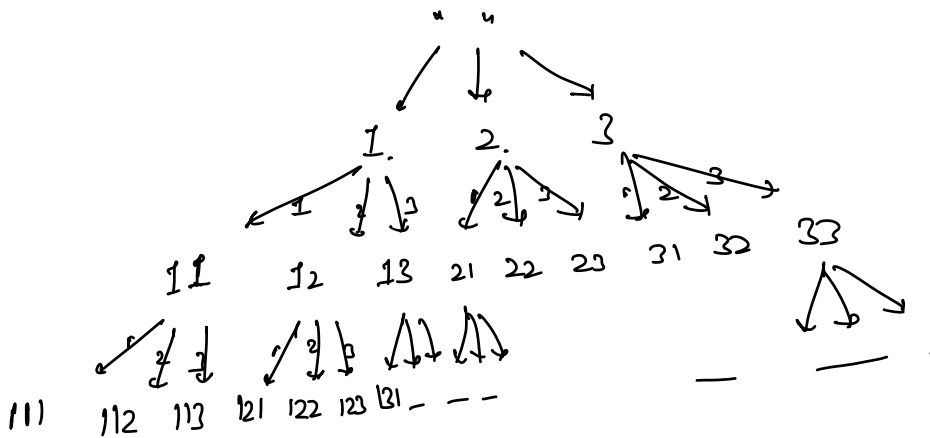
no
pop from st2.

→ Nth number

1, 2, 3.

N=1	1		31
N=2	2		32
N=3	3		33
N=4	11		111
N=5	12		112
N=6	13		113
	21		121
	22		122
	23		123
			1

$\left[\begin{array}{l} N=13, \text{ans} = 111 \\ N=17, \text{ans} = 122 \\ N=6, \text{ans} = 13 \end{array} \right]$



N=17

~~1~~ | ~~2~~ | ~~3~~ | ~~11~~ | ~~12~~ | 13 | 21 | 22 | 23 | 31 | 32 | 33 | 111 | 112 | 113 | 121 | 122 |

count = 2, 4, 8, 16, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17

```
Queue<int> q;
```

```
q.enqueue(1);
```

```
q.enqueue(2);
```

```
q.enqueue(3);
```

```
while(count <= N){
```

```
    x = q.front();
```

```
    q.dequeue();
```

```
    q.enqueue(x*10+1);
```

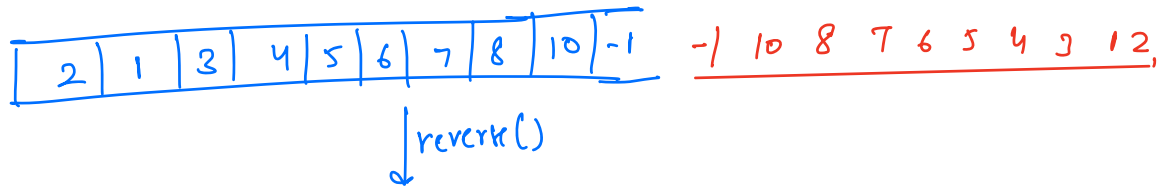
```
    q.enqueue(x*10+2);
```

```
    q.enqueue(x*10+3);
```

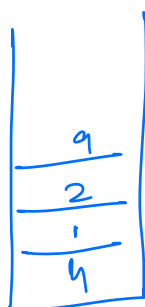
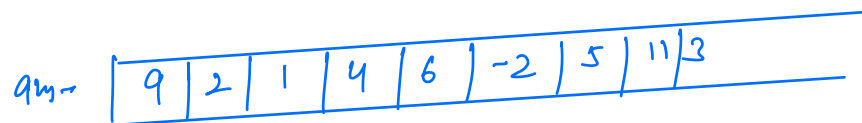
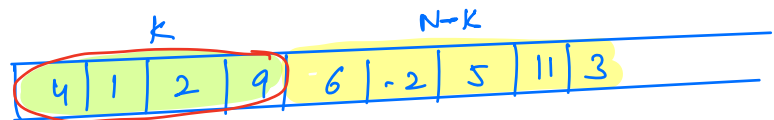
```
}
```

$T.C \rightarrow O(N)$
 $S.C \rightarrow O(N)$

• Reverse a Queue -



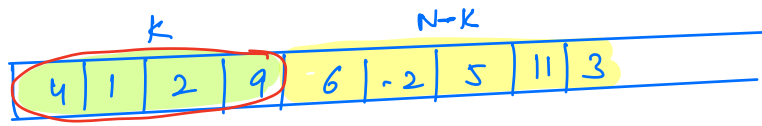
Reverse first k elements of the queue.



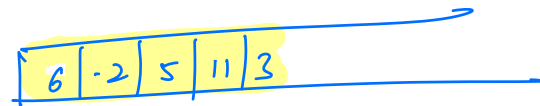
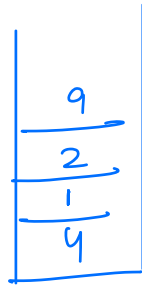
1 stack



1 Queue.



① push first k elements in the stack



② pop all elements from stack & enqueue them in the same queue.



③ Dequeue $\overset{(5)}{n-k}$ elements from the queue & enqueue them in the same queue.

