# Todoy's Quote -

# YOU DON'T HAVE TO BE GREAT TO START, BUT YOU HAVE TO START TO BE GREAT

[Arrays.]

## Today's.

$$arr = \begin{cases} 3, 1, 2, 5, 7, 4, 9 \end{cases}$$
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# How to print all the elements of array.

Void print Arr ( aux) 
$$\S$$

Tic  $\to$  o(N)

Sign (  $i = 0$  to  $N-1$ )  $\S$ 

print ( aux(i));

 $\S$ 
 $\S$ 

N iterations.

Qi) Given N array elements, count no of elements, having at least one element greater than itself.

observation]: for maximum element, we can't have any non greater than itself.

Jind ans. → total nor of elements - C.

no of times maxim no. is appearing.

```
pseudocode. \rightarrow

int count liveater (arr, N) {

max = arr, [0];

for (i \rightarrow 1 to n-1) {

if (arr[i] > max) {

max = arr, [i]

y

c = 0;

for (i \rightarrow 0 to n-1) {

if (arr, [i] == max) {

if (arr, [i] == max) {

c = c+1;

}

return N-C;
```

to Do: fry to solve it by iterating on the array only once.

### Python

# Java.

```
int n = arr.length;
int max = arr[0];
for(int i = 1; i < n; i++){
    if(arr[i] > max){
        max = arr[i];
    }
}
int count = 0;
for(int i = 0; i < n; i++){
    if(arr[i] == max){
        count++;
    }
}
System.out.println(n - count);</pre>
```

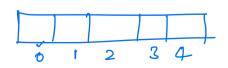
Q:) Given N array elements, check if there exists a poir(i,j) such that am(i] + am(j] == k & !!= j'

Note:-, i and j are index value, K is given sum.

arr(); {3 -2 1 4 3 6 8 3 : true

arr(7: \$2 4 -3 73! falke

aux(i) + aux(j) => 2 + 3 = 6

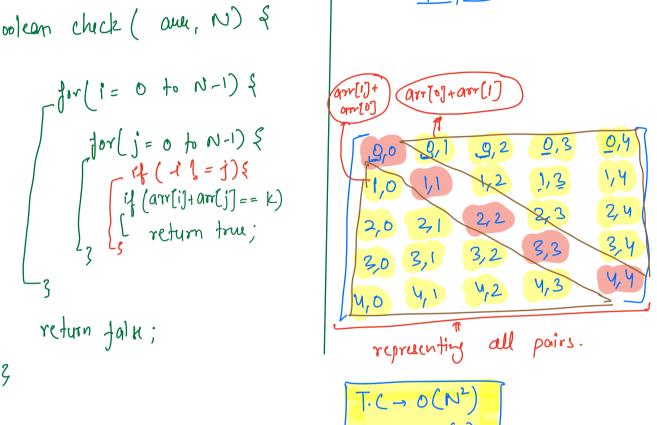


idea-1 - Check all the pairs. If any pair (i,j) is having sum = K -> retyrn true.

boolean check ( and, N) &

2

All pairs



$$\begin{array}{c} T. (\rightarrow O(N^2) \\ S. (\rightarrow O(1) \\ \end{array}$$

boolean check Pair (au, N, K)  $\S$ for (i = 0 to N-1)  $\S$ for (j = i+1 to N-1)  $\S$ if  $[arr[i] + arr[j] = = K) \S$ return true  $\S$ return false;

Htable-		
20	[i+l, N-	iferations
0	[1,10-1]	N-1
1	[2,11.	N-2
2	[3,N-1]	N-3
3	[4,N-1]	N- 4
1	1	1
N-1 (0	, , ,	<b>,</b>
10 1 10	ا (۱۰ سار	0

# total no. of iteration = 
$$(N-1)+(N-2)+(N-3)+---3+2+1+0$$
  
=  $N(N-1)$ 

$$\begin{array}{c} T. \left( \rightarrow 0 \left( N^2 \right) \right) \\ S. \left( \rightarrow 0 \left( 1 \right) \right) \end{array}$$

(1) Criven an array: Reverse entire array. [s.c.-0(1)]
Note - array itself should change.

aun [8]: {-1 4 7 6 -2 17 8 10 }

arr[7 - \$10, 8, 17, -2, 6, 7, 4, -13

= 2 1 6 2 5 4 3 We have to break.

aunt]: 10 9 8 2 3 6 -1

We have to break bseudo-code.

Void reverse ( and, n) 
$$\S$$
 $i = 0$ ,  $j = n-1$ 

while (  $i = j$ )  $\S$ 
 $femp = and (i)$ 
 $and (i) = and (j)$ 
 $and (j) = femp$ 
 $i = i+1$ 
 $j = j-1$ 
 $3$ 

The objective of the series o

# Todo", try to solve this question using for loop.

Storting index ending index

(1) Given NI array elements and is it et.

Reverse array from si to ei. Note si == ei

aux [7 = 9 =

pseudo-code-

٤

reverse Array Part (am, N, si, ei)  $\xi$  i = si, j = eiwhile (i < j)  $\xi$ femp = am(i) swap arr(i) with am(i) = am(j) am(j) = temp

# no. of iteration = ei-si/2

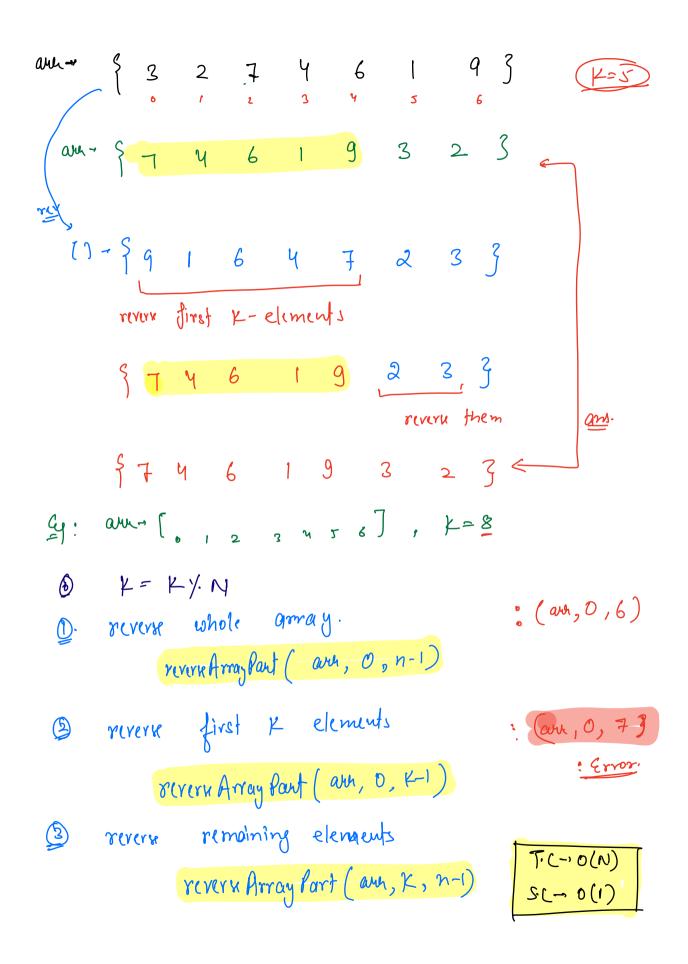
T.(->0(ei-si), s.(->0(1)

O(N)

Q1 Cliven N array elements, Rotate array from last to first by K times. \{ Google, Amazon }

arr 
$$[T]$$
:  $\begin{cases} 3 & -2 & 1 & 4 & 6 & 9 & 8 & 3 \\ 1 & 2 & 3 & 4 & 5 & 6 \end{cases}$ 

$$\underbrace{k=1}_{a} \cdot \begin{cases} 8 & 3 & -2 & 1 & 4 & 6 & 9 & 3 \\ 8 & 3 & -2 & 1 & 4 & 6 & 9 & 3 \\ 8 & 3 & -2 & 1 & 4 & 6 & 9 & 3 \\ 8 & 3 & -2 & 1 & 4 & 6 & 9 & 3 \\ 8 & 3 & -2 & 1 & 4 & 6 & 9 & 3 \\ 8 & 3 & 3 & -2 & 1 & 4 & 3 & 3 \\ 8 & 3 & 3 & 3 & 3 & 4 & 4 & 4 & 4 \\ 8 & 3 & 3 & 3 & 3 & 4 & 4 & 4 & 4 \\ 8 & 3 & 3 & 3 & 3 & 4 & 4 & 4 & 4 \\ 8 & 3 & 3 & 3 & 3 & 4 & 4 & 4 & 4 \\ 8 & 3 & 3 & 3 & 4 & 4 & 4 & 4 \\ 8 & 3 & 3 & 3 & 4 & 4 & 4 & 4 \\ 8 & 3 & 3 & 3 & 4 & 4 & 4 & 4 \\ 8 & 3 & 3 & 3 & 4 & 4 & 4 & 4 \\ 8 & 3 & 3 & 3 & 4 & 4 & 4 & 4 \\ 8 & 3 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 & 4 \\ 8 & 3 & 3 & 4 \\ 8 & 3 & 3 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 \\ 8 & 3 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 & 4 \\ 8 & 3 & 4 &$$



$$\frac{N - \text{Doub} ds}{\text{for}(i = 1 ; i < = N ; i + t)}$$

$$\frac{\text{for}(i = 1 ; i < = N ; i + t)}{\text{for}(j = 1 ; j < = N ; j + t)}$$

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