

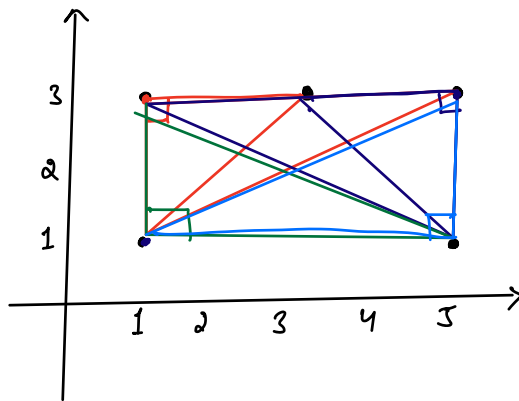
Agenda: -

→ Questions on Hashing

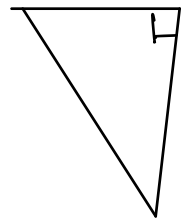
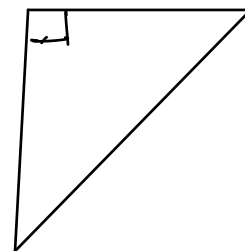
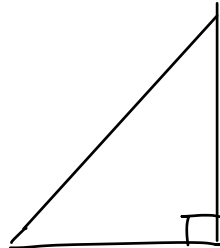
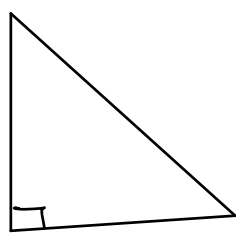
→ Rolling Hash

→ Rabin Karp Algorithm

Q) Given co-ordinates of N distinct points on a 2D plane. Count the no. of right angled triangle using the given set of points such that two small sides of a \triangle should be parallel to x-axis & y-axis. (P and B)



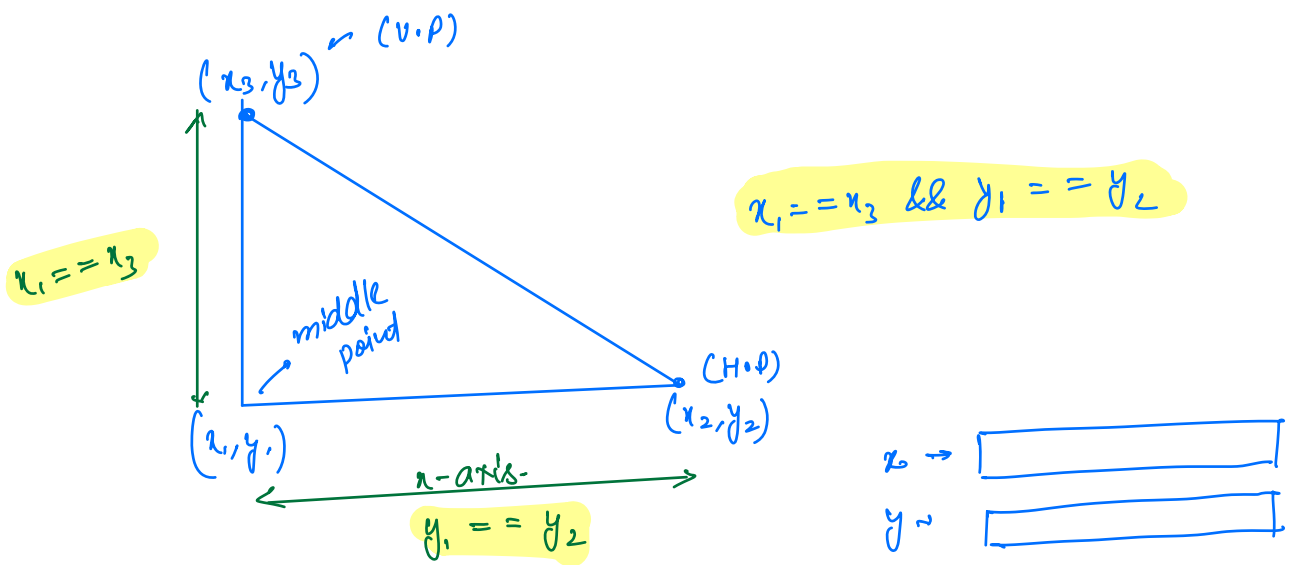
ans=6.



triangle \rightarrow triplet of co-ordinates is required.

idea-1 \rightarrow Consider all the triplets & check if they can form a right angled triangle or not.

two-sides \rightarrow $\begin{cases} \parallel x \text{ axis} \\ \parallel y \text{ axis} \end{cases}$



// All triplets.

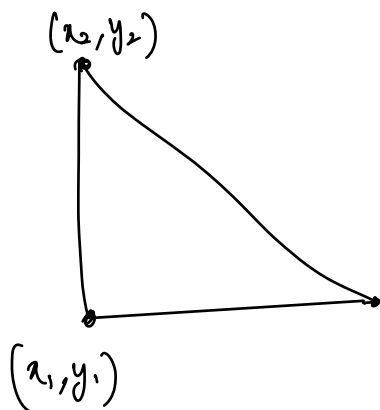
```

for( i=0; i < N; i++) { // i -> middle point
    for( j=0; j < N; j++) { // j -> vertical point
        if( i == j ) { continue; }
        for( k=0; k < N; k++) { // k -> horizontal point
            if( k == i || k == j ) { continue; }
            if( x[i] == x[j] && y[i] == y[k] ) {
                count++;
            }
        }
    }
}
return count;

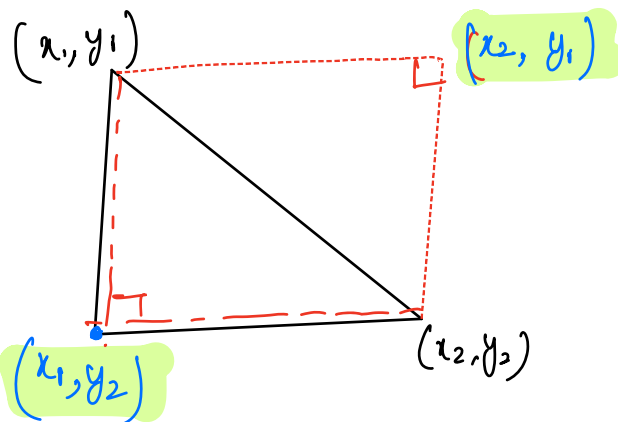
```

$T.C \rightarrow O(N^3)$
 $S.C \rightarrow O(1)$

$O(N^2)$?



? , y1



idea. \rightarrow using Hashset / Hashmaps.

// Add all the points in the hashset.

for(i = 0 ; i < N ; i++) {

for(j = i + 1 ; j < N ; j++) {

// handle edge cases.

check(x[i], y[j]) { count++ }

check(x[j], y[i]) { count++ }

T.C $\rightarrow O(1)$.

return count;

x \rightarrow

2	4	4	3	4
---	---	---	---	---

y \rightarrow

5	3	1	7	7
---	---	---	---	---

[2-5 , 2-3 , 4-1 , 3-7 , 4-7]

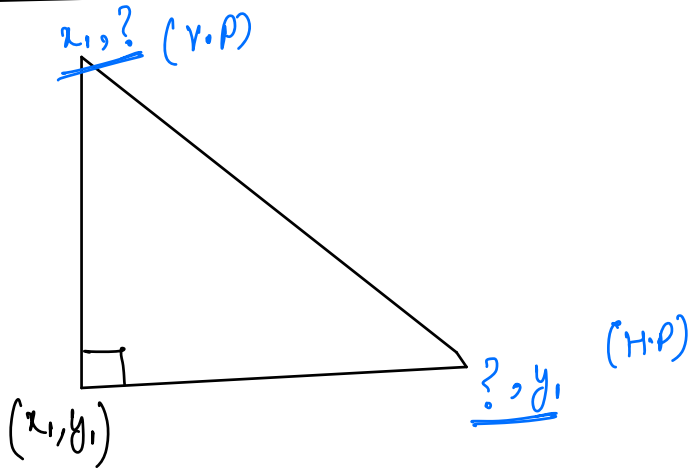
Key \rightarrow string \rightarrow x - y

T.C $\rightarrow O(N^2)$
S.C $\rightarrow O(N)$

$O(N)$

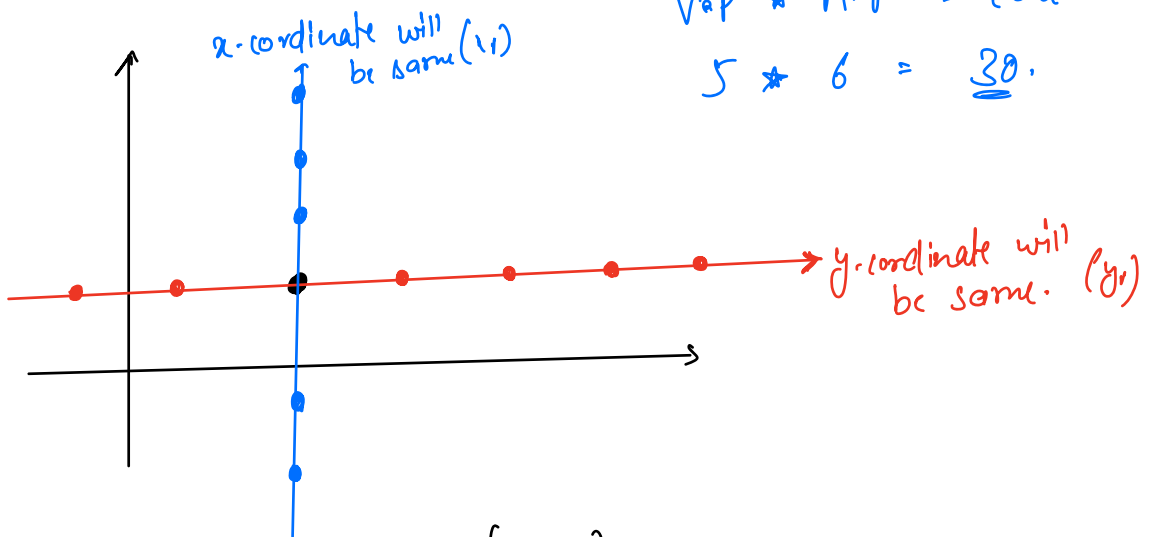
Syntax.

hs.contains(x[i] + " - " + y[j]) { count++ }



$$V.P * H.P = \text{count.}$$

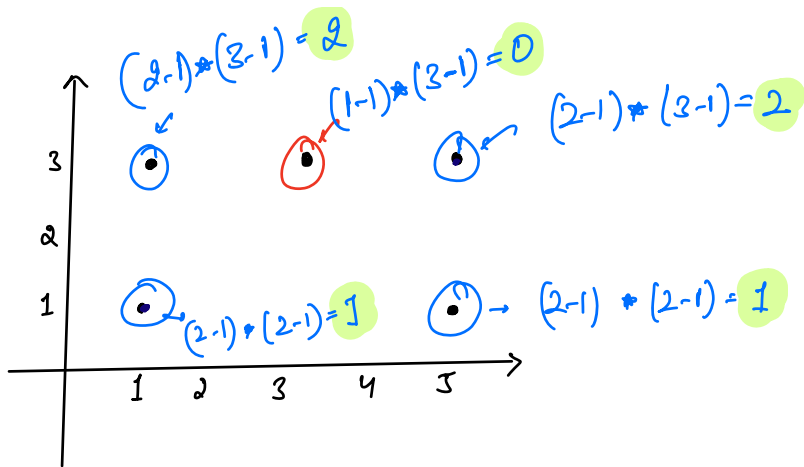
$$5 * 6 = \underline{30.}$$



$$\left\{ \begin{array}{l} \text{Count of co-ordinates} \\ \text{with x-coordinate as } x_1 \\ - 1 \end{array} \right\} * \left\{ \begin{array}{l} \text{count of co-ordinates} \\ \text{with y-coordinate as } y_1 \\ - 1 \end{array} \right\}$$

2 frequency Hashmaps.

\downarrow
 \downarrow
 x-coordinates y-coordinates



ans = 6.

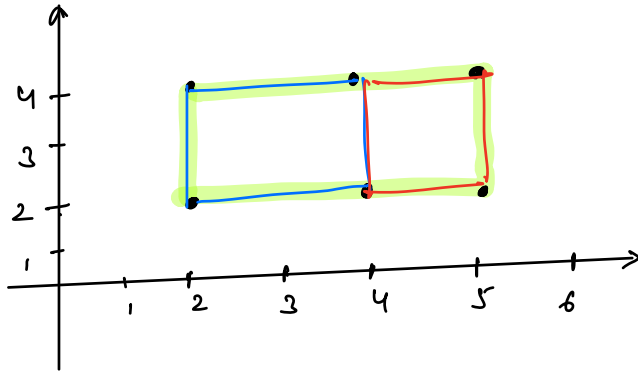
//pseudo-code.

- ① Take 2 hashmap, $hm1^{(x)}$ and $hm2^{(y)}$
- ② Insert all x-coordinates in $hm1$ & all y-coordinates in $hm2$.

③ for ($i = 0$; $i < N$; $i++$) {
 $c_1 = hm1[x(i)] - 1$
 $c_2 = hm2[y(i)] - 1$
 $ans += (c_1 * c_2);$
 }
 return ans;

T.C $\rightarrow O(N)$
 S.C $\rightarrow O(N)$

Q Given N points on a 2-D plane. Find count of rectangles we can form such that sides are parallel to x -axis and y -axis.

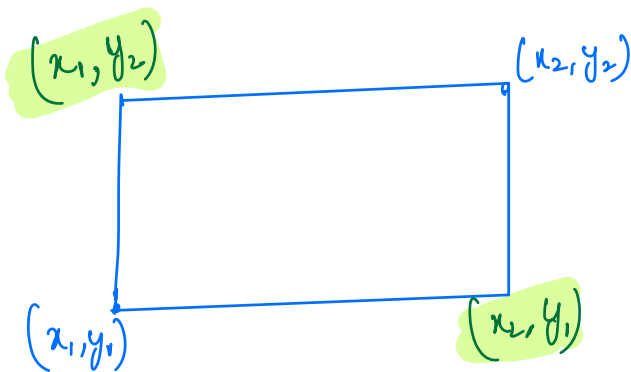


[ans=3.]

idea.1 - Consider all the quadruplets & check if they are satisfying the condition.

T.C $\rightarrow O(N^4)$

$\swarrow \searrow$
 $\binom{N}{4}$ $\binom{N}{4}$



$\Rightarrow \{ \# \text{ todo } \}$

Text : S : a b c b a x b a (N)
 pattern : T : c b a (K)

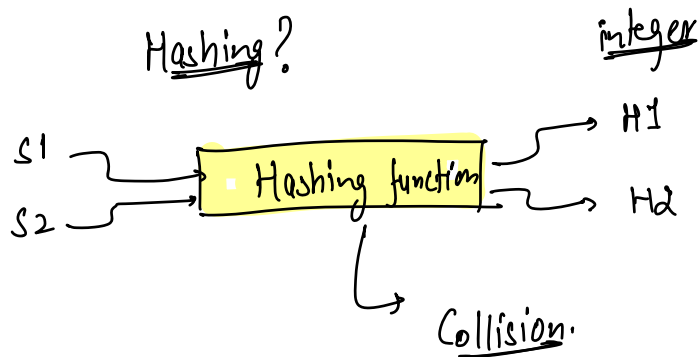


Check if pattern is present as a substring in text?

idea-1 Consider all the substring of size K (length of pattern)

T.C $\rightarrow O(N^2)$

{ KMP
 Z-Algo
 Rabin Karp }



$H1 = H2$

① Sum of ASCII values.

a b c
 ↓ ↓ ↓
 0 1 2 $\Rightarrow 3$

{ a b c
 a c b
 b a c
 b c a
 c a b
 c b a }

{ lot of collision will be there. }

②. ASCII values with indices.

$$\sum \text{ch}(i) * (i+1)$$

$$\begin{array}{ccc} a & b & c \\ 0 & 1 & 2 \end{array} \Rightarrow (97*1) + (98*2) + (99*3)$$

\Rightarrow collision.

③. Decimal no.

4 1 3 9 6 2

$$\begin{array}{c} \Downarrow \\ \textcircled{4} \times 10^{\textcircled{5}} + \textcircled{1} \times 10^{\textcircled{4}} + \textcircled{3} \times 10^{\textcircled{3}} + \textcircled{9} \times 10^{\textcircled{2}} + \textcircled{6} \times 10^{\textcircled{1}} + \textcircled{2} \times 10^{\textcircled{0}} \end{array}$$

\uparrow
[0-9]

b a d e f
0 1 2 3 4

$$[b*26^4 + a*26^3 + d*26^2 + e*26^1 + f*26^0]$$

$$\sum s(i) * p^{N-i-1}$$

p=26.

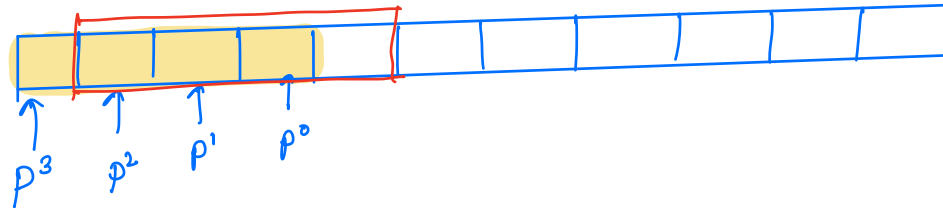
\Downarrow
very large.

{ overflow }

Polynomial
Rolling
Hash

$$\sum s[i] * p^{N-i-1} \% m$$

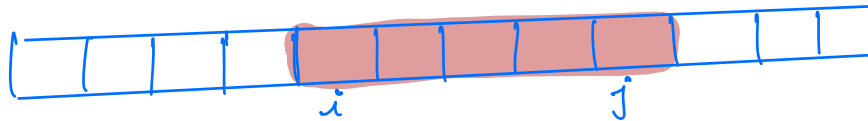
$$[m \approx 10^9]$$



$$H_1 = s[0] * p^3 + s[1] * p^2 + s[2] * p^1 + s[3] * p^0$$

$$H_2 = s[1] * p^3 + s[2] * p^2 + s[3] * p^1 + s[4] * p^0$$

$$H_2 = (H_1 - s[0] * p^3) * p + s[4] \Rightarrow T.C \rightarrow O(1)$$



$$H_{old} = i-1 \text{ to } j-1$$

$$H_{new} = i \text{ to } j = ?$$

$$H_{new} = (H_{old} - s[i-1] * p^{N-i+1}) * p + s[j]$$

$$\begin{matrix} T.C \rightarrow O(N) \\ S.C \rightarrow O(1) \end{matrix}$$

\Rightarrow Rabin Karp Algorithm
uses Rolling Hash.

\leftarrow (99% accurate)

{ Plagiarism checker }

Disadvantage? \Rightarrow collision

$$\sum \text{sig} \star p^{n-i-1} \% m \Rightarrow \begin{matrix} m = 10^9 \\ \underline{0 \text{ to } m-1} \end{matrix}$$

$$\begin{array}{ccccccc} & & H_1 & & & & \\ \text{---} & \text{---} & \text{---} & \text{---} & \text{---} & \text{---} & \end{array}$$

0 n-1

$$S_1 = H_1 = 0$$

$$S_2 = H_2 = 1/m$$

$$S_2 = H_2 = 2/m$$

$$S_4 = H_2 = 3/m$$

$$S_N = H_N = \frac{N-1}{n}$$

$$N = 10^5$$

$$M = 10^9$$

$$\frac{N}{n} = \frac{10^5}{10^9} = 10^{-4}$$

$$\left\{ \% \text{ of collision} = 0.0001 * 100 = 0.01\% \right\}$$

