

Today's content

- ① Given N elements . find first missing +ve no.
- ② Search in row-wise & column-wise sorted matrix.
- ③ Merge Intervals

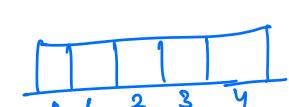
Q) find the first missing natural number.
[1, 2, 3, ...]

Eg₁: arr[5]: [3, -1, 1, 2, 7] : 4

Eg₂: arr[7]: [-9, 2, 6, 4, -8, 1, 3] : 5

Q.: arr[6]: [1, 2, 5, 6, 4, 3] : 7

Q.: arr[6]: [1, 0, -5, -6, 4, 2] : 3.

idea-1 arr[5] - 

1 →
2 →
3 →
4 →
5 →
6 → ans.

Obs. min: 1
max: N+1

idea-1.

for(i=1; i<=N; i++) {
 // check if ith. No. is present in
 // arr?

for(j=0; j<N; j++) {
 if(arr[j] == i) {
 }
 }
}

// if ith. no. is not present in arr()
// return i

return N+1;

T.C → O(N²)
S.L → O(1)

Idea-2. Using HashSet

1. // Insert all the elements in HashSet (hs)

```

for( i=1; i <= N; i++) {
    if(i is not present in hs) {
        return i;
    }
}
return N+1

```

T.C $\rightarrow O(N)$
S.C $\rightarrow O(N)$

Condition \rightarrow No extra space is allowed.

Idea-3 Sort & then iterate to get first missing natural number.

① arr[7]: [-9, 2, 6, 4, -8, 1, 3]

arr[7]: [-9, -8, 1, 2, 3, 4, 6] \rightarrow [ans = 5]

X X ✓ ✓ ✓ ✓ 5

② arr[8]: [-3, 1, 5, 8, 14, 2, 7, 3]

T.C $\rightarrow O(N \log N)$
S.C $\rightarrow O(1)$

After sorting
arr[8]: [-3, 1, 2, 3, 5, 7, 8, 14] \rightarrow [ans = 4]

X ✓ ✓ ✓ 4

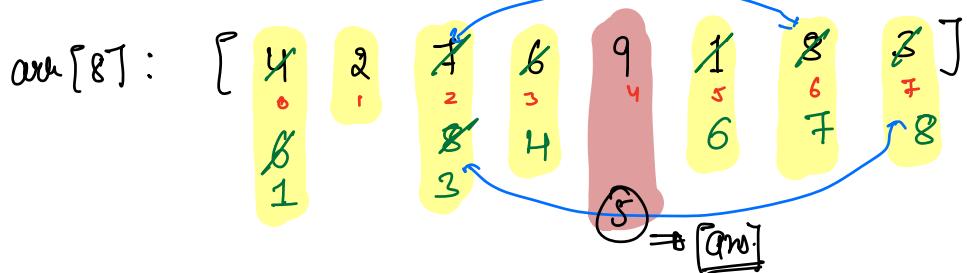
Ignore -ve number & duplicates

Implementation
(todo)

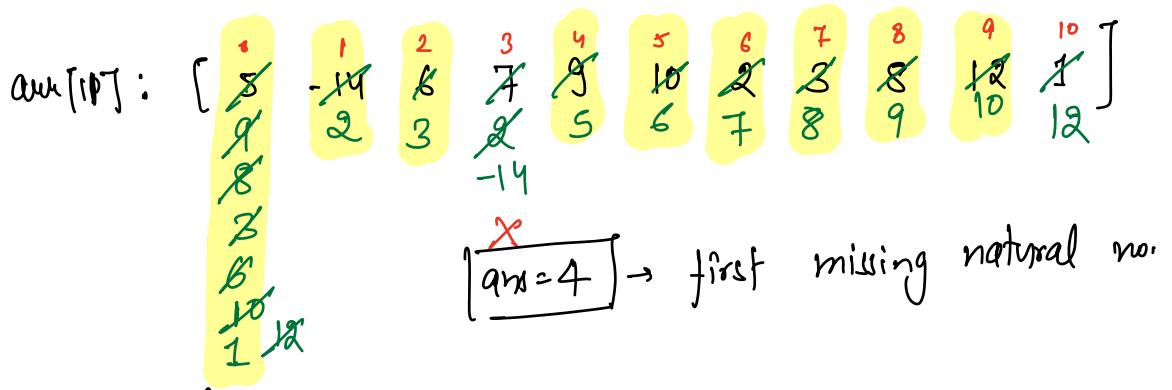
Optimization: $O(N \log N) \rightarrow O(\underline{N})$, S.C $\rightarrow O(1)$

idea-4 Keep the element at its correct position.

	<u>$N=5$</u>	<u>Generalize</u>	
$i=0$	$\text{val}=1$	<u>val</u>	<u>idx</u>
$i=1$	$\text{val}=2$	1	0
$i=2$	$\text{val}=3$	2	1
$i=3$	$\text{val}=4$	3	2
$i=4$	$\text{val}=5$	4	3
		⋮	⋮
		x	($i-1$)
		⋮	($N-1$)



- | <u>i</u> | <u>val</u> | |
|----------|--|--------------|
| <u>0</u> | <u>swap(arr[0], arr[3]) , swap(arr[0], arr[5])</u> | <u>stop.</u> |
| <u>1</u> | <u>2 (nothing)</u> | |
| <u>2</u> | <u>swap(arr[2], arr[6]) , swap(arr[2], arr[7])</u> | <u>stop.</u> |
| <u>3</u> | <u>4 ✓</u> | |
| <u>4</u> | <u>9 [ignore]</u> | |
| <u>5</u> | <u>✓</u> | |
| <u>6</u> | <u>✓</u> | |
| <u>7</u> | <u>✓</u> | |



i : val
 $\underline{\text{swap}}(\text{arr}[0], \text{arr}[1])$, $\underline{\text{swap}}(\text{arr}[0], \text{arr}[8])$, $\underline{\text{swap}}(\text{arr}[0], \text{arr}[7])$
 $\underline{\text{swap}}(\text{arr}[0], \text{arr}[2])$, $\underline{\text{swap}}(\text{arr}[0], \text{arr}[5])$, $\underline{\text{swap}}(\text{arr}[0], \text{arr}[9])$

(stop)

<u>1</u> :	<u>ignore</u>
<u>2</u> :	<u>3 ✓</u>
<u>3</u> :	<u>swap(arr[3], arr[6])</u> , <u>swap(arr[3], arr[1])</u> <u>stop</u> .
<u>4</u> :	<u>✓</u>
<u>5</u> :	<u>✓</u>
<u>6</u> :	<u>✓</u>
<u>7</u> :	<u>✓</u>
<u>8</u> :	<u>✓</u>
<u>9</u> :	<u>✓</u>
<u>10</u> :	<u>swap(arr[10], arr[0])</u> <u>stop</u> .

observation: In single swap, at least one element will be shifted at its correct position.
At max, how many swaps are possible? $\Rightarrow \underline{N}$ swaps.

pseudo-code

```
for( i=0; i < N; i++) {
```

while($\underbrace{\text{arr}[i] > 0 \text{ && arr}[i] \leq N}_{\text{relevant data}} \text{ && } \underbrace{\text{arr}[i] \neq i+1}_{\text{not satisfied.}}) {$

val = arr[i]

if (arr[i] == arr[val-1]) { break }

swap arr[i] with arr[val-1]

T.C $\rightarrow O(N)$
S.C $\rightarrow O(1)$

}

// iterate again & then find first missing Natural no.

```
for ( i= 0 ; i < N; i++) {
```

if (arr[i] != i+1) {

return (i+1);

}

return N+1

i	swap.count
0	s_0^+
1	s_2^+
2	s_3^+
1	\vdots
N-1	s_n^+

$$= s_1 + s_2 + s_3 + \dots + s_n$$

iterations. = $\underline{\Omega}N$

Eg \rightarrow arr[5] : [~~1~~
~~2~~
~~3~~
~~4~~
~~5~~] ! [~~1~~
~~2~~
~~3~~
~~4~~] 2]

i
val:
① swap(arr[0] and arr[3]), swap(arr[0] and arr[2])
swap(arr[0] and arr[2]), swap(arr[0] & arr[2]), --
- - - T.L.E. (infinite loop)

if (arr[i] == arr[val-1])
we need to break
otherwise, infinite loop will be there.

Q Given a matrix where every row & column are sorted.
find an element \underline{k} .

	0	1	2	3	4	5
0	-1	2	4	5	9	11
1	1	4	7	8	10	14
2	3	7	9	10	12	18
3	6	10	12	14	16	20
4	9	13	16	19	22	24
5	11	15	19	21	24	27
6	14	20	25	29	31	39
7	18	24	29	32	34	42

$\underline{k=15}$

idea1. traverse the whole matrix &
check if $\text{arr}[i][j] == k$ or not
 $T.C \rightarrow O(N \cdot M)$, $S.C \rightarrow O(1)$

Binary Search → search for any
element in sorted array.
 $\hookrightarrow \log N$.

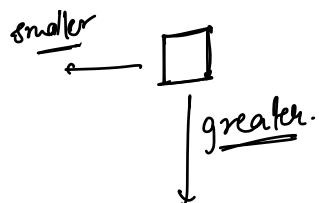
idea2. In every row, apply binary search.
 $T.C \rightarrow O(N \log M)$

idea3. In every col, apply binary search
 $T.C \rightarrow O(M \log N)$

	0	1	2	3	4	5
0	-1	2	4	5	9	11
1	1	4	7	8	10	14
2	3	7	9	10	12	18
3	6	10	12	14	16	20
4	9	13	16	19	22	24
5	11	15	19	21	24	27
6	14	20	25	29	31	39
7	18	24	29	32	34	42

$\underline{k=15}$

\underline{k}



PF26

	0	1	2	3	4	5
0	1	2	4	5	9	11
1	1	4	7	8	10	14
2	3	7	9	10	12	18
3	6	10	12	14	16	20
4	9	13	16	19	22	24
5	11	15	19	21	24	27
6	14	20	25	29	31	39
7	18	24	29	32	34	42

[Todo # try with bottom-left corner.]

pseudo-code:

$i = 0, j = M-1$

while ($i < N \text{ and } j \geq 0$) {

 if ($\text{arr}[i][j] < k$) {

$i++$

 } else if ($\text{arr}[i][j] > k$) {

$j--$

 } else {

 return true;

}

observation
At every step we are either skipping a row or a column.

no. of steps = $N + M$

$T.C \rightarrow O(N+m)$
$S.C \rightarrow O(1)$

Merge Intervals ↳ range / boundary / $[a, b]$.

	I_1	I_2
	$[2, 6]$	$[3, 7]$
	$[2, 8]$	$[4, 6]$
	$[3, 7]$	$[4, 10]$
	$[3, 6]$	$[6, 10]$
	$[2, 5]$	$[8, 10]$
	$[5, 8]$	$[1, 3]$

Merged interval

$[2, 7]$

$[2, 8]$

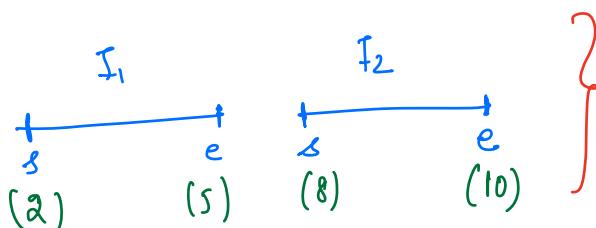
$[3, 10]$

$[3, 10]$

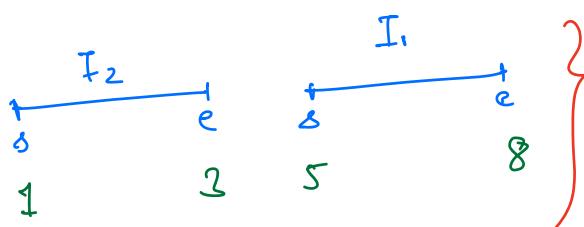
{No overlapping}

{No overlapping}

Non-overlapping.

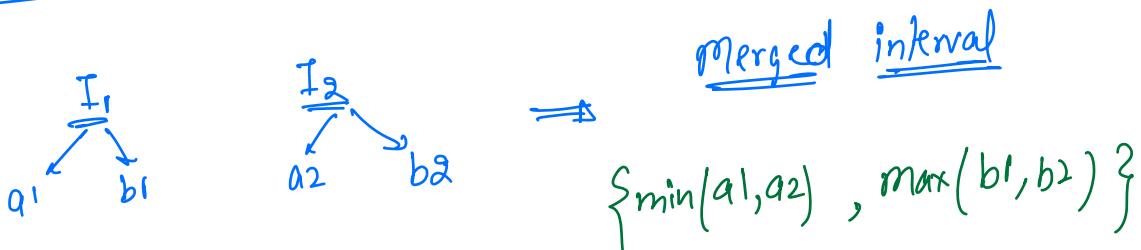


$$I_2.e > I_1.e$$



$$I_2.e < I_1.e$$

if I_1 & I_2 are overlapping. (Generalisation for merged interval)



Given N non-overlapping intervals & they are sorted based on start-time. You are given a new interval.
Merge all the intervals.

Eg. $N=8$

$[1, 3]$

$[4, 7]$

$$[10, 14] + [10, 22] \Rightarrow [10, 22]$$

$$[16, 19] + [10, 22] \Rightarrow [10, 22]$$

$$[21, 24] + [10, 22] \Rightarrow [10, 24]$$

$$[27, 30] \quad [10, 24]$$

$[32, 35]$

$[38, 41]$

New interval

New non-overlapping intervals

$[1, 3]$

$[4, 7]$

$[10, 24]$

$[27, 30]$

$[32, 35]$

$[38, 41]$

final
req.
ans.

N=5

[1, 5]

[8, 10]

$$[11, 14] + [12, 22] \Rightarrow [11, 22]$$

$$[15, 20] + [11, 22] \Rightarrow [11, 22]$$

$$[20, 24] + [11, 22] \Rightarrow [11, 24]$$

Interval

New non-overlapping intervals

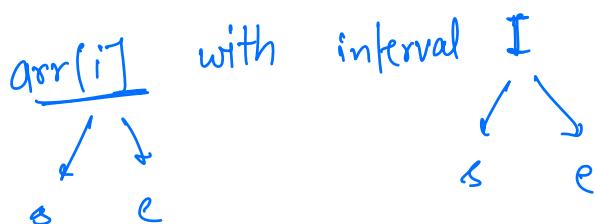
[1, 5]

[8, 10]

[11, 24]



final
req.
ans.



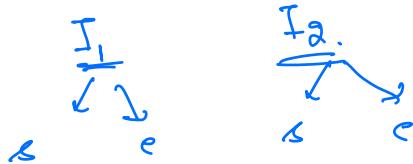
Pseudo-code

```
Intervals[] mergeIntervals( arr[], Interval I ) {
```

```
    for( i=0 ; i < N ; i++ ) {  
        // arr[i] → interval in the array.  
        if( arr[i].e < I.e ) {  
            // they are non-overlapping, update  
            ans.insert( arr[i] )  
        }  
        else if( arr[i].s > I.e ) {  
            // they are non-overlapping  
            ans.insert( I );  
            for( j=i ; j < N ; j++ ) {  
                ans.insert( arr[j] );  
            }  
            return ans  
        }  
        else {  
            // overlapping  
            // merge arr[i] & Ith  
            I.s = min( arr[i].s, I.s )  
            I.e = max( arr[i].e, I.e )  
        }  
    }  
    ans.insert( I );  
    return ans;
```

T.C → O(N)
S.C → O(1)

Doubts :



$$\left[\begin{array}{cc} \frac{I_1}{I_2} & \frac{F_2}{F_1} \\ \frac{I_2}{I_1} & \frac{F_1}{F_2} \end{array} \right]$$

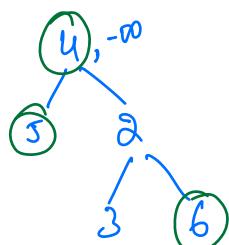
non-overlapping. ① $I_1 \cdot s < I_2 \cdot s$.

check if $(I_1 \cdot e < I_2 \cdot s)$

② $I_2 \cdot s < I_1 \cdot s$

check if $(I_2 \cdot e < I_1 \cdot s)$

Count of nodes with more value than all its ancestors



[Keep an track of maximum element]