Trie.

① Spell-checker

playgriund.
loyal
life
lyfe.

word → l.

T.C → O(N*l)

↓

HashMap / Hashset.

② Auto-complete.

play
place
plate
playground.

Trie. (prefix-tree)

→ hierarchical data structure.

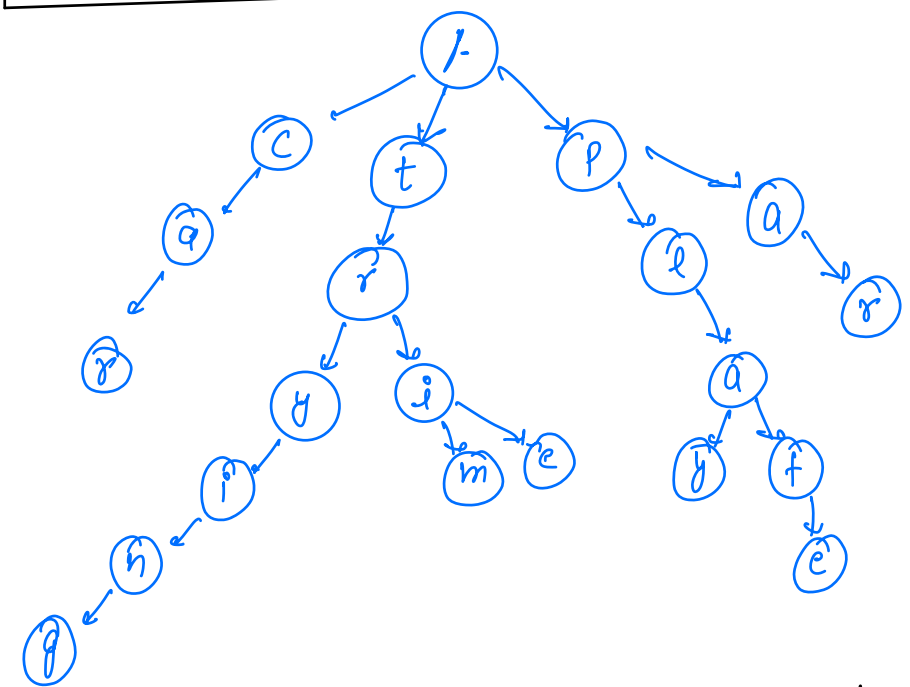
→ N-children (any) tree

→ It is used for information retrieval.

→ It is a data-structure which stores the information from top to down.

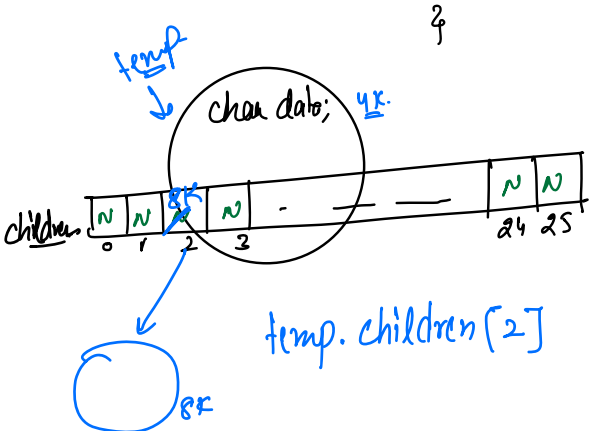
dict →

try	trim	trie	play	trying
plate	car	par	~~trimmer~~	pla



```
class Node {
    char data;
    Node children[26]
}
```

```
class Node {
    char data;
    Node a;
    Node b
    Node c;
        )
        )
    Node z;
}
```

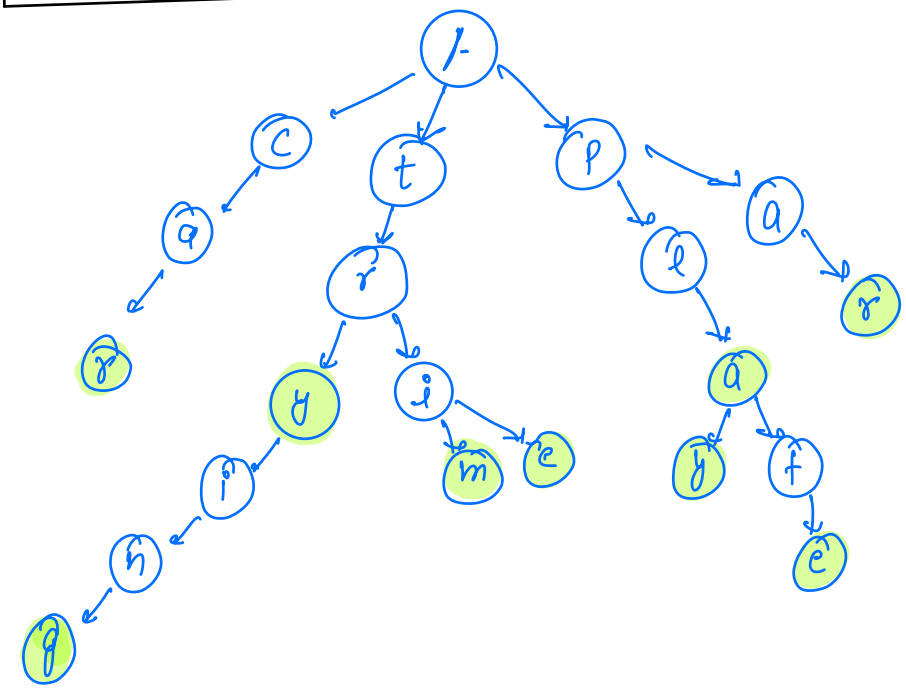


temp.children[2] = new Node('c');

a → '0'
b → '1'
c → '2'
'd' → '3'
'z' → '25'

dict →

try	trim	trie	play	trying
plate	car	par	~~trimmer~~	pla



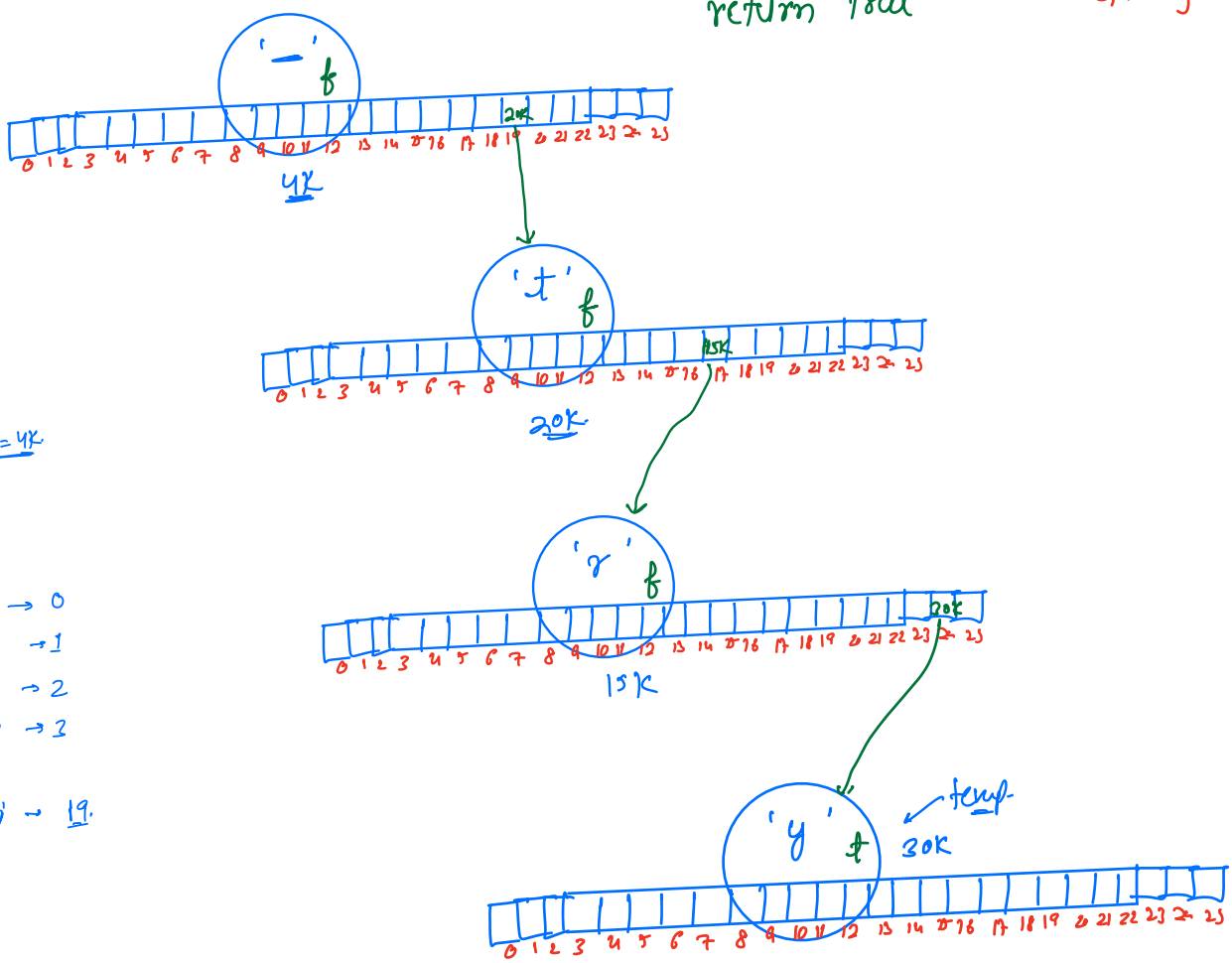
search (trie)
search (tri)
search (try)
search (tr)

There should be a marker whether the current node is denoting the end of word or not.

```
class Node {
    char data;
    Node children[26]
    boolean isEnd;
}
```

traverse on the string and for every character, we need to traverse tries.

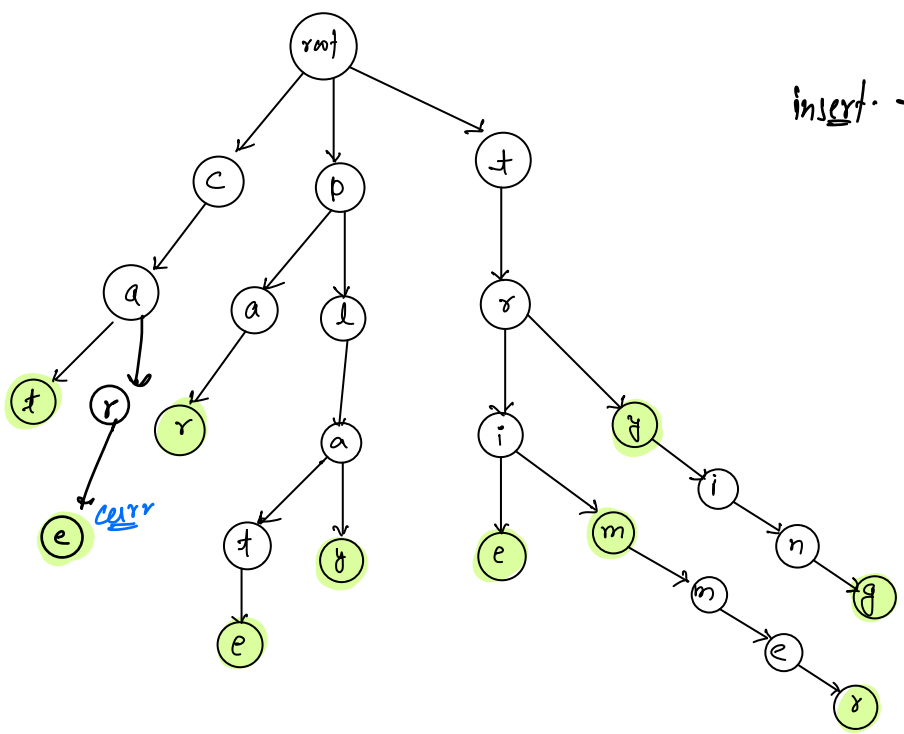
- if any character of string is not found in trie.
  - return false
- all characters are there
  - last char of the string.
    - ch.isEnd == true
    - return true
  - last char of the string
    - ch.isEnd == false
    - return false;



insert.

```
void insert( Node root, String word ){
    Node curr = root;
    for( i=0; i< word.length(); i++){
        idx = word[i] - 'a';
        if( curr.children[idx] == NULL){
            curr.children[idx] = new Node(word[i]);
        }
        curr = curr.children[idx]
    }
    curr.isEnd = true;
}
```

{T.C → O(L)
S.C → O(L)}



insert. → care.

Search.

```
boolean search ( Node root , String word ){
    Node curr = root ;
    for( i=0; i< word.length(); i++){
        idx = word[i] - 'a' ;
        if ( curr.children[idx] == NULL){
            return false;
        }
        curr = curr.children[idx]
    }
    return curr.isEnd ;
}
```

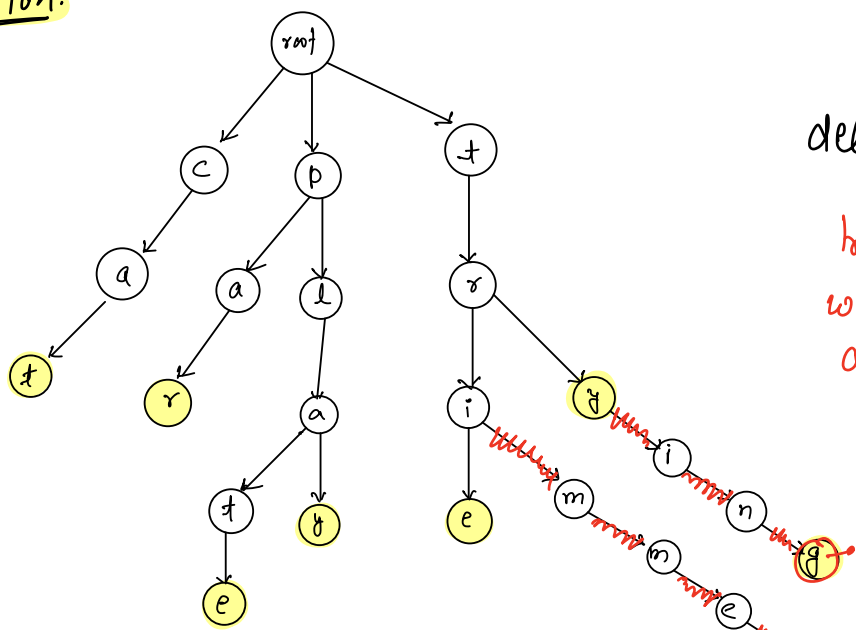
T.C → O(l)

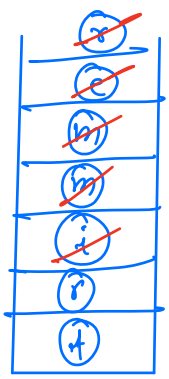
deletion.

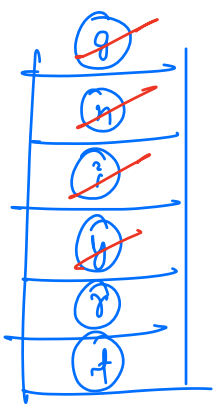
delete (trying) ✓

delete (trimmer) ✓

traverse for current word & mark isEnd as false for the last node.







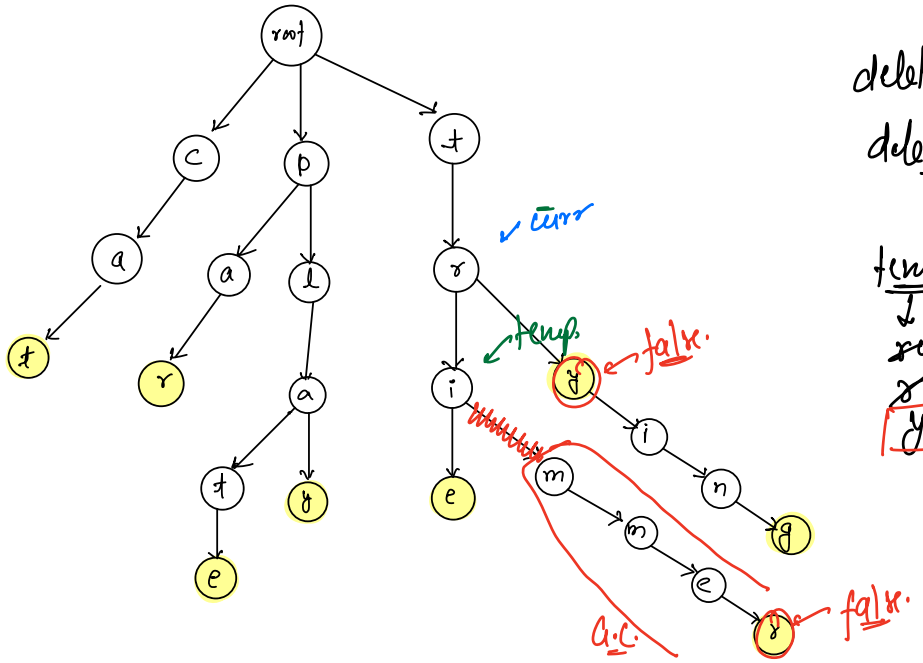
T.C → O(l)
S.C → O(l)

nodes that can't be deleted.

→ nodes where isEnd is marked as true.

→ nodes which are having > 1 children.

idea → Find the last node that can't be deleted.



delete(trimmer)
delete (try)

temp.
~~root~~
r
y

nextchar
t
y

Last node that can't be deleted.

~~root~~
r
i

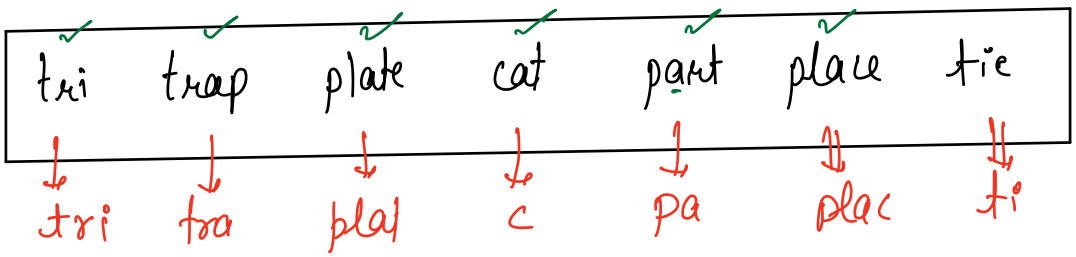
temp.children [m - a] = NULL;

pseudo-code

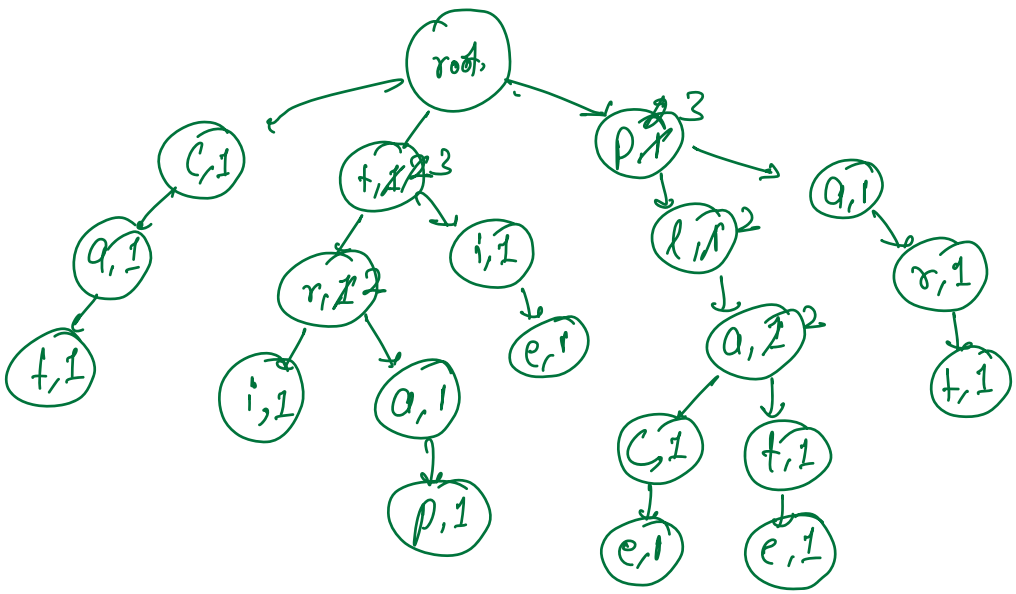
```
void DeleteWord(Node root, String word){
    Node curr = root, temp = null, nextChar = '-'
    for(i=0; i < word.length()-1; i++){
        int count = 0;
        for(i = 0; i < 25; i++){
            if(curr.children[i] != NULL){
                count++
            }
        }
        if(count > 1 || curr.isEnd == true){
            // we can't delete this node
            temp = curr, nextChar = word[i]
        }
        idx = word[i] - 'a';
        curr = curr.children[idx];
    }
    idx = word[N-1] - 'a';          // 'z' - 'a'
    curr = curr.children[idx];
    curr.isEnd = false;
    int count = 0;
    for(i = 0; i < 25; i++){
        if(curr.children[i] != NULL){
            count++
        }
    }
    if(count != 0) { // do nothing }
    else{
        temp.children[nextChar - 'a'] = NULL;
    }
}
```

Q.1 Find shortest unique prefix to represent each word.

Note: Assume that no word is prefix of another word.
In other words, the representation is always possible.



✗ trip, tole.

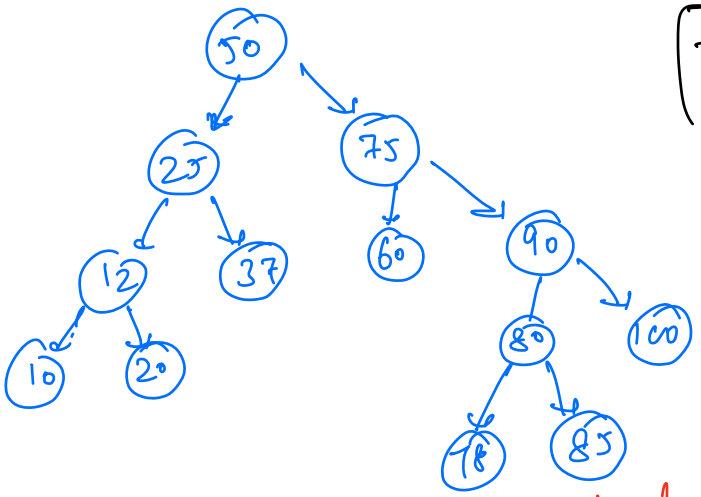


→ Small variation { including frequencies of characters }

```
Node {
    char data;
    Node children[26];
    int freq;
}
```

$|\text{height of l.st} - \text{height of r.s.t}| \leq 1 \Rightarrow$

[height balanced binary tree



→ false.

boolean isHb = true;

```
int height (Node root){
    if (root == NULL) {return -1}
    lh = height (root.left);
    rh = height (root.right);
    if ( |lh-rh| > 1 ) { isHb = false}
    return max (lh, rh)+1
}
```