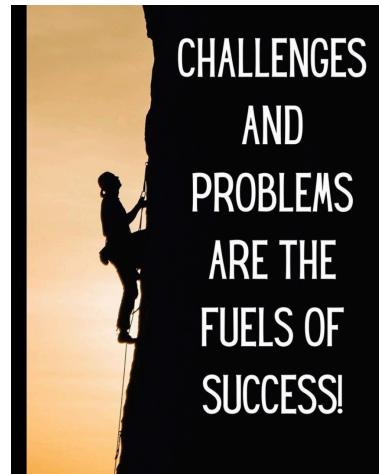


## Today's Quote



## Today's content

- understand sorting
- few problems on sorting
- 1 sorting algo ↪
- Comparator ↪

Sorting  $\Rightarrow$  Arranging of data in inc|dec order based on a parameter

Eg:  $\{2, 4, 7, 11, 15, 20\}$   $\rightarrow$  inc order, par $\rightarrow$  value

Eg:  $\{15, 10, 9, 6, -2\}$   $\rightarrow$  dec. order, par $\rightarrow$  value

Eg:  $\{1, \underline{2}, \underline{3}, \underline{7}, \underline{4}, \underline{9}, \underline{6}\}$   $\rightarrow$  sorted in inc. order  
based on factors count.

factors: 1 2 2 2 3 3 4

Inbuilt library functions  $\rightarrow$

→ sort(), in every language.

→ How? logic?  $\Rightarrow \{\text{In advance module}\}$

T.C  $\Rightarrow N \log N$ , if N elements are present in the array.

## Q.) Elements Removal

Given N elements, at every step remove an array element.

[Cost to remove element = Sum of array elements present in array]

Find min cost to remove all elements.

Note: first add the cost of removal & then remove it.

Eg:  $\text{arr}[3] = \{ \underline{\underline{2}}, \underline{1}, \underline{\underline{4}} \}$

remove 2 :      cost  
                    7  
  
remove 1 :      5  
  
remove 4 :      4  
  
total cost = 16

remove 4 :      cost  
                    7  
  
remove 2 :      3  
  
remove 1 :      1  
  
total cost = 11

Eg:  $\text{arr}[4] : \{ \underline{3}, \underline{6}, \underline{2}, \underline{\underline{4}} \}$

cost:  
  
remove 6 :      15  
  
remove 4 :      9  
  
remove 3 :      5  
  
remove 2 :      2  
  
total cost = 31

observation  $\Rightarrow$  Delete elements in decreasing order to get min cost ?

arr = { a b c d }  
cost:

remove a  $\rightarrow$  a + b + c + d

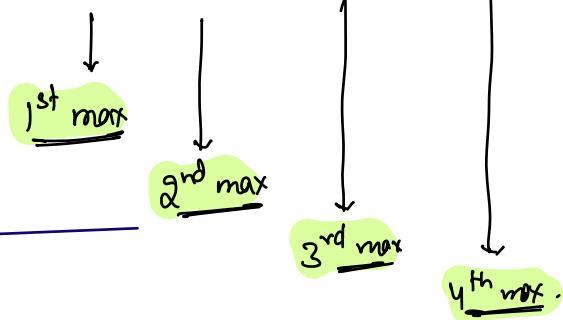
remove b  $\rightarrow$  b + c + d

remove c  $\rightarrow$  c + d

remove d  $\rightarrow$  d

Total cost:

$$a + \cancel{2}b + 3c + 4d$$



Pseudo-code

```
int minCost ( arr, N ) {
    sort ( arr, desc ) // sort the array in dec. order.
    cost = 0
    for ( i = 0 ; i < N ; i++ ) {
        // How many times arr[i] will be present in total cost
        cost += arr[i] * (i+1)
    }
    return cost;
}
```

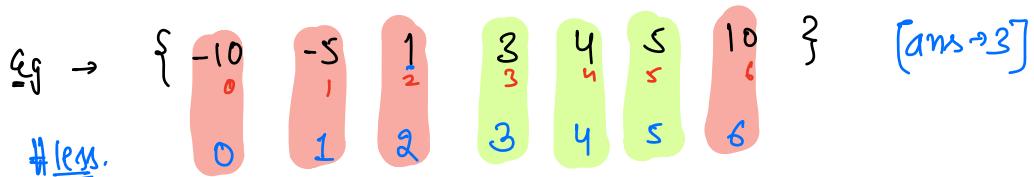
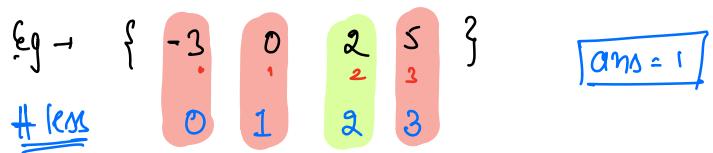
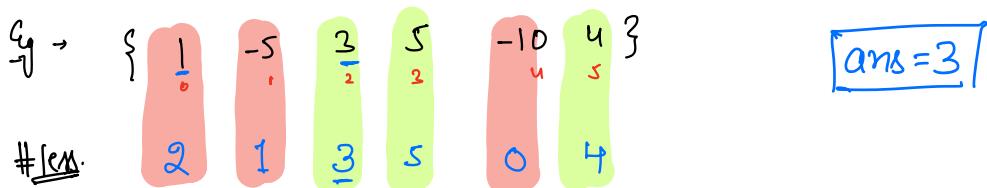
T.C $\rightarrow O(N \log N)$
S.C $\rightarrow O(1)$

## Q) Noble Integers {distinct data?}

Given N array element, calculate no. of noble integers.

An element ele in arr[] is said to be noble if

{No. of elements < ele = ele itself?}



observation:  
-ve integer can't  
be noble integers

Brute force solution: For every element, find count of smaller elements & then check if [count == current element].

```
int nobleIntegers (arr, N) {
    ans = 0
    for( i=0 ; i < N ; i++ ) {
        count = 0
        for( j=0 ; j < N ; j++ ) {
            if( arr[j] < arr[i] ) { count += 1 }
        }
        if( count == arr[i] ) { ans += 1 }
    }
    return ans;
}
```

T.C  $\rightarrow O(N^2)$   
S.C  $\rightarrow O(1)$

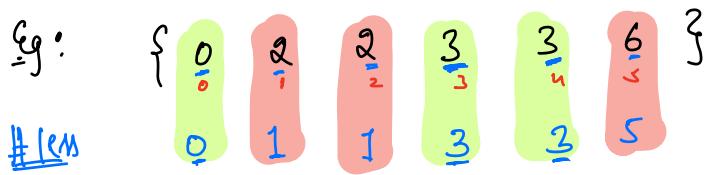
## Optimisation

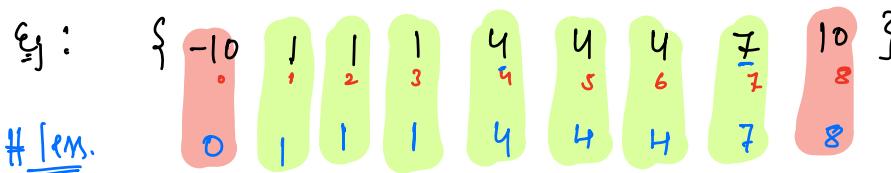
# sort the array in increasing order

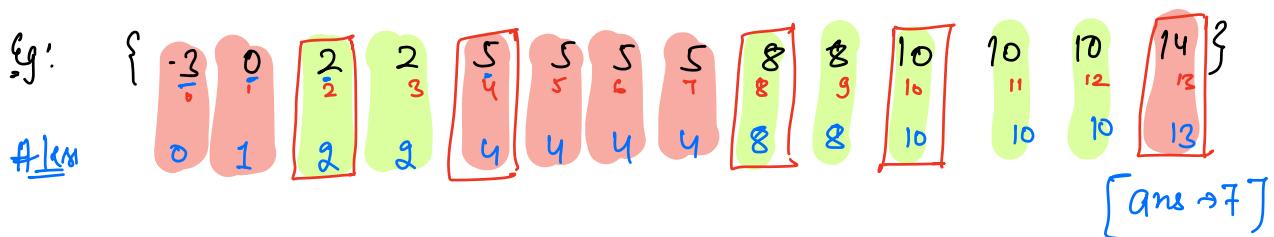
```
int nobleIntegers ( arr, N ) {  
    sort ( arr, asc. )  
    ans = 0;  
    for ( i = 0 ; i < N : i++ ) {  
        if ( arr[i] == i ) { ans += 1 }  
    }  
    return ans;  
}
```

T.C  $\rightarrow O(N \log N)$   
S.C  $\rightarrow O(1)$

Q: Noble integers : {data can repeat}

Eg: { } [ans → 3]

Eg: { } [ans → 7]

Eg: { } [ans → 7]

Observations:

- if element is repeating, ( $\text{arr}[i] == \text{arr}[i-1]$ )  
count of smaller elements is going to remain same.
- if element is coming for the first time ( $\text{arr}[i] != \text{arr}[i-1]$ )  
count of smaller elements =  $i$

→ idea 1 of previous question ✓

→ idea 2 [we need to make some changes]

## pseudo code

```

int nobleIntegers ( arr, N ) {
    sort ( arr, asc. )
    ans = 0 , count = 0
    if ( arr[0] == 0 ) { ans += 1 }   ⇒ Edge case.
    for ( i = 1 ; i < N ; i++ ) {
        // for every arr[i], we need to find count of smaller
        // elements
        if ( arr[i] != arr[i-1] ) {
            count = i
        }
        else {
            // that element is repeating
        }
        if ( arr[i] == count ) { ans += 1 }
    }
    return ans;
}

```

T.C →  $O(N \log N)$   
S.C →  $O(1)$

$$\text{arr}[i] == \text{arr}[i-1]$$

Ex:

$\left\{ \frac{2}{0}, \frac{3}{1}, \frac{3}{2}, \frac{3}{3}, \frac{4}{4}, \frac{4}{5}, \frac{6}{6} \right\}$

$$\text{count} = 1, \text{ans} = 1, 2, 3$$

- Give enough time while reading the qn
- take some exs.

Sorting algo

arr → { 3 6 8 2 8 4 1 }

Bubble Sort

iteration 1 : { 3 6 8 2 4 8 1 } : 8 at correct position.

iteration 2 : { 2 3 4 6 8 } : 6, 8 at correct position.

iteration 3 : { 2 3 4 6 8 } : 4, 6, 8 at correct position  
[whole array is sorted].

observation

After 1 iteration → last 1 element at correct position

After 2 iteration → last 2 elements " " "

After 3 " → last 3 elements " " "

After N-1 iteration → last N-1 elements at correct position

↓  
last N elements at correct position

idea & pseudo-code

We need N-1 iterations →  
in every iteration, traverse from left to right  
and check the adjacent elements.

```

bubbleSort ( arr, N ) {
    for( i = 1 ; i < N ; i++ ) {
        for ( j = 0 ; j < N-1 ; j++ ) {
            if ( arr[j] > arr[j+1] ) {
                // swap arr[j] with arr[j+1]
            }
        }
    }
}
  
```

T.C → O(N<sup>2</sup>)  
S.C → O(1)

// Decreasing order.

```
- bubbleSort ( arr, N ) {  
    for( i = 1 ; i < N ; i++ ) {  
        for ( j = 0 ; j < N-1 ; j++ ) {  
            if ( arr[j] < arr[j+1] ) {  
                // swap arr[j] with arr[j+1]  
            }  
        }  
    }  
}
```

has power to control  
the inc/dec. order.

idea -

```
- bubbleSort ( arr, N ) {  
    for( i = 1 ; i < N ; i++ ) {  
        for ( j = 0 ; j < N-1 ; j++ ) {  
            if ( comp(arr[j],arr[j+1]) ) {  
                // swap arr[j] with arr[j+1]  
            }  
        }  
    }  
}
```

bool can comp(int a, int b) {  
 // Just implement this  
 // function to get your  
 // own sorted order  
 // if a is going to come  
 // before b or vice-versa.  
}

## Comparator in Java / Python

```
Comparator customComparator = (Integer a, Integer b) {
```

{

- if you want a to come before b : { return -1 }
- if you want a & b same : { return 0 }
- if you want b to come before a : { return 1 }

}

```
def compare (a, b) :
```

{

- if you want a to come before b : { return -1 }
- if you want a & b same : { return 0 }
- if you want b to come before a : { return 1 }

}

```
ar.sort (key = cmp_to_key(compare))
```

⇒ {check the syntax in your language.}

Q1 Given N elements. Sort them in increasing order of their no. of factors.

Note: If two elements have same no. of factors, element with less value should come first.  
→ No extra space is allowed.

Ex- arr: { 9, 3, 4, 8, 16, 37, 6, 13, 15 }  
# factors : { 3, 2, 3, 4, 5, 2, 4, 2, 4 }

sorted. order → { 3, 13, 37, 4, 9, 6, 8, 15, 16 }

Idea:

→ sort(arr, comp)

↙

{ your own  
sorted order }

{ T.C → in adv. }

```
int comp ( int a, int b ) {
    int fa = factor(a)
    int fb = factor(b)
    if ( fa < fb ) {
        // a should come first
        return -1
    }
    else if ( fa == fb && a < b ) {
        // a should come first
        return -1
    }
    else {
        // b should come first
        return 1
    }
}
```

Doubts :

arr- { ~~3~~ ~~6~~ ~~4~~ ~~2~~ }

Cost of removal = sum of all  
the elements  
before removing

Cost:

remove 6 :  $6 + 4 + 3 + 2$

remove 4 :  $4 + 3 + 2$

remove 3 :  $3 + 2$

$[6 \cdot 1 + 4 \cdot 2 + 3 \cdot 3 + 2 \cdot 4]$

remove 2 : 2

arr  $\rightarrow [6 \downarrow \quad 4 \quad 3 \quad 2 \quad 2] \downarrow$   
 $\downarrow \quad \downarrow \quad \downarrow \quad \downarrow$   
 $6 \cdot 1 \quad 4 \cdot 2 \quad 3 \cdot 2 \quad 2 \cdot 4$

$[ans += arr[i] * (i+1)]$

✓ Arrays.

✓ Modular Arithmetic

B.M.

- watch recording →  $1.25x$   
→  $1.5x$
- try to attempt H.W.
- 25-30 minutes to every question
- connect with T.A.