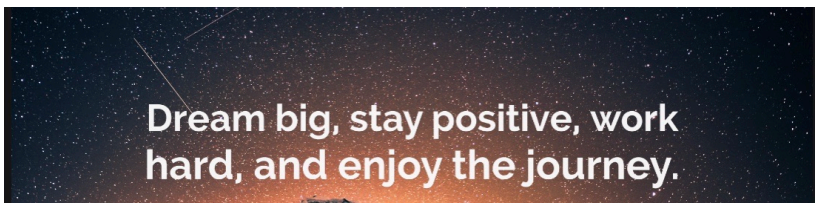
Today's Quote



Today's content

- Intro
- flip
- sort ch[]
- Reverse string
- Longest palindromic substring

String. → array of characters
→ sequence of characters
→ bunch of characters → {a b c} ≠ {b a c}

* Not same.
order is also important.

characters → ASCII value.

'A' - 65 ⇄ 'a' - 97 (+32 / -32)
'B' - 66 ⇄ 'b' - 98 (+32 / -32)
'C' - 67 'c' - 99
'D' - 68 'd' - 100
|
'Z' - 90 'z' - 122

'0' - 48
'1' - 49
'2' - 5

$[32 \rightarrow 2^5]$

'g' - 57

'10' → It is not a single character

char ch = 'q'
2 bytes.
ASCII = 57

ch = 'g'
ch = ch + 8
print(ch) Ascii → 65
'A'

string. → array of characters.

String s = "a b c d"

s → | a | b | c | d |
0 1 2 3

→ print(s[2])
→ c will get printed.

Q) Given a char[], toggle every character.

↳ Capital ⇌ Small

Note → Input contains only small & capital characters.

Eg → AnaConDa

o/p → aNAcONdA

```
toggleCharacters( s[], N){
    for( i=0; i<N; i++){
        if (s[i] >=65 && s[i] <=90){
            // s[i] is capital
            s[i] += 32;
        }
        else{ // s[i] is small
            s[i] -= 32;
        }
    }
}
```

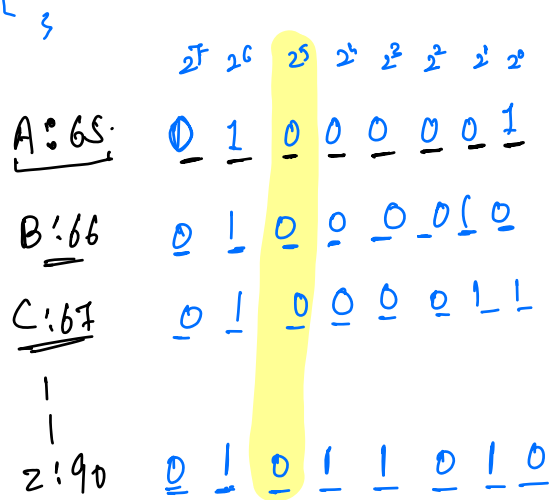
s[i] = s[i] ^ 32

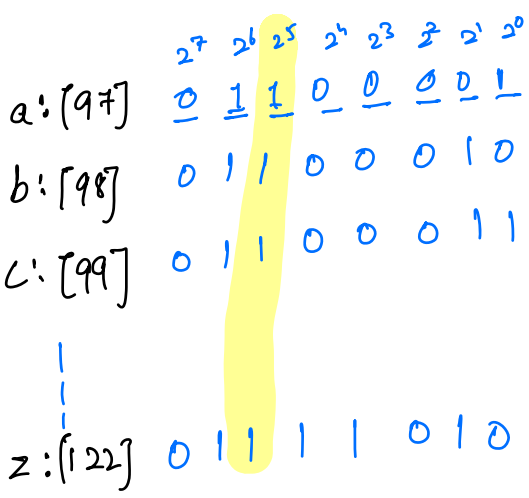
or

s[i] = s[i] ^ (1<<5)

T.C → O(N)

S.C → O(1)





Q) Given a char array, which contains only lower-case alphabets. Sort given ch[] in alphabetical order.

constraints

$1 <= N <= 10^5$

'a' <= ch[i] <= 'z'

Ex: s = d a b a c d b

After sorting ⇓

s = a a b b c d d

ideas

① Sort char[] using bubble sort.

T.C → $O(N^2)$

S.C → $O(1)$

② Use inbuilt sort function + custom comparator. {if needed}

T.C → O(nlogn).

③ s = d a b a c d b

'a' → 2

'b' → 2

'c' → 1

'd' → 2

s = a a b b c d d

'a' - 'a' = 0

'b' - 'a' = 1

'c' - 'a' = 2

⋮

'z' - 'a' = 25

count[26] →

2	2	1	2	0	— — —	0	0
0	1	2	3	4	— — —	24	25
'a'	'b'	'c'	'd'	'e'		'y'	'z'

count[0] → frequency of 'a'
count[1] → frequency of 'b'
count[2] → frequency of 'c'
⋮ ⋮
count[25] → frequency of 'z'

```
sortString ( char[] s , N ) {
    count[26] ;
    for ( i = 0 ; i < N ; i++) {
        idx = s[i] - 'a'
        count[idx] += 1 ;
    }                                   } N
    K = 0
    for( i = 0 ; i <= 25 ; i++){
        ch = i + 'a' ;
        for( j = 1 ; j <= count[i]; j++){
            s[K] = ch
            K += 1
        }                               } n
}
```

T.C → O(N)
S.C → O(1)

Eg : s → b b c c d d
0 1 2 3 4 5 6

count →

0	1	2	3	4	5	—	—	25
0	2	2	2	1	0	—	—	0

i	j	iteration
0	[1, c[0]]	c[0]
1	[1, c[1]]	c[1]
2	[1, c[2]]	c[2]
⋮		⋮
25	[1, c[25]]	c[25]

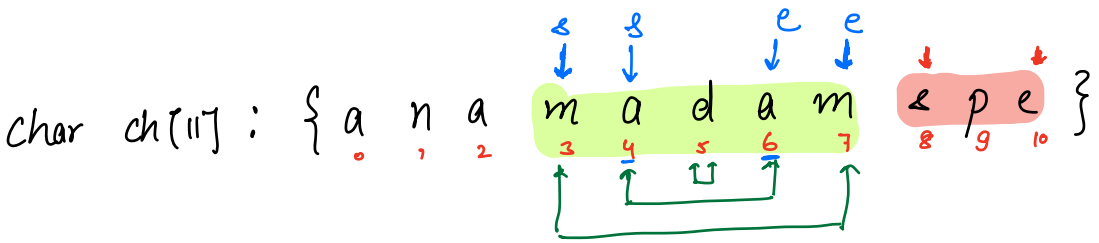
= n

Substring concept is same as sub-array.

- 1) continuous part of a string
- 2) full string can be sub-string.
- 3) A single character can be a substring.

Q) check if given substring is palindrome or not.

Eg: [madam nayan level
mam civic malayalam
dad radar]



start index of the substring

end index of the substring.

```
boolean isPalindrome ( ch[] str, s , e ){
    while ( s < e) {
        if (ch[s] != ch[e]){
            return false;
        }
        s++ , e--
    }
    return true;
}
```

T.C → O(N)
S.C → O(1)

Q: Given a string, calculate length of the longest palindromic substring.

Eg: a b a c a b
0 1 2 3 4 5
[ans = 5]

Eg: a b c d e
0 1 2 3 4
[ans = 1]

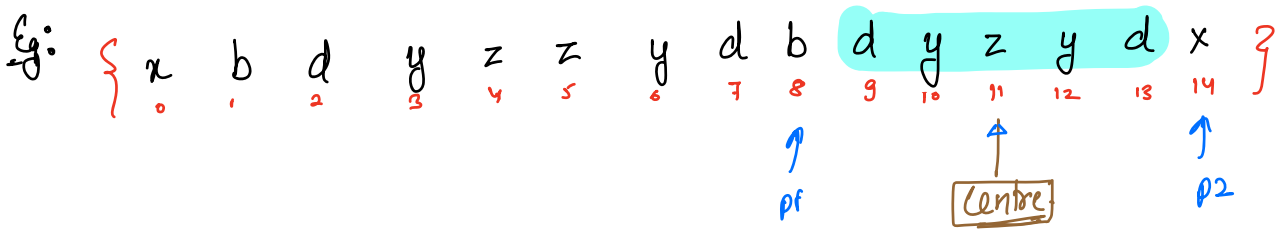
idea-1: For all the substrings, check if it is palindrome or not & get the max-length.

$\frac{n(n+1)}{2} * n \Rightarrow$ T.C = $O(N^3)$, S.C = $O(1)$

constraints $1 <= N <= 3 \times 10^3$

```
int longestPal ( char[] s , N) {
    ans = 0;
    for( i = 0 ; i < N ; i++) {          // i → start index of the substring
        for( j = i ; j < N : j++) {      // j → end index of the substring.
            // substring [i, j]
            if ( isPalindrome ( s, i , j ) ){
                ans = max (ans, j - i + 1);
            }
        }
    }
    return ans
}
```

Eg:



length = (P_1, P_2)

$= P_2 - P_1 - 1$.

idea:

Take any idx as the centre and try to expand to find the length of longest palindromic substring. ⇒ max odd length palindrome.

T.C → $O(N^2)$

max length palindrome.

Take any two consecutive characters as the centre & try to expand to find the length of longest palindromic substring. ⇒ max even length palindrome.

T.C → $O(N^2)$

T.C → $O(N^2)$

S.C → $O(1)$

```
int expand ( char s[] , p1 , p2 ){
    while ( p1 >= 0 && p2 < N && s[p1] == s[p2] ){
        p1--
        p2++
    }
    return p2 - p1 - 1 ;
}
```

```
int longestPal ( char[] s , n ){
    ans = 0/1

    for( i = 0 ; i < n ; i++ ){   // max odd-length palindrome.
        //Centre -> s[i]
        p1 = i , p2 = i
        ans = max (ans , expand ( s , p1 , p2) ) ;
    }

    for( i = 0 ; i < n-1 ; i++ ){   // i = N-1 , p1 = N-1 , p2 = N
        //centres -> s[i], s[i+1]
        p1 = i , p2 = i+1
        ans = max ( ans , expand(s , p1 , p2) ) ;
    }

    return ans ;
}
```

Doubts →

Eg: { x b d y z z y d b d y z y d x }
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14

$i = 4$

$P_1 = \cancel{4}\ \cancel{3}\ 2$

$P_2 = \cancel{8}\ \cancel{6}\ 7$

```
for( i=0 ; i < n-1 ; i++) {
    if (a[i] > a[i+1]) {
        swap a[i], a[i+1]
        i = -1
    }
}
```

$1+2+3+\cdots+n$

$\Rightarrow \frac{n(n+1)}{2} = n^2$

arr =>

3	4	2	5	
~~x~~	~~y~~	~~z~~	~~z~~	1
0	1	2	3	4

$i = \cancel{0}\ -1\ 0\ \cancel{1}\ -1\ 0\ -1\ 0\ X\ \cancel{2}\ 0$

$X = A \& B$. proof

CustomComparator.

$X \wedge A + X \wedge B$

$(A \& B) \wedge A + (A \& B) \wedge B$