Learn to sell. Learn to build.
If you can do both, you will be unstoppable.

- The Almanack of Naval Ravikant

Today's content

→ LRU Cache

→ Is linked list palindrome?

→ Intersection of two linked-lists.

L.R.U. Cache → temporary memory. → small.

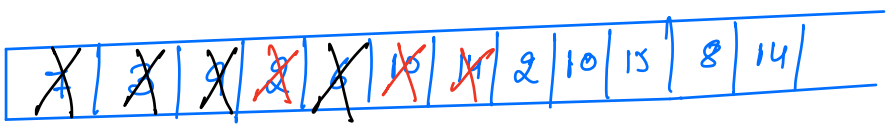
↓

Least Recently Used.

nos - 7 3 9 2 6 10 14 2 10 15 8 14

Capacity → 5

size → ~~1 2 3 4 5~~

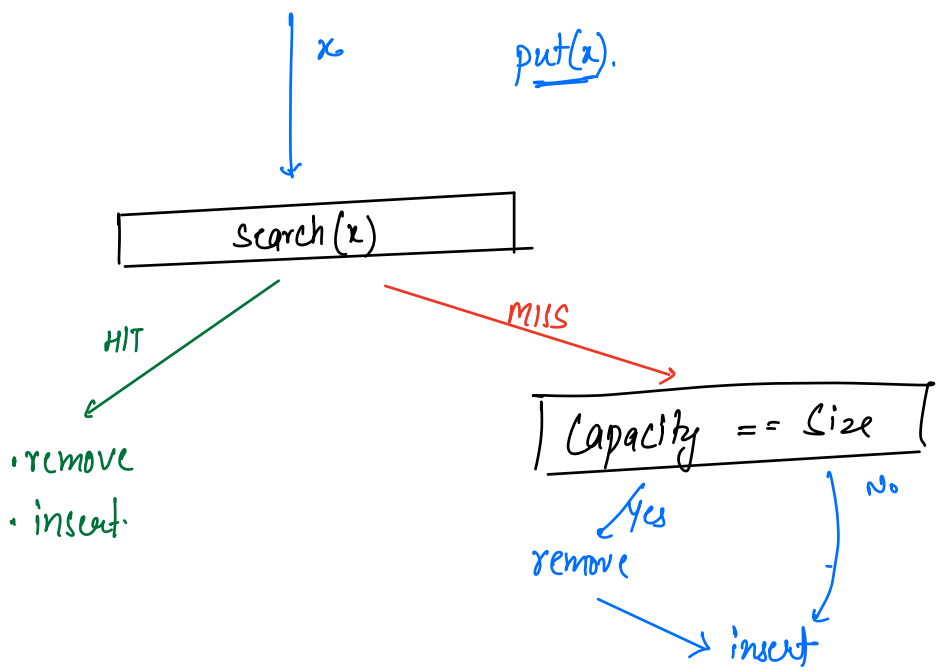


- search
- insert
- remove

HIT. Miss.

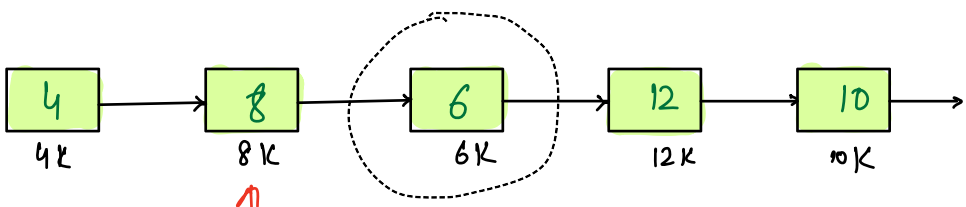
⇓

[already present in cache.]



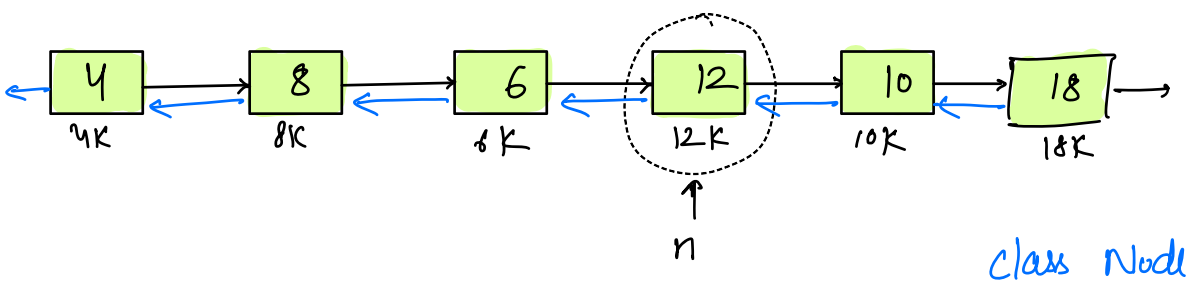
	array	L.L	Hashmap	LL + Hashmap	DLL + Hashmap
search	O(N)	O(N)	O(1)	O(1)	O(1)
insert	O(1)	O(1)	order is not maintained.	O(1)	O(1)
remove	O(N)	O(1)		O(N)	O(1)
		⇓			

[traversal is already happening in searching part.]



prev.

[prev.next = node.next]



D.L.L.

[6K.next = n.next
10K.prev = n.prev]

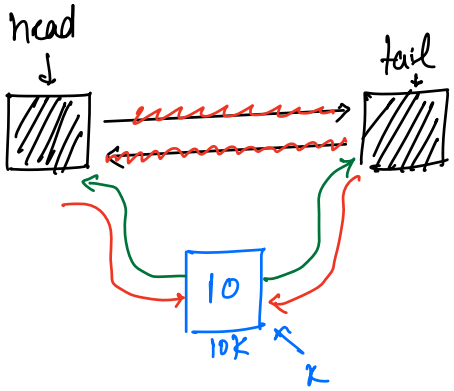
```
class Node {
    int val;
    Node next;
    Node prev;
```

Q. - 10 15 19 20 15 18 23 20 19

head → ~~null~~
tail → ~~null~~

Capacity → 5.

Size → ~~0~~ 2.
3 4
5

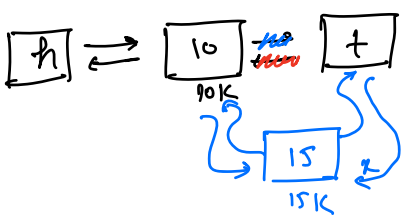


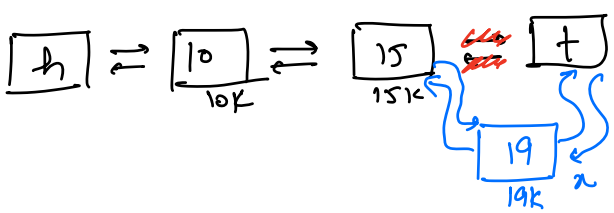
Hashmap.

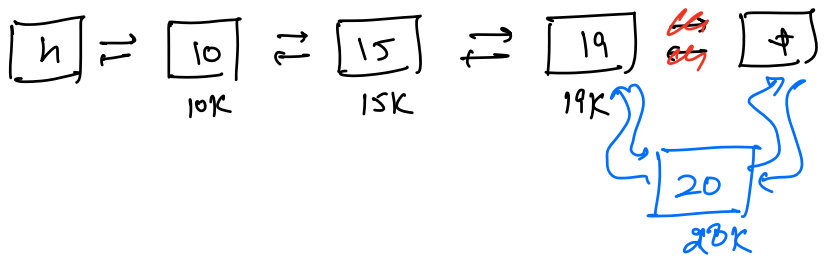
<int , reference of Node>

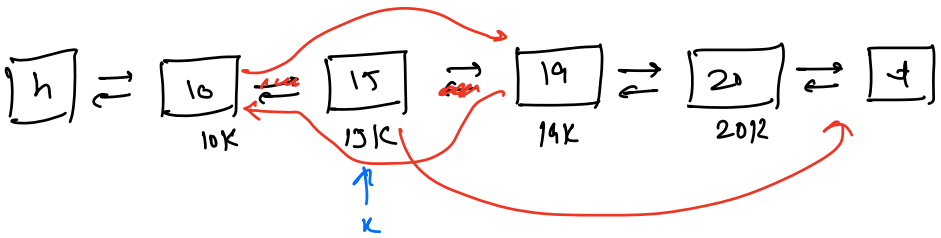
~~< 10, 10K >~~
< 15 , 15K >
< 19 , 19K >
< 20 , 20K >
< 18 , 18K >
< 23 , 23K >

```
void addToTail (Node x) {
    x.next = tail          ①
    x.prev = tail.prev     ②
    tail.prev = x          ③
    x.prev.next = x        ④
}
```



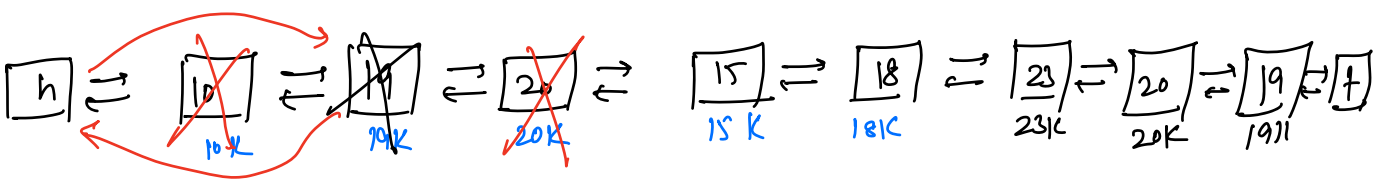






addToTail(x):

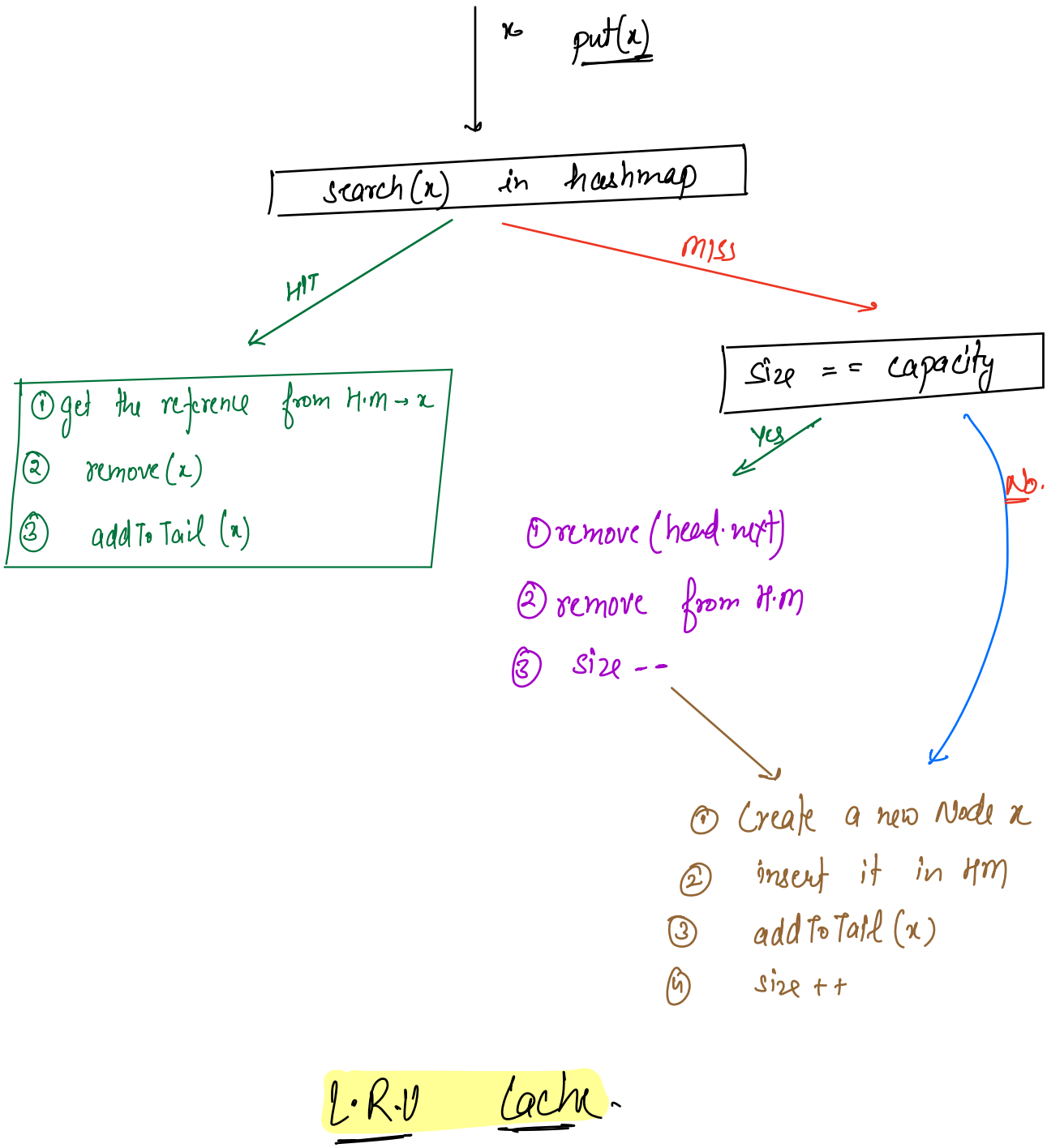
```
remove (Node x) {
    x.prev.next = x.next
    x.next.prev = x.prev
}
```



→ insert

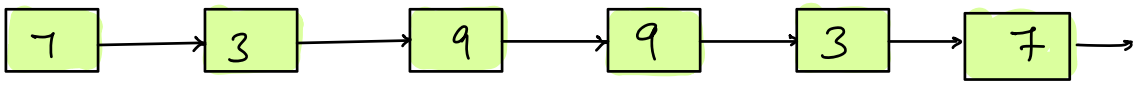
→ search

→ remove


x
put(x)
search(x) in hashmap
HIT
miss
① get the reference from H.M → x
② remove(x)
③ addToTail(x)
size == capacity
Yes
No.
① remove(head.next)
② remove from H.M
③ size --
① Create a new Node x
② insert it in HM
③ addToTail(x)
④ size ++

L·R·U Cache.

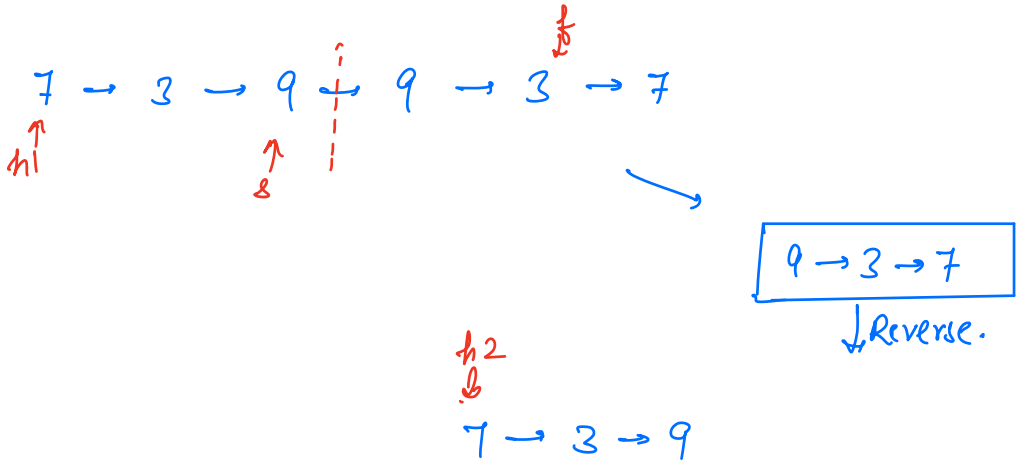
② Is linked list a palindrome ?



① extra space — ll to array.

7	3	9	9	3	7

② No extra space.



① Find middle node. [using slow & fast]

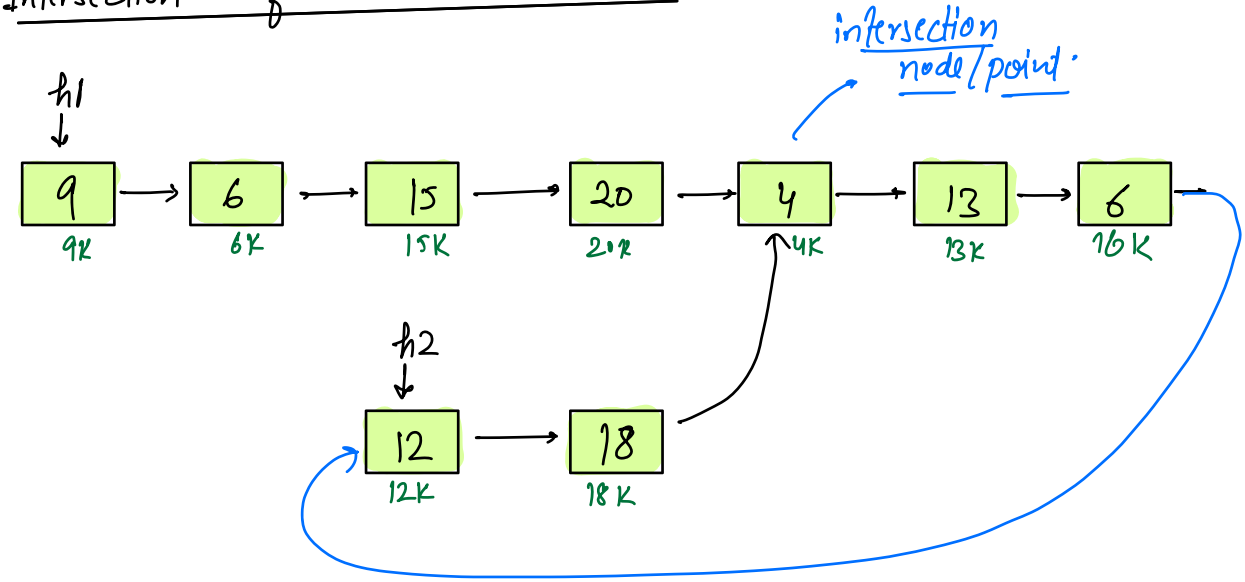
② h1 → head, h2 = slow.next [slow.next = null]

③ Reverse the second linked list. [ReverseLL(h2)]

④ Compare two linked-lists.

⑤ Re-construct the linked-list. → [Reverse LL(h2) slow.next = h2]

Intersection Of Two Linked-Lists →



idea-1. Using Hashset/hashmap.

T.C → O(N)
S.C → O(N)

9K
6K
15K
20K
4K → intersection node.
13K
16K

idea-2.

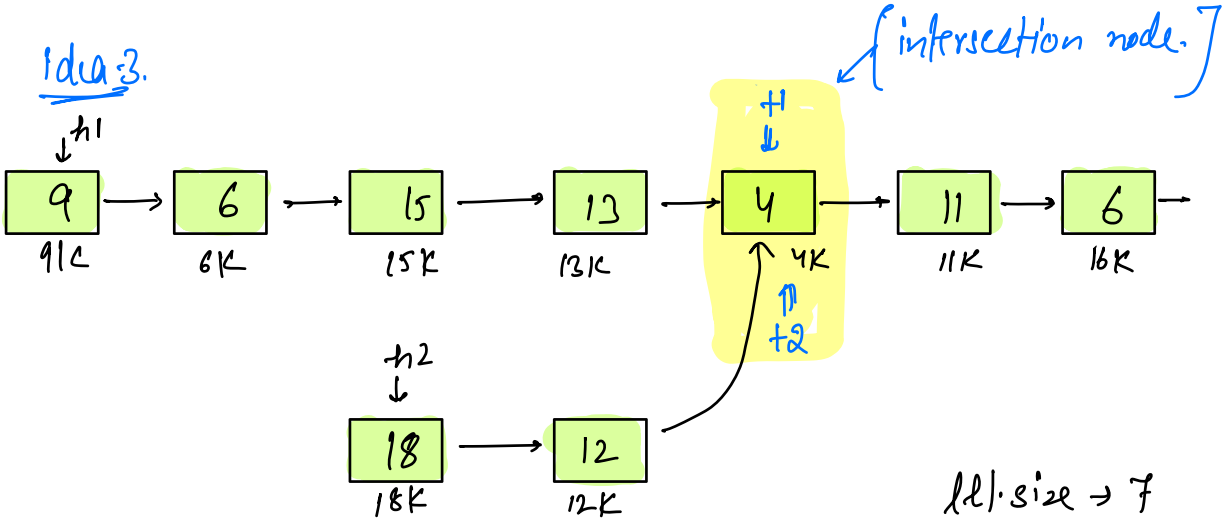
① Create a connection b/w tail of one linked-list & head of another linked-list. (loop will be created)

② find the starting point of loop. ⇒ Intersection point of 2 L.L's.

③ Break that connection (built in ①)

T.C → O(N)
S.C → O(1)

Idea:3.



ll1.size → 7

ll2.size → 5

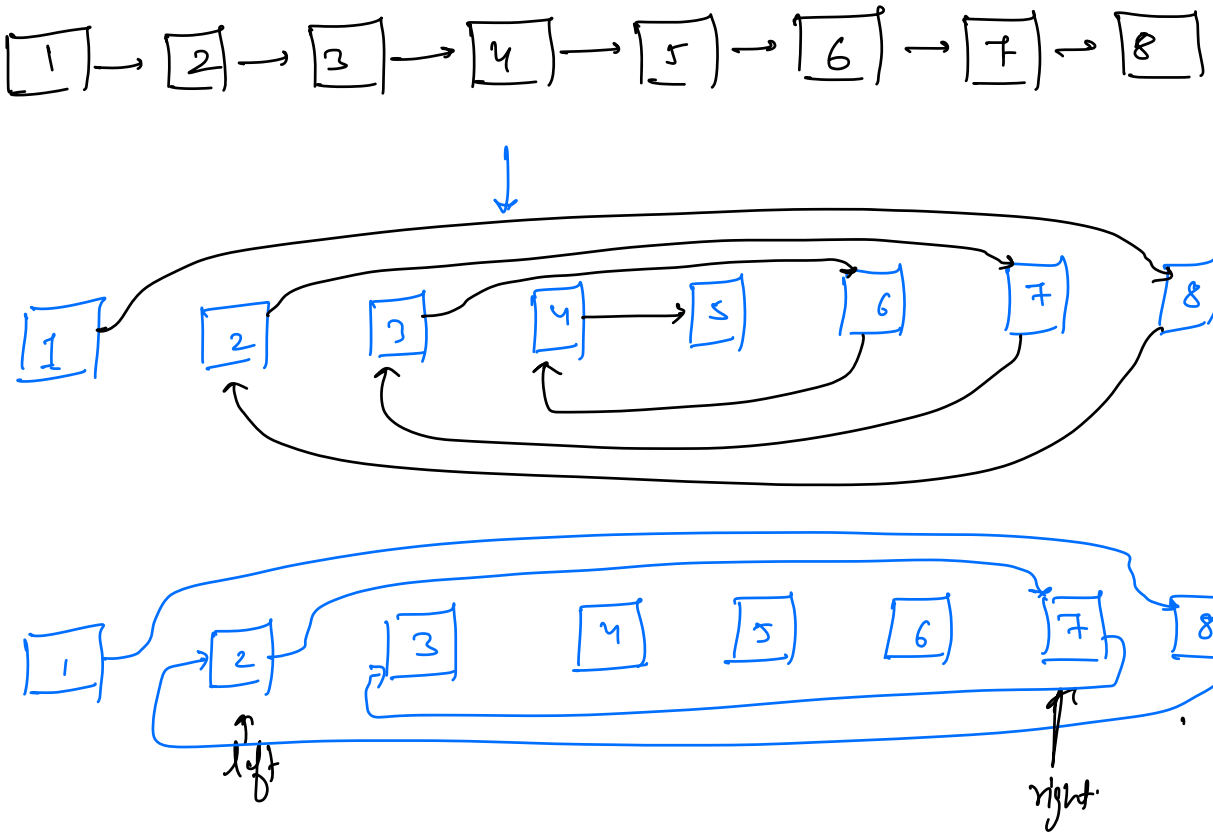
diff → 2.

idea:3

① Find the size of two linked-lists.

② Take t1 = h1, t2 = h2
and move the pointer corresponding to longer linked-list by diff = | ll1.size - ll2.size | steps.

③ Move t1 & t2 simultaneously by one step at a time till they are pointing to same node reference.


1
2
3
4
5
6
7
8
left
right