

2.4 – 2.5

Recursion 1

```
public int factorial(int n) {           //c1
    if(n==0){                          //c2
        return 1;                      //c3
    }else{
        return n*(factorial(n-1));    //c4 + T(n-1)
    }
}
```

$$T(n) = T(n-1) + c$$

$$\underline{T(n) = O(n)}$$

```
public static int fibonacci(int n) {           //c1
    if(n<=1){                              //c2
        return n;                          //c3
    }else{
        return (fibonacci(n-2)+ fibonacci(n-1));    // c4 + T(n-2) + T(n-1)
    }
}
```

$$T(n) = c4 + T(n-2) + T(n-1)$$

$$\underline{T(n) = O(2^n)}$$

```

public static int bunnyEars(int bunnies) { // c1
    if(bunnies == 0){           //c2
        return 0;               //c3
    }else{
        return 2+(bunnyEars(bunnies-1)); //c4 + T(2 + (n-1))
    }
}

```

$T(n) = c4 + T(2 + (n-1))$

$T(n) = c + T(n + 1)$

$T(n) = O(n)$

```

public int bunnyEars2(int bunnies) {
    if(bunnies == 0){           //c1
        return 0;               //c2
    }else if(bunnies % 2 == 1){ //c3
        return 2 + bunnyEars2(bunnies-1); //c4 + T(2 + (n-1))
    }else {
        return 3 + bunnyEars2(bunnies-1); //c5 + T(3 + (n-1))
    }
}

```

$T(n) = c4 + T(2 + (n-1)) + c5 + T(3 + (n-1))$

$T(n) = O(n)$

```

public int triangle(int rows) {
    if(rows == 0){        //c1
        return 0;        //c2
    }else{
        return rows + 1 * (triangle(rows-1)); //c3 + T(n -1)
    }
}

```

$T(n) = c3 + T(n -1)$

$T(n) = O(n)$

Recursión 2

```

public boolean groupSum6(int start, int[] nums, int target){
    if (start >= nums.length)        //c1
        return (target == 0);        //c2
    if (groupSum6(start+1, nums, target - nums[start])) return true;    //c3 + T(n -1)
    if (nums[start] != 6 && groupSum6(start+1, nums, target)) return true; //c4 + T(n-1)
    return false;                    //c5
}

```

$T(n) = T(n-1) + T(n+1) + c$

$n = \text{start}$

$T(n) = O(2^n)$

```

public boolean groupNoAdj(int start, int[] nums, int target) {
    if (start >= nums.length)      //c1
        return (target == 0);      //c2
    if (groupNoAdj(start+1, nums, target)) return true;    //c3 + T(n-1)
    if (groupNoAdj(start+2, nums, target-nums[start])) return true; //c4 + T(n-2)
    return false; //c5
}

```

$T(n) = T(n-1) + T(n-2) + c$

$n = \text{start}$

$T(n) = (2^n)$

```

private boolean checkOne(int start, int[] nums) {
    if (start == 0) return true;    //c1
    if (start > 0 && nums[start-1] % 5 == 0 && nums[start] == 1) //c2 + T(n-1)
        return false; //c3
    else
        return true; //c4
}

```

$T(n) = T(n-1) + c$

$T(n) = c \cdot n + c1$

$n = \text{start}$

$T(n) = O(n)$

```

public boolean groupSum5(int start, int[] nums, int target) {
    if (start >= nums.length) return (target == 0);    //c1
    if (groupSum5(start+1, nums, target-nums[start]) && checkOne(start, nums)) // c2 + T(n-1)
        return true;    //c3
    if (nums[start] % 5 != 0 && groupSum5(start+1, nums, target)) return true; // c4 + T(n-1)
    return false;    //c5
}

```

$T(n) = T(n-1) + T(n-1)$

$T(n) = O(n) + O(2^n)$

$n = \text{start}$

$T(n) = O(2^n)$

```
private void altArray(int[] nums) {  
    for (int i = 0; i < nums.length; i++) { //c.n  
        if (i > 0 && nums[i] == nums[i-1]) { //c2  
            nums[i-1] += nums[i]; //c3  
            if (i+1 < nums.length && nums[i] != nums[i+1])  
                nums[i] = 0; //c4  
            else if (i == nums.length-1) // c5  
                nums[i] = 0; //c6  
        }  
    }  
}
```

$T(n) = O(1) + O(n)$

$n = \text{nums}$

$T(n) = O(n)$

```
public boolean groupSumClump(int start, int[] nums, int target) {  
    altArray(nums); //O(2^n)  
    if (start >= nums.length) return target == 0; //c2  
    if (groupSumClump(start+1, nums, target-nums[start])) return true; //c3 + T(n-1)  
    if (groupSumClump(start+1, nums, target)) return true; //c4 + T(n-1)  
    else return false; //c5  
}
```

$T(n) = T(n-1) + T(n-1)$

$n = \text{start}$

$$T(n) = O(2^n) + O(n)$$

$$\underline{T(n) = O(2^n)}$$

```
private boolean recArray ( int[] nums, int index, int sum1, int sum2 ) {
    if ( index >= nums.length ) {    //c1
        return sum1 == sum2;        //c2
    }
    int value = nums[index];    //c3
    return (recArray(nums, index + 1, sum1 + value, sum2) ||
        recArray(nums, index + 1, sum1, sum2 + value)); //c4 + T(n -1) + T(n-1)
}
```

$$T(n) = c4 + T(n -1) + T(n-1)$$

n = index

$$\underline{T(n) = O(2^n)}$$

```
public boolean splitArray(int[] nums) {
    int index = 0;
    int sum1 = 0;
    int sum2 = 0;
    return recArray(nums, index, sum1, sum2); // O(2^n)
}
```

$$\underline{T(n) = O(2^n)}$$

2.3

Dado un arreglo de enteros, retorna true o false si es posible escoger un subconjunto de esos enteros, de tal manera que la suma de los elementos de ese subconjunto sea igual a target. El parámetro start funciona como un contador y representa un índice en el arreglo de números nums.

```
public boolean SumaGrupo(int start, int[] nums, int target) {  
    if (start >= nums.length) return target == 0;  
        return SumaGrupo(start + 1, nums, target - nums[start]) || SumaGrupo(start + 1, nums, target);  
}
```

Se llama recursivamente incrementando los valores de start y de target hasta que start se hace mayor o igual a el tamaño del arreglo y ahí la función empieza a retornar los valores a los que equivale cada cambio de valor cuando se llamaba recursivamente al método si este alguna vez es igual al target que se ingresó como parámetro retorna true