

# Functional Data Structures

## Exercise Sheet 10

### Exercise 10.1 Tries with 2-3-trees

In this exercise, you shall develop a trie data structure for keys of type  $'a \text{ list}$  (instead of  $\text{bool list}$ ).

Thus, a node needs to store a map from  $'a$  to the next trie.

In a first step, we encode the map as  $'a \Rightarrow 'b \text{ option}$

**datatype**  $'a \text{ trie} = \text{Leaf} \mid \text{Node } \text{bool} \text{ ``}'a \multimap 'a \text{ trie}``$

Define and prove correct membership, insertion and deletion (without shrinking the trie).

**fun**  $\text{isin} :: \text{``}'a \text{ trie} \Rightarrow 'a \text{ list} \Rightarrow \text{bool}``$

**fun**  $\text{ins} :: \text{``}'a \text{ list} \Rightarrow 'a \text{ trie} \Rightarrow 'a \text{ trie}``$

**lemma**  $\text{ins\_correct}: \text{``} \text{isin} (\text{ins } as \ t) \ bs = (as = bs \vee \text{isin } t \ bs) \text{``}$

**fun**  $\text{delete} :: \text{``}'a \text{ list} \Rightarrow 'a \text{ trie} \Rightarrow 'a \text{ trie}``$  **where**

**lemma**  $\text{delete\_correct}: \text{``} \text{isin} (\text{delete } as \ t) \ bs = (as \neq bs \wedge \text{isin } t \ bs) \text{``}$

Now refine the trie data structure to use 2-3-trees for the map. Note: To make the provided interface more usable, we introduce some abbreviations here:

**abbreviation**  $\text{``empty23} \equiv \text{Tree23.Leaf}``$

**abbreviation**  $\text{``inv23 } t \equiv \text{bal } t \wedge \text{sorted1 } (\text{inorder } t) \text{``}$

**lemmas**  $\text{map23\_thms}[\text{simp}] = \text{map\_empty } \text{map\_update } \text{map\_delete}$   
 $\text{invar\_empty } \text{invar\_update } \text{invar\_delete}$

The refined trie datatype

**datatype**  $'a \text{ trie}' = \text{Leaf}' \mid \text{Node}' \text{ bool} \text{ ``} ('a \times 'a \text{ trie}') \text{ tree23}``$

Define an invariant for  $\text{trie}'$  and an abstraction function to  $\text{trie}$ . Then define membership, insertion, and deletion, and show that they behave correctly wrt. the abstract  $\text{trie}'$ . Finally, combine the correctness lemmas to get a set interface based on 2-3-tree tries.

```

fun trie'_inv :: "'a::linorder trie'  $\Rightarrow$  bool"
fun trie'_ $\alpha$  :: "'a::linorder trie'  $\Rightarrow$  'a trie'"

```

```

fun isin' :: "'a::linorder trie'  $\Rightarrow$  'a list  $\Rightarrow$  bool"
fun ins' :: "'a::linorder list  $\Rightarrow$  'a trie'  $\Rightarrow$  'a trie'"
fun delete' :: "'a::linorder list  $\Rightarrow$  'a trie'  $\Rightarrow$  'a trie'"

```

**lemma** ins'\_correct: "trie'\_inv t  $\implies$  (isin' (ins' xs t) ks  $\longleftrightarrow$  xs=ks  $\vee$  isin' t ks)  $\wedge$  trie'\_inv (ins' xs t)"

**lemma** delete'\_correct: "trie'\_inv t  $\implies$  (isin' (delete' xs t) ks  $\longleftrightarrow$  xs $\neq$ ks  $\wedge$  isin' t ks)  $\wedge$  trie'\_inv (delete' xs t)"

## Exercise 10.2 Union Function on Tries

Define a function to merge two tries and show its correctness

```

fun union :: "trie  $\Rightarrow$  trie  $\Rightarrow$  trie"
lemma "isin (union a b) x = isin a x  $\vee$  isin b x"

```

## Homework 10.1 Tries with Same-Length Keys

Submission until Friday, 22. 6. 2018, 11:59am.

Consider the following trie datatype:

```

datatype trie = LeafF | LeafT | Node "trie * trie"

```

It is meant to store keys of the same length only. Thus, the *Node* constructor stores inner nodes, and there are two types of leaves, *LeafF* if this path is not in the set, and *LeafT* if it is in the set.

Define an invariant *is\_trie* *N* *t* that states that all keys in *t* have length *N*, and that there are no superfluous nodes, i.e., no nodes of the form *Node* (*LeafF*, *LeafF*).

```

fun is_trie :: "nat  $\Rightarrow$  trie  $\Rightarrow$  bool"

```

Hint: The following should evaluate to true!

```

value "is_trie 42 LeafF"
value "is_trie 2 (Node (LeafF, Node (LeafT, LeafF)))"

```

Whereas these should be false

```

value "is_trie 42 LeafT" — Wrong key length
value "is_trie 2 (Node (LeafT, Node (LeafT, LeafF)))" — Wrong key length
value "is_trie 1 (Node (LeafT, Node (LeafF, LeafF)))" — Superfluous node

```

Define membership, insert, and delete functions, and prove them correct!

```

fun isin :: "trie  $\Rightarrow$  bool list  $\Rightarrow$  bool"
fun ins :: "bool list  $\Rightarrow$  trie  $\Rightarrow$  trie"

```

```

lemma isin_in:
  assumes "is_trie n t" and "length as = n"
  shows "isin (ins as t) bs = (as = bs  $\vee$  isin t bs)"
  and "is_trie n (ins as t)"

fun delete2 :: "bool list  $\Rightarrow$  trie  $\Rightarrow$  trie" where
lemma
  assumes "is_trie n t"
  shows "isin (delete2 as t) bs = (as  $\neq$  bs  $\wedge$  isin t bs)"
  and "(is_trie n (delete2 as t))"

```

Hints:

- Like in the *delete2* function for standard tries, you may want to define a "smart-constructor" *node* :: *trie*  $\times$  *trie*  $\Rightarrow$  *trie* for nodes, that constructs a node and handles the case that both successors are *Leaf*.
- Consider proving auxiliary lemmas about the smart-constructor, instead of always unfolding it with the simplifier.

## Homework 10.2 Enumeration of Keys in Tries

*Submission until Friday, 22. 6. 2018, 11:59am.*

Write a function that enumerates all keys in a trie, in lexicographic order! Prove it correct.

```

fun enum :: "trie  $\Rightarrow$  bool list list"

lemma enum_correct:
  "set (enum t) = { xs. isin t xs }" and "sorted_wrt op< (enum t)"

```

Note that Booleans are ordered by *False* < *True*, and that we imported *List\_lexord*, which defines a lexicographic ordering on lists, if the elements are ordered.

```

value "[True, True, False] < [True, True, True, True]"

```

## Homework 10.3 Be Original!

*Submission until Friday, 13. 7. 2017, 11:59am.*

Develop a nice Isabelle formalization yourself!

- This homework goes in parallel to other homeworks for the rest of the lecture period. From next sheet on, we will reduce regular homework load a bit, such that you have a time-frame of 3 weeks with reduced regular homework load.

- This homework will yield 15 points (for minimal solutions). Additionally, up to 15 bonus points may be awarded for particularly nice/original/etc solutions.
- You may develop a formalization from all areas, not only functional data structures.
- Document your solution, such that it is clear what you have formalized and what your main theorems state!
- Set yourself a time frame and some intermediate/minimal goals. Your formalization needs not be universal and complete after 3 weeks.
- You are welcome to discuss the realizability of your project with the tutor!
- In case you should need inspiration to find a project: Sparse matrices, skew binary numbers, arbitrary precision arithmetic (on lists of bits), interval data structures (e.g. interval lists), spatial data structures (quad-trees, oct-trees), Fibonacci heaps, prefix tries/arrays and BWT, etc.