

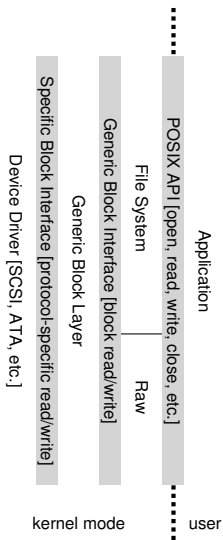
# CSL 301

## OPERATING SYSTEMS

### Lecture 24

#### I/O

Instructor  
Dr. Dhiman Saha

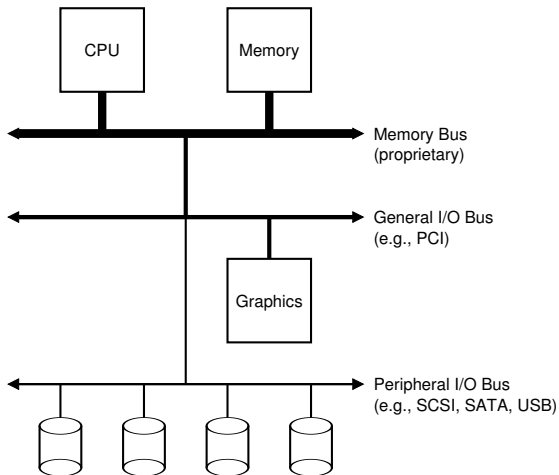


The File System Stack

# Persistence

- ▶ How should I/O be integrated into systems?
- ▶ What are the general mechanisms?
- ▶ How can we make them efficient?

# I/O devices connect to the CPU and memory via a bus



- ▶ Block device e.g disk
- ▶ Character device e.g keyboard
- ▶ The device interface (Abstracts out the internal details)
  - ▶ Commands supported
  - ▶ Current status of the device e.g busy
  - ▶ Data that needs to be transferred

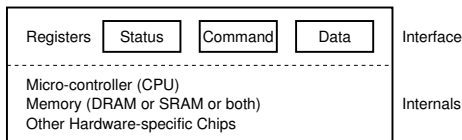


Figure 36.3: **A Canonical Device**

# How to deal with read/write to device

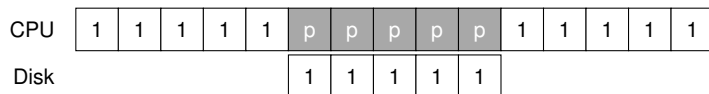
## Two Approaches

- ▶ Explicit instructions
  - ▶ e.g., `in` and `out` in x86
  - ▶ r/w to specific registers like status and command
- ▶ Memory Mapped I/O
  - ▶ Registers are abstracted as memory locations
  - ▶ Certain parts of memory are reserved
  - ▶ OS r/w to those specific locations
  - ▶ Memory hardware does the translation

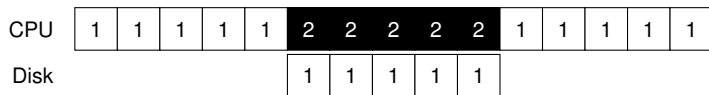
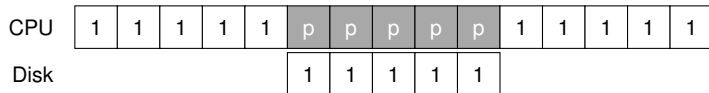
```
While (STATUS == BUSY)
    ; // wait until device is not busy
Write data to DATA register
Write command to COMMAND register
    (Doing so starts the device and executes the command)
While (STATUS == BUSY)
    ; // wait until device is done with your request
```

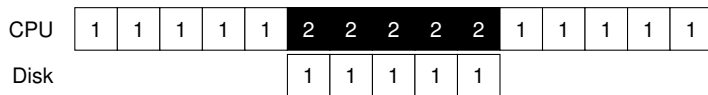
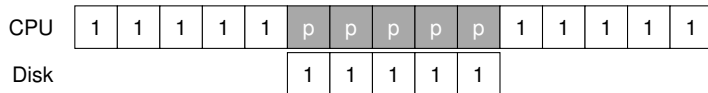
- ▶ Idea of polling
- ▶ Notion of Programmed I/O
- ▶ Both waste CPU cycles

How can the OS check device status without frequent polling?







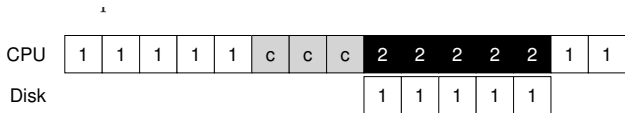


Issue

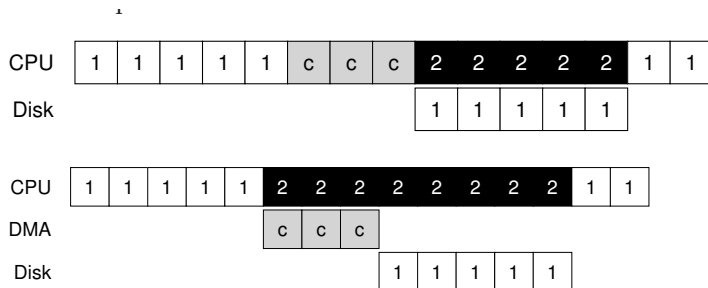
Interrupts

Livelock

- ▶ Already discussed
- ▶ Interrupt Descriptor Table
- ▶ Interrupt Number
- ▶ Interrupts allow for overlap of computation and I/O,



- CPU cycles still wasted in copying data to/from device registers.



# The Need for Device Drivers

- ▶ Devices have specific interfaces.
- ▶ The OS needs to be general and device-neutral.
- ▶ Device drivers encapsulate the details of device interaction.

- ▶ A piece of software in the OS must know in detail how a device works.
- ▶ This piece of software is called a device driver,
- ▶ Any specifics of device interaction are encapsulated within

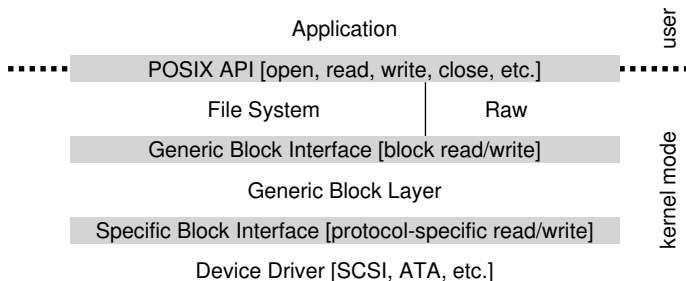
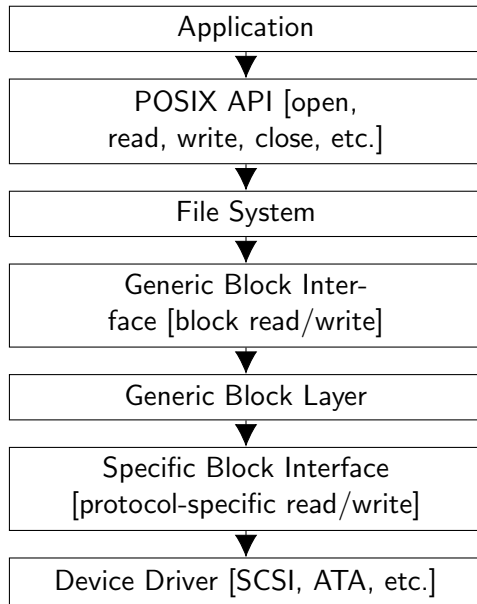


Figure 36.4: **The File System Stack**

# The File System Stack





# Summary

- ▶ Two techniques to improve device efficiency:
  - ▶ Interrupts
  - ▶ Direct Memory Access (DMA)
- ▶ Two approaches to accessing device registers:
  - ▶ Explicit I/O instructions
  - ▶ Memory-mapped I/O
- ▶ Device drivers encapsulate low-level details to make the OS device-neutral.