

# **Hash Functions and Chaos**

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**What is a hash function?**

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- Converts input of arbitrary length into fixed-length output
- Frequently used in hash tables

Hash function:  $\text{hash}(\text{key}) = \text{key} \% 17$

key	hash(key)	value
5	5	"hello"
8	8	"goodbye"
36	2	"compression"
22	5	"collision"

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- Resistant to collisions (given  $\text{hash}(a) == X$ , it is difficult to find  $b \neq a$  such that  $\text{hash}(b) == X$ )
- Avalanche effect (given  $a$  and  $a'$  are slightly different, even if  $\text{hash}(a) == X$  is known, it is hard to predict  $\text{hash}(a')$ )

# Desirable properties of hash functions

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Sensitive dependence on initial conditions

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Chaos?

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Sensitive dependence on initial conditions



Chaos?

(transitivity helps)

The big question:

**Can chaotic discrete maps be used as effective  
hash functions?**

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  - Rolling hash
  - Tent map
  - Dyadic map
  - “Multi-state” hash

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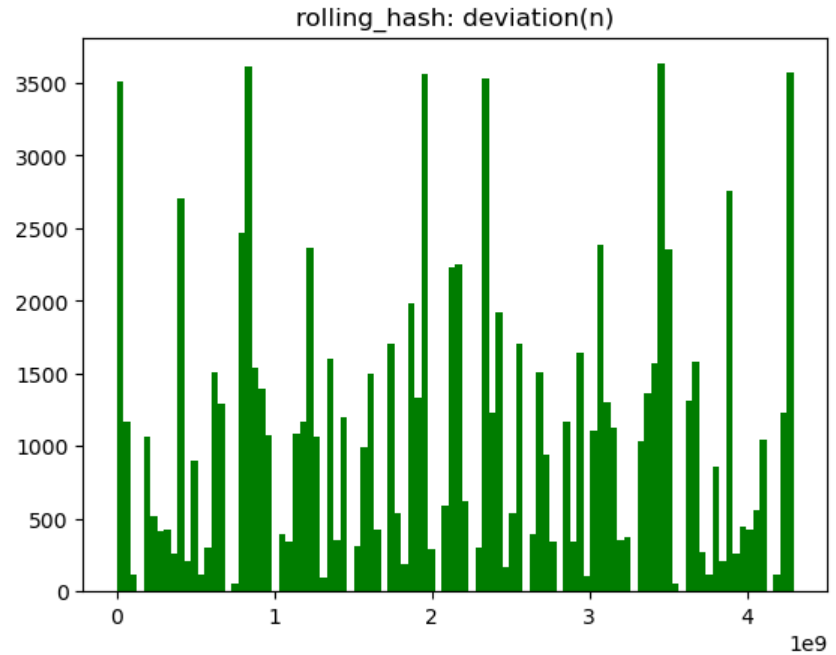
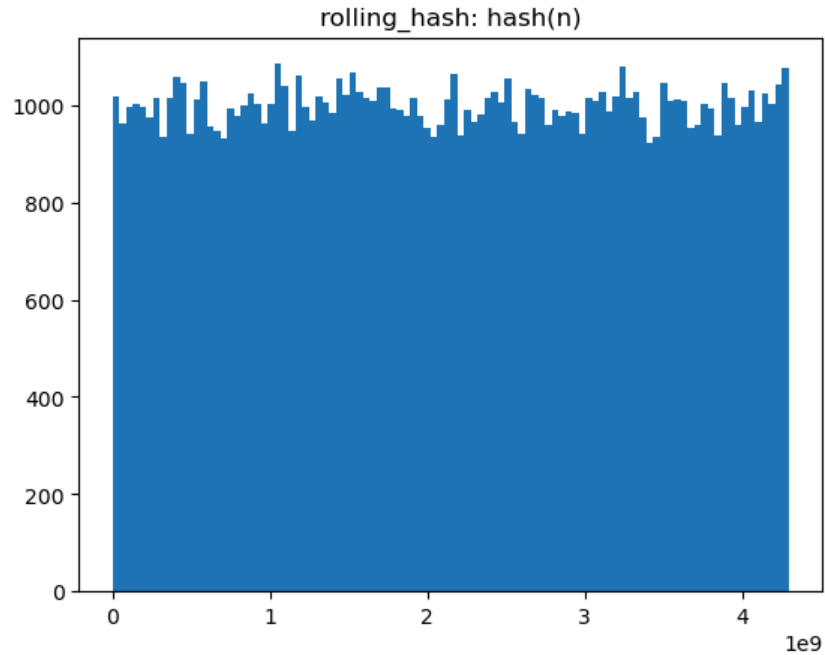
- Hash strings (`char[ ]`) to 32-bit unsigned integers (`unsigned int`)
- 4 different hashing algorithms
  - Rolling hash
  - Tent map
  - Dyadic map
  - “Multi-state” hash
- 3 different metrics
  - Distribution of hashed values
  - Number of collisions
  - Avalanche effect

# Rolling hash

```
unsigned int rolling_hash(char *n) {  
    unsigned int hash = 0;  
    unsigned int curr_pow = 1;  
    for (unsigned long i = 0; i < strlen(n); i++) {  
        hash += n[i] * curr_pow;  
        curr_pow *= 7919;  
    }  
    return hash;  
}
```

(Interpret the string as an integer in base 7919.)

# Rolling hash



Collision percentage: 0.003%

# Tent map hash

$$T(x) = \begin{cases} 2x & \text{if } x < \frac{2^{32}-1}{2} \\ 2 \cdot ((2^{32} - 1) - x) & \text{otherwise} \end{cases}$$

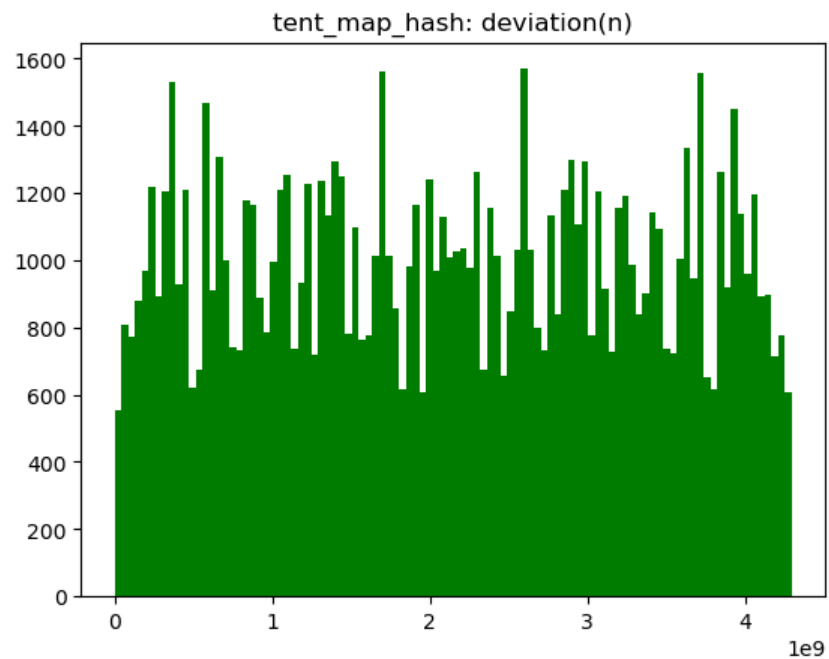
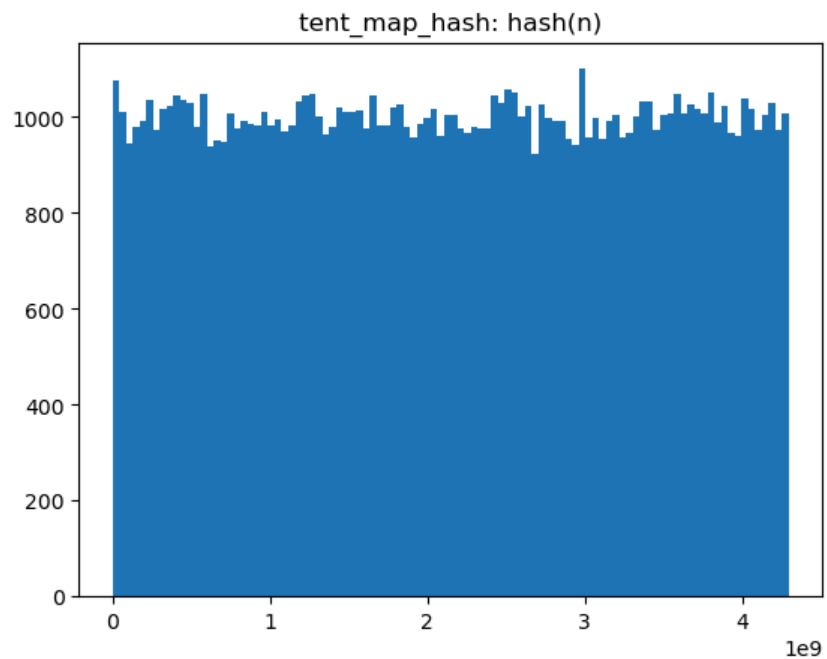
```
unsigned int tent_map_hash(char *n) {  
    unsigned int hash = rolling_hash(n);  
    int iterations = 53;  
}
```

(Find the rolling hash value of the string.)

```
    for (int i = iterations; i > 0; i--) {  
        if (hash < UINT_MAX / 2) {  
            hash = 2 * hash;  
        } else {  
            hash = 2 * (UINT_MAX - hash);  
        }  
    }  
    return hash;  
}
```

(Perform 53 iterations of the tent map on the hash value.)

# Tent map hash



Collision percentage: 0.001%

# Dyadic map hash

$$D(x) = \begin{cases} 2x & \text{if } x < \frac{2^{32}-1}{2} \\ 2 \cdot \left(x - \frac{2^{32}-1}{2}\right) & \text{otherwise} \end{cases}$$

```
unsigned int dyadic_map_hash(char *n) {  
    unsigned int hash = rolling_hash(n);  
    int iterations = 53;  
}
```

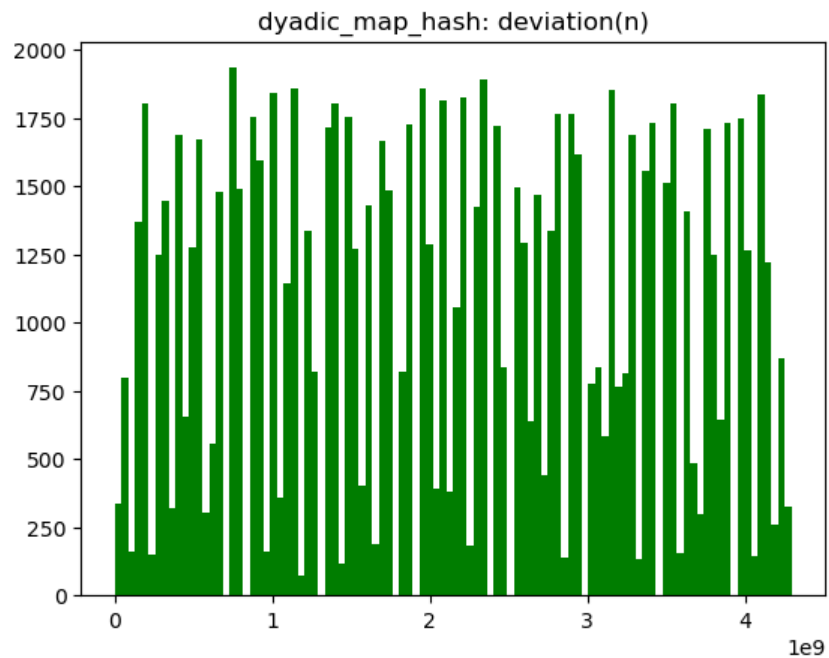
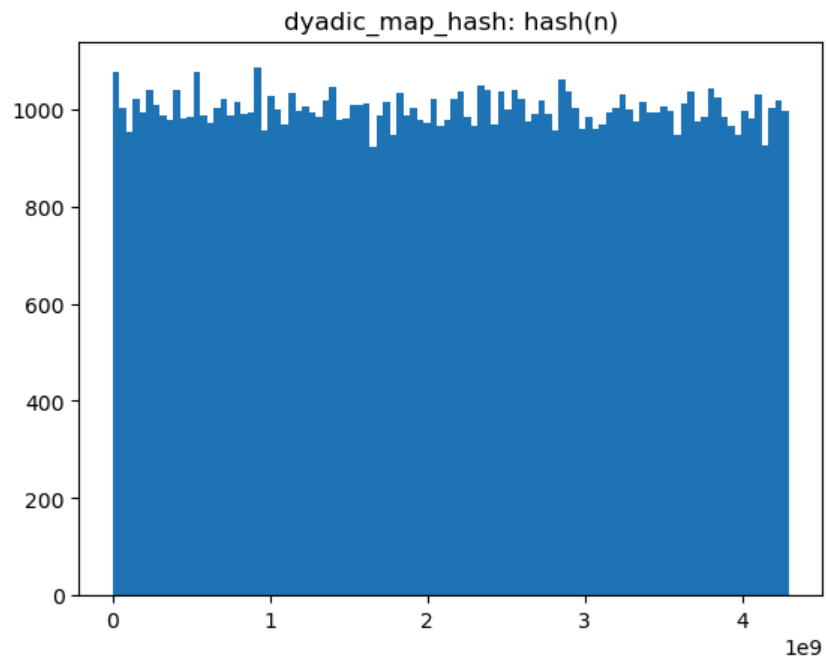
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    } else {  
        hash = 2 * (hash - UINT_MAX / 2);  
    }  
}  
return hash;  
}
```

(Perform 53 iterations of the dyadic map on the hash value.)

# Dyadic map hash



Collision percentage: 0.002%

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tent\_map\_hash with input "asdfqwer": 914563734

tent\_map\_hash with input "asdfqwgr": 1406886406

tent\_map\_hash with input "asdfqwer.": 4153872816

dyadic\_map\_hash with input "asdfqwer": 914564396

dyadic\_map\_hash with input "asdfqwgr": 1406886924

dyadic\_map\_hash with input "asdfqwer.": 141098142

# The dyadic map and tent map

- Topological conjugacy:

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where

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$$C(x) = \sin(\pi x)$$

- If  $C(x) = \sin(\pi x) \approx x$ , then  $T(x) \approx D(x)$

**What can we improve?**



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- Rolling hash is weak
- Easy to find collisions
  - Degrades performance due to lots of accidental collisions
- Not chaotic, so small perturbations  $\rightarrow$  small changes in hash
  - Transfers over to tent map and dyadic map hashes
- **Solution:** remove dependence on rolling hash

# MD5 algorithm

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  - Apply a (chaotic) map on (B, C, D) and store it in A
  - Rotate all values ( $A \rightarrow B, B \rightarrow C, \dots$ )
  - After the chunk is processed, add the values of A, B, C, D to the global state



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- Return A concat B concat C concat D

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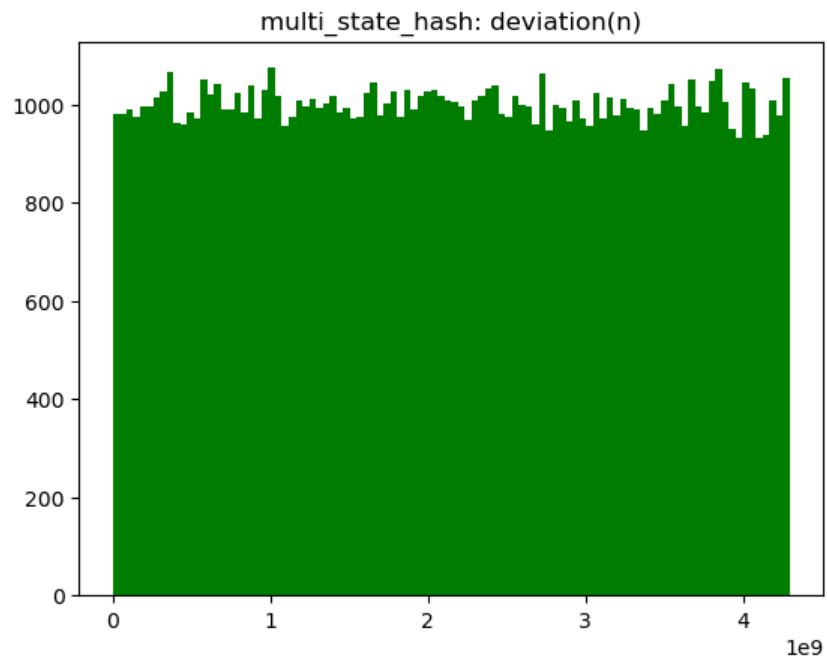
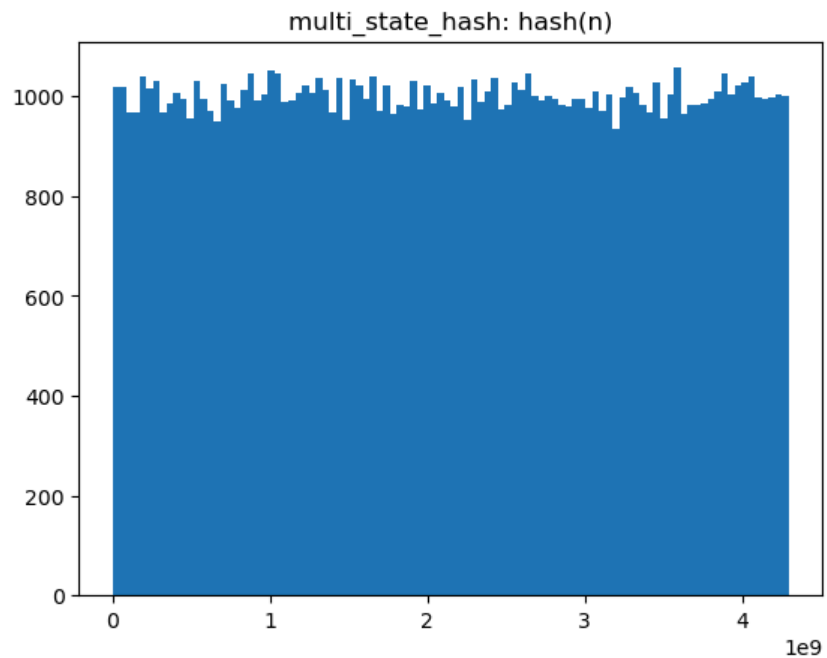
# Multi-state hash

- Maintain four state variables: A, B, C, D (each an 8-bit integer)
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- Add the next 8-bit chunk of input to A, apply the tent map to A, then rotate state values
  - Repeat 32 times per 8-bit chunk

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# Multi-state hash



Collision percentage: 1.162%

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Things that worked:

- Determinism
- Hash distribution with multi-state hash
- Avalanche effect

Things that didn't really work:

- Collision resistance

## **Potential future research:**

- How can collision resistance for the multi-state hash be improved?
- How much padding is really necessary for the multi-state hash to be effective?

**Thanks for listening!**

Source code for hashes + ipynb of analysis available at  
[github.com/ambareesh1510/chaotic-hash-functions](https://github.com/ambareesh1510/chaotic-hash-functions)