High-level signatures and initial semantics

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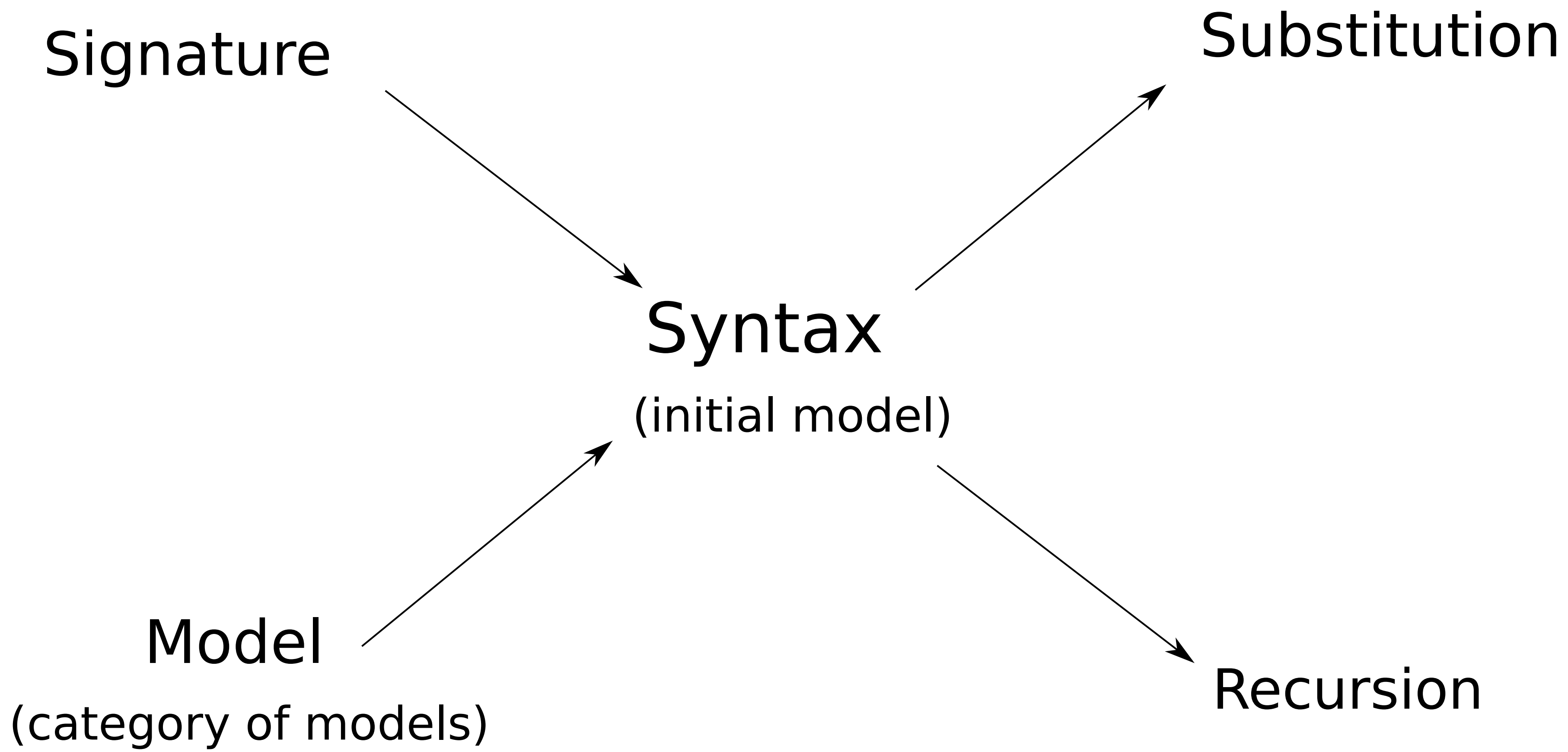
Overview

Purpose of our work: to specify and construct untyped syntaxes with variables and a well-behaved substitution (e.g. lambda-calculus).

More specifically (terms in italics will be explained):

1. we work with a general notion of *signature.* Not all of them do *generate a syntax*
2. classical *binding signatures* embed into our signatures as *algebraic signatures,* and indeed generate a syntax*.*
3. our main result: any *quotient* of algebraic signatures also generates a syntax

What is a syntax?



Signatures which we care about: those whose category of models have an *initial object*, i.e., generate a syntax.

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Example: 0, ★

Consider the syntax generated by a binary operation ★ and a constant **0** (and variables):

expr ::=	x	*(variable)*
	| t_1 ★ t_2	*(binary operation)*
	| 0	*(constant)*

The syntax can be considered as the endofunctor B (on Set):

$$B(X) = \text{expressions over } X$$

$$B(\emptyset) = \{0, 0 \star 0, \ldots\}$$

$$B(\{x, y\}) = \{0, 0 \star 0, \ldots, x, y, x \star y, \ldots\}$$

Example: 0, ★

The binary operation ★ induces a natural transformation:

$$B \times B \to B$$

The constant **0** induces a natural transformation:

$$1 \to B$$

Variables induce a natural transformation

$$\mathrm{Id}_{\mathrm{Set}} \to B$$

They gather into a single natural transformation:

$$B \times B + 1 + \mathrm{Id}_{\mathrm{Set}} \to B$$

i.e. B is an algebra for the endofunctor $F \mapsto F \times F + 1 + \mathrm{Id}_{\mathrm{Set}}$ on the category $\mathrm{End}_{\mathrm{Set}}$.

Actually, B can be defined to be the initial algebra.

Binding Signatures

Definition

Binding signature = a family of lists of natural numbers.

Each list specifies an operation in the syntax:

- length of the list = number of arguments of the operation
- natural number in the list = number of bound variables in the corresponding argument

Syntax with 0, ⋆: $\big((), (0, 0)\big)$

nullary operation **0** binary operation ⋆

Lambda calculus: $\big((0\,,\,0)\,,\,(1)\big)$

application lambda-abstraction

Signatures and endofunctors

In the same spirit as in the first example $(0,\star)$, any binding signature gives rise to an endofunctor Σ on the category $\mathrm{End}_{\mathrm{Set}}$.

A notion of model: $\Sigma + \mathrm{Id}_{\mathrm{Set}}$ -algebra

The initial $\Sigma + \mathrm{Id}_{\mathrm{Set}}$-algebra of a binding signature Σ always exists.

Does this initial algebra come with a well-behaved substitution?

Substitution [Fiore-Plotkin-Turi 1999]

The endofunctor Σ induced by a binding signature comes with a *strength* which allows [FPT] to refine the notion of model:

Σ-monoid

$\Sigma + \mathrm{Id}_{\mathrm{Set}}$-algebra **equipped with a well-behaved substitution**.

Σ-monoid morphisms

algebra morphisms commuting with substitution.

Theorem [FPT]:

The initial $\Sigma + \mathrm{Id}_{\mathrm{Set}}$-algebra of a binding signature Σ comes with a well-behaved substitution that makes it initial in the category of **Σ-monoids**.

This suggests defining signatures to be endofunctors on $\mathrm{End}_{\mathrm{Set}}$ *with strength* (as in [Matthes-Uustalu 2004]).

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Our signatures and models

Binding signatures $\hookrightarrow$ Endofunctors with strength $\hookrightarrow$ Our signatures

A **signature** Σ is a functorial assignment:

monad module over R

$$R \mapsto \Sigma(R)$$

A **model of** $\boldsymbol{\Sigma}$ is a pair:

$$(R, \quad \sigma : \Sigma(R) \to R)$$

monad module morphism

monad	:=	endofunctor with substitution
module over a monad	:=	endofunctor with substitution
module morphism	:=	natural transformation preserving substitution

Monads

Reminder:

- $B(X)$ = expressions built out of 0, ★ and variables taken in X
- Variables induce a natural transformation $\eta : \mathrm{Id}_{\mathrm{Set}} \to B$

In addition, there is a **substitution** $\mathrm{bind} : B(X) \to (X \to B(Y)) \to B(Y)$ required to satisfy some equations.

A triple (B, η, bind) is called a **monad**.

A **monad morphism** between two monads R and S is a family of maps $(f_X : R(X) \to S(X))_X$ preserving variables and substitution.

Operations are module morphisms

In the $(0, \star)$ language,

$$(t \star u)[x \mapsto v_x] = t[x \mapsto v_x] \star u[x \mapsto v_x]$$

$\star$ commutes with substitution

In the right hand side, substitution acts on a pair of expressions.

We abstract this situation as follows:

- pairs of expressions form a **module** $B \times B$ over the monad B,

- $\star$ yields a **module morphism** from $B \times B$ to B

Module over a monad

$B \times B$ comes with a substitution with B-expressions required to satisfy some equations:

$$(B \times B)(X) \to (X \to B(Y)) \to (B \times B)(Y)$$

Such a functor with B-substitution is called a **module over the monad B**.

Examples of modules over a monad

Some examples of **modules over a monad $\boldsymbol{R}$**:

- R itself
- M x N for any modules M and N (in particular, R x R)
- Given a module M, the **derivative of $\boldsymbol{M}$** is the module M' defined by M'(X) = M(X + $\{x\}$).

The derivative is used to model an operation binding a variable

(Cf next slide).

Examples of module morphisms

A **module morphism** between two modules M and N on the same monad R is a family of maps $(f_X{:}M(X) \to N(X))_X$ commuting with substitution.

$id_M : M \to M$

the family of identity maps $(id_{M(X)}{:}M(X) \to M(X))_X$ for any module M

$\star : B \times B \to B$

$app : L \times L \to L$

the application operation of the lambda calculus monad L.

$abs : L' \to L$

Indeed, in $\lambda x.t$, the term *t* depends on an additional free variable x:
If $t \in L(Y + \{x\}) = \boldsymbol{L'(Y)}$, then $\lambda x.t \in L(Y)$

Signatures and models

A **signature** Σ is a functorial assignment:

monad　　　　module over R

$$R \mapsto \Sigma(R)$$

A **model of Σ** is a pair:

$$(R, \quad \rho : \Sigma(R) \to R)$$

monad　　　　module morphism

A **model morphism** $m : (R,\rho) \to (S,\sigma)$ is a monad morphism commuting with the module morphism:

$$\begin{tikzcd}
\Sigma(R) \arrow[r, "\rho"] \arrow[d, "\Sigma(m)"'] & R \arrow[d, "m"] \\
\Sigma(S) \arrow[r, "\sigma"'] & S
\end{tikzcd}$$

Existence of syntax?

Notion of signature too general: existence of the syntax (**= initial model**) ?

Counter-example: the signature $R \mapsto \mathscr{P} \circ R$

powerset endofunctor on $\mathbb{S}\mathrm{et}$

Examples of signatures with syntax

- $R \mapsto 1 + R \times R$

  Models are monads R equipped with module morphisms $1 \to R$ and $R \times R \to R$.

  The syntax is the language B generated by a constant **0** and binary operation ★.

- $R \mapsto R \times R + R'$

  Models are monads R equipped with two modules morphisms: $R \times R \to R$ and $R' \to R$.

  The syntax is the lambda calculus.

Algebraic signatures

More generally, the syntax exists for any signature induced by a disjoint sum of products of finite derivatives of the monad ($R \mapsto R' \times R'' \times R''' + R \times R'' \times R''' \times R + \ldots$).

We call such a signature an **algebraic signature**. They correspond to binding signatures through the inclusion:

$$\text{Binding signatures} \hookrightarrow \text{Endofunctors with strength} \overset{\mathcal{I}}{\hookrightarrow} \text{Our signatures}$$

Our main result: Quotients of algebraic signatures generate a syntax.

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Quotient of a signature

Quotient of a set:

A quotient of a set X is a set Y together with a surjection $p : X \to Y$.

$$x \sim x' \qquad \Longleftrightarrow \qquad p(x) = p(x')$$

Quotient of a signature:

A quotient of a signature Σ is a signature Ψ together with a (natural) family of module morphisms $(f_R : \Sigma(R) \to \Psi(R))_R$ that is pointwise surjective.

$$R \mapsto \begin{tikzcd} \Sigma(R) \arrow[d, two heads, "f_R"] \\ \Psi(R) \end{tikzcd}$$

Syntax for presentable signatures

A **presentable signature** is a quotient of a binding signature.

Main Theorem: For any presentable signature, there is a syntax.

We now give examples of new kinds of operations specified by presentable signatures (more can be found in the article).

Example 1: Symmetric operations

Binary commutative operation +:

$$t + u = u + t$$

As a quotient of an algebraic signature:

$$R \mapsto \begin{tikzcd} R \times R \arrow[d, two heads] \\ R \times R / \{(x,y) \sim (y,x)\} \end{tikzcd}$$

This generalizes to **n-ary permutation invariant operations**.

Example 2: Explicit substitution

An operation $\_\langle x_i \mapsto t_i \rangle$ that mimics the behavior of the substitution in the sense that it enjoys some of its coherences:

- invariance under **permutation**

$$F(x,y)\langle x \mapsto t, y \mapsto u\rangle = F(y,x)\langle x \mapsto u, y \mapsto t\rangle$$

- invariance under **weakening**

$$F(x)\langle x \mapsto t, y \mapsto u\rangle = F(x)\langle x \mapsto u\rangle$$

- invariance under **contraction**

$$F(x,y)\langle x, y \mapsto t\rangle = F(x,x)\langle x \mapsto t\rangle$$

Example 2: Explicit substitution

Explicit substitution as a quotient of the algebraic signature:

$$\Sigma(R) := R'\times R \;+\; R''\times R\times R \;+\; R'''\times R\times R\times R \;+\; \ldots$$

$$t\langle x \mapsto u\rangle \qquad\qquad t\langle x \mapsto u, y \mapsto v\rangle \qquad\qquad t\langle x \mapsto u, y \mapsto v, z \mapsto w\rangle$$

$$R \mapsto \begin{tikzcd} \Sigma(R) \arrow[d, two heads] \\ \Sigma(R)\,/\sim \end{tikzcd}$$

- **permutation**: $t\langle x \mapsto u, y \mapsto v\rangle \sim t[x \rightleftarrows y]\langle x \mapsto v, y \mapsto u\rangle$
- **weakening**: $t\langle x \mapsto u\rangle \sim t\langle x \mapsto u, y \mapsto v\rangle$
- **contraction**: $t\langle x \mapsto u, y \mapsto u\rangle \sim t[y := x]\langle x \mapsto u\rangle$

Conclusion

We have given a criterion for signatures to generate a syntax. This criterion encompasses the classical binding signatures. It also allows new operations that satisfy some equations.

Our main theorem has been formalized using the Coq library UniMath.

Future work:

- extend our framework to encompass general equations (e.g. associative binary operation, lambda-calculus modulo beta/eta equivalence);

- extend our framework to simply typed syntaxes.

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- extend our framework to encompass general equations (e.g. associative binary operation, lambda-calculus modulo beta/eta equivalence);

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Thank you!

FIN PROVISOIRE

Ne pas lire les slides qui suivent (ce sont des anciennes slides que je garde au cas où).

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Our signatures

In the next slides, we present the notion of signature and models we work with.

$$\text{Binding signatures} \hookrightarrow \text{Endofunctors with strength} \overset{\mathcal{I}}{\hookrightarrow} \text{Our signatures}$$

Theorem (Zsido): for any endofunctor with strength Σ, there is an adjunction between our category of models and the [FPT] one:

$$\begin{tikzcd} \mathrm{Models}(\mathcal{I}(\Sigma)) \arrow[rr, bend left] \arrow[rr, phantom, "\top"] & & \Sigma\text{-monoids} \arrow[ll, bend left] \end{tikzcd} \quad .$$

Fixpoint operator with coherences

But we would like to encode some of the expected behaviour of such a fixed point:

- invariance under permutation
- invariance under weakening
- invariance under contraction. Roughly:

```
let rec f₁ = F(f₁,f₂)
    and f₂ = F(f₁,f₂)
in f₁
```

=

```
let rec f = F(f,f)
in f
```

A construction satisfying these invariances can be specified by quotienting the naive algebraic signature.

Fixpoint operator

A fixpoint operator:

A language with (mutual) fixpoints comes with a construction

```
let rec f1 = t1
    and f2 = t2
         ...
    and fn = tn
in fi
```

where each $\mathbf{f}_j$ may appear as a variable in each expression $\mathbf{t}_i$.

Thus, it takes **n** expressions $\mathbf{t}_1,..,\mathbf{t}_n$ depending on **n** fresh variables $\mathbf{f}_1,..,\mathbf{f}_n$ and produces an expression which no longer depend on them.

As such, it can be specified by a binding signature.

Our work

We work with a notion of signature, and associated models, based on the notion of module over a monad.

Goal of our work: To identify a large subclass of these signatures whose category of models have an initial object.

Our main result: Quotients of "algebraic signatures" generate a syntax.

Example 2: Syntactic closure operator

Syntactic closure operator $\forall$

$\forall xyz.t$ binds the variables x, y and z in the term t

Example of an operation invariant under **permutation** and **weakening**:

- permutation: $\forall xy.t = \forall yx.t$

- weakening: $\forall x.t = \forall xy.t$ if t does not depend on y