High-level signatures and initial semantics

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Overview

Topic: specification and construction of untyped syntaxes with variables and a well-behaved substitution (e.g. lambda calculus).

Our work:

1. general notion of ***signature*** based on ***monads*** and ***modules***.

   - *Caveat:* Not all of them do ***generate a syntax***

   - special case: classical **binding signatures**

2. our main result: any ***quotient*** of algebraic signatures generates a syntax.

This talk: explain the words in bold

Operations covered by our result

Some examples:

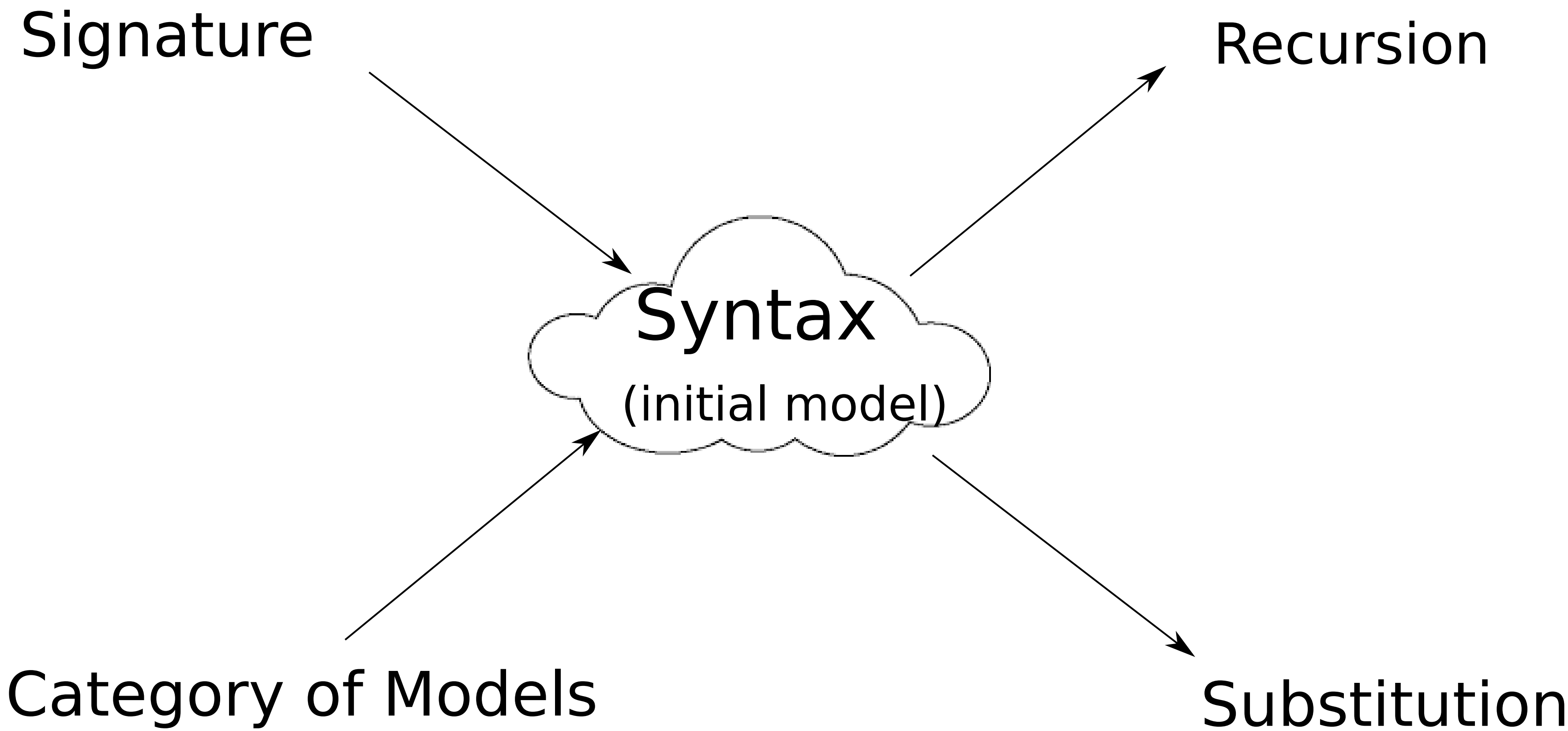
- Symmetric operations

- Explicit substitution

- Coherent fixed point operation

- Syntactic closure operator

What is a syntax?



generates a syntax = existence of the initial model

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Categorical formulation of a term language

Example: syntax with a binary operation, a constant, and variables

$$\begin{aligned} \text{expr} ::= {} & x && \textit{(variable)} \\ \mid {} & t_1 \bigstar t_2 && \textit{(binary operation)} \\ \mid {} & 0 && \textit{(constant)} \end{aligned}$$

The syntax can be considered as the endofunctor B (on Set):

$$B : X \mapsto \{\text{expressions over } X\}$$

For example:

$$B(\emptyset) = \{0, 0 \star 0, \ldots\}$$
$$B(\{x, y\}) = \{0, 0 \star 0, \ldots, x, y, x \star y, \ldots\}$$

Categorical formulation of a term language

The binary operation $\star$ induces a natural transformation:

$$B \times B \to B$$

The constant 0 induces a natural transformation:

$$1 \to B$$

Variables induce a natural transformation:

$$\mathrm{Id}_{\mathrm{Set}} \to B$$

They gather into a single natural transformation:

$$B \times B + 1 + \mathrm{Id}_{\mathrm{Set}} \to B$$

i.e. B is an algebra for the endofunctor $F \mapsto F \times F + 1 + \mathrm{Id}_{\mathrm{Set}}$ on the category $\mathrm{End}_{\mathrm{Set}}$.

Actually, B can be **defined** to be the initial algebra.

Binding Signatures

Definition

Binding signature = a family of lists of natural numbers.

Each list specifies one operation in the syntax:

- length of the list = number of arguments of the operation
- natural number in the list = number of bound variables in the corresponding argument

Syntax with 0, ★:

$$\Big(\underbrace{()}_{0},\ \underbrace{(0,\ 0)}_{\star}\Big)$$

Lambda calculus:

$$\Big(\underbrace{(0\ ,\ 0)}_{app}\ ,\ \underbrace{(1)}_{abs}\Big)$$

Initial semantics for binding signatures

In the same spirit as in the first example $(0, \star)$, any binding signature gives rise to an endofunctor Σ on the category $\mathrm{End}_{\mathrm{Set}}$.

A notion of model: $\Sigma + \mathrm{Id}_{\mathrm{Set}}$-algebra

The initial $\Sigma + \mathrm{Id}_{\mathrm{Set}}$-algebra of a binding signature Σ always exists.

Does this initial algebra come with a well-behaved substitution?

Classical results on initial semantics

The endofunctor Σ induced by a binding signature comes with a *strength* which allows [FPT] to refine the notion of model:

Σ-monoid:

$\Sigma+\mathrm{Id}_{\mathrm{Set}}$-algebra **equipped with a well-behaved substitution**.

Σ-monoid morphisms:

algebra morphisms commuting with substitution.

Theorem [FPT]:

The initial $\Sigma + \mathrm{Id}_{\mathrm{Set}}$-algebra of a binding signature comes with a well-behaved substitution that makes it initial in the category of **Σ-monoids**.

This suggests defining signatures to be endofunctors on $\mathrm{End}_{\mathrm{Set}}$ *with strength* (as in [Matthes-Uustalu 2004]).

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The Big Picture of signatures and models

Binding signatures $\hookrightarrow$ Endofunctors with strength $\hookrightarrow$ Our signatures

A **signature** Σ is a functorial assignment:

$$R \mapsto \Sigma(R)$$

A **model of** $\boldsymbol{\Sigma}$ is a pair: $(R, \quad \rho : \Sigma(R) \to R)$

monad	:=	endofunctor with substitution
module over a monad	:=	endofunctor with substitution
module morphism	:=	natural transformation preserving substitution

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Substitution and monads

Reminder:

- $B(X)$ = expressions built out of 0, ★ and variables taken in X
- Variables induce a natural transformation $\eta : \mathrm{Id}_{\mathrm{Set}} \to B$

substitution $\mathrm{bind} : B(X) \to (X \to B(Y)) \to B(Y)$ subject to satisfy some equations.

A triple (B, η, bind) is called a **monad**.

A **monad morphism** between two monads R and S is a family of maps $(f_X : R(X) \to S(X))_X$ preserving variables and substitution.

Preview: Operations are module morphisms

$\star$ commutes with substitution

$$(t \star u)[x \mapsto v_x] = t[x \mapsto v_x] \star u[x \mapsto v_x]$$

In the right hand side, substitution acts on a pair of expressions.

Categorical formulation

$B \times B$ supports B-substitution $\rightsquigarrow$ $B \times B$ is a **module over $\boldsymbol{B}$**

$\star$ commutes with substitution $\rightsquigarrow$ $\star : B \times B \to B$ is a **module morphism**

Building blocks for binding signatures

Essential constructions of **modules over a monad $\boldsymbol{R}$**:

- R itself
- M x N for any modules M and N (in particular, R x R)
- The **derivative of a module $\boldsymbol{M}$** is the module M' defined by $M'(X) = M(X + \{\bullet\})$.

The derivative is used to model an operation binding a variable

(Cf next slide).

Syntactic operations are module morphisms

A **module morphism** between two modules M and N on the same monad R is a family of maps $(f_X{:}M(X) \to N(X))_X$ commuting with substitution.

$id_M : M \to M$

the family of identity maps $(id_{M(X)}{:}M(X) \to M(X))_X$ for any module M

$\star : B \times B \to B$

$app : L \times L \to L$

the application operation of the lambda calculus monad L.

$abs : L' \to L$

Indeed, in $\lambda x.t$, the term *t* depends on an additional free variable x:

If $t \in L(Y + \{x\})$ **= $\boldsymbol{L'(Y)}$**, then $abs(t){=}\lambda x.t \in L(Y)$

The Big Picture again

A **signature** Σ is a functorial assignment:

$$\underset{\text{monad}}{R} \mapsto \Sigma(\underset{\text{module over } R}{R})$$

A **model of Σ** is a pair:

$$(\underset{\text{monad}}{R}, \quad \underset{\text{module morphism}}{\rho} : \Sigma(R) \to R)$$

A **model morphism** $m : (R,\rho) \to (S,\sigma)$ is a monad morphism commuting with the module morphism:

$$\begin{tikzcd} \Sigma(R) \arrow[r, "\rho"] \arrow[d, "\Sigma(m)"'] & R \arrow[d, "m"] \\ \Sigma(S) \arrow[r, "\sigma"'] & S \end{tikzcd}$$

Syntax

Definition

Given a signature Σ, its **syntax** is an initial object in its category of models.

Question: Does the syntax exist for every signature?

Answer: No.

Counter-example: the signature $R \mapsto \mathscr{P} \circ R$

powerset endofunctor on $\mathbb{S}\mathrm{et}$

Examples of signatures generating syntax

- $R \mapsto 1 + R \times R$

  Model: (R equipped with module morphisms $1 \to R$

  and $R \times R \to R$.

  The syntax is our previous $(0,\bigstar)$ language.

- $R \mapsto R \times R + R'$

  Models are monads R equipped with two modules morphisms:

  $R \times R \to R$ and $R' \to R$.

  The syntax is lambda calculus.

Algebraic signatures

More generally, the syntax exists for any signature induced by a disjoint sum of products of finite derivatives of the monad ($R \mapsto R' \times R'' \times R''' + R \times R'' \times R''' \times R + \ldots$).

We call such a signature an **algebraic signature**. They correspond to binding signatures through the inclusion:

Binding signatures $\hookrightarrow$ Endofunctors with strength $\hookrightarrow$ Our signatures

Our main result: Quotients of algebraic signatures generate a syntax.

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1. Binding signatures and their models

2. Signatures and models based on monads and modules

3. Our main result

- Definition of presentable signatures
- Generated syntax for presentable signatures
- Examples of presentable signatures

Quotient of a signature

Quotient of a set:

A quotient of a set X is a set Y together with a surjection $p : X \to Y$.

$$x \sim x' \qquad \Longleftrightarrow \qquad p(x) = p(x')$$

Quotient of a signature:

A quotient of a signature Σ consists of:

- a signature Ψ

- a (natural) family of surjective module morphisms $(f_R : \Sigma(R) \to \Psi(R))_R$

$$R \mapsto \begin{array}{c} \Sigma(R) \\ \Big\downarrow f_R \\ \Psi(R) \end{array}$$

Syntax for presentable signatures

Definition

A **presentable signature** is a quotient of an algebraic signature.

Theorem

Any presentable signature generates a syntax.

Question: Are there interesting examples of presentable signatures?

Answer:

- Symmetric operations
- Explicit substitution
- Coherent fixed point operation
- ...

Example 1: Symmetric operations

Binary commutative operation +:

$$t + u = u + t$$

As a quotient of an algebraic signature:

$$R \mapsto \begin{array}{c} R \times R \\ \Big\downarrow\!\!\downarrow \\ R \times R \,/\, \{(x,y) \sim (y,x)\} \end{array}$$

This generalizes to **n-ary permutation invariant operations**.

Example 2: Explicit substitution

An operation $\_\langle x_i \mapsto t_i \rangle$ satisfying coherence equations:

- invariance under **permutation**

$$F(x,y)\langle x \mapsto t, y \mapsto u\rangle = F(y,x)\langle x \mapsto u, y \mapsto t\rangle$$

- invariance under **weakening**

$$F(x)\langle x \mapsto t, y \mapsto u\rangle = F(x)\langle x \mapsto u\rangle$$

- invariance under **contraction**

$$F(x,y)\langle x, y \mapsto t\rangle = F(x,x)\langle x \mapsto t\rangle$$

Example 2: Explicit substitution

Signature of explicit substitution as a quotient of the algebraic signature Σ:

$$\Sigma(R) := R'\times R \;+\; R''\times R\times R \;+\; R'''\times R\times R\times R \;+\; \ldots$$

$$t\langle x \mapsto u\rangle \qquad\qquad t\langle x \mapsto u, y \mapsto v\rangle \qquad\qquad t\langle x \mapsto u, y \mapsto v, z \mapsto w\rangle$$

$$R \mapsto \begin{tikzcd} \Sigma(R) \arrow[d, two heads] \\ \Sigma(R)\,/\sim \end{tikzcd}$$

- **permutation**: $t\langle x \mapsto u, y \mapsto v\rangle \sim t[x \rightleftarrows y]\langle x \mapsto v, y \mapsto u\rangle$
- **weakening**: $t\langle x \mapsto u\rangle \sim t\langle x \mapsto u, y \mapsto v\rangle$
- **contraction**: $t\langle x \mapsto u, y \mapsto u\rangle \sim t[y := x]\langle x \mapsto u\rangle$

Conclusion

Summary of the talk:

- presented a notion of signature and models
- identified a class of signatures that generate a syntax
  - encompasses the classical binding signatures
  - encompasses operations satisfying some equations

Future work:

- add equations (e.g. lambda calculus modulo beta/eta equivalence);
- extend our framework to simply typed syntaxes.

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Thank you!