A categorical diagram editor to help formalising commutation proofs

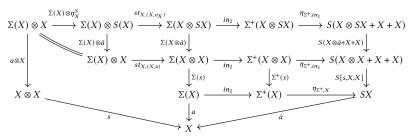
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Motivation

• Diagrammatic reasoning is useful



- × Theorem provers (e.g., Coq) are text-based
- ⇒ This talk: a tool to bridge the gap

Yet Another Diagram Editor (YADE)

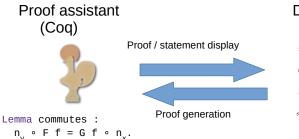


Diagram editor (YADE)



More about YADE

- Available online¹: run by your browser
- Written in **Elm** (~6000 LoC)



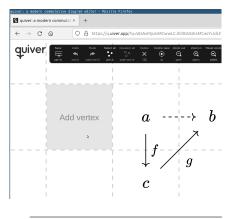
"A delightful language for reliable web applications."

- Functional PI
- Compiles to js

¹https://amblafont.github.io/graph-editor/index.html

Comparison with Quiver

"quiver¹ is a modern, graphical editor for commutative and pasting diagrams, capable of rendering high-quality diagrams for screen viewing, and exporting to LaTeX via tikz-cd."



Quiver features:

- More display options (colors, ...)
- Rigid grid layout
- Export to LaTeX
- × No proof generation

Note: YADE exports to quiver

¹https://q.uiver.app/

This talk

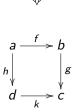
- Short presentation of YADE's features that help formalisation
- Demo

- Parse equations
- 2 Forward reasoning
- Backward reasoning
- Demo

- Parse equations
- 2 Forward reasoning
- Backward reasoning
- 4 Demo

A simple format for equations

$$a -- f -> b -- g -> c = a -- h -> d -- k -> c$$

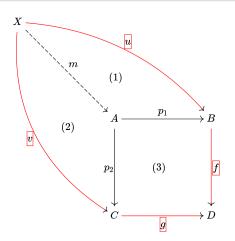


Load an equation from Coq

Specific Coq notations to comply to the above format

- Parse equations
- 2 Forward reasoning
- Backward reasoning
- 4 Demo

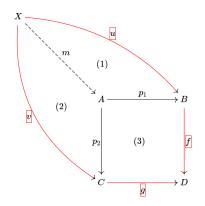
Diagrams are proofs



Generated formal proof

If inner subdiagrams (1)-(3) commute, then the outer diagram commutes.

Algorithm for proof generation



- Identify all subdiagrams
- Start from one branch of the outer diagram
- Repeatedly "apply" subdiagrams to progress, until reaching the other branch.

Target: the UniMath mathematical library

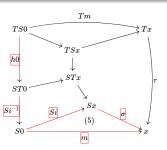
History

- started in 2014 with Voevodsky's repository Foundations
- Large Coq mathematical library (~300 000 lines)
 - ~ two thirds on (bi)category theory
 - Verbose style
- Inconsistent! (Type: Type)

Why targeting this library?

- I use UniMath in my formalisation projects.
- The implementation could be easily adapted to another target.

Commutation proofs in UniMath

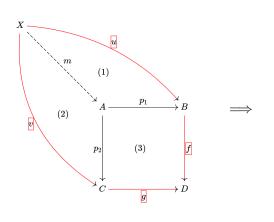


Applying¹ (5) to the bottom left branch

$$\begin{array}{c} (h_0 \cdot Si^{-1}) \cdot m \\ h_0 \cdot (Si^{-1} \cdot m) \\ Si^{-1} \cdot m \\ m \\ Si \cdot \sigma \end{array} \begin{array}{c} rewrite \ assoc' \\ apply \ cancel_precomposition \\ apply \ cancel_precomposition \\ apply \ (5) \end{array}$$

¹rewrite (5) sometimes works but is hard to maintain.

Generated Coq script



```
Goal \{ u \cdot f = v \cdot g \}.
assert(eq : { u = m \cdot p_1 }).
etrans.
  apply cancel_postcomposition.
  apply eq.
clear eq.
assert(eq : \{ p_1 \cdot f = p_2 \cdot g \}).
etrans.
  repeat rewrite assoc'.
  apply cancel precomposition.
  repeat rewrite assoc.
  apply eq.
repeat rewrite assoc.
clear eq.
assert(eq : { m \cdot p_2 = v }).
etrans.
  apply cancel_postcomposition.
  apply eq.
clear eq.
apply idpath.
Qed.
```

Commutation of subdiagrams are explicitly asserted and admitted.

- Parse equations
- 2 Forward reasoning
- Backward reasoning
- 4 Demo

Forward vs Backward reasoning in Coq

 $n: G \Rightarrow F$ natural.

How to rewrite $n_x \cdot Ff$ into $Gf \cdot n_y$?

Forward reasoning in Coq

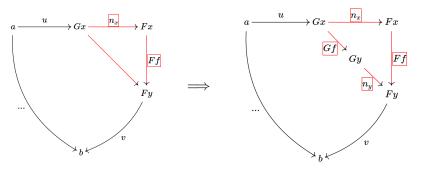
- Explicitly state the equation
- 2 Prove it by naturality of n
- Apply the equation

Backward reasoning in Coq

- Focus on $n_X \cdot Ff$
- Apply naturality of n
- \Rightarrow No need to provide the r.h.s $Gf \cdot n_y$

Backward reasoning with YADE

- Generate Coq script that focuses on the subterm
- Manually apply the lemma, in Coq
- Open and paste the result in YADE
- ⇒ The diagram is completed automatically



- Parse equations
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Demo: Associativity of the "composite" multiplication

Let

- $\delta: ST \to TS$ be a monadic distributive law.
- R := TS
- $\mu': RR \to R$ defined in the only sensible way.

Then,

$$\begin{array}{c} RRR \xrightarrow{R\mu'} RR \\ \mu'R \downarrow & \downarrow \mu' \\ RR \xrightarrow{\mu'} R \end{array}$$