

# A diagram editor to mechanise categorical proofs

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# Naming convention

## Yet Another Diagram Editor

I will refer to my editor as **YADE**, or Coreact-**YADE**.

ANR Project<sup>1</sup> (2023 - 2027): Coq-based Rewriting: Towards Executable Applied Category Theory

<sup>1</sup> <https://coreact.wiki/>

# How to test the editor

A web app that runs locally in your browser

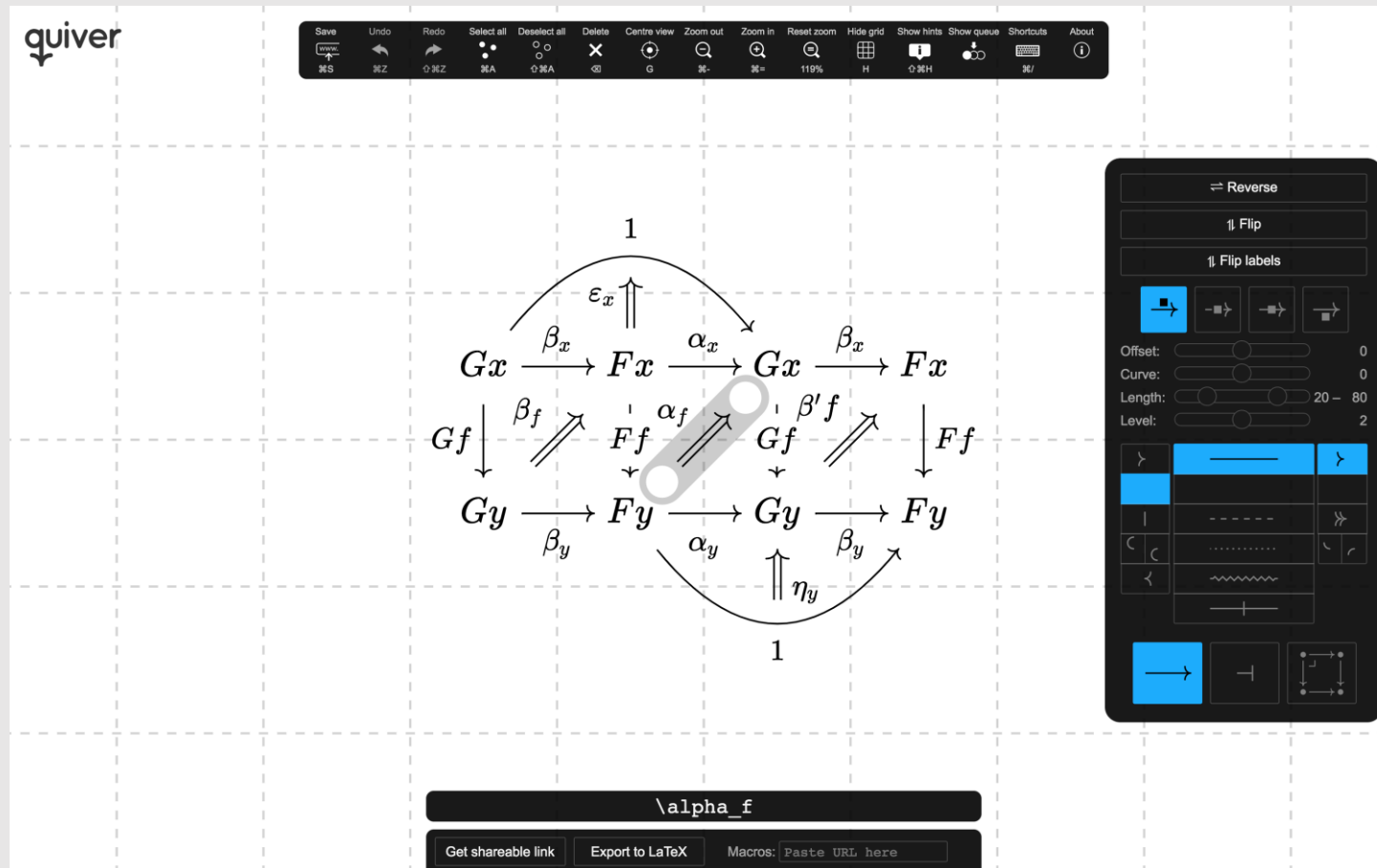
<https://amblafont.github.io/graph-editor/index.html>

A standalone desktop program

- Embeds the web app using electron
- Additional features (mechanisation)

# Related software: Quiver

*“a modern, graphical editor for commutative and pasting diagrams, capable of rendering high-quality diagrams for screen viewing, and exporting to LaTeX via tikz-cd.”*



# Comparison with quiver

About the same size (around 10k of LoC)

	Quiver	YADE
Programming Languages	<div><div>Languages</div><div><div></div><div></div><div></div><div></div></div><div>JavaScript 90.7%   CSS 5.3%</div><div>TeX 2.4%   Other 1.6%</div></div>	<div><div>Languages</div><div><div></div><div></div><div></div><div></div><div></div><div></div></div><div>Elm 72.6%   HTML 12.7%</div><div>TeX 8.5%   TypeScript 4.7%</div><div>JavaScript 1.5%</div></div>
Styling options	+	-
User-friendly	+	-
Editing features	-	+ Tabs, copy & paste, find & replace, expand selection to connected components, ...
LaTeX export	yes	yes <sup>1</sup>

<sup>1</sup> Implemented by Tom Hirschowitz

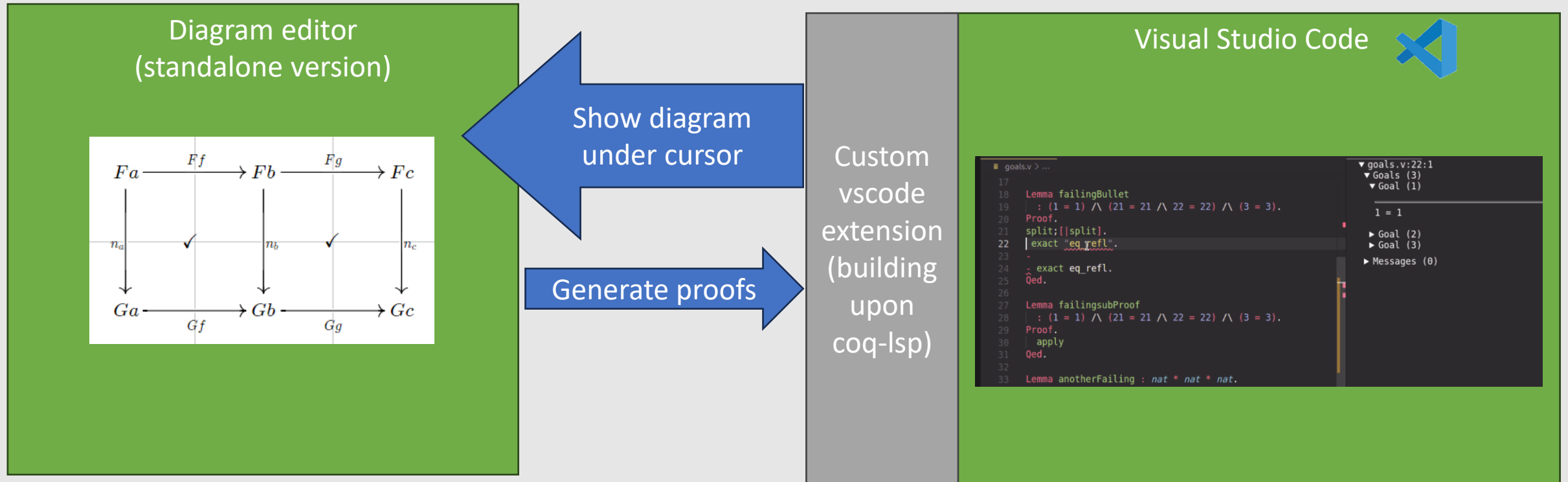
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Mechanisation features	-	+

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# Architecture



(+ Coq library for custom notations)

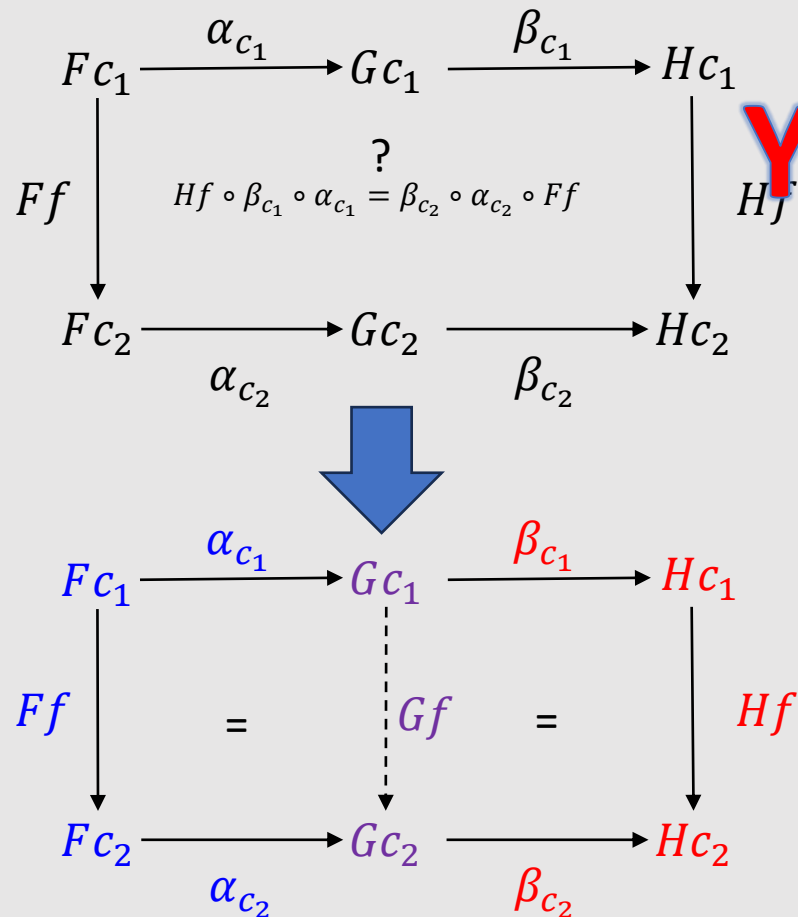
# Natural transformations compose:

If  $\alpha: F \Rightarrow G$  and  $\beta: G \Rightarrow H$  are natural, then so is  $(\beta_c \circ \alpha_c: Fc \rightarrow Hc)_c$

Diagrammatic proof

Automatic generation?

Computer-friendly proof



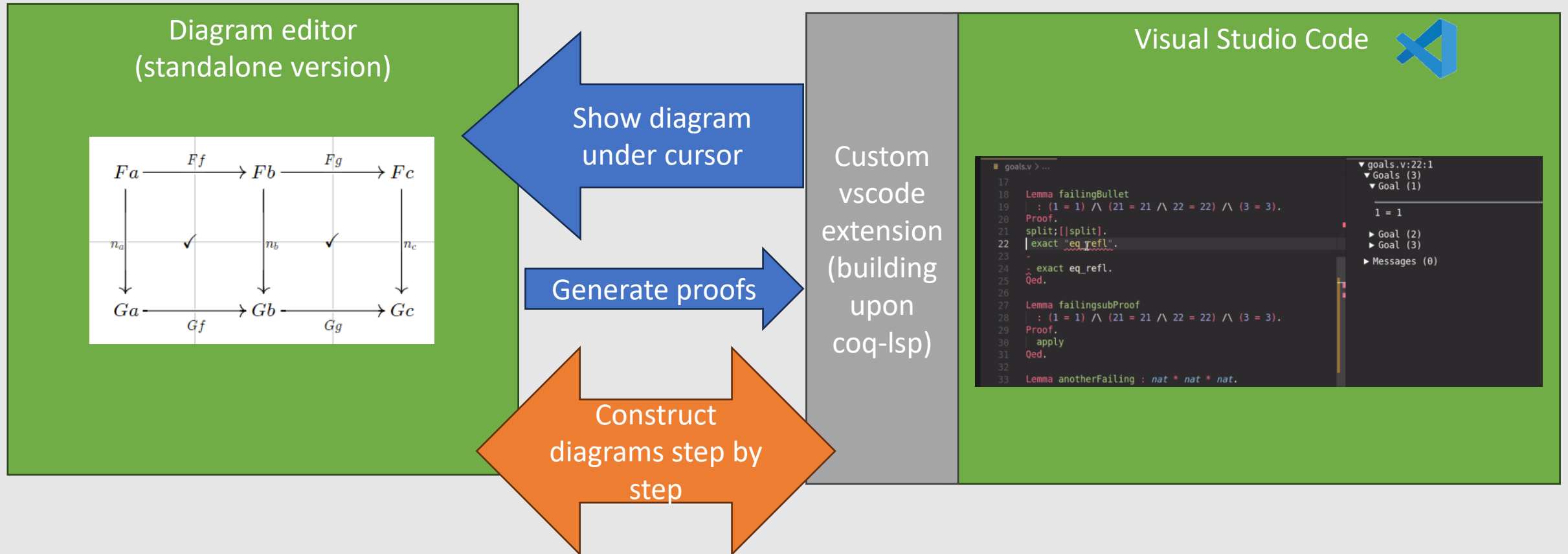
YADE



$$\begin{aligned}
 & Hf \circ \beta_{c_1} \circ \alpha_{c_1} \\
 &= \beta_{c_2} \circ Gf \circ \alpha_{c_1} \\
 &= \beta_{c_2} \circ \alpha_{c_2} \circ Ff
 \end{aligned}$$



# Architecture



(+ Coq library for custom notations)

# Building the diagrammatic proof interactively

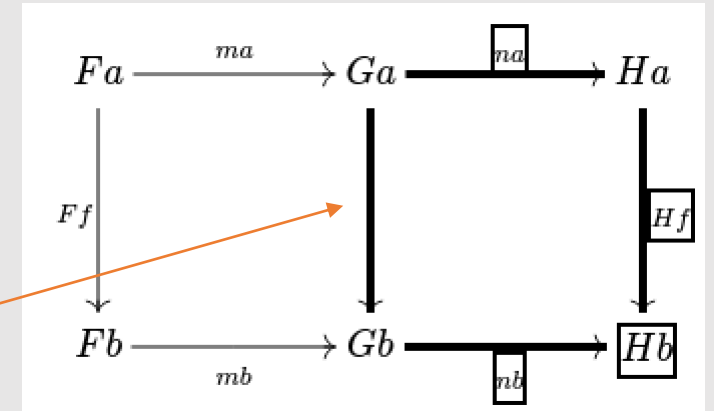
1) Select a subdiagram

2) Create a proof node, labelled with the Coq tactic `naturality`.

⇒ Coq (in vscode) checks that this tactic solves the goal:

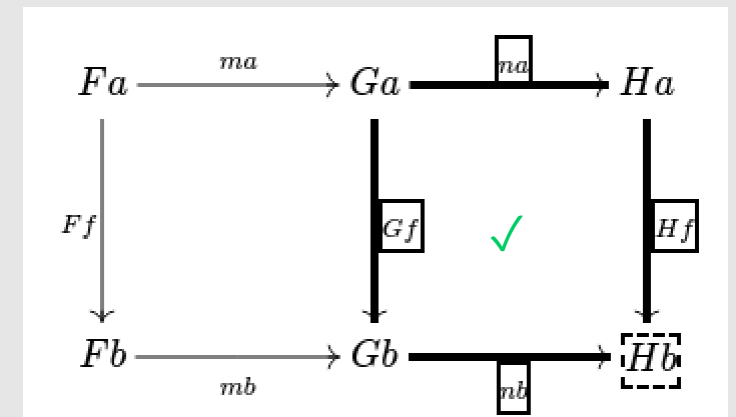
$$Hf \circ n_a = nb \circ _$$

*unnamed arrow*

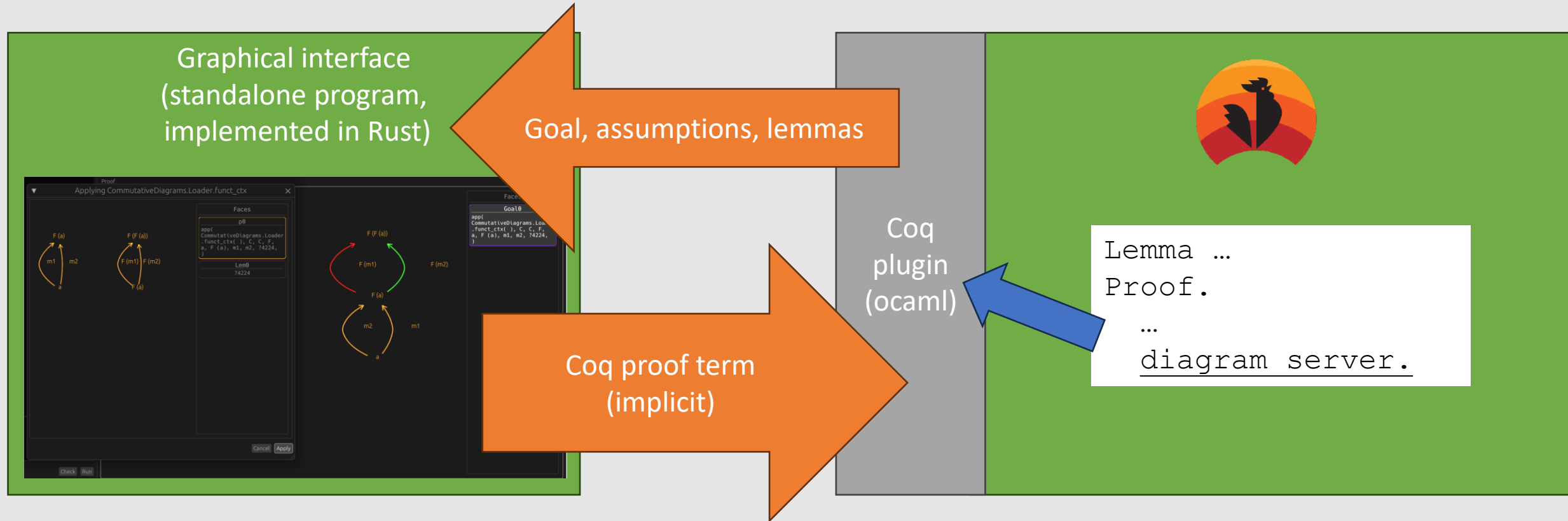


⇒ The diagram gets completed in YADE:

- The *unnamed arrow* is instantiated by Coq's inferred instantiation
- The proof node is marked as validated (indicated by a green ✓)



# Another related software for mechanisation: Luc Chabassier's interface<sup>1</sup> for diagrammatic proofs



- Automatic layout (limited manual editing)
- Proof tactics (including invoking a lemma)

<sup>1</sup><https://github.com/dwarfmaster/commutative-diagrams>

# Demo<sup>1</sup> of YADE

(Based on the category theory library of Hierarchy Builder + custom tactics & notations)

A distributive law  $\delta: TS \Rightarrow ST$  between two monads  $S$  and  $T$  induces a monad structure on  $ST$ .

Let us show that the induced multiplication  $STST \xrightarrow{S\delta T} SSTT \xrightarrow{\mu^S \mu^T} ST$  is associative.

<sup>1</sup> <https://github.com/amblafont/vscode-yade-example>

# Proof generation: sketch of the algorithm

1) Save all subdiagrams as rewrite rules: “top right branch” → “bottom left branch”

$$(1) u \rightarrow p_1 \circ m$$

$$(2) p_2 \circ m \rightarrow v$$

$$(3) f \circ p_1 \rightarrow g \circ p_2$$

2) Identify the top right branch of the outer diagram.

$$f \circ u$$

3) Use greedily the rewrite rules until reaching the bottom left branch

$$f \circ u \xrightarrow{(1)} f \circ p_1 \circ m \xrightarrow{(3)} g \circ p_2 \circ m \xrightarrow{(2)} g \circ v$$

4) Remember the rewrite steps and generate the coq proof script accordingly.

If a Coq proof is provided inside a subdiagram, use it to justify the rewrite step.

