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pyADCG

A preliminary implementation of a new parallel solver
for nonconvex MINLPs in Pyomo/Python

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Introduction

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- standard global solution approach for MIPs and MINLPs
- **However:** many branching steps necessary for large MINLPs
(Benchmark Vigerske 2015, most problems have < 1000 variables)

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ADCG (Alternating Direction Column Generation Method):

- exact parallel MINLP decomposition method
- combines AD and CG without using branching

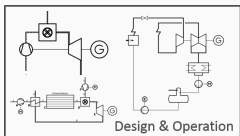
ADCG solves modular optimization problems

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Planning (energy systems)



Engineering design (glider, Airbus)

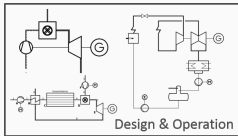


sub-problems: MINLPs

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sub-problems: MINLPs

Transport optimization (crew scheduling)



sub-problems: constrained shortest path

The quasi-separable MINLP

Nonconvex **quasi-separable** MINLP:

$$\min\{c^T x : x \in P \cap G\} \quad \text{with} \quad G := \bigtimes_{k \in K} G_k$$

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where G_k defines the feasible set of a (MINLP) **sub-problem**:

$$G_k := \{y \in [\underline{x}_{I_k}, \bar{x}_{I_k}] : y_i \in \{0, 1\}, i \in I_k^{\text{int}}, g_j(y) \leq 0, j \in J_k\}$$

and $P := \{x \in [\underline{x}, \bar{x}] : Ax \leq b\}$ defines the **coupling** constraints

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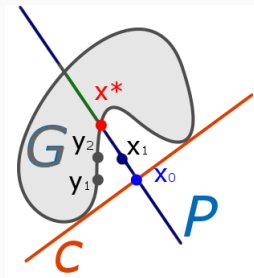
(MINLPs can be **transformed** into quasi-separable MINLPs by adding new variables and constraints)

Assumption: MINLP **sub-problems** can be solved **quickly**

The Alternating Direction Method

Basic steps of an ADM for solving $\min\{c^T x : x \in P \cap \bigcap_{k \in K} G_k\}$:

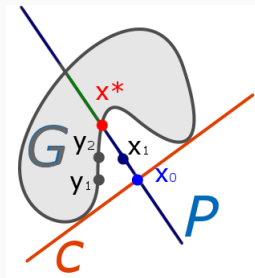
1. $y^{i+1} \leftarrow \text{G-project } x^i \in P$
regarding $c + A^T \lambda^i$
2. $x^{i+1} \leftarrow \text{P-project } y^{i+1} \in G$
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where

- **G-project** by (parallel) solving MINLP sub-problems
- **P-project** by solving a (large) QP master-problem
- the starting point $x^0 \in P$ is the solution of a **convex relaxation**

Traditional ADMs do **not converge** always towards the solution point x^* .

The Alternating Direction Column Generation Method

(Global optimization without branching)

Penalty function, full-projection problem, global convergence

Motivation: Houska, Frasch and Diehl, *An Augmented Lagrangian based Algorithm for Distributed Non-Convex Optimization* (2014), **ALADIN**:

- Check quality of x^i using the **penalty function**

$$\text{PenF}(x^i) := c^T x^i + \gamma \cdot \text{Viol}(x^i, P \cap G)$$

$$(\rightarrow \min \text{PenF}(x) \Leftrightarrow \text{MINLP})$$

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- If $\text{PenF}(x^i) - \text{PenF}(x^{i+1})$ is not large enough:
solve approximately (via a dual line search)
a **full-projection problem** ($x^i \in P \rightarrow P \cap G$)

$$\begin{aligned} x^{i+1} = \text{approx argmin} \quad & c^T x + \rho \sum_{k \in K} \|x_{l_k} - x_{l_k}^i\|^2 \\ \text{s.t.} \quad & x \in G \cap P \end{aligned}$$

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Theorem: If ρ is large, then the full-projection problem has a **zero duality gap** \rightarrow **global convergence** of ALADIN

ADCG: target constraint and CG-based full-projection

- Use a target constraint: $P' = \{x \in P : c^T x \leq \text{tarVal}\}$
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CG generates **sample points** $S_k \subseteq G_k$ (inner approximation)
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- ADCG performs **CG**-steps, until
 - (i) **PenF** is **reduced** sufficiently or
 - (ii) $s + \text{redCost}(\lambda) > 0 \quad \Rightarrow P' \cap G = \emptyset \Leftrightarrow \text{tarVal} < v^*$

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Theorem: x^i converges towards a **global solution** of the MINLP
(**without branching**)

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For details of ADCG see:

1. Nowak. *Column Generation based Alternating Direction Methods for solving MINLPs*. 2015 (Optimization Online)
2. Nowak, Breielfeld, *The Alternating Direction Column Generation Method (Global Optimization without Branching)* (in preparation)

Preliminary results with pyADCG

Implementation:

- object-oriented implementation in Python/Pyomo
- development started in beginning of 2016
- not finished, currently only for nonconvex QQPs

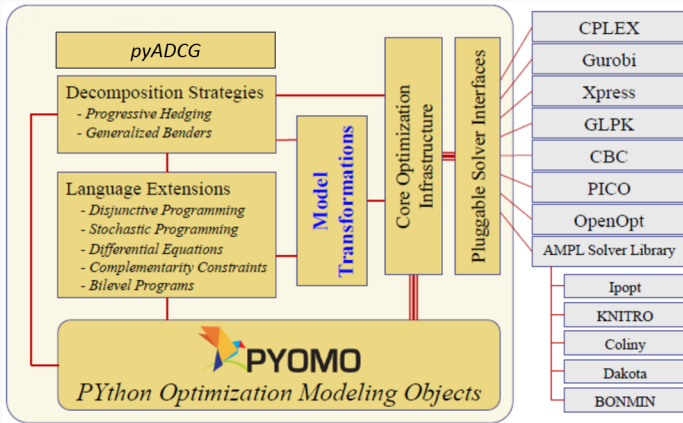
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pyADCG is a [meta-solver](#) consisting of the [sub-solvers](#):

- [IpOpt](#) for master-problems and local optimization
- [SCIP](#) for MINLP sub-problems

What is Pyomo?



Pyomo provides a **modeling language** and supports the development of **meta-solvers** based on sub-solvers via **Python** (→ fast development)

Preliminary numerical results with random QQPs

(without parallelization):

- very few ADKG iterations (maybe because duality gap is small)
- time for solving master-problems 1%
- time for solving sub-problems 94%

⇒ significant potential for [parallel solving](#) the sub problems

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Planned:

- solving [MINLPlib2](#) models (translated to Pyomo) in development
- automatic decomposition and other [enhancements](#)
- [application](#) to energy conversion systems

Questions?