
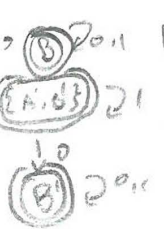


NFA \rightarrow  lang $(a^*b)^*$
 DFA \rightarrow  lang $(a^*b)^*$
 $NFA = Q$
 $DFA = Q = P(Q)$
 $\delta(CB, 1) = \delta(Cb, 1) = B'$
 $\delta(CAB', 1) = \delta(CA, 1) \cup \delta(CB, 1)$
 $PUB' = B'$

CS Theory cheat sheet Exam 1

Finite Automata $(Q, \Sigma, \delta, q_0, F)$ - FA

Q - Set of states

Σ - Alphabet

δ - Transition function $Q \times \Sigma \rightarrow Q$ NFA: $Q \times \Sigma \rightarrow P(Q)$ $\Sigma_a = \Sigma \cup \{ \epsilon \}$

q_0 - Start state

F - Accept state

Reg Lang - Lang Desc by FA - RL

$A \cup B = \{ x \mid x \in A \text{ or } x \in B \}$

$A \cdot B = \{ xy \mid x \in A \text{ and } y \in B \}$

If A_1 and A_2 Reg then $A_1 \cup A_2$ is Reg

Deterministic: given current state and input know next state

Lang is Reg iff Desc. by Reg X

Gen NFA ① Arrows desc w/ string ② Start \rightarrow any other state, no \rightarrow to start

③ Accept state \leftarrow from all other states no outgoing \rightarrow ④ except $\neq q_{start}$

⑤ as a graph NFA is complete

PFA \rightarrow GNFA \rightarrow Reg Ex

\rightarrow  \Rightarrow  \Rightarrow  \Rightarrow  \Rightarrow \Rightarrow