University of Toronto
Faculty of Arts and Science

CSC165H1S Term Test 1, Version 3

Date: February 13, 2023 **Duration**: 75 minutes **Instructor(s)**: G. Baumgartner, T. Fairgrieve

No Aids Allowed

Given Name(s):

Family Name(s):

Student Number (9 or 10 digits):

UTORid (e.g., `pitfra12`):

- This test has **4** questions. There are a total of **10 pages, DOUBLE-SIDED**.
- Final expressions in predicate logic must have negation symbols ($\neg$) applied **only** to predicates or propositional variables, e.g., $\neg Prime(x)$ or $\neg p$.
- You may not define your own propositional operators, predicates, or sets unless asked for in the question.
- In your proofs, you may always use definitions we have covered in this course. However, you may **not** use any external facts about these definitions unless they are given in the question.
- You may **not** use proofs by induction on this test.
- A list of standard equivalences, predicate definitions and sets is given on Page 2.

Question	Q1	Q2	Q3	Q4	**Total**
Mark					
Out of	10	8	5	11	**34**

On this test you may use the following standard equivalences, predicate definitions and sets. Read $\iff$ as 'is equivalent to'.

Standard Equivalences (p, q, r, $P(x)$, $Q(x)$, *etc.* are arbitrary sentences. D is an arbitrary set)

- *Commutativity*
  $p \land q \iff q \land p$
  $p \lor q \iff q \lor p$
  $p \Leftrightarrow q \iff q \Leftrightarrow p$
- *Associativity*
  $p \land (q \land r) \iff (p \land q) \land r$
  $p \lor (q \lor r) \iff (p \lor q) \lor r$
- *Identity*
  $p \land (q \lor \neg q) \iff p$
  $p \lor (q \land \neg q) \iff p$
- *Absorption*
  $p \land (q \land \neg q) \iff q \land \neg q$
  $p \lor (q \lor \neg q) \iff q \lor \neg q$
- *Idempotency*
  $p \land p \iff p$
  $p \lor p \iff p$
- *Double Negation*
  $\neg\neg p \iff p$
- *Negation rules*
  $\neg(p \land q) \iff \neg p \lor \neg q$
  $\neg(p \lor q) \iff \neg p \land \neg q$
  $\neg(p \Rightarrow q) \iff p \land \neg q$
  $\neg(p \Leftrightarrow q) \iff ((p \land \neg q) \lor (\neg p \land q))$
- *Distributivity*
  $p \land (q \lor r) \iff (p \land q) \lor (p \land r)$
  $p \lor (q \land r) \iff (p \lor q) \land (p \lor r)$
- *Contrapositive*
  $p \Rightarrow q \iff \neg q \Rightarrow \neg p$
- *Implication*
  $p \Rightarrow q \iff \neg p \lor q$
- *Biconditional*
  $p \Leftrightarrow q \iff (p \Rightarrow q) \land (q \Rightarrow p)$
- *Renaming*
  (where $P(x)$ does not contain variable y)
  $\forall x \in D, P(x) \iff \forall y \in D, P(y)$
  $\exists x \in D, P(x) \iff \exists y \in D, P(y)$
- *Quantifier Negation*
  $\neg\forall x \in D, P(x) \iff \exists x \in D, \neg P(x)$
  $\neg\exists x \in D, P(x) \iff \forall x \in D, \neg P(x)$
- *Quantifier Commutativity*
  $\forall x \in D, \forall y \in D, S(x,y) \iff$
  $\quad \forall y \in D, \forall x \in D, S(x,y)$
  $\exists x \in D, \exists y \in D, S(x,y) \iff$
  $\quad \exists y \in D, \exists x \in D, S(x,y)$
- *Quantifier Distributivity*
  (where S does not contain variable x)
  $S \land \forall x \in D, Q(x) \iff \forall x \in D, S \land Q(x)$
  $S \lor \forall x \in D, Q(x) \iff \forall x \in D, S \lor Q(x)$
  $S \land \exists x \in D, Q(x) \iff \exists x \in D, S \land Q(x)$
  $S \lor \exists x \in D, Q(x) \iff \exists x \in D, S \lor Q(x)$
  $(\forall x \in D, P(x)) \land (\forall x \in D, Q(x)) \iff$
  $\quad \forall x \in D, P(x) \land Q(x)$
  $(\exists x \in D, P(x)) \lor (\exists x \in D, Q(x)) \iff$
  $\quad \exists x \in D, P(x) \lor Q(x)$

Standard Predicate Definitions

- $Even(n)$: "$\exists k \in \mathbb{Z},\ n = k \cdot 2$", where $n \in \mathbb{Z}$
- $Odd(n)$: "$\exists k \in \mathbb{Z},\ n = k \cdot 2 + 1$", where $n \in \mathbb{Z}$
- $d \mid n$: "$\exists k \in \mathbb{Z},\ n = k \cdot d$", where $d, n \in \mathbb{Z}$
- $Prime(n)$: "$n > 1 \land \big(\forall d \in \mathbb{N},\ d \mid n \Rightarrow (d = 1 \lor d = n)\big)$", where $n \in \mathbb{N}$
- $Atomic(n)$: "$\forall a, b \in \mathbb{N},\ (n \nmid a \land n \nmid b) \Rightarrow n \nmid ab$", where $n \in \mathbb{N}$

Standard Sets

- The set of natural numbers, $\mathbb{N} = \{0, 1, 2, \ldots\}$.
- The set of integers, $\mathbb{Z} = \{\ldots, -2, -1, 0, 1, 2, \ldots\}$. The set of positive integers, $\mathbb{Z}^+ = \{1, 2, 3, \ldots\}$.
- The set of real numbers, $\mathbb{R}$. The set of positive real numbers, $\mathbb{R}^+$.

1. **[10 marks] Short answer questions.**

   (a) **[2 marks]** Let A, B and C be different **non-empty** sets and consider the (incorrect) claim:

   $$(A \cap B) \cup C = (A \cap C) \cup B$$

   Show that this claim is not correct. Use suitably constructed sets A, B and C.

   Solution

   Let $A = \{1\}$, $B = \{2\}$ and $C = \{3\}$.
   Then $A \cap B = \varnothing$ and $(A \cap B) \cup C = \{3\}$.
   But $A \cap C = \varnothing$ and $(A \cap C) \cup B = \{2\}$, which is different from $\{3\}$.
   And so, $(A \cap B) \cup C \neq (A \cap C) \cup C$.

   (b) **[3 marks]** Using a truth table, prove the equivalence rule:

   $$\neg(p \Rightarrow q) \text{ is equivalent to } (p \wedge \neg q)$$

   You must include columns with intermediate results to demonstrate how you determined the rows of the table.

   Solution

p	q	$(p \Rightarrow q)$	$\neg(p \Rightarrow q)$	$\neg q$	$p \wedge \neg q$	$\neg(p \vee q)$ is equivalent to $\neg(p \vee q)$
True	True	True	False	False	False	True
True	False	False	True	True	True	True
False	True	True	False	False	False	True
False	False	True	False	True	False	True

   Note that since the fourth and sixth columns are the same, it follows that $\neg(p \Rightarrow q)$ is equivalent to $(p \wedge \neg q)$.

(c) **[1 mark]** Rewrite the summation $\sum_{i=3}^{13}(3i+1)$ as an equivalent summation where the index of the summation has lower bound 0. Do **not** determine the numerical value of the resulting sum.

Solution

$\sum_{j=0}^{10}(3j+10)$.

(d) **[2 marks]** Is the statement below True or False? Justify your response. You may make use any facts that you have seen in the course about the predicate $\mid$.

$$\exists m \in \mathbb{N}, \forall n \in \mathbb{N},\ n > 1 \Rightarrow (m < n \wedge m \mid n)$$

Solution

This statement is True. Consider $m = 1$. For any $n \in \mathbb{N}$ that is greater than 1, we know that $1 < n$ and that 1 divides n, as we know that 1 divides every natural number.

(e) **[2 marks]** In the following, $f : \mathbb{N} \to \mathbb{R}^+$ and $g : \mathbb{N} \to \mathbb{R}^+$ are functions. The set $\mathbb{R}^+$ is the set of positive real numbers.

Write the negation of the following statement:

$$\exists c \in \mathbb{R}^+, \exists n \in \mathbb{R}^+, \forall m \in \mathbb{N},\ m \geq n \Rightarrow g(m) \leq c \cdot f(m)$$

Solution

$$\forall c \in \mathbb{R}^+, \forall n \in \mathbb{R}^+, \exists m \in \mathbb{N},\ m \geq n \wedge g(m) > c \cdot f(m)$$

2. [**8 marks**] **Translations.** Let W be the set of all positive integers, and recall the following terms:

- A positive integer is said to be even when it is divisible by 2.
- A positive integer is said to be odd when it is not divisible by 2.
- A positive integer that can be formed by multiplying two smaller positive integers is said to be a composite number.

Suppose we define the following predicates:

- $Even(x)$: "x is even", where $x \in W$.
- $Odd(x)$: "x is odd", where $x \in W$.
- $Composite(x)$: "x is a composite number", where $x \in W$.

In Parts (a) – (e), translate each of the following statements from English into symbolic predicate logic. No explanation is necessary. Do not define any of your own predicates or sets. You may use the predicates $<, \leq, =, >, \geq$ and $\neq$ to compare two numbers.

(a) [**1 mark**] The positive integer 165 is an odd composite number.

Solution

$Composite(165) \wedge Odd(165)$

It is okay to write $165 \in W \wedge Composite(165) \wedge Odd(165)$ but it can be taken as a given that 165 is a positive integer.

The statement $\forall x \in W, x = 165 \Rightarrow (Composite(x) \wedge Odd(x))$ is also okay but not concise.

(b) [**1 mark**] Positive integers are either odd, composite or are the number 2.

Solution

$\forall x \in W, Odd(x) \vee Composite(x) \vee x = 2$

(It would be okay to also express the exclusivity of these cases.)

(c) [**2 marks**] Some odd positive integer is composite.

Solution

$\exists x \in W, Odd(x) \wedge Composite(x)$

(d) [**2 marks**] Even positive integers that are composite are greater than 2.

Solution

$\forall x \in W, (Even(x) \wedge Composite(x)) \implies x > 2$

(e) [**2 marks**] The positive integer 9 is the smallest odd composite number.

Solution

$Odd(9) \land Composite(9) \land \forall x \in W, Odd(x) \land Composite(x) \implies x \geq 9$

3. **[5 marks] A proof about numbers.**

(a) **[1 mark]** Translate the following statement into predicate logic:

"There is a positive integer c such that for every positive integer x,
$100x^2 + x + 148 \leq c \cdot (x^6 + 1)$."

Use $\mathbb{Z}^+$ to denote the set of positive integers.

Solution

$\exists c \in \mathbb{Z}^+, \forall x \in \mathbb{Z}^+,\ 100x^2 + x + 148 \leq c \cdot (x^6 + 1)$

(b) **[4 marks]** Prove the statement from Part (a). We have left you space for rough work here but write your formal proof in the box below.

Solution

Proof. Let $c = 1000$ (or some other number greater than or equal to both 100 and 148). Let $x \in \mathbb{Z}^+$.
Then

$$\begin{aligned}
c \cdot (x^6 + 1) &= 1000(x^6 + 1) \\
&= 1000x^6 + 1000 \\
&= 999x^6 + x^6 + 1000 \\
&\geq 999x^6 + x^6 + 148 && \text{(since } 1000 \geq 148\text{)} \\
&\geq 999x^6 + x + 148 && \text{(since } x \geq 1, x^6 \geq x\text{)} \\
&\geq 100x^6 + x + 148 && \text{(since } 999 \geq 100\text{),}
\end{aligned}$$

as required.
(There are many other valid sequences of steps.) □

4. **[11 marks] A proof about multiples of 165.**

(a) **[1 mark]** Define a unary predicate $Multiple165$ with domain $\mathbb{N}$ so that $Multiple165(n)$ means:

n is a multiple of 165.

(That is, n is equal to 165 times a natural number.)

Solution

$Multiple165(n)$: " $\exists k \in \mathbb{N},\ n = 165k$ ", where $n \in \mathbb{N}$.

(b) **[1 mark]** Translate the following statement into predicate logic:

"If n is a natural number and n^2 is **not** a multiple of 165 then n is **not** a multiple of 165."

Do not use the predicate $Multiple165$ in your final statement; use its definition instead.

Solution

For convenience during the later proofs we'll use different variable names here for the two local existentially quantified variables.

$$\forall n \in \mathbb{N},\ \neg(\exists k \in \mathbb{N},\ n^2 = 165k) \Rightarrow \neg(\exists j \in \mathbb{N},\ n = 165j)$$

which is logically equivalent to

$$\forall n \in \mathbb{N},\ (\forall k \in \mathbb{N},\ n^2 \neq 165k) \Rightarrow (\forall j \in \mathbb{N},\ n \neq 165j)$$

(c) **[1 mark]** Write the contrapositive of the statement from Part (b):

Solution

$$\forall n \in \mathbb{N},\ (\exists j \in \mathbb{N},\ n = 165j) \Rightarrow (\exists k \in \mathbb{N},\ n^2 = 165k)$$

(d) [**4 marks**] Write a direct proof of the statement from Part (**c**). (That is, write a direct proof of the contrapositive of the original statement.) We have left you space for rough work here, but write your formal proof in the box below.

Solution

Recall that the statement is:

$$\forall n \in \mathbb{N},\ (\exists j \in \mathbb{N},\ n = 165j) \Rightarrow (\exists k \in \mathbb{N},\ n^2 = 165k)$$

Proof.
Let $n \in \mathbb{N}$.
Assume there is a $j \in \mathbb{N}$ with $n = 165j$.
Let $k = 165 \cdot j^2$, which is in $\mathbb{N}$.
Then $n^2 = (165j)^2 = 165 \cdot 165 \cdot j^2 = 165k$. $\square$

(e) [**4 marks**] Write a direct proof of the statement from Part (**b**).
This proof can be tricky, so be sure to write out the direct proof outline and any comments (e.g. what you want to show in a part of it, what you may use in a part of it) for partial credit.

We have left you space for rough work here and on the last page, but write your formal proof in the box below.

Solution

Recall that the statement is:

$$\forall n \in \mathbb{N},\ (\forall k \in \mathbb{N},\ n^2 \neq 165k) \Rightarrow (\forall j \in \mathbb{N},\ n \neq 165j)$$

Proof.
Let $n \in \mathbb{N}$.
Assume $\forall k \in \mathbb{N},\ n^2 \neq 165k$.
(WTS $\forall j \in \mathbb{N},\ n \neq 165j$)
Let $j \in \mathbb{N}$.
Let $k = 165 \cdot j^2$, which is in $\mathbb{N}$.
Then $n^2 \neq 165k$ (by the assumption), i.e. $n^2 \neq (165j)^2$.
So $n \neq 165j$. $\square$