

## Lab 4: Implementation of Dynamic Programming using C++

### Objective:

To implement and analyze two classic greedy algorithms, 0/1 Knapsack Algorithm and Single Source Shortest Path Problem. These algorithm are chosen to demonstrate the efficiency and practical applications of dynamic programming approach for solving problems.

### Theory:

Dynamic Programming (DP) is an algorithmic technique that solves complex problems by breaking them down into simpler subproblems. Its particularly useful for optimization problems.

1. 0/1 Knapsack Problem: The 0/1 Knapsack Problem is a classic optimization problem where we need to select items to maximize value while staying within a weight limit. Unlike the Fractional Knapsack Problem, items cannot be divided, they are either excluded or included.

Algorithm:

Input : array of items with value and weight

Output: Maximum value

- a.  $N :=$  Length of items.
  - b. Create a 2d array  $dp[n+1][max\_weight+1]$  initialized with all 0.
  - c. For  $i$  from 1 to  $n$ 
    - i. For  $w$  from 0 to  $max\_weight$ 
      1. If  $items[i-1].weight \leq w$  then
      2.  $Dp[i][w] := \max(dp[i-1][w], dp[i-1][w-items[i-1].weight] + items[i-1].value)$
      3. Else
      4.  $dp[i][w] := dp[i-1][w]$
  - d. return  $dp[n][W]$
2. Single Source Shortest path Problem: This problem involves finding the shortest path from a source node to every other node in a weighted graph, where the weight can also be negative. It detects negative cycles and stops the algorithm.

Algorithm:

Input : number of edges (E), number of vertices (V), set of edges where edge has source, destination and weight.

Output: vertices and minimum distance to each vertex

- a. Initialize  $\text{dist}[V]$  with  $\text{dist}[\text{src}] = 0$  and rest Infinity
- b. For  $l$  in 0 to  $v-1$ 
  - i. For  $j$  in 0 to  $E$ 
    1.  $u = \text{edges}[j].\text{src}$
    2.  $v = \text{edges}[j].\text{dest}$
    3.  $\text{weight} = g.\text{edges}[j].\text{weight};$
    4. if  $\text{dist}[u] \neq \text{Infinity}$  and  $\text{dist}[u] + \text{weight} < \text{dist}[v]$
    5.  $\text{dist}[v] = \text{dist}[u] + \text{weight}$
- c. For  $i$  in 0 to  $e$ 
  - i.  $u = \text{edges}[i].\text{src}$
  - ii.  $v = \text{edges}[i].\text{dest}$
  - iii.  $\text{weight} = \text{edges}[i].\text{weight}$
  - iii. if  $\text{dist}[u] \neq \text{Infinity}$  and  $\text{dist}[u] + \text{weight} < \text{dist}[v]$
  - iv. print "Negative weight cycle"
  - v. exit
  - vi. for  $i$  in 0 to
  - vii. print "vertex"  $i$  "distance = "  $\text{dist}[i]$
  - viii. exit
- d. for  $i$  in 0 to  $v$ 
  - i. print "vertex"  $i$  "distance = "  $\text{dist}[i]$
- e. exit

## Observation:

### 1. 0/1 Knapsack Problem:

```
#include <iostream>
#include <functional>

using namespace std;

class KnapSack{
public:

    int knapSackRec(int W, int wt[], int val[], int index, int** dp){
        if(index < 0 || W == 0){
            return 0;
        }

        if(dp[index][W] != -1){
            return dp[index][W];
        }

        if(wt[index] > W){
            return dp[index][W] = knapSackRec(W, wt, val, index - 1, dp);
        }else{
            dp[index][W] = max(val[index] + knapSackRec(W - wt[index], wt, val, index - 1,
dp), knapSackRec(W, wt, val, index - 1, dp));
            return dp[index][W];
        }
    }

    int knapSack(int W, int wt[], int val[], int n){
        int** dp;
        dp = new int*[n];

        for (int i = 0; i < n; i++){
            dp[i] = new int[W + 1];
        }

        for (int i = 0; i < n; i++){
            for (int j = 0; j < W + 1; j++){
                dp[i][j] = -1;
            }
        }
        return knapSackRec(W, wt, val, n - 1, dp);
    }
};
```

```

int main(){
    int size = 1000;
    for (int i = 0; i < 5; i++){
        int profit[size];
        int weight[size];
        for (int i = 0; i < size; i++){
            auto v = rand() % 10;
            auto w = rand() % 10;
            profit[i] = v;
            weight[i] = w;
        }
        int n = sizeof(profit) / sizeof(profit[0]);
        int W = 10;
        KnapSack knapSack;
        auto knap = [&]() {
            knapSack.knapSack(W, weight, profit, n);
        };
        cout << "Time taken to process " << size << " elements: " << getTime(knap) <<
        "ns" << endl;
        size += 500;
    }
}

```

Output:

```

• → lab4 git:(main) x g++ Knapsack.cpp -o Knapsack.out
• → lab4 git:(main) x ./Knapsack.out
Time taken to process 1000 elements: 261ns
Time taken to process 1500 elements: 344ns
Time taken to process 2000 elements: 488ns
Time taken to process 2500 elements: 611ns
Time taken to process 3000 elements: 697ns

```

## 2. Single source shortest problem:

```
#include <iostream>
#include <bits/stdc++.h>

using namespace std;

struct Edge {
    int src, dest, weight;
};

struct Graph {
    int V, E;
    struct Edge* edge;
};

struct Graph* createGraph(int V, int E)
{
    struct Graph* graph = new Graph;
    graph->V = V;
    graph->E = E;
    graph->edge = new Edge[E];
    return graph;
}

void printArr(int dist[], int n)
{
    printf("Vertex  Distance from Source\n");
    for (int i = 0; i < n; ++i)
        printf("%d \t\t %d\n", i, dist[i]);
}

void BellmanFord(struct Graph* graph, int src)
{
    int V = graph->V;
    int E = graph->E;
    int dist[V];

    for (int i = 0; i < V; i++)
        dist[i] = INT_MAX;
    dist[src] = 0;

    for (int i = 1; i <= V - 1; i++) {
        for (int j = 0; j < E; j++) {
            int u = graph->edge[j].src;
```

```

        int v = graph->edge[j].dest;
        int weight = graph->edge[j].weight;
        if (dist[u] != INT_MAX
            && dist[u] + weight < dist[v])
            dist[v] = dist[u] + weight;
    }
}

for (int i = 0; i < E; i++) {
    int u = graph->edge[i].src;
    int v = graph->edge[i].dest;
    int weight = graph->edge[i].weight;
    if (dist[u] != INT_MAX
        && dist[u] + weight < dist[v]) {
        printf("Graph contains negative weight cycle");
        return;
    }
}

printArr(dist, V);

return;
}
int main()
{
    int V = 5;
    int E = 8;
    struct Graph* graph = createGraph(V, E);

    graph->edge[0].src = 0;
    graph->edge[0].dest = 1;
    graph->edge[0].weight = -1;

    graph->edge[1].src = 0;
    graph->edge[1].dest = 2;
    graph->edge[1].weight = 4;

    graph->edge[2].src = 1;
    graph->edge[2].dest = 2;
    graph->edge[2].weight = 3;

    graph->edge[3].src = 1;
    graph->edge[3].dest = 3;
    graph->edge[3].weight = 2;

    graph->edge[4].src = 1;

```

```

graph->edge[4].dest = 4;
graph->edge[4].weight = 2;

graph->edge[5].src = 3;
graph->edge[5].dest = 2;
graph->edge[5].weight = 5;

graph->edge[6].src = 3;
graph->edge[6].dest = 1;
graph->edge[6].weight = 1;

graph->edge[7].src = 4;
graph->edge[7].dest = 3;
graph->edge[7].weight = -3;

auto bellman = [&]() {
    BellmanFord(graph, 0);
};

cout << getTime(bellman) << "ns time taken" << endl;
return 0;
}

```

Output:

```

• → lab4 git:(main) x g++ Bellman_ford.cpp -o Bellman_ford.out
• → lab4 git:(main) x ./Bellman_ford.out
Vertex    Distance from Source
0          0
1         -1
2          2
3         -2
4          1
32ns time taken

```

Conclusion:

We solved fractional knapsack problem and single source shortest algorithm by applying dynamic programming in C++.