Proportional Logic

- The field of Natural Language Processing (NLP) is advancing at a high speed, with devices such as Amazon Alexa, Google Home etc. using the underlying technology to turn spoken language into commands. This technology is based on proportional logic, because it tries to immitate the human language -> the same sentence can be said in multiple ways, but always keeping the same meaning.
- A statement is either **true** or **false**, it can not be both.

Compound statements (connectives)

Name	Term	Symbol
NOT	Negation	\sim
AND	Conjunction	\wedge
OR	Disjunction	V

- Variables in proportional logic can only have two values, either true or false. This is why they are called boolean variables.
- Negation has presedence over conjunction and disjunction.

Name	Term
"It is hot"	\overline{p}
"It is sunny"	q
"It is not hot but it is sunny"	$\sim p \wedge q$
"It is neither hot nor sunny"	$\sim (p \wedge q)$
"It is not true that it is hot and sunny"	$\neg(p \land q)$

Truth tables

Negation

$\sim p$
F
Τ

${\bf Conjunction}$

\overline{p}	q	$(p \wedge q)$
$\overline{\mathbf{T}}$	Τ	${ m T}$
Τ	\mathbf{F}	\mathbf{F}
F	\mathbf{T}	\mathbf{F}
F	F	\mathbf{F}

Disjunction

\overline{p}	q	$(p \vee q)$
\mathbf{T}	\mathbf{T}	${ m T}$
\mathbf{T}	\mathbf{F}	${ m T}$
F	Τ	${ m T}$
F	F	\mathbf{F}

Exclusive or (XOR)

$$p \oplus q \equiv (p \vee q) \wedge \sim (p \wedge q)$$

\overline{p}	q	$p \vee q$	$p \wedge q$	$\sim (p \wedge q)$	$p \oplus q$
$\overline{\mathrm{T}}$	Τ	Т	Т	F	F
${ m T}$	\mathbf{F}	${ m T}$	\mathbf{F}	${ m T}$	${ m T}$
\mathbf{F}	${ m T}$	${ m T}$	\mathbf{F}	${f T}$	${ m T}$
F	F	\mathbf{F}	\mathbf{F}	${ m T}$	\mathbf{F}

Three input statements With three statements, it makes sense to evaluate the each of the statements individually and then put the stements together.

\overline{p}	q	r	$p \wedge q$	$\sim r$	$(p \wedge q) \vee \sim r$
Τ	Т	Т	Т	F	
${\rm T}$	${\rm T}$	\mathbf{F}	${ m T}$	${ m T}$	${ m T}$
${\rm T}$	\mathbf{F}	\mathbf{T}	\mathbf{F}	F	\mathbf{F}
${\rm T}$	\mathbf{F}	\mathbf{F}	\mathbf{F}	${ m T}$	${ m T}$
\mathbf{F}	${\rm T}$	${\rm T}$	\mathbf{F}	\mathbf{F}	\mathbf{F}
\mathbf{F}	${\rm T}$	\mathbf{F}	\mathbf{F}	${ m T}$	${ m T}$
\mathbf{F}	\mathbf{F}	\mathbf{T}	\mathbf{F}	F	\mathbf{F}
\mathbf{F}	F	F	\mathbf{F}	\mathbf{T}	${ m T}$

De Morgan's law When the negation sign is applied to a parantheses, the expressions "and" as well as "or" are interchanged.

$$\sim (p \vee q) \equiv \sim p \wedge \sim q$$

$$\sim (p \land q) \equiv \sim p \lor \sim q$$

Tautology and contradiction A tautology (denoted by the symbol t) is a statement that is always true regardless of the truth values of its component statements

$$p \wedge t \equiv \mathbf{t}$$

A contradiction (denote by the symbol c) is a statement that is always false regardless of the truth values of its component statements

$$p \wedge c \equiv \mathbf{c}$$

Basic logical equivalences

Name	Laws
Commutative	$p \land q \equiv q \land p$
Associative	$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$
Distributive	$p \land (q \lor r) \equiv (p \land q) \lor (p \land r)$
Identity	$p \wedge t \equiv p$
Negation	$p \lor \sim p \equiv \mathbf{t}$
Double negative	$\sim (\sim p) \equiv p$
Idempotent	$p \wedge p \equiv p$
Universal bound	$pee\mathbf{t}\equiv\mathbf{t}$
De Morgan's	$\sim (p \land q) \equiv \sim p \lor \sim q$
Absorption	$p \vee (p \wedge q) \equiv p$
Negations of t and c	$\sim {f t} \equiv {f c}$

Name	Laws
Commutative	$p \vee q \equiv q \vee p$
Associative	$(p \vee q) \vee r \equiv p \vee (q \vee r)$
Distributive	$p \lor (q \land r) \equiv (p \lor q) \land (p \lor r)$
Identity	$p \lor c \equiv p$
Negation	$p \wedge \sim p \equiv \mathbf{c}$
Double negative	N. A.
Idempotent	$p \wedge p \equiv p$
Universal bound	$p \wedge \mathbf{c} \equiv \mathbf{c}$
De Morgan's	$\sim (p \vee q) \equiv \sim p \wedge \sim q$
Absorption	$p \land (p \lor q) \equiv p$
Negations of t and $\mathbf c$	$\sim {f c} \equiv {f t}$

Condtional statements

If-then conditional statements, if is the hypothesis and q is the conclusion

p	q	$p \to q$	$\sim p \vee q$
\overline{T}	Т	${ m T}$	Τ
\mathbf{T}	\mathbf{F}	\mathbf{F}	F
\mathbf{F}	${\rm T}$	${ m T}$	F
\mathbf{F}	\mathbf{F}	${ m T}$	${ m T}$

If-and only if means p implies q and q implies p.

$$(p \to q) \land (q \to p)$$

$$\begin{array}{c|cccc} \hline p & q & p \leftrightarrow q \\ \hline T & T & T \\ T & F & F \\ F & T & F \\ F & F & T \\ \end{array}$$