

# Algorithms, Design & Analysis

## Lecture 12: **Articulation points and Topological sort**

Muhammad Ibrahim & Rana Hammad Ahmad

Information Technology University

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# About Your Fellows

- Hi there! We are **Ibrahim** and **Hammad**.
- We are Associate Students at ITU.

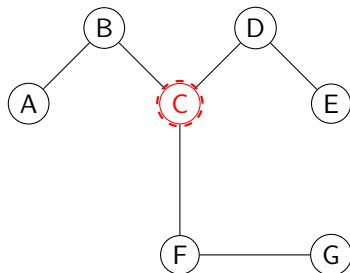
# Articulation Points

# Articulation Points

## Definition:

- Articulation point is a node which if removed increase the number of connected components.

# Articulation Point Example



**Explanation:** The vertex **C** is an articulation point because its removal disconnects the graph into three separate components.

# Algorithm For Articulation Points

# Articulation Points Algorithm

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**Algorithm 1** Finding Articulation Points

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```
1: for each  $x \in V$  do
2:    $flag \leftarrow \text{False}$ 
3:    $t \leftarrow 1$ 
4: end for
5: for each  $v \in N(x)$  do
6:   if  $v$  is unvisited then
7:     DFS( $v$ )
8:   end if
9:   if  $flag = \text{True}$  then
10:     $v$  is an Articulation Point (ARTP)
11:   end if
12:    $flag \leftarrow \text{True}$ 
13: end for
```

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# Depth-First Search Algorithm For Articulation Points



# DFS Algorithm

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## Algorithm 2 DFS( $v$ )

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```

1:  $v.d = t, v.l = t, t \leftarrow t + 1$ 
2:  $v.visited \leftarrow \text{True}$ 
3: for each  $u \in N(v)$  do
4:   if  $u$  is unvisited then
5:     DFS( $u$ )
6:   end if
7:   if  $u.d < v.d$  and  $u \neq v.\pi$  then
8:      $v.l \leftarrow \min(u.d, v.l)$ 
9:   end if
10:  if  $u \neq v.\pi$  then
11:     $v.l \leftarrow \min(v.l, u.l)$  checking if there is a path to ancestors via child
12:  end if
13:  if  $u.l > v.d$  then
14:     $v$  is an ARTP
15:  end if
16: end for

```

# Complexity

- Runtime complexity of this algorithm is  $\Theta(|V| + |E|)$ .

# Topological Sort

# Topological Sort

## Definition:

- It is linear ordering of graph vertices such that for every directed edge  $uv$  from vertex  $u$  to vertex  $v$ ,  $u$  comes before  $v$  in the ordering.

# Topological Sort

## Conditions:

- Tree: Minimally connected graph (there exist a path between every pair vertices).
- Directed Acyclic Graph (DAG)
- It checks indegrees and outdegrees

# Topological Sort

## Example:

- **Task Scheduler**

- A Task Scheduler is a system that manages the execution order of tasks while considering dependencies. This is a perfect real-world application of Topological Sorting, which is used in Directed Acyclic Graphs (DAGs) to order tasks such that each task appears before any tasks that depend on it.

# How Topological Sorting Works in a Task Scheduler

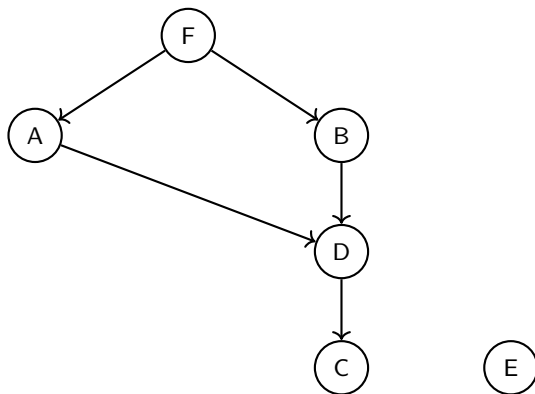
- **Tasks as Graph Nodes:**

- Each task is represented as a **node** in a directed graph.

- **Dependencies as Directed Edges:**

- A directed edge  $A \rightarrow B$  means **task A must be completed before task B**.
- Check if in-degree is zero. If it is remove it and put it in the list.

# Task Dependency Graph (Topological Sorting)



**Valid Execution Order:**  $F \rightarrow A \rightarrow B \rightarrow D \rightarrow C \rightarrow E$



# Topological Sorting Algorithm (1/2)

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## Algorithm 3 Topological Sorting (TS)

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```
1: Given a directed graph  $G = (V, E)$ 
2: Compute in-degrees of all vertices
3: Initialize an empty queue  $W \leftarrow []$ 
4: for each  $u \in V$  do
5:   if  $d(u) = 0$  then
6:     Add  $u$  to  $W$ 
7:   end if
8: end for
```

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# Topological Sorting Algorithm (2/2)

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## Algorithm 4 Topological Sorting (TS) - Continued

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```
1: while  $W$  is not empty do  
2:   Remove a vertex  $u$  from  $W$   
3:   for each outgoing edge  $(u, v)$  do  
4:     Remove edge  $(u, v)$  from  $G$   
5:     Decrease in-degree of  $v$   
6:     if  $d(v) = 0$  then  
7:       Add  $v$  to  $W$   
8:     end if  
9:   end for  
10: end while
```

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# Topological Sort(Home Work)

- Prove this Topological Sort in Linear time  $O(V + E)$ .