

# Small omega and Small o Notations

In this lesson, we discuss notations which imply loose bounds.

The small o and small  $\omega$  are complementary notations to the big O and big  $\Omega$  notations. For algorithm analysis, the most important notation is the big O. For the sake of completeness, we mention the small o and small  $\omega$  notations too.

## Small o

The small o is **not** an asymptotically tight upper bound. The formal definition is similar to big O, with one important difference. A function  $f(n)$  belongs to the set  $o(g(n))$ , if the following condition is satisfied.

$$0 \leq f(n) < cg(n)$$

*for any positive constant  $c$ , there exists some constant  $n_0$  which if  $n$  surpasses, the above inequality holds*

Note that the inequality should not hold for *some* positive constant  $c$ , as is the case for big O; rather, it should hold for all positive constants.

Unlike big O, which may or may not be tight, the small o notation is necessarily not tight.

## Small $\omega$

Small  $\omega$  is similarly **not** a tight lower bound. Big  $\Omega$ , on the contrary, may or may not be a tight bound. A function  $f(n)$  belongs to the set  $\omega(g(n))$  if the following inequality holds:

$$0 \leq g(n) < cf(n)$$

For **any** positive constant  $c$ , there exists some constant  $n_0$  which if  $n$  surpasses, the above inequality holds. The above inequality should hold for **all constants**. Note that in the case of big omega, the inequality was required to hold for *some constant*.

### Explanation

We can see that the small case notations are, in a sense, *relaxed* compared to their upper case notations.

- $2n^2$  is equal to  $\omega(n)$  but  $2n^2 \neq \omega(n^2)$ . To understand the first claim consider the inequality:

$$cn \leq 2n^2$$

Let's pick a really big number for  $c=1000$ . To make the inequality hold I can set  $n_0=c=1000$ , and because of  $n$  being squared the right side will be bigger for every value of  $n$  greater than  $n_0$ .

Now, to understand why  $2n^2 \neq \omega(n^2)$ , consider the inequality:

$$cn^2 < 2n^2$$

If I pick  $c = 3$  then - no matter what value of  $n_0$  I choose - I can never satisfy the above inequality.

- For small  $\omega$ ,  $2n^2$  is  $o(n^3)$  but  $2n^2 \neq o(2n^2)$

1

If  $f(n)$  is  $\Theta(n^3)$ , is  $f(n)$  also  $o(n^3)$  and  $\omega(n^3)$  ?

2

If  $f(n)$  is  $O(n^3)$ , is  $f(n)$  also  $o(n^3)$ ?

Check Answers