

Score for this attempt: **19** out of 20

Submitted Mar 15 at 10:20am

This attempt took 75 minutes.

Question 1

1 / 1 pts

Is the following statement true or false?

Propositional logic cannot deal with modifiers like "there exists," "all," and "only," while First-Order Logic can.



Correct!

☒ True

☐ False

Question 2

1 / 1 pts

Consider the following formula, in which the occurrences of variables are numbered from 1 to 6 for reference:

$$\exists x(P(x, y) \rightarrow \forall y P(y, x))$$

1 2 3 4 5 6

Question 2

1 / 1 pts

Consider the following formula, in which the occurrences of variables are numbered from 1 to 6 for reference:

$$\underset{1}{\exists} \underset{2}{x} (\underset{3}{P}(\underset{4}{x}, \underset{5}{y}) \rightarrow \underset{6}{\forall} \underset{y}{P}(\underset{y}{y}, \underset{x}{x}))$$

Which occurrence of variables denotes a free occurrence of a variable in this formula?

Correct!

☒ 3

☐ 2

☐ 5

☐ 6

6

Question 3

1 / 1 pts

Is the following statement true or false?

Consider 2 sets: $A = \{1, 2\}$ and $B = \{a, b, c\}$. There are 0 onto-functions from A to B.

Correct!

☒ True

☐ False



Question 4

1 / 1 pts

Consider the graph below, noting that each arrow is labeled with a number in red.

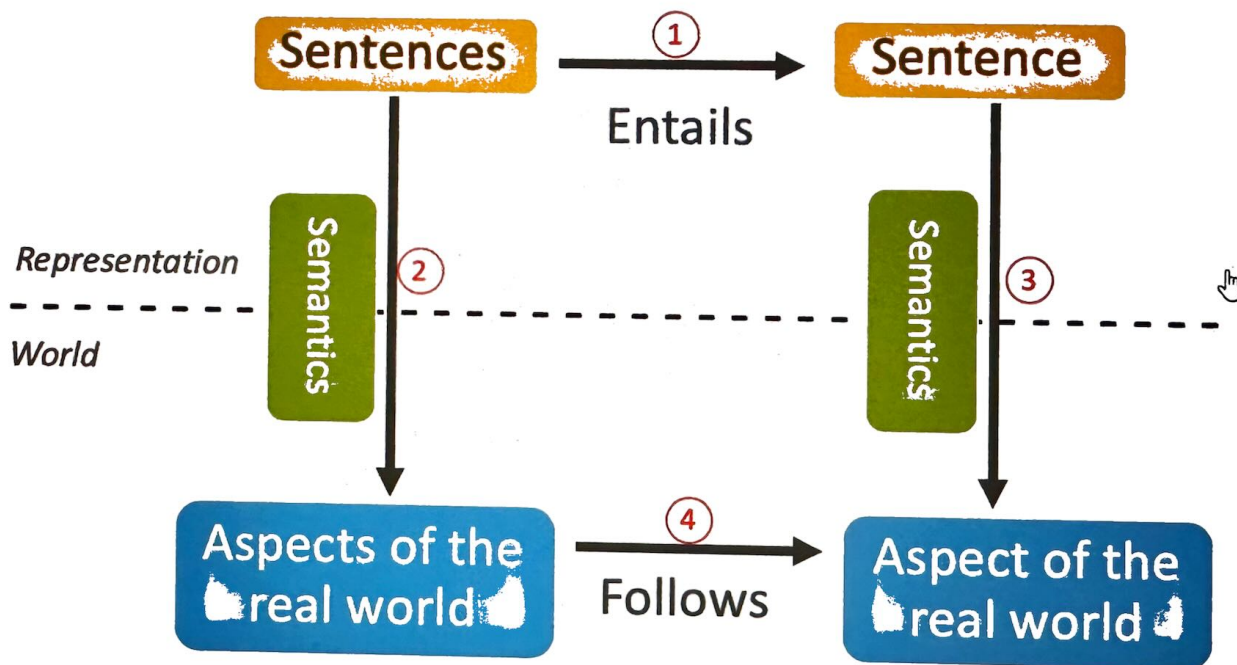
Which arrow denotes the step in which we write a clingo program to represent the Sudoku problem?



Question 4

Consider the graph below, noting that each arrow is labeled with a number in red.

Which arrow denotes the step in which we write a clingo program to represent the Sudoku problem?



Correct!

☒ 2

☐ 3

☐ 1

1

4

Question 5

1 / 1 pts

Let $\Gamma = \{p \vee q, q\}$ be a set of formulas where p and q are atoms. Which of the following formulas does Γ entail?

☐ $p \wedge q$

☒ $p \rightarrow q$

☐ p

☐ $q \rightarrow p$

Correct!

Question 6

0 / 1 pts

How many stable model have the following program:

$\{p(X): X=3..5\}.$

p

$q \rightarrow p$

Question 6

0 / 1 pts

How many stable model have the following program:

$\{p(X): X=3..5\}.$

$\vdash p(3).$

$\vdash \text{not } p(5).$



You Answered

☒ 4

Correct Answer

☐ 2

☐ 5

☐ 3

Question 7

Original Score: 1 / 1 pts Regraded Score: 1 / 1 pts

Question 7

Original Score: 1 / 1 pts Regraded Score: 1 / 1 pts

⚠ This question has been regraded.

Recall the definition of propositional formula, provided below.

A propositional formula of signature σ is defined recursively as follows:

- every atomic formula of σ is a formula
- both 0-place connectives \top, \perp are formulas
- if F is a formula, then $\neg F$ is a formula
- For any binary connective \odot , if F and G are formulas then $(F \odot G)$ is a formula
- if F is a formula then $\forall x F$ and $\exists x F$ are formulas

Which of the options listed below is a propositional formula strictly according to the above definition?
(Select all that apply.)

Let $\sigma = \{a, P, Q\}$, where a is an object constant, P is a unary and Q, S are binary predicate constant. x, y, z are general variables.

☐ $p \leftarrow \perp$

You Answered

☒ q

It is not in signature.

Correct Answer

⓪ p

Question 8

1 / 1 pts

What is the correct translation of the following statement into mathematical logic?

"Some integers are negative".

☐ $\exists x (\text{integer}(x) \vee \text{negative}(x))$

☐ $\forall x (\text{integer}(x) \rightarrow \text{negative}(x))$

Correct!

☒ $\exists x (\text{integer}(x) \wedge \text{negative}(x))$

☐ $\exists x (\text{integer}(x) \rightarrow \text{negative}(x))$

Question 9

1 / 1 pts

Which option is the value of the following aggregate?

$\# \text{sum}\{M * N, M, N : M = 1..3, N = 1..3\}$

25

$\exists x (\text{integer}(x) \rightarrow \text{negative}(x))$

Question 9

1 / 1 pts

Which option is the value of the following aggregate?

$\# \text{sum}\{M*N, M, N : M=1..3, N=1..3\}$

☐ 25

☐ 6

☐ 9

☒ 36

Correct!

Question 10

1 / 1 pts

Which answer option is equivalent to the following Clingo program?

$p(X**Y, X*|Y*X|) :- X=-2..0, Y=-5..-2, X \leq Y.$

☒ p(0, -8)

☐ p(0, 8)

Correct!

Question 11

1 / 1 pts

Which of the formulas listed below has a model but not a Herbrand model? (Select all that apply.)

$$P(a) \wedge \neg P(a)$$

$$\neg P(a) \rightarrow P(a)$$

Correct!

☒ $a=b$

Correct!

☒ $P(a) \wedge P(b) \wedge (\forall x P(x) \rightarrow x=c)$



Question 12

1 / 1 pts

Which statement listed below correctly describes the stable models of the following propositional rule?

$$r \leftarrow \neg r$$

Correct!

☒ It has no stable model.

☐ Its stable model is \emptyset .

☐ Its stable models are \emptyset and $\{p\}$.

☐ Its stable model is $\{p\}$.

Question 13

1 / 1 pts

Is the following statement true or false?

If a propositional program contains a rule of form $F \leftarrow \perp$ then removing that rule doesn't change its stable models.

Correct!

☒ True

☐ False



Question 14

1 / 1 pts

Which option listed below shows all of (and only) the local variables in the following clingo program?

$\{q(I,J): J=1..2\} :- I = 1..3.$

☐ Both I and J

☒ J

☐ I

☐ Neither I nor J

Correct!

Question 15

1 / 1 pts

What is the value of the following aggregate when $n = 6$?

$\#count\{X*Y:X=2..n, Y=2..n, X*Y \leq n\}$

☐ 3

☐ 16

☐ 25

☒ 2

Correct!

By placing $n = 6$ into the rule, we can get $\#count\{X*Y:X=2..6, Y=2..6, X*Y \leq 6\}$. The potential value of X could be the integer in the interval of 2 to 6, and the same for Y , however, not all of the them are the correct values because of the condition $X*Y \leq 6$.

Enumerate the value of X and Y to satisfy the condition and we can get (X, Y) to be $(2,2), (2, 3), (3, 2)$. However, the value of $X*Y$ when (X, Y) is $(3, 2)$ is equal to $(2,3)$, so we only need to count one of them, and so we get the return value to be 2.

Question 16

1 / 1 pts

Assuming the signature consists of p, q, r . Which of the options are **models** of the following program?

$p \leftarrow q \wedge r$

Question 16

1 / 1 pts

Assuming the signature consists of p, q, r . Which of the options are **models** of the following program?

$p \leftarrow q \wedge r$

$q \leftarrow p$

$r \leftarrow p$

Choose all that apply.

Correct!

☒ $\{p, q, r\}$



Correct!

☒ $\{q\}$

☐ $\{p, q\}$

Correct!

☒ $\{r\}$

☐ $\{p\}$

☐ $\{q, r\}$

Correct!

☒ \emptyset (empty set)

Question 17

1 / 1 pts

Correct!

0 (empty set)

Question 17

1 / 1 pts

Which option is equivalent to the following clingo program?

$p(1,1). p(1,2). p(2,1). p(2,2).$

☐ $p(1;2, 1;2).$

☐ $p(X, X^{**}|X-Y|) :- X=1..2, Y=2..3.$

☐ $p(X,Y) :- X=1..2, Y=1..X.$

☒ $p(1..2, 1..2).$

Correct!

Question 18

1 / 1 pts

Find the reduct of (the propositional image of) the following clingo program relative to the interpretation $\{q(a)\}.$

$\{p(a)\}.$

$q(X) :- p(X).$

$\perp \vee T$

Question 18

1 / 1 pts

Find the reduct of (the propositional image of) the following clingo program relative to the interpretation $\{q(a)\}$.

$\{p(a)\}$.

$q(X) \text{ :- } p(X)$.

☐ $\perp \vee \top$
 $q(a) \leftarrow \perp$

☐ $p(a) \vee \top$
 $q(a) \leftarrow \perp$

☐ $p(a) \vee \neg p(a)$
 $q(a) \leftarrow p(a)$

☒ $p(a) \vee \top$
 $q(a) \leftarrow p(a)$

Correct!

Question 19

1 / 1 pts

Question 19

1 / 1 pts

The following clingo program represents what will happen when we roll a dice. However, not all stable models (e.g., {roll_dice, get(1), get(2)}) of this program represent solutions of rolling a dice. Choose the option such that: adding the option to the clingo program can weed out the elements of the search space that do not represent solutions.

roll_dice.

{get(N): N=1..6} :- roll_dice.

☐ :- roll_dice, get(1), get(2).



☐ :- get(N1), get(N2), N1 != N2.

☒ :- not get(1), not get(2), not get(3), not get(4), not get(5), not get(6).

☐ :- get(N1), get(N2), N1 != N2.

☐ :- 1{get(N): N=1..6}1.

Correct!

Question 20

1 / 1 pts

☐ :- not get(1), not get(2), not get(3), not get(4), not get(5), not get(6).

☐ :- get(N1), get(N2), N1 != N2.

☐ :- 1{get(N): N=1..6}1.

Question 20

1 / 1 pts

In Hamiltonian Cycle problem, we use "edge(X,Y)" to represent that "there is a directed edge from X to Y" and use "in(X,Y)" to represent that "edge from X to Y is in the Hamiltonian cycle". Then which option is represented by the following clingo rule?

$1\{in(X,Y): edge(X,Y)\}1 :- node(Y).$

- ☐ Every node is reachable.
- ☐ For each node Y, we choose exactly 1 edge of it to be in the Hamiltonian cycle.
- ☐ Every node in the graph has exactly 1 outgoing edge.
- ☒ Every node in the graph has exactly 1 incoming edge.

Correct!

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