

# CHAOS IN CRYPTOGRAPHY

AMEY P GAIKWAD

RAHUL DANDWATE

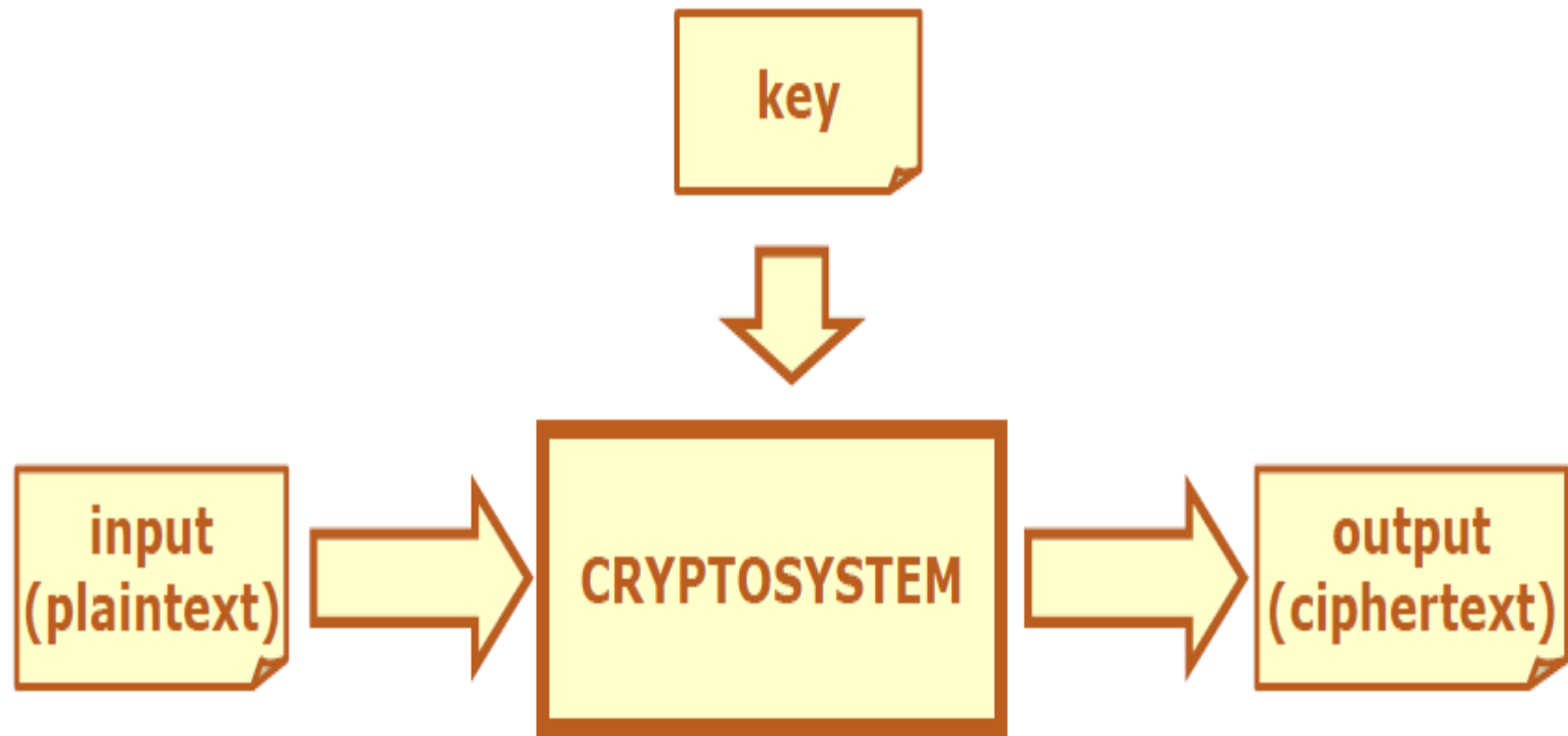
SUMUKH VAIDYA

SIDDHESH PEDNEKAR

PALLAV JOSHI

# Cryptography

- ▶ Cryptography is the study of methods of converting messages into a disguised form so that only the intended recipient can remove the disguise and read the message
- ▶ Cryptography - The art of making
- ▶ Cryptanalysis - The art of breaking



► Mathematically ,

A Cryptosystem is a five tuple (  $P, C, K, E, D$  ) satisfying the following conditions;

1.  $P$  is a set of possible plaintexts
2.  $C$  is set of possible ciphertexts
3.  $K$ , the key space, is a finite set of possible keys
4. For each  $k \in K$ , there is an encryption rule  $e_k \in E$  and a corresponding decryption rule  $d_k \in D$  where

$$e_k: P \rightarrow C \quad \text{and} \quad d_k: C \rightarrow P$$

are functions such that  $d_k(e_k(x)) = x$  for every plaintext  $x \in P$

# Kerkchoff Shannon Principle

- ▶ A cryptosystem must be secure even if everything about the system , except the key is public knowledge.

Shannon's Principle:

The enemy knows the system i.e the ALGORITHM

# Chaos in Cryptography

- ▶ Chaotic cryptography is the application of the mathematical chaos theory to the practice of the cryptography, the study or techniques used to privately and securely transmit information with the presence of a third-party or adversary. The use of chaos or randomness in cryptography has long been sought after by entities wanting a new way to encrypt messages.
- ▶ The mapping between these two fields will be highlighted later on.

# Principles of Modern Ciphers

- ▶ Maximum entropy of the cipher
- ▶ Maximum diffusion of the key
- ▶ Long cycle length of cipher ( to avoid repetition )

Diffusion - Transposition

+

Confusion - Substitution



# Applications Of Chaos In Cryptography

- ▶ Pseudo Random Number Generators ( PRNGs ) (An overview)
- ▶ Standalone algorithms developed from chaotic maps
- ▶ Chaos synchronization as means to secure communications
- ▶ Cellular Automata cryptography



# RNGs and PRNGs

## RNGs

- ▶ Truly random
- ▶ Generated through natural processes ( fluctuations ) feeding the black box to generate the random numbers.(Hardware)
- ▶ Completely random and nondeterministic.

## PRNGs

- ▶ As the name suggests pseudo random
- ▶ Generated through a mathematical algorithm with the help of an initial seed.
- ▶ Deterministic ( know the seed , know the random number )

# TRNGs and PRNGs

- ▶ These random number generators are used in variety of cryptosystems (eg. The Vernam Cipher )
- ▶ For absolute secrecy the random numbers generated need to be TRNGs. However due to computational limitations PRNGs are used.
- ▶ These PRNGs are going to result in **practical rather than perfect secrecy** ( due to computational limitations ).
- ▶ The **seed space** of TRNGs is much larger than the seed space of PRNGs

# Chaos and PRNGs

- ▶ The pseudo random number generators are deterministic. (given the initial seed since the algorithm is a mathematical computation )
- ▶ If we change the initial seed value an entirely different random number is generated.
- ▶ One method of producing PRNGs is **the middle squares method** developed by Neumann.
- ▶ The random number generated is not due to the error propagation but rather due to the inherent mathematical formulation of the algorithm.
- ▶ For the middle squares method taking an  $n$  digit seed atmost  $10^n$  numbers are generated before reusing the seed.
- ▶ Prevent reusing the seed (OTP).

# One Time Pads ( OTP )

- ▶ Before moving to the actual presence of chaos in an actual cryptosystem an example of an ideal cipher design is mentioned for comparisons with the future designs shown
- ▶ Perfect secrecy
- ▶ Claude Shannon - 1) impossible to break 2) Vernam cipher
- ▶ Assumptions :
  1. One time use
  2. Truly random numbers are needed
  3. As long as the plain text ( the key )
  4. The key is kept complete secret

# OTP Cipher Design

- ▶ Generally modulo arithmetic is used
- ▶ In the simplest possible case, the alphabets are numbered from 0 to 25.
- ▶ A TRN is generated and added to the plain text
- ▶ The sum is calculated modulo 26.
- ▶ The ciphertext is obtained.

	H	E	L	L	O	message
	7 (H)	4 (E)	11 (L)	11 (L)	14 (O)	message
+	23 (X)	12 (M)	2 (C)	10 (K)	11 (L)	key
=	30	16	13	21	25	message + key
=	4 (E)	16 (Q)	13 (N)	21 (V)	25 (Z)	(message + key) mod 26
	E	Q	N	V	Z	→ ciphertext

Source : Wikipedia

# Cipher Evaluation

- ▶ The entropy of the cipher text is very large and hence this cryptosystem is a very good and an efficient one.
- ▶ Given a plain text the a priori and the posteriori probabilities of the plain text are equal.
- ▶ In other words the entropy of a plaintext is equal to the entropy of the plaintext given the ciphertext. Its impossible to predict the plaintext from the ciphertext.
- ▶  $H(M)=H(M|C)$
- ▶ The keyspace consists of just one key vector

# Security Analysis

- ▶ If TRNG is used then the OTP is unbreakable. As proven by Shannon **its information theoretically secure** and has the property which he termed as absolute secrecy.
- ▶ Given a cipher text **any plain text can be concocted from it** by using random numbers and using the modulo in a reverse manner.
- ▶ Even if a PRNG with a very long seed is used then even with the present supercomputers it would take years to break it down.( This is possible since after some finite value the seed is reused ).

	E	Q	N	V	Z	ciphertext
	4 (E)	16 (Q)	13 (N)	21 (V)	25 (Z)	ciphertext
-	23 (X)	12 (M)	2 (C)	10 (K)	11 (L)	key
=	-19	4	11	11	14	ciphertext - key
=	7 (H)	4 (E)	11 (L)	11 (L)	14 (O)	ciphertext - key (mod 26)
	H	E	L	L	O	→ message

	4 (E)	16 (Q)	13 (N)	21 (V)	25 (Z)	ciphertext
-	19 (T)	16 (Q)	20 (U)	17 (R)	8 (I)	possible key
=	-15	0	-7	4	17	ciphertext-key
=	11 (L)	0 (A)	19 (T)	4 (E)	17 (R)	ciphertext-key (mod 26)



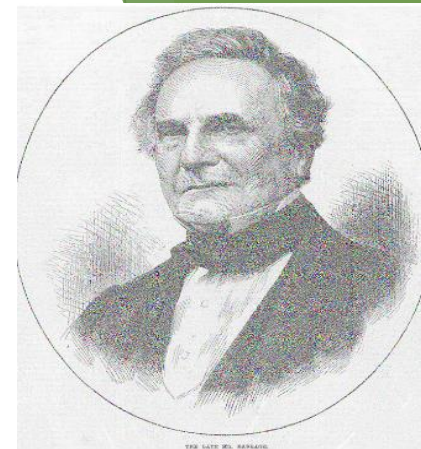
# The Vernam Cipher Design

- ▶ Convert the plain text and the random number of the same length as the plain text into binary.
- ▶  $c = p \oplus n$  (for encryption )
- ▶  $p = n \oplus c$  (for decryption )
- ▶ **Why XOR?**

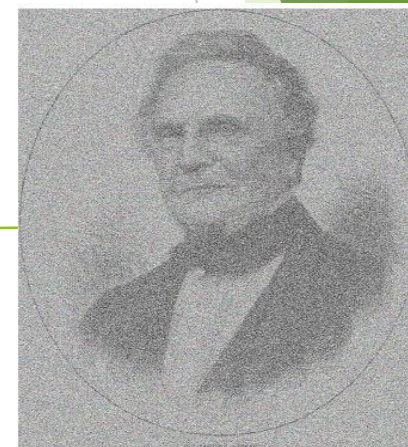
Take an image. Suppose each pixel is coded 0 or 1. The XOR function converts the code into the ciphertext with no statistical bias unlike the OR and the AND operators.



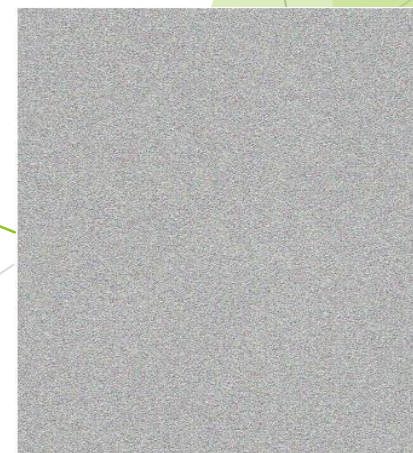
OR



AND



XOR



Courtesy: Khan Academy

# Cipher Evaluation

- ▶ Depending on the length of the plain text cryptanalysis of the Vernam cipher varies.
- ▶ The key space of the Vernam cipher is one ( just one random number is used for converting plain to cipher text.)
- ▶ The length of the key will be equal to the length of the plain text.
- ▶ It is an example of a **symmetric cryptosystem**.

# Security Analysis

- ▶ Even when using a PRNG for generating the random numbers if the seed is long enough a sufficiently random number can be generated which leave the cryptosystem as practically secure rather than perfectly secure. ( a random number with no statistical bias ).
- ▶ But **generally practically secure** is more than sufficient for protection of data.
- ▶ An important condition for chaos is also present.
  1. Random number with no statistical bias(maximum confusion)
  2. Ultrasensitivity to the initial seed K ( maximum diffusion.)

Chaos is beautifully manifested fundamentally in the pseudo random number generation algorithm which carries over to this very simple but strong and efficient cryptosystem.

The example show the action of the Vernam cipher.

Everything depends upon the random number generator.

Change the seed, get a different random number and hence a different cipher text.

Plaintext:

Attack at 03:00am.

8-bit ASCII code	8-bit cipher	XOR
01000001	00000101	01000100
01110100	00000000	01110100
01110100	01000010	00110110
01100001	00000001	01100000
01100011	00001000	01101011
01101011	00000110	01101101
00100000	00001101	00101101
01100001	00011011	01111010
01110100	01000011	00110111
00100000	00011111	00011111
00110000	00001001	00111001
00110011	00010000	00100011
00111010	00010001	00101011
00110000	00001001	00111001
00110000	00001010	00111010
01100001	00010010	01110011
01101101	10000000	01101111
00101110	00001001	00100111

Ciphertext:

Dt6'km-z7&9#+9:so'

Courtesy : Prof. JM Blackledge

# An Improvisation

- ▶ The PRNG described earlier shows the problems of periodicity.
- ▶ After a large number of cycles the seed used initially is re-used. This leads to periodicity and finite but long cycle lengths of the random numbers
- ▶ Something better can be done by changing from periodic to aperiodic (though it will still be a PRNG )
- ▶ This is the problem which is addressed by CHAOS

# A Chaotic Map (As a PRNG)

- ▶ As mentioned earlier of the presence of chaos in the inherent process of random number generation, here we aim to change the periodicity of the random number generated.
- ▶ Use of one dimensional maps to generate random numbers proves to be the method for doing so.
- ▶ The random number generated depends not only on the initial number generated but also on the parameters involved in the one dimensional map.
- ▶ Care must be taken to choose the parameters where the map would show chaos.
- ▶ Actual implementation of this method and its security will be discussed later.

# Cryptographical Perspective

1. The key space for the cryptosystem has increase from just one ( the one random number ) to  $N+1$  where  $N$  is the number of parameters of the one dimensional map.
2. Comparing to conventional random number generators, the numbers generated, the 1D map generator is infinite, aperiodic and uncorrelated.

The increase in key space is significant making the cryptanalysts job more difficult. Now he has to not only worry about the initial seed but also about the parameter space. **This can also help in multiagorithmicity. Use two or more more algorithms in conjunction and in a periodic manner to further complicate the random number generation.**



# The Logistic Map

- ▶ To this end a logistic map can be used as a one dimensional map as described.
- ▶ The system show chaotic regime for the parameter values of  $r > 3.57$  but used value is  $r=4$ .
- ▶ So choosing a value of  $r$  while avoiding the gaps in the middle ( for saddle node bifurcations ) and maintaining a closed state phase space values between 0 and 1 can lead to a long number of iterates leading to a random number formation.

# Problems with this method

- ▶ Though this method scores much higher than the earlier implementation, it is not completely devoid of its own drawbacks.
- 1. Truncation error in computers :
  - i. Single precision ( 32 bits ) - on an average 5000 iterations
  - ii. Double precision ( 64 bits ) - on an average  $10^9$  iterations
- 2. Correlations between the successive iterates. (Will be mentioned briefly in Simple and Advanced ciphers .)

# Conclusion

- ▶ Chaos is beautifully manifested fundamentally in the pseudo random number generation algorithm which carries over to this very simple but strong and efficient cryptosystem
- ▶ **The entire cryptosystem breaks if the initial seed is known.** The system is completely deterministic.
- ▶ Since we using binary or in other cases modulo the **entire ciphertext space is bounded** ( another important condition for chaos to be shown.
- ▶ Chaos is entwined closely with the cryptographic necessity of **maximum confusion + diffusion**

# Summary

- ▶ As a conclusion to the application of chaos for random number generation, it is important to note the **way chaos maps itself to cryptographic demands**.
- ▶ The details of the mapping will be highlighted in due course.

The background features abstract, overlapping green geometric shapes, primarily triangles and polygons, in various shades of green, creating a modern and dynamic visual effect.

# Chaos based ciphers

## Chaotic Maps

# Introduction

- ▶ Chaos- best known as sensitivity to initial conditions exhibited by dynamical systems described by differential equations or iterated mappings.
- ▶ We will be aiming to present the connection between these two fields theoretically and practically through the implementation of two simple ciphers.
- ▶ Symmetric Key cryptosystems

# Chaos Dynamics

- ▶ Mainly iterated mappings will be used
- ▶ Elements of the orbit determine chaotic behavior.
- ▶ The system must be ergodic and have a continuous invariant density.
- ▶ Topological mixing (homogenous invariant measure )
- ▶ Positive Lyapunov exponents
- ▶ Attractors
- ▶ Robust chaos
- ▶ Topological transitivity

# Mapping chaos to Cryptography

Properties	Chaos theory	Cryptography
Sensitive dependence on initial conditions	Characteristic chaos concept	Key dependent confusion and diffusion
Discrete time	Time based iteration in case of chaotic maps	Time based encryption rounds.
The security manifesto	The system parameters and the initial condition	The cryptographic keys (the number of keys and how each key adds to the security )
Type	True chaotic system	Pseudo chaotic system
Type of system	Deterministic	Deteministic



# Chaos and Pseudo Chaos

- ▶ True chaos occurs in the chaotic maps and is successfully seen in the Differential equations as shown by Lorenz.
- ▶ Pseudo chaos is the best we can get from chaos for application in cryptography.
- ▶ This Pseudo chaotic nature which is the limit to which we can tap the true potential of chaos is due to our limitations and not due to the nature of the chaotic maps which are going to be used for encryption.
- ▶ Floating Point approximation is the major factor contributing to this/

# Chaos and Pseudo Chaos

## Chaos

- ▶ Infinite number of states
- ▶ Over the real numbers
- ▶ Completely chaotic
- ▶ Infinite Orbits

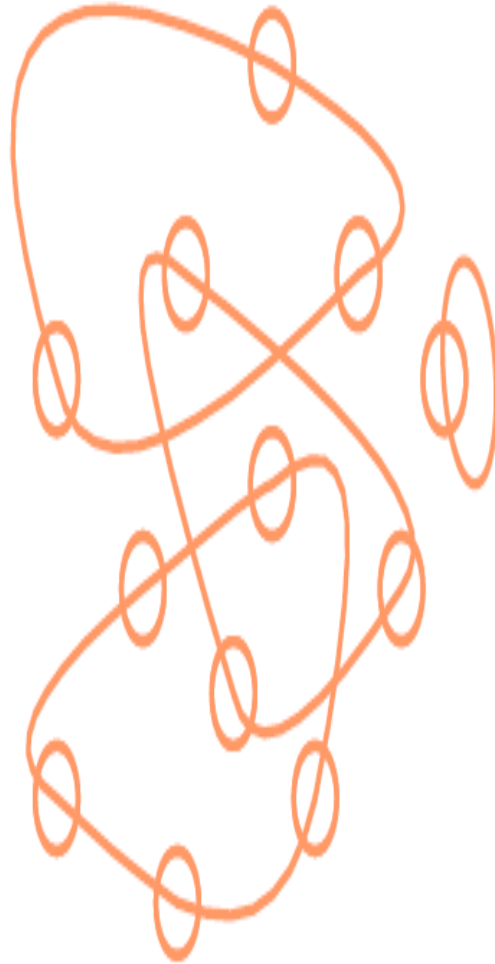
## Pseudo chaos

- ▶ Finite number of states
- ▶ Limit on precision
- ▶ Low cycle lengths
- ▶ Finite Orbits

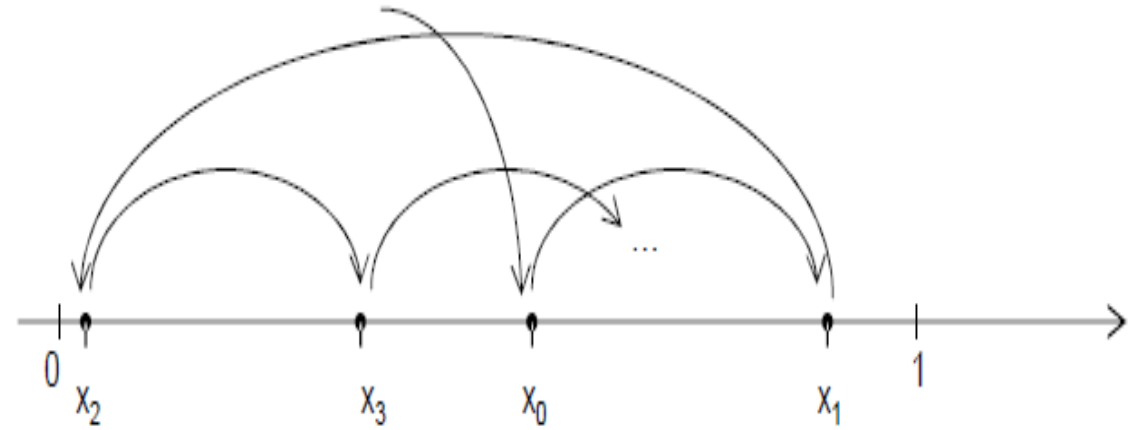
Infinite orbit (Chaos)



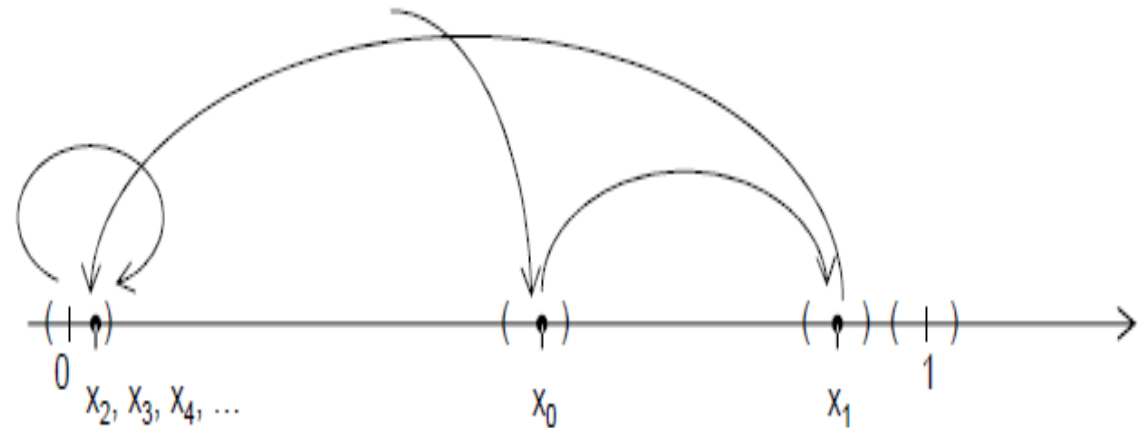
Finite orbits (Pseudo-Chaos)



Continuous Chaos



Floating-point Approximation



# Chaotic Cipher Design

- ▶ Use space discretized versions of chaotic maps that approximate real valued systems
- ▶ Discretization is necessary (why?)
- ▶ Must limit the numbers we can represent within an interval over a given set of numbers.
- ▶ Performance
- ▶ Approximation for chaotic maps is required
- ▶ Transformations should generate values far apart in Euclidian space and also in bit values.

# Cipher Evaluation

- ▶ Measure of entropy of the underlying chaotic maps.
  1. Modify the order of the key. ( Entropy  $\propto \log_2 k$  )
  2. Initial conditions- Chaos is attained and its properties expressed
  3. Randomizing not only the initial conditions but also the number of iterations
  4. Size of data blocks-
    - a) Difficulty in decrypting
    - b) Time consuming
  5. Set of elements are over the real domain

# Security Analysis

- ▶ At a theoretical level cryptographic primitives are deemed if they possess the following characteristics
  1. Randomness increasing
  2. Computationally unpredictable
  3. The cryptographical primitive must increase the entropy of the system over which it operates
  4. Successful cryptanalytic attack depends on the statistical attributes of the cipher.
  5. Randomness increasing is a necessary condition for security but its not sufficient.
  6. The most important signature of security lies in “randomness” and not in its increase.

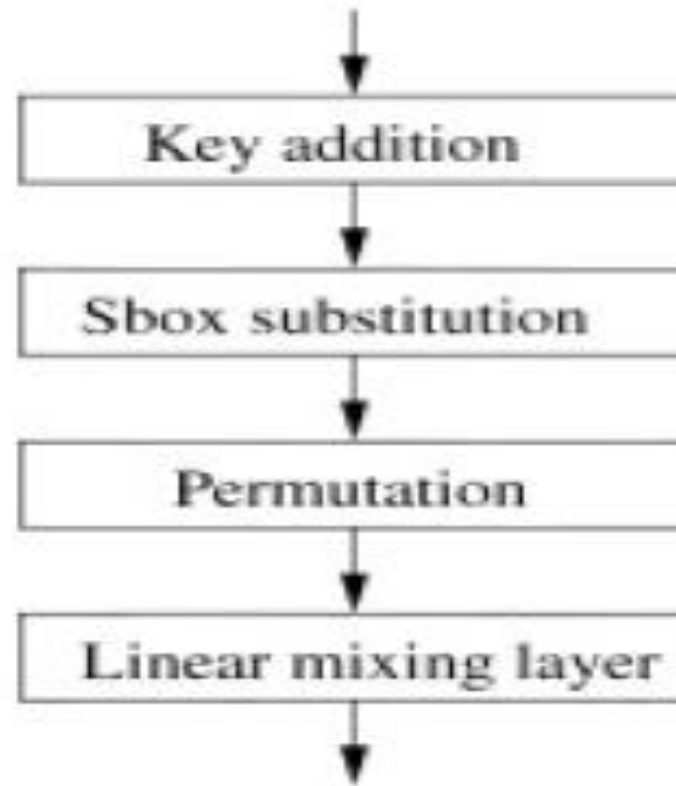


Fig. 4. A typical symmetric key block cipher structure based on the substitution-permutation network design. It contains the four fundamental operations of key addition, S-box substitution, permutation, and linear mixing.

# Conclusion

- ▶ The chaotic maps are one way to implement chaos in the field of security analysis. As can be seen the method is not at all efficient. Much more simple, efficient and robust algorithms exist for encryption.
- ▶ A major setback is that the security properties are not thorough nor are they provable given the complexity involved in the algorithms.
- ▶ But these algorithms prove the immense scope of chaos in this field.
- ▶ Specific case studies have been performed by us which will be taken up by the other members. Strict adherence to the three main points of design, evaluation and security analysis have been made.



# Case Study

1. Simple Cipher
2. Advanced Cipher
3. Baptista Method
4. Rabbit Cipher

# Simple Cipher

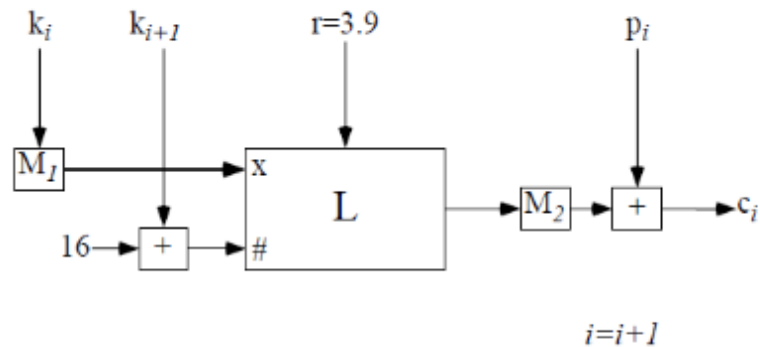
- ▶ Proposed by Roskin and Casper.
- ▶ It is a Block cipher based on unpredictability of logistic map.
- ▶ Symmetric Key Encryption Algorithm.

# Simple Cipher

## ► Keys

1. Array of 50 randomly generated numbers
2. Initial condition
3. Parameter of the logistic map

# Simple Cipher - Block Diagram



## ► Encryption

- Keys mapped from  $[0, 256]$  to range  $[0, 1]$
- The key is iterated in Block L
- Output mapped from  $[0, 1]$  to  $[0, 256]$
- Final output added to plain text

## ► Decryption

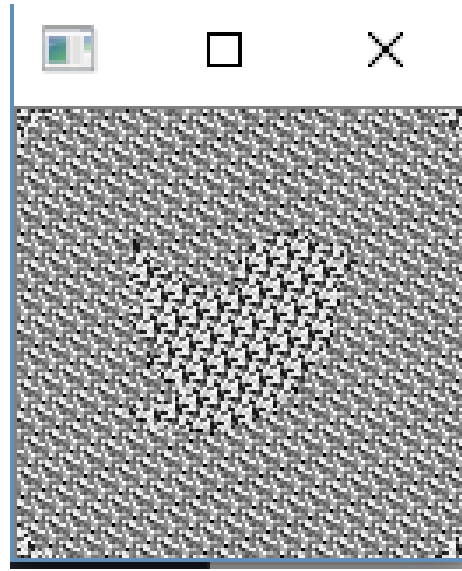
- Given the keys and cipher text, receiver can completely determine the message.

# Simple Cipher - Test Case

Input Image

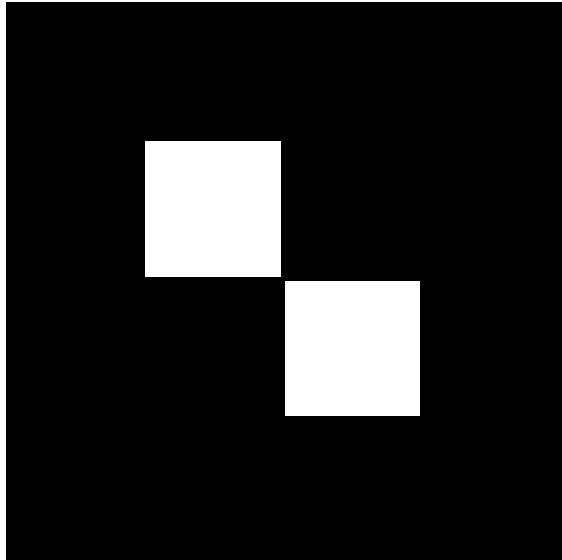


Encrypted Image

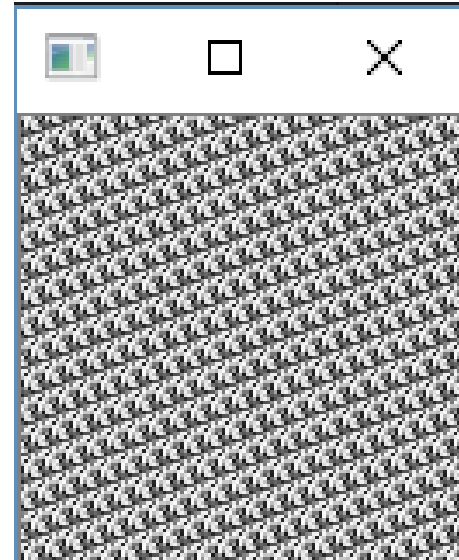


# Simple Cipher - Test case

Input Image



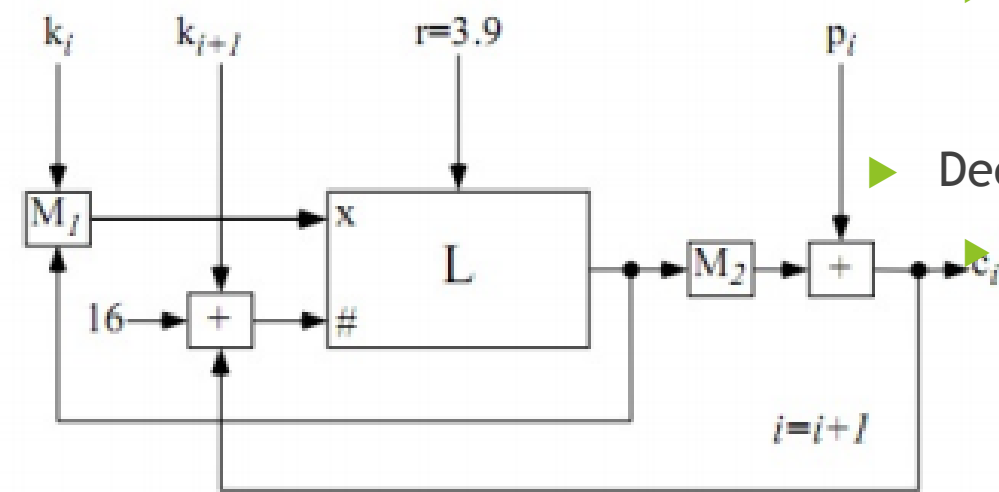
Encrypted Image



# Advanced Cipher

- ▶ Proposed by Roskin and Casper.
- ▶ Extended version of simple cipher with a feedback to reduce the periodicity in the cipher text.

# Advanced Cipher - Block Diagram



## ► Encryption

- Similar to simple cipher but with a additional feedback in the initial value and the number of iteration

## ► Decryption

- Similar to simple cipher



# Advanced Cipher - Test Case

Input Image

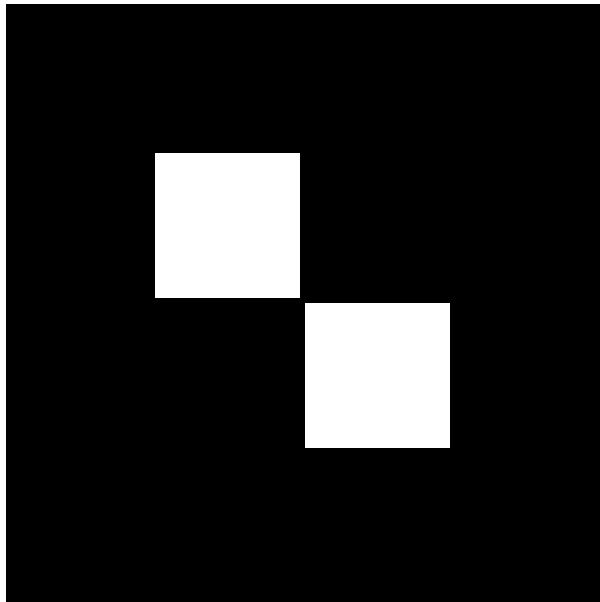


Encrypted Image

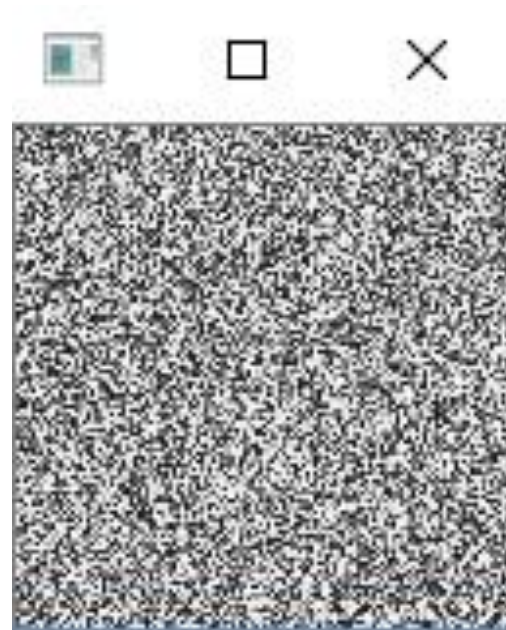


# Advanced Cipher - Test Case

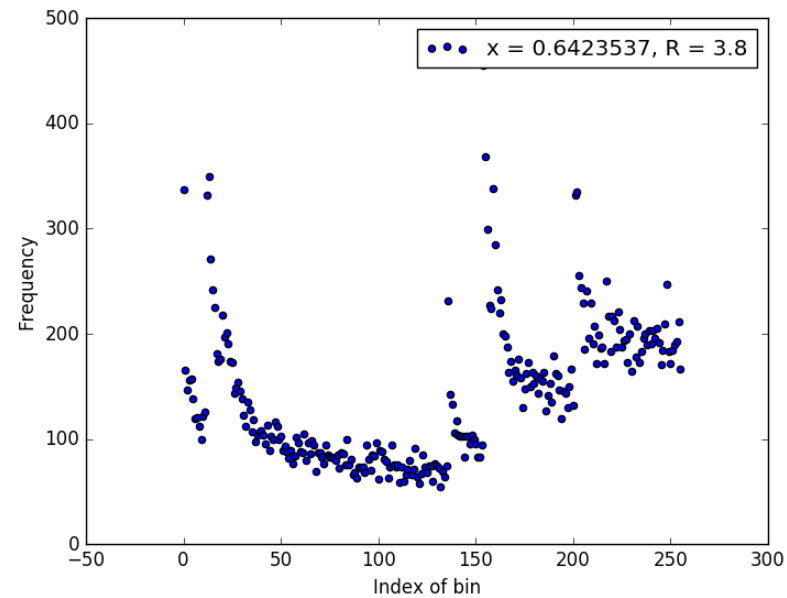
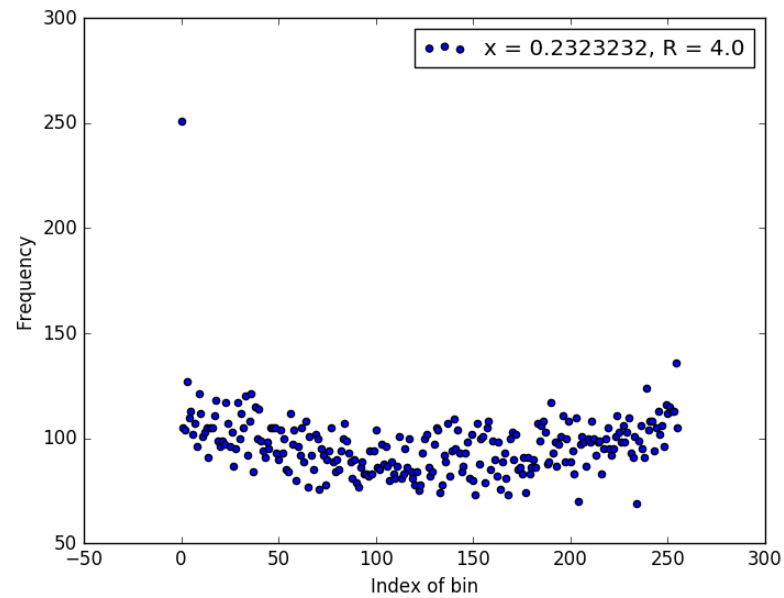
Input Image



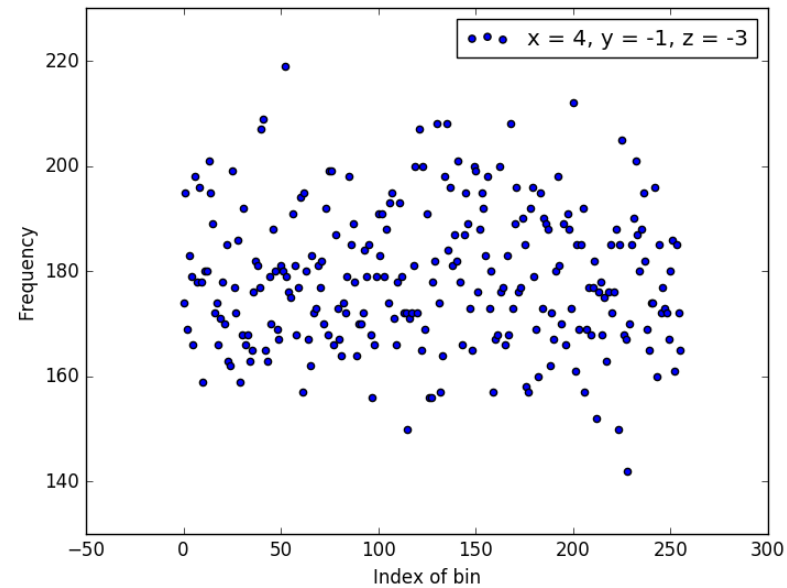
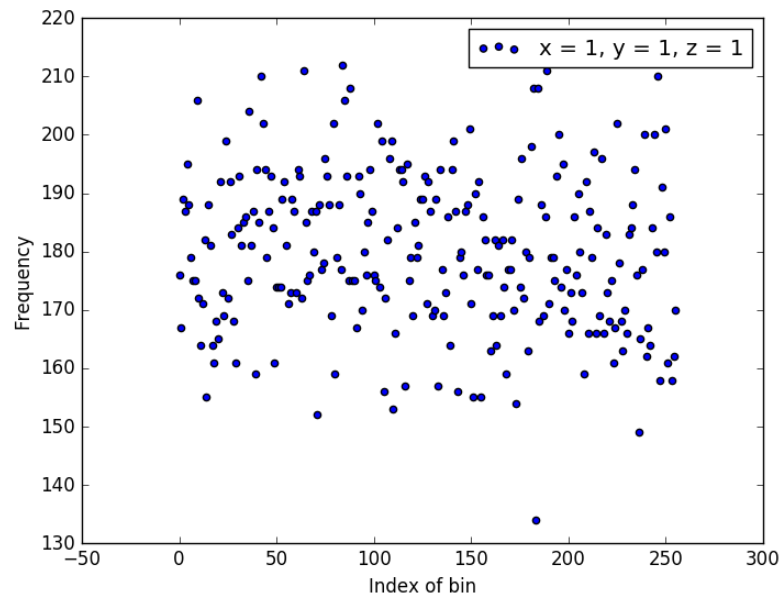
Encrypted Image



# Baptista Method - Logistic Map

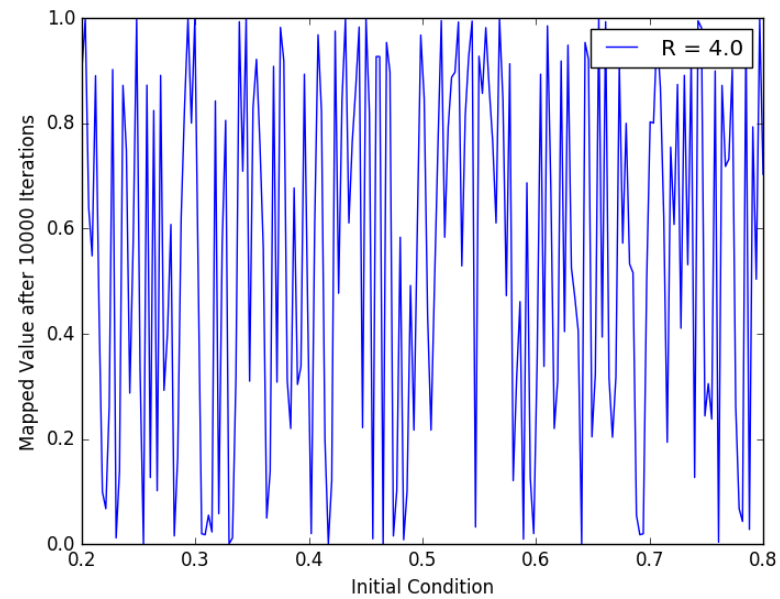


# Baptista Method - Lorenz System

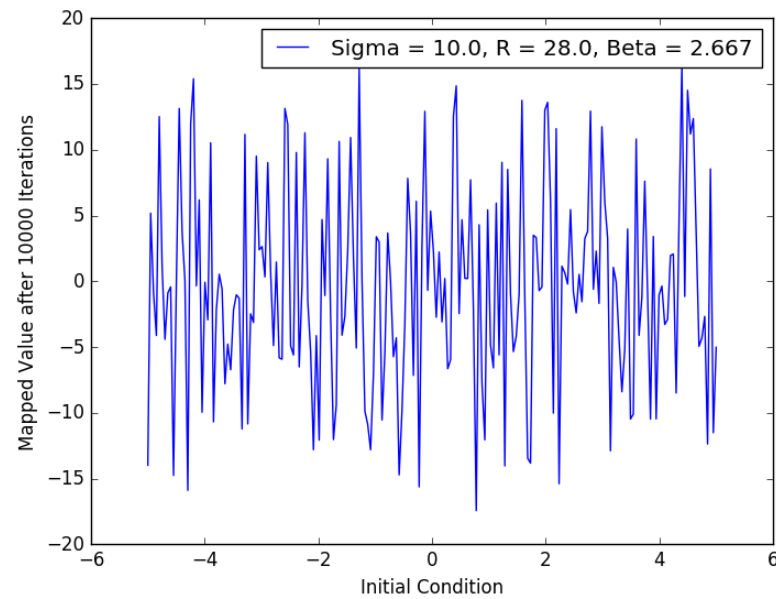


# Baptista Method

## Logistic Map



## Lorenz System



# Baptista Method

- ▶ It is a block cipher.
- ▶ Symmetric Key cipher.

# Baptista Method

## ► Keys

### ► Using Logistic Map

1. Initial value or the seed of the map
2. Parameter of logistic map

### ► Using Lorenz System

1. Initial value or the seed of the map
2. Parameter of Lorenz System
3. Variable which is used for encryption i.e. X or Y or Z
4. Bound of the phase space used

# Baptista Method

## ► Encryption

- Phase Space is divided into 256 bins.
- Bin index is the value for the particular pixel / plaintext.
- Cipher text = Number of iterations to reach a bins

## ► Decryption

- The receiver has both the cipher text and the key.
- The receiver can determine the plain text using the number of iteration.

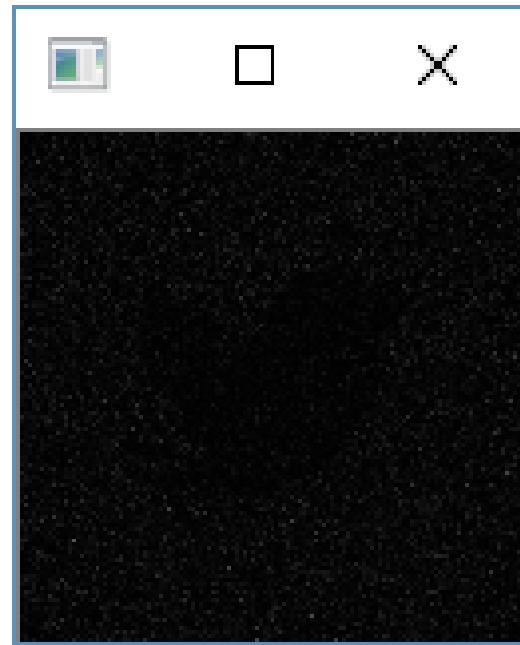


# Baptista Method - Test Case (Logistic Map)

Input Image



Encrypted Image

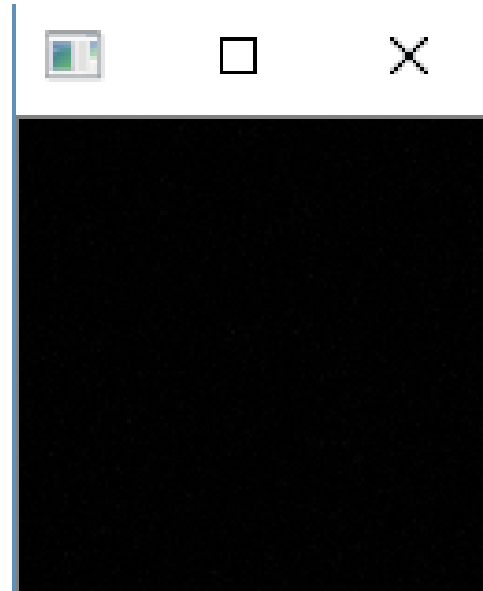


# Baptista Method - Test Case (Lorenz System)

Input Image



Encrypted Image



# Rabbit Cipher

- ▶ Inspired by “Random” behaviour of the chaotic maps.
- ▶ It takes advantage of the random-like properties of real valued chaotic maps & at the same time secures optimal cryptographic properties when discretizing them.
- ▶ Designed to work with 128 - bit data sizes, as both the key and the output are 128 - bit in length.
- ▶ Internal data structure consists of 8 state variables and 8 counters.
- ▶ Rabbit Cipher takes a 128-bit key as input and generates for each iteration an output block 128 pseudo random bits from a combination of internal state bits

# Algorithm

- ▶ The algorithm for the cipher can be broken down into the
- ▶ following four main components :
- ▶ 1) Key setup
- ▶ 2) Next state (round) function
- ▶ 3) Counter system
- ▶ 4) Extraction scheme

# Cipher Evaluation

- ▶ The design of this system is initiated by constructing a chaotic system of non-linear maps, then the system is restricted to fixed-point valued, i.e. each variable is represented by an integer type number, where a virtual decimal point is introduced manually.
- ▶ What is unique is that instead of operating over the domain of real numbers (meaning that implementation requires floating point operations), the domain is scaled up by  $2^{32}$  to translate decimal numbers into integers. This key modification to the phase space maintains the same precision as in the real interval while at the same time improving implementation performance.

# Key Setup

- ▶ The size of Internal state is 513 bits divided between 8 32- bit counters & 32-bit state variables and one counter carry bit.
- ▶ Algorithm is initialized by expanding 128-bit key into both 8 state variables and the 8 counters.
- ▶ There is one-to-one correspondence between the between the key and the initial state variables  $x_{j,0}$  and the initial counter  $c_{j,0}$

The key,  $K^{[127..0]}$ , is divided into eight subkeys:  $k_0 = K^{[15..0]}$ ,  $k_1 = K^{[31..16]}$ , ...,  $k_7 = K^{[127..112]}$ . The state and counter variables are initialized from the subkeys as follows:

$$x_{j,0} = \begin{cases} k_{(j+1 \bmod 8)} \diamond k_j & \text{for } j \text{ even} \\ k_{(j+5 \bmod 8)} \diamond k_{(j+4 \bmod 8)} & \text{for } j \text{ odd} \end{cases} \quad (1)$$

and

$$c_{j,0} = \begin{cases} k_{(j+4 \bmod 8)} \diamond k_{(j+5 \bmod 8)} & \text{for } j \text{ even} \\ k_j \diamond k_{(j+1 \bmod 8)} & \text{for } j \text{ odd.} \end{cases} \quad (2)$$

The system is iterated four times, according to the next-state function defined below, to diminish correlations between bits in the key and bits in the internal state variables. Finally, the counter values are re-initialized according to:

$$c_{j,4} = c_{j,4} \oplus x_{(j+4 \bmod 8),4} \quad (3)$$

to prevent recovery of the key by inversion of the counter system.



# Next State Function

$$x_{0,i+1} = g_{0,i} + (g_{7,i} \lll 16) + (g_{6,i} \lll 16)$$

$$x_{1,i+1} = g_{1,i} + (g_{0,i} \lll 8) + g_{7,i}$$

$$x_{2,i+1} = g_{2,i} + (g_{1,i} \lll 16) + (g_{0,i} \lll 16)$$

$$x_{3,i+1} = g_{3,i} + (g_{2,i} \lll 8) + g_{1,i}$$

$$x_{4,i+1} = g_{4,i} + (g_{3,i} \lll 16) + (g_{2,i} \lll 16)$$

$$x_{5,i+1} = g_{5,i} + (g_{4,i} \lll 8) + g_{3,i}$$

$$x_{6,i+1} = g_{6,i} + (g_{5,i} \lll 16) + (g_{4,i} \lll 16)$$

$$x_{7,i+1} = g_{7,i} + (g_{6,i} \lll 8) + g_{5,i}$$

$$g_{j,i} = ((x_{j,i} + c_{j,i})^2 \oplus ((x_{j,i} + c_{j,i})^2 \ggg 32)) \bmod 2^{32}$$

# Counter System

$$\begin{aligned}c_{0,i+1} &= c_{0,i} + a_0 + \phi_{7,i} \mod 2^{32} \\c_{1,i+1} &= c_{1,i} + a_1 + \phi_{0,i+1} \mod 2^{32} \\c_{2,i+1} &= c_{2,i} + a_2 + \phi_{1,i+1} \mod 2^{32} \\c_{3,i+1} &= c_{3,i} + a_3 + \phi_{2,i+1} \mod 2^{32} \\c_{4,i+1} &= c_{4,i} + a_4 + \phi_{3,i+1} \mod 2^{32} \\c_{5,i+1} &= c_{5,i} + a_5 + \phi_{4,i+1} \mod 2^{32} \\c_{6,i+1} &= c_{6,i} + a_6 + \phi_{5,i+1} \mod 2^{32} \\c_{7,i+1} &= c_{7,i} + a_7 + \phi_{6,i+1} \mod 2^{32}\end{aligned}$$

where the counter carry bit,  $\phi_{j,i+1}$ , is given by

$$\phi_{j,i+1} = \begin{cases} 1 & \text{if } c_{0,i} + a_0 + \phi_{7,i} \geq 2^{32} \wedge j = 0 \\ 1 & \text{if } c_{j,i} + a_j + \phi_{j-1,i+1} \geq 2^{32} \wedge j > 0 \\ 0 & \text{otherwise.} \end{cases}$$

Furthermore, the  $a_j$  constants are defined as:

$$\begin{array}{ll} a_0 = 0x4D34D34D & a_1 = 0xD34D34D3 \\ a_2 = 0x34D34D34 & a_3 = 0x4D34D34D \\ a_4 = 0xD34D34D3 & a_5 = 0x34D34D34 \\ a_6 = 0x4D34D34D & a_7 = 0xD34D34D3. \end{array}$$

# Extraction Scheme

After each iteration 128 bits of output are generated as follows:

$$\begin{array}{ll} s_i^{[15..0]} &= x_{0,i}^{[15..0]} \oplus x_{5,i}^{[31..16]} & s_i^{[31..16]} &= x_{0,i}^{[31..16]} \oplus x_{3,i}^{[15..0]} \\ s_i^{[47..32]} &= x_{2,i}^{[15..0]} \oplus x_{7,i}^{[31..16]} & s_i^{[63..48]} &= x_{2,i}^{[31..16]} \oplus x_{5,i}^{[15..0]} \\ s_i^{[79..64]} &= x_{4,i}^{[15..0]} \oplus x_{1,i}^{[31..16]} & s_i^{[95..80]} &= x_{4,i}^{[31..16]} \oplus x_{7,i}^{[15..0]} \\ s_i^{[111..96]} &= x_{6,i}^{[15..0]} \oplus x_{3,i}^{[31..16]} & s_i^{[127..112]} &= x_{6,i}^{[31..16]} \oplus x_{1,i}^{[15..0]} \end{array}$$

where  $s_i$  is the 128-bit keystream block at iteration  $i$ .

# Encryption/Decryption Scheme

- ▶ Extracted bits are XOR'ed with plaintext/ciphertext to encrypt/decrypt

$$c_i = p_i \oplus s_i,$$

$$p_i = c_i \oplus s_i,$$

- ▶ Where  $c_i$  and  $p_i$  denote the  $i^{\text{th}}$  ciphertext and plaintext, respectively.

# Security Analysis

- ▶ One-to-one correspondence between the key, the state and the counter prevents key redundancy. It also distributes the key bits in an optimal way to prepare for the the system iteration.
- ▶ After one iteration of the next state function, each key bit has affected all eight state variables.
- ▶ The key expansion scheme ensures that after two iterations of the next state function, all state bits are affected by all key bits with a measured probability of 0.5. A safety margin is provided by iterating the system 4 times.
- ▶ The counter modification makes it hard to recover the key by inverting the counter system, as this would require additional knowledge of the state variables. It also destroys the one-to-one correspondence between key and counter.
- ▶ All the output bytes of the next - state function depend on the maximal 12 input bytes. Consequently, removing any of those input bytes will result in nearly maximal entropy of error of the output bytes.

# Resulting Attacks

## ► Attack on the Key Step Function

1) Counter bits and state bits depend strongly and highly non-linearly on the bits. This makes attacks made on guessing parts of the key difficult.

2) Although counter modification destroys one-to-one correspondence between the key and counters, but the probability of two randomly selected keys to have the same counters is

$$1/2^{128}.$$

## ► Distinguishing and Correlation Attack

1) Attacker tries to distinguish a sequence generated by the cipher from a sequence of truly random numbers.

2) This type of attack is impossible because, in order to observe such a correlation, output from approximately  $2^{114}$  iterations will have to be generated.

## ► Divide and Conquer Attack

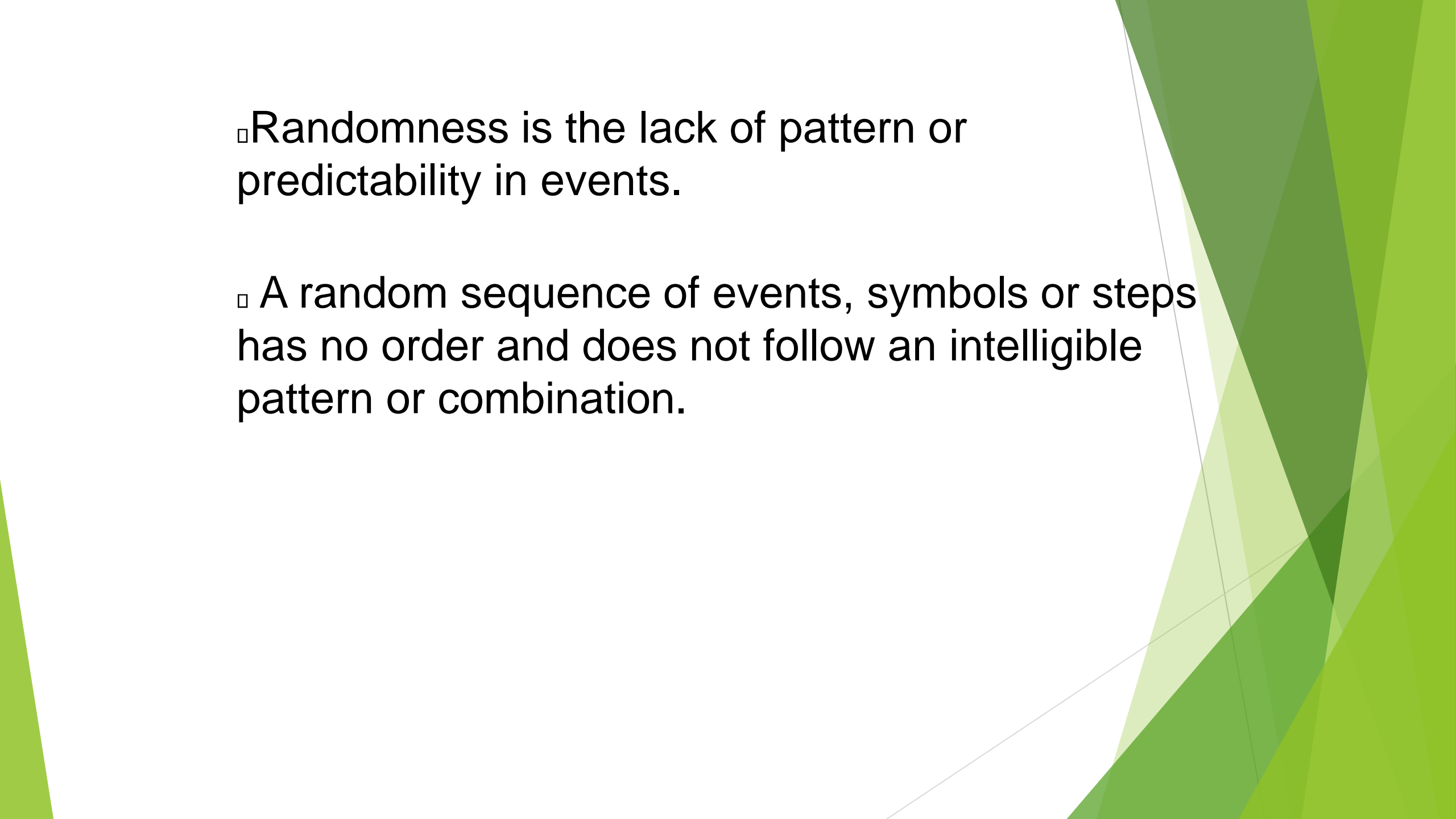
1) Attack is feasible only if a fraction of state variable is known, in order to predict a significant fraction of output bit.

2) Attacker will need to guess each of the 192 bits of input correctly.

key

# Random Numbers and their Analysis



The background of the slide features abstract, overlapping green geometric shapes, primarily triangles and polygons, in various shades of green, creating a modern and dynamic visual effect.

▣ Randomness is the lack of pattern or predictability in events.

▣ A random sequence of events, symbols or steps has no order and does not follow an intelligible pattern or combination.

- A numeric sequence is said to be statistically random when it contains no recognizable patterns or regularities.
- Sequences such as the results of an ideal dice roll, or the digits of  $\pi$  exhibit statistical randomness.

- Statistical randomness does not necessarily imply "true" randomness, i.e., objective unpredictability.
- Pseudorandomness is sufficient for many uses, such as statistics, hence the name statistical randomness.

# Generating Pseudorandom Numbers

- ▣ In our analysis, we used :
  - ▣ 1. Lorenz System
  - ▣ 2. Logistic Map
  - ▣ 3. A Simple cipher using Baptista method
  - ▣ 4. Advanced cipher using Baptista method
- ▣ to encrypt and decrypt an image.

# Lorenz System

$$\begin{pmatrix} \dot{x} \\ \dot{y} \\ \dot{z} \end{pmatrix} = \begin{pmatrix} -\sigma(x - y) \\ (r - z)x - y \\ xy - bz \end{pmatrix}$$

# Map

□

$$X_n = r * X_{n-1} * (1 - X_{n-1})$$

# Logistic

# Simple cipher

- ▣ Based on a different implementation of logistic map

# Advanced cipher

- ▣ Another, more complex implementation of logistic map.



# Testing Pseudorandom numbers

- ▣ Diehard suite of tests:
- ▣ Designed to test for quality of random number generators
- ▣ Consists of the following tests:

- Birthday spacings: Choose random points on a large interval. The spacings between the points should be asymptotically exponentially distributed.
- Overlapping permutations: Analyze sequences of five consecutive random numbers. The 120 possible orderings should occur with statistically equal probability.

▫Ranks of matrices: Select some number of bits from some number of random numbers to form a matrix over  $\{0,1\}$ , then determine the rank of the matrix. Count the ranks.

▫Monkey tests: Treat sequences of some number of bits as "words". Count the overlapping words in a stream. The number of "words" that do not appear should follow a known distribution.

▣ Count the 1s: Count the 1 bits in each of either successive or chosen bytes. Convert the counts to "letters", and count the occurrences of five-letter "words" .

▣ Parking lot test: Randomly place unit circles in a  $100 \times 100$  square. A circle is successfully parked if it does not overlap an existing successfully parked one. After 12,000 tries, the number of successfully parked circles should follow a certain normal distribution.

▣ Minimum distance test: Randomly place 8,000 points in a  $10,000 \times 10,000$  square, then find the minimum distance between the pairs. The square of this distance should be exponentially distributed with a certain mean.

▣ Random spheres test: Randomly choose 4,000 points in a cube of edge 1,000. Center a sphere on each point, whose radius is the minimum distance to another point. The smallest sphere's volume should be exponentially distributed with a certain mean.

▫ The squeeze test: Multiply  $2^{31}$  by random floats on  $(0,1)$  until you reach 1. Repeat this 100,000 times. The number of floats needed to reach 1 should follow a certain distribution.

▫ Overlapping sums test: Generate a long sequence of random floats on  $(0,1)$ . Add sequences of 100 consecutive floats. The sums should be normally distributed with characteristic mean and variance.

# Results of Tests

▣ Results for Logistic Map:





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dieharder-3.31.0: make

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🕒

🏠 Home	sts_serial	5	100000	100	0.82728997	PASSED			
📁 Downloads	sts_serial	6	100000	100	0.82911383	PASSED			
📁 Desktop	sts_serial	6	100000	100	0.56967810	PASSED			
📁 Documents	sts_serial	7	100000	100	0.60889433	PASSED			
📁 Music	sts_serial	7	100000	100	0.81407600	PASSED			
📁 Pictures	sts_serial	8	100000	100	0.31184296	PASSED	🖼️ Screenshot from 2016-10-22 20:43.png	🖼️ Screenshot from 2016-10-23 10:22:44.png	🖼️ Screenshot from 2016-10-23 12:42:25.png
📁 Videos	sts_serial	8	100000	100	0.99843546	WEAK			
📁 Other	sts_serial	9	100000	100	0.46645155	PASSED			
📁 Templates	sts_serial	9	100000	100	0.02931867	PASSED			
📁 Trash	sts_serial	10	100000	100	0.80945147	PASSED			
📁 Desktop	sts_serial	10	100000	100	0.83545328	PASSED			
📁 Downloads	sts_serial	11	100000	100	0.50491188	PASSED			
📁 Documents	sts_serial	11	100000	100	0.45548374	PASSED			
📁 Music	sts_serial	12	100000	100	0.81931992	PASSED			
📁 Pictures	sts_serial	12	100000	100	0.87411097	PASSED			
📁 Videos	sts_serial	13	100000	100	0.52285048	PASSED			
📁 Other	sts_serial	13	100000	100	0.92375324	PASSED			
📁 Templates	sts_serial	14	100000	100	0.78060314	PASSED			
📁 Trash	sts_serial	14	100000	100	0.98061154	PASSED			
📁 Desktop	sts_serial	15	100000	100	0.96488043	PASSED			
📁 Downloads	sts_serial	15	100000	100	0.65170909	PASSED			
📁 Documents	sts_serial	16	100000	100	0.61512873	PASSED			
📁 Music	sts_serial	16	100000	100	0.23950135	PASSED			
📁 Pictures	rgb_bitdist	1	100000	100	0.80197860	PASSED			
📁 Videos	rgb_bitdist	2	100000	100	0.96807445	PASSED			
📁 Other	rgb_bitdist	3	100000	100	0.77039678	PASSED			
📁 Templates	rgb_bitdist	4	100000	100	0.77033890	PASSED			
📁 Trash	rgb_bitdist	5	100000	100	0.33439519	PASSED			
📁 Desktop	rgb_bitdist	6	100000	100	0.44010926	PASSED			
📁 Downloads	rgb_bitdist	7	100000	100	0.27018817	PASSED			
📁 Documents	rgb_bitdist	8	100000	100	0.10182237	PASSED			
📁 Music	rgb_bitdist	9	100000	100	0.51095835	PASSED			
📁 Pictures	rgb_bitdist	10	100000	100	0.18759793	PASSED			
📁 Videos	rgb_bitdist	11	100000	100	0.18249321	PASSED			
📁 Other	rgb_bitdist	12	100000	100	0.91505496	PASSED			
📁 Templates	rgb_minimum_distance	2	10000	1000	0.57844191	PASSED			
📁 Trash	rgb_minimum_distance	3	10000	1000	0.86502037	PASSED			
📁 Desktop	rgb_minimum_distance	4	10000	1000	0.63070560	PASSED			
📁 Downloads	rgb_minimum_distance	5	10000	1000	0.86282200	PASSED			
📁 Documents	rgb_permutations	2	100000	100	0.62102183	PASSED			

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rgb_permutations	3	100000	100 0.36547008	PASSED				
rgb_permutations	4	100000	100 0.68097404	PASSED				
rgb_permutations	5	100000	100 0.95081949	PASSED				
rgb_lagged_sum	0	1000000	100 0.94104638	PASSED				
rgb_lagged_sum	1	1000000	100 0.64081754	PASSED				
rgb_lagged_sum	2	1000000	100 0.97265792	PASSED				
rgb_lagged_sum	3	1000000	100 0.45521773	PASSED				
rgb_lagged_sum	4	1000000	100 0.26352220	PASSED				
rgb_lagged_sum	5	1000000	100 0.12913947	PASSED				
rgb_lagged_sum	6	1000000	100 0.60681017	PASSED				
rgb_lagged_sum	7	1000000	100 0.77460412	PASSED				
rgb_lagged_sum	8	1000000	100 0.76285066	PASSED				
rgb_lagged_sum	9	1000000	100 0.86790613	PASSED				
rgb_lagged_sum	10	1000000	100 0.39853351	PASSED				
rgb_lagged_sum	11	1000000	100 0.74681952	PASSED				
rgb_lagged_sum	12	1000000	100 0.80728999	PASSED				
rgb_lagged_sum	13	1000000	100 0.30710347	PASSED				
rgb_lagged_sum	14	1000000	100 0.95507870	PASSED				
rgb_lagged_sum	15	1000000	100 0.01989391	PASSED				
rgb_lagged_sum	16	1000000	100 0.47859195	PASSED				
rgb_lagged_sum	17	1000000	100 0.19696157	PASSED				
rgb_lagged_sum	18	1000000	100 0.93003082	PASSED				
rgb_lagged_sum	19	1000000	100 0.80033529	PASSED				
rgb_lagged_sum	20	1000000	100 0.18421376	PASSED				
rgb_lagged_sum	21	1000000	100 0.61937336	PASSED				
rgb_lagged_sum	22	1000000	100 0.42586385	PASSED				
rgb_lagged_sum	23	1000000	100 0.80266953	PASSED				
rgb_lagged_sum	24	1000000	100 0.46501454	PASSED				
rgb_lagged_sum	25	1000000	100 0.85802887	PASSED				
rgb_lagged_sum	26	1000000	100 0.03269628	PASSED				
rgb_lagged_sum	27	1000000	100 0.09438973	PASSED				
rgb_lagged_sum	28	1000000	100 0.00972809	PASSED				
rgb_lagged_sum	29	1000000	100 0.46737811	PASSED				
rgb_lagged_sum	30	1000000	100 0.61454382	PASSED				
rgb_lagged_sum	31	1000000	100 0.86000314	PASSED				
rgb_lagged_sum	32	1000000	100 0.00132700	WEAK				
rgb_kstest_test	0	10000	1000 0.13488329	PASSED				
dab_bytedistrib	0	5120000	1 0.45230409	PASSED				
dab_dct	256	50000	1 0.57223651	PASSED				

Preparing to run test 207. ntuple = 0

▣ Results for lorenz map

```
Applications Sun Oct 23 5:52 PM
x Downloads: dieharder
+ /: apt-get x Downloads: dieharder
sumukhvaidya@sumukhvaidya:~/Downloads$ python lorenz.py
sumukhvaidya@sumukhvaidya:~/Downloads$ dieharder -a -f lorenzx.txt
#####
# dieharder version 3.31.0 Copyright 2003 Robert G. Brown #
#####
# rng_name | filename | [rands/second] |
# mt19937 | lorenzx.txt | 1.69e+08 |
#####
# test_name | [ntup] | [tsamples] | [psamples] | [p-value] | [Assessment] |
#####
# diehard_birthdays | 0 | 100 | 100 | 0.61272691 | PASSED |
# diehard_operm5 | 0 | 1000000 | 100 | 0.05529783 | PASSED |
# diehard_rank_32x32 | 0 | 40000 | 100 | 0.99968166 | WEAK |
# diehard_rank_6x8 | 0 | 100000 | 100 | 0.08246443 | PASSED |
# diehard_bitstream | 0 | 2097152 | 100 | 0.91627639 | PASSED |
# diehard_opso | 0 | 2097152 | 100 | 0.62513041 | PASSED |
# diehard_oqso | 0 | 2097152 | 100 | 0.18202942 | PASSED |
# diehard_dna | 0 | 2097152 | 100 | 0.39249496 | PASSED |
# diehard_count_1s_str | 0 | 256000 | 100 | 0.05002979 | PASSED |
# diehard_count_1s_byt | 0 | 256000 | 100 | 0.99989878 | WEAK |
# diehard_parking_lot | 0 | 12000 | 100 | 0.63285114 | PASSED |
# diehard_2dsphere | 2 | 8000 | 100 | 0.14012094 | PASSED |
# diehard_3dsphere | 3 | 4000 | 100 | 0.31318825 | PASSED |
# diehard_squeeze | 0 | 100000 | 100 | 0.81408529 | PASSED |
# diehard_sums | 0 | 100 | 100 | 0.68959521 | PASSED |
# diehard_runs | 0 | 100000 | 100 | 0.88975596 | PASSED |
# diehard_runs | 0 | 100000 | 100 | 0.45741447 | PASSED |
# diehard_craps | 0 | 200000 | 100 | 0.01663107 | PASSED |
# diehard_craps | 0 | 200000 | 100 | 0.04462043 | PASSED |
# marsaglia_tsang_gcd | 0 | 10000000 | 100 | 0.58286364 | PASSED |
# marsaglia_tsang_gcd | 0 | 10000000 | 100 | 0.16645147 | PASSED |
# sts_monobit | 1 | 100000 | 100 | 0.99145561 | PASSED |
# sts_runs | 2 | 100000 | 100 | 0.99929866 | WEAK |
# sts_serial | 1 | 100000 | 100 | 0.56825194 | PASSED |
# sts_serial | 2 | 100000 | 100 | 0.49834722 | PASSED |
# sts_serial | 3 | 100000 | 100 | 0.81960283 | PASSED |
# sts_serial | 3 | 100000 | 100 | 0.24333729 | PASSED |
# sts_serial | 4 | 100000 | 100 | 0.45103118 | PASSED |
# sts_serial | 4 | 100000 | 100 | 0.05502749 | PASSED |
# sts_serial | 5 | 100000 | 100 | 0.43110219 | PASSED |
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sts_serial	5	100000	100	0.96248916	PASSED
sts_serial	6	100000	100	0.82279427	PASSED
sts_serial	6	100000	100	0.64626745	PASSED
sts_serial	7	100000	100	0.44386384	PASSED
sts_serial	7	100000	100	0.33224742	PASSED
sts_serial	8	100000	100	0.06754073	PASSED
sts_serial	8	100000	100	0.75305109	PASSED
sts_serial	9	100000	100	0.28202103	PASSED
sts_serial	9	100000	100	0.89684405	PASSED
sts_serial	10	100000	100	0.40022811	PASSED
sts_serial	10	100000	100	0.37405366	PASSED
sts_serial	11	100000	100	0.13164364	PASSED
sts_serial	11	100000	100	0.88359805	PASSED
sts_serial	12	100000	100	0.11869649	PASSED
sts_serial	12	100000	100	0.32217210	PASSED
sts_serial	13	100000	100	0.28710664	PASSED
sts_serial	13	100000	100	0.92561194	PASSED
sts_serial	14	100000	100	0.57673816	PASSED
sts_serial	14	100000	100	0.48259175	PASSED
sts_serial	15	100000	100	0.70060894	PASSED
sts_serial	15	100000	100	0.25288565	PASSED
sts_serial	16	100000	100	0.78111931	PASSED
sts_serial	16	100000	100	0.84811160	PASSED
rgb_bitdist	1	100000	100	0.87480860	PASSED
rgb_bitdist	2	100000	100	0.78345241	PASSED
rgb_bitdist	3	100000	100	0.74002059	PASSED
rgb_bitdist	4	100000	100	0.12950856	PASSED
rgb_bitdist	5	100000	100	0.95022943	PASSED
rgb_bitdist	6	100000	100	0.00354592	WEAK
rgb_bitdist	7	100000	100	0.23097805	PASSED
rgb_bitdist	8	100000	100	0.46117498	PASSED
rgb_bitdist	9	100000	100	0.23153263	PASSED
rgb_bitdist	10	100000	100	0.60864526	PASSED
rgb_bitdist	11	100000	100	0.37296738	PASSED
rgb_bitdist	12	100000	100	0.34141264	PASSED
rgb_minimum_distance	2	10000	1000	0.70213145	PASSED
rgb_minimum_distance	3	10000	1000	0.84312334	PASSED
rgb_minimum_distance	4	10000	1000	0.53365538	PASSED
rgb_minimum_distance	5	10000	1000	0.74222616	PASSED
rgb_permutations	2	100000	100	0.75148530	PASSED

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poorly." In other words, we shouldn't conclude an idea is

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didn't want a Q&A service. It was more appropriate to

Answers just implemented Q&A the wrong way.

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Google's sudden change and how it changed

Google's sudden change and how it changed

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rgb_minimum_distance	3	10000	1000	0.84312334	PASSED
rgb_minimum_distance	4	10000	1000	0.53365538	PASSED
rgb_minimum_distance	5	10000	1000	0.74222616	PASSED
rgb_permutations	2	100000	100	0.75148530	PASSED
rgb_permutations	3	100000	100	0.95940183	PASSED
rgb_permutations	4	100000	100	0.75203253	PASSED
rgb_permutations	5	100000	100	0.74861822	PASSED
rgb_lagged_sum	0	1000000	100	0.80599468	PASSED
rgb_lagged_sum	1	1000000	100	0.61975360	PASSED
rgb_lagged_sum	2	1000000	100	0.28459587	PASSED
rgb_lagged_sum	3	1000000	100	0.31499871	PASSED
rgb_lagged_sum	4	1000000	100	0.99385383	PASSED
rgb_lagged_sum	5	1000000	100	0.00889537	PASSED
rgb_lagged_sum	6	1000000	100	0.33464318	PASSED
rgb_lagged_sum	7	1000000	100	0.76349269	PASSED
rgb_lagged_sum	8	1000000	100	0.47527029	PASSED
rgb_lagged_sum	9	1000000	100	0.60892043	PASSED
rgb_lagged_sum	10	1000000	100	0.12309924	PASSED
rgb_lagged_sum	11	1000000	100	0.18134351	PASSED
rgb_lagged_sum	12	1000000	100	0.96160533	PASSED
rgb_lagged_sum	13	1000000	100	0.93478146	PASSED
rgb_lagged_sum	14	1000000	100	0.89259251	PASSED
rgb_lagged_sum	15	1000000	100	0.67233860	PASSED
rgb_lagged_sum	16	1000000	100	0.98523017	PASSED
rgb_lagged_sum	17	1000000	100	0.42998003	PASSED
rgb_lagged_sum	18	1000000	100	0.39700656	PASSED
rgb_lagged_sum	19	1000000	100	0.52849105	PASSED
rgb_lagged_sum	20	1000000	100	0.62726546	PASSED
rgb_lagged_sum	21	1000000	100	0.65816634	PASSED
rgb_lagged_sum	22	1000000	100	0.99549280	WEAK
rgb_lagged_sum	23	1000000	100	0.85966971	PASSED
rgb_lagged_sum	24	1000000	100	0.99654605	WEAK
rgb_lagged_sum	25	1000000	100	0.39080301	PASSED
rgb_lagged_sum	26	1000000	100	0.55165994	PASSED
rgb_lagged_sum	27	1000000	100	0.58093275	PASSED
rgb_lagged_sum	28	1000000	100	0.13345599	PASSED
rgb_lagged_sum	29	1000000	100	0.04414923	PASSED
rgb_lagged_sum	30	1000000	100	0.35220794	PASSED
rgb_lagged_sum	31	1000000	100	0.08103711	PASSED
rgb_lagged_sum	32	1000000	100	0.49071791	PASSED



▣ Results for Logistic map: Baptista method

Applications

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```
sumukhvaidya@sumukhvaidya:~/Downloads$ dieharder -a -f logisticbaptista.txt
#####
#               dieharder version 3.31.0 Copyright 2003 Robert G. Brown               #
#####
#
#  rng_name | All | Map | filename | rands/second | Search tools |
#  mt19937 | | | logisticbaptista.txt | 1.70e+08 | |
#
#  test_name | ntup | tsamples | psamples | p-value | Assessment |
#####
#
#  diehard_birthdays | 0 | 100 | 100 | 0.80617875 | PASSED |
#  diehard_operm5 | 0 | 1000000 | 100 | 0.73292464 | PASSED |
#  diehard_rank_32x32 | 0 | 40000 | 100 | 0.31575637 | PASSED |
#  diehard_rank_6x8 | 0 | 100000 | 100 | 0.22319294 | PASSED |
#  diehard_bitstream | 0 | 2097152 | 100 | 0.17648766 | PASSED |
#  diehard_opso | 0 | 2097152 | 100 | 0.54921589 | PASSED |
#  diehard_oqso | 0 | 2097152 | 100 | 0.18112667 | PASSED |
#  diehard_dna | 0 | 2097152 | 100 | 0.48719625 | PASSED |
#  diehard_count_1s_str | 0 | 256000 | 100 | 0.84178050 | PASSED |
#  diehard_count_1s_byt | 0 | 256000 | 100 | 0.99954981 | WEAK |
#  diehard_parking_lot | 0 | 12000 | 100 | 0.61642584 | PASSED |
#  diehard_2dsphere | 0 | 8000 | 100 | 0.81208401 | PASSED |
#  diehard_3dsphere | 0 | 4000 | 100 | 0.78742498 | PASSED |
#  diehard_squeeze | 0 | 100000 | 100 | 0.69321972 | PASSED |
#  diehard_sums | 0 | 100 | 100 | 0.00140798 | WEAK |
#  diehard_runs | 0 | 100000 | 100 | 0.72107212 | PASSED |
#  diehard_runs | 0 | 100000 | 100 | 0.67394043 | PASSED |
#  diehard_craps | 0 | 200000 | 100 | 0.00876260 | PASSED |
#  diehard_craps | 0 | 200000 | 100 | 0.22240892 | PASSED |
#  marsaglia_tsang_gcd | 0 | 10000000 | 100 | 0.30360953 | PASSED |
#  marsaglia_tsang_gcd | 0 | 10000000 | 100 | 0.99998518 | WEAK |
#  sts_monobit | 1 | 100000 | 100 | 0.81063273 | PASSED |
#  sts_runs | 2 | 100000 | 100 | 0.65215572 | PASSED |
#  sts_serial | 1 | 100000 | 100 | 0.22523046 | PASSED |
#  sts_serial | 2 | 100000 | 100 | 0.93843923 | PASSED |
#  sts_serial | 3 | 100000 | 100 | 0.68374098 | PASSED |
#  sts_serial | 3 | 100000 | 100 | 0.43292104 | PASSED |
#  sts_serial | 4 | 100000 | 100 | 0.34156656 | PASSED |
#  sts_serial | 4 | 100000 | 100 | 0.91352431 | PASSED |
#  sts_serial | 7 | 100000 | 100 | 0.80202549 | PASSED |
```



Applications

Sat Oct 29 12:26 PM

(61%)

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Downloads: diehard

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Downloads: dieharder

sts_serial	5	100000	100	0.29094613	PASSED
sts_serial	6	100000	100	0.56672255	PASSED
sts_serial	6	100000	100	0.93511319	PASSED
sts_serial	7	100000	100	0.55289871	PASSED
sts_serial	7	100000	100	0.98467944	PASSED
sts_serial	8	100000	100	0.85629227	PASSED
sts_serial	8	100000	100	0.58451139	PASSED
sts_serial	9	100000	100	0.84295915	PASSED
sts_serial	9	100000	100	0.56955261	PASSED
sts_serial	10	100000	100	0.64618846	PASSED
sts_serial	10	100000	100	0.02416346	PASSED
sts_serial	11	100000	100	0.51349405	PASSED
sts_serial	11	100000	100	0.75513379	PASSED
sts_serial	12	100000	100	0.66121688	PASSED
sts_serial	12	100000	100	0.68665862	PASSED
sts_serial	13	100000	100	0.19165981	PASSED
sts_serial	13	100000	100	0.04663361	PASSED
sts_serial	14	100000	100	0.02770660	PASSED
sts_serial	14	100000	100	0.81964488	PASSED
sts_serial	15	100000	100	0.26478771	PASSED
sts_serial	15	100000	100	0.83693656	PASSED
sts_serial	16	100000	100	0.89493989	PASSED
sts_serial	16	100000	100	0.66881497	PASSED
rgb_bitdist	1	100000	100	0.70013328	PASSED
rgb_bitdist	2	100000	100	0.64597393	PASSED
rgb_bitdist	3	100000	100	0.65209182	PASSED
rgb_bitdist	4	100000	100	0.36269831	PASSED
rgb_bitdist	5	100000	100	0.75566936	PASSED
rgb_bitdist	6	100000	100	0.50338195	PASSED
rgb_bitdist	7	100000	100	0.86000575	PASSED
rgb_bitdist	8	100000	100	0.22875648	PASSED
rgb_bitdist	9	100000	100	0.60561112	PASSED
rgb_bitdist	10	100000	100	0.14021067	PASSED
rgb_bitdist	11	100000	100	0.20685415	PASSED
rgb_bitdist	12	100000	100	0.46436677	PASSED
rgb_minimum_distance	2	10000	1000	0.00991504	PASSED
rgb_minimum_distance	3	10000	1000	0.62322203	PASSED
rgb_minimum_distance	4	10000	1000	0.40281444	PASSED
rgb_minimum_distance	5	10000	1000	0.45416185	PASSED
rgb_permutations	1	100000	100	0.52329514	PASSED

ApplicationsSat Oct 29 12:26 PM(62%)Downloads: dieharder

Downloads: dieharder

rgb_permutations	3	100000	100	0.29906561	PASSED
rgb_permutations	4	100000	100	0.12650274	PASSED
rgb_permutations	5	100000	100	0.03527228	PASSED
rgb_lagged_sum	0	1000000	100	0.83069177	PASSED
rgb_lagged_sum	1	1000000	100	0.94365878	PASSED
rgb_lagged_sum	2	1000000	100	0.44801797	PASSED
rgb_lagged_sum	3	1000000	100	0.07545888	PASSED
rgb_lagged_sum	4	1000000	100	0.33952920	PASSED
rgb_lagged_sum	5	1000000	100	0.93922510	PASSED
rgb_lagged_sum	6	1000000	100	0.12595126	PASSED
rgb_lagged_sum	7	1000000	100	0.86603541	PASSED
rgb_lagged_sum	8	1000000	100	0.21981053	PASSED
rgb_lagged_sum	9	1000000	100	0.06793234	PASSED
rgb_lagged_sum	10	1000000	100	0.89019152	PASSED
rgb_lagged_sum	11	1000000	100	0.83419457	PASSED
rgb_lagged_sum	12	1000000	100	0.18632533	PASSED
rgb_lagged_sum	13	1000000	100	0.46281806	PASSED
rgb_lagged_sum	14	1000000	100	0.59220384	PASSED
rgb_lagged_sum	15	1000000	100	0.44433466	PASSED
rgb_lagged_sum	16	1000000	100	0.99210887	PASSED
rgb_lagged_sum	17	1000000	100	0.96144658	PASSED
rgb_lagged_sum	18	1000000	100	0.88379198	PASSED
rgb_lagged_sum	19	1000000	100	0.45584143	PASSED
rgb_lagged_sum	20	1000000	100	0.68922359	PASSED
rgb_lagged_sum	21	1000000	100	0.79246981	PASSED
rgb_lagged_sum	22	1000000	100	0.35085446	PASSED
rgb_lagged_sum	23	1000000	100	0.58415349	PASSED
rgb_lagged_sum	24	1000000	100	0.38476404	PASSED
rgb_lagged_sum	25	1000000	100	0.95054937	PASSED
rgb_lagged_sum	26	1000000	100	0.30076039	PASSED
rgb_lagged_sum	27	1000000	100	0.93854100	PASSED
rgb_lagged_sum	28	1000000	100	0.01542562	PASSED
rgb_lagged_sum	29	1000000	100	0.75605344	PASSED
rgb_lagged_sum	30	1000000	100	0.74029289	PASSED
rgb_lagged_sum	31	1000000	100	0.24319889	PASSED
rgb_lagged_sum	32	1000000	100	0.35886430	PASSED
rgb_kstest_test	0	10000	1000	0.05716786	PASSED
dab_bytedistrib	5120000	5120000	100	0.97988149	PASSED
dab_dct	256	50000	100	0.55797791	PASSED

Preparing to run test 207: ntuple = 0

▣ Results for simple cipher

ApplicationsSat Oct 29 1:42 PM

Downloads: dieharder

Downloads: dieharder...anced\_cipher\_data.txt

sumukhvaidya@sumukhvaidya:~/Downloads\$ dieharder -a -f simple\_cipher\_data.txt

```
#####
#
#       dieharder version 3.31.0 Copyright 2003 Robert G. Brown
#
#####
#
#  rng_name | filename | rands/second |
#-----|-----|-----|
#  mt19937 | simple_cipher_data.txt | 1.33e+08 |
#-----|-----|-----|
#
#  test_name | tntup | tsamples | psamples | p-value | Assessment
#-----|-----|-----|-----|
#
#  diehard_birthdays | 0 | 100 | 100 | 0.33347960 | PASSED
#  diehard_operm5 | 0 | 1000000 | 100 | 0.96241696 | PASSED
#  diehard_rank_32x32 | 0 | 40000 | 100 | 0.86682159 | PASSED
#  diehard_rank_6x8 | 0 | 100000 | 100 | 0.75849261 | PASSED
#  diehard_bitstream | 0 | 2097152 | 100 | 0.58330877 | PASSED
#  diehard_opso | 0 | 2097152 | 100 | 0.84448682 | PASSED
#  diehard_oqso | 0 | 2097152 | 100 | 0.89814812 | PASSED
#  diehard_dna | 0 | 2097152 | 100 | 0.27349596 | PASSED
#  diehard_count_1s_str | 0 | 256000 | 100 | 0.33582784 | PASSED
#  diehard_count_1s_byt | 0 | 256000 | 100 | 0.51635808 | PASSED
#  diehard_parking_lot | 0 | 12000 | 100 | 0.03198680 | PASSED
#  diehard_2dsphere | 2 | 8000 | 100 | 0.46363630 | PASSED
#  diehard_3dsphere | 3 | 100000 | 100 | 0.31041624 | PASSED
#  diehard_squeeze | 0 | 100000 | 100 | 0.14050698 | PASSED
#  diehard_sums | 0 | 100 | 100 | 0.00008866 | WEAK
#  diehard_runs | 0 | 100000 | 100 | 0.96827983 | PASSED
#  diehard_runs | 0 | 100000 | 100 | 0.08285685 | PASSED
#  diehard_craps | 0 | 200000 | 100 | 0.30752619 | PASSED
#  diehard_craps | 0 | 200000 | 100 | 0.31221833 | PASSED
#  marsaglia_tsang_gcd | 0 | 10000000 | 100 | 0.73358049 | PASSED
#  marsaglia_tsang_gcd | 0 | 10000000 | 100 | 0.76384898 | PASSED
#  sts_monobit | 1 | 100000 | 100 | 0.22101010 | PASSED
#  sts_runs | 2 | 100000 | 100 | 0.01363567 | PASSED
#  sts_serial | 1 | 100000 | 100 | 0.95851300 | PASSED
#  sts_serial | 2 | 100000 | 100 | 0.02583550 | PASSED
#  sts_serial | 3 | 100000 | 100 | 0.28068035 | PASSED
#  sts_serial | 3 | 100000 | 100 | 0.74400261 | PASSED
#  sts_serial | 4 | 100000 | 100 | 0.55363454 | PASSED
#  sts_serial | 4 | 100000 | 100 | 0.70402853 | PASSED
#  sts_serial | 5 | 100000 | 100 | 0.06201125 | PASSED
#  sts_serial | 5 | 100000 | 100 | 0.19647496 | PASSED
#####
```

Applications

Sat Oct 29 1:42 PM

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Downloads: dieharder

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...anced\_cipher\_data.txt

sts_serial	5	100000	100	0.06201125	PASSED
sts_serial	5	100000	100	0.19647496	PASSED
sts_serial	6	100000	100	0.93760755	PASSED
sts_serial	6	100000	100	0.17565414	PASSED
sts_serial	7	100000	100	0.92802973	PASSED
sts_serial	7	100000	100	0.86099413	PASSED
sts_serial	8	100000	100	0.66531821	PASSED
sts_serial	8	100000	100	0.86460283	PASSED
sts_serial	9	100000	100	0.71843302	PASSED
sts_serial	9	100000	100	0.83373548	PASSED
sts_serial	10	100000	100	0.35077366	PASSED
sts_serial	10	100000	100	0.06022224	PASSED
sts_serial	11	100000	100	0.79816977	PASSED
sts_serial	11	100000	100	0.49728003	PASSED
sts_serial	12	100000	100	0.25900552	PASSED
sts_serial	12	100000	100	0.86416769	PASSED
sts_serial	13	100000	100	0.74620850	PASSED
sts_serial	13	100000	100	0.73845307	PASSED
sts_serial	14	100000	100	0.88514615	PASSED
sts_serial	14	100000	100	0.99175669	PASSED
sts_serial	15	100000	100	0.25713341	PASSED
sts_serial	15	100000	100	0.96908365	PASSED
sts_serial	16	100000	100	0.29225435	PASSED
sts_serial	16	100000	100	0.67724159	PASSED
rgb_bitdist	1	100000	100	0.96807716	PASSED
rgb_bitdist	2	100000	100	0.77578280	PASSED
rgb_bitdist	3	100000	100	0.04842969	PASSED
rgb_bitdist	4	100000	100	0.15862382	PASSED
rgb_bitdist	5	100000	100	0.01055362	PASSED
rgb_bitdist	6	100000	100	0.19360240	PASSED
rgb_bitdist	7	100000	100	0.65988300	PASSED
rgb_bitdist	8	100000	100	0.85552077	PASSED
rgb_bitdist	9	100000	100	0.22872912	PASSED
rgb_bitdist	10	100000	100	0.61919611	PASSED
rgb_bitdist	11	100000	100	0.58364289	PASSED
rgb_bitdist	12	100000	100	0.88953751	PASSED
rgb_minimum_distance	2	10000	1000	0.96428893	PASSED
rgb_minimum_distance	3	10000	1000	0.36372510	PASSED
rgb_minimum_distance	4	10000	1000	0.44955668	PASSED
rgb_minimum_distance	5	10000	1000	0.54372169	PASSED

advcipher\_data.txt

advcipher.py

advcipher.txt

Anaconda3-4.1.1-Linux-x86\_64.sh

Churchill.djvu

Cryptography.pptx

hs101-COURSE-2016\_2.pdf

hs101-COURSE-2016\_2(1).pdf

libpepflashplayer.so

logistic.txt

logisticaptista.txt

lorenz.py

PassportApplicationForm\_Main\_English\_V3.0.pdf

pimean.png

Poisson.png

pycharm-edu-3.0.1.tar.gz

simple\_cipher\_data.txt

SolnRotn4.12.21.pdf

Sumukh.JPG

Applications

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...anced\_cipher\_data.txt

rgb_bitdist	10	100000	100	0.61919611	PASSED
rgb_bitdist	11	100000	100	0.58364289	PASSED
rgb_bitdist	12	100000	100	0.88953751	PASSED
rgb_minimum_distance	2	10000	1000	0.96428893	PASSED
rgb_minimum_distance	3	10000	1000	0.36372510	PASSED
rgb_minimum_distance	4	10000	1000	0.44955668	PASSED
rgb_minimum_distance	5	10000	1000	0.54372169	PASSED
rgb_permutations	2	100000	100	0.96288421	PASSED
rgb_permutations	3	100000	100	0.47646958	PASSED
rgb_permutations	4	100000	100	0.65000037	PASSED
rgb_permutations	5	100000	100	0.67181661	PASSED
rgb_lagged_sum	0	1000000	100	0.38770099	PASSED
rgb_lagged_sum	1	1000000	100	0.66589189	PASSED
rgb_lagged_sum	2	1000000	100	0.40318201	PASSED
rgb_lagged_sum	3	1000000	100	0.76688365	PASSED
rgb_lagged_sum	4	1000000	100	0.17580288	PASSED
rgb_lagged_sum	5	1000000	100	0.98375846	PASSED
rgb_lagged_sum	6	1000000	100	0.61895122	PASSED
rgb_lagged_sum	7	1000000	100	0.00499714	WEAK
rgb_lagged_sum	8	1000000	100	0.35141900	PASSED
rgb_lagged_sum	9	1000000	100	0.70201257	PASSED
rgb_lagged_sum	10	1000000	100	0.84073477	PASSED
rgb_lagged_sum	11	1000000	100	0.72417336	PASSED
rgb_lagged_sum	12	1000000	100	0.77524426	PASSED
rgb_lagged_sum	13	1000000	100	0.84896209	PASSED
rgb_lagged_sum	14	1000000	100	0.84977466	PASSED
rgb_lagged_sum	15	1000000	100	0.96568395	PASSED
rgb_lagged_sum	16	1000000	100	0.75286635	PASSED
rgb_lagged_sum	17	1000000	100	0.46369400	PASSED
rgb_lagged_sum	18	1000000	100	0.86166809	PASSED
rgb_lagged_sum	19	1000000	100	0.81119460	PASSED
rgb_lagged_sum	20	1000000	100	0.96406188	PASSED
rgb_lagged_sum	21	1000000	100	0.74030396	PASSED
rgb_lagged_sum	22	1000000	100	0.69331757	PASSED
rgb_lagged_sum	23	1000000	100	0.92129308	PASSED
rgb_lagged_sum	24	1000000	100	0.06453867	PASSED
rgb_lagged_sum	25	1000000	100	0.69922107	PASSED
rgb_lagged_sum	26	1000000	100	0.48365384	PASSED
rgb_lagged_sum	27	1000000	100	0.59755302	PASSED
rgb_lagged_sum	28	1000000	100	0.70529601	PASSED

▣ Results for Advanced cipher









Applications

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dieharder -a -f advanced\_cipher\_data.txt

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Downloads: dieharder

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...anced\_cipher\_data.txt

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rgb_bitdist	10	100000	100 0.26928932	PASSED						
rgb_bitdist	11	100000	100 0.28662366	PASSED						
rgb_bitdist	12	100000	100 0.49474232	PASSED						
rgb_minimum_distance	2	10000	1000 0.80741898	PASSED						
rgb_minimum_distance	3	10000	1000 0.62614165	PASSED						
rgb_minimum_distance	4	10000	1000 0.20455962	PASSED	dv2.png	datasim1.png	datasim2.png	datasim3.png	datasim4.png	logis1.png
rgb_minimum_distance	5	10000	1000 0.91074928	PASSED						
rgb_permutations	2	100000	100 0.92821662	PASSED						
rgb_permutations	3	100000	100 0.87499334	PASSED						
rgb_permutations	4	100000	100 0.61250580	PASSED						
rgb_permutations	5	100000	100 0.42157327	PASSED						
rgb_lagged_sum	0	1000000	100 0.32170641	PASSED						
rgb_lagged_sum	1	1000000	100 0.22831682	PASSED	logis4.png	logisbapt1.png	logisbapt2.png	logisbapt3.png	logisbapt4.png	lorenz1.png
rgb_lagged_sum	2	1000000	100 0.89962694	PASSED						
rgb_lagged_sum	3	1000000	100 0.92341344	PASSED						
rgb_lagged_sum	4	1000000	100 0.69377453	PASSED						
rgb_lagged_sum	5	1000000	100 0.03729818	PASSED						
rgb_lagged_sum	6	1000000	100 0.99530514	WEAK						
rgb_lagged_sum	7	1000000	100 0.77654511	PASSED	enxz4.png	Screenshot from 2016-10-06 00:37:02.png	Screenshot from 2016-10-22 11:20:43.png	Screenshot from 2016-10-23 10:22:44.png		
rgb_lagged_sum	8	1000000	100 0.11601075	PASSED						
rgb_lagged_sum	9	1000000	100 0.79708076	PASSED						
rgb_lagged_sum	10	1000000	100 0.16836223	PASSED						
rgb_lagged_sum	11	1000000	100 0.97368300	PASSED						
rgb_lagged_sum	12	1000000	100 0.24687878	PASSED						
rgb_lagged_sum	13	1000000	100 0.04995688	PASSED						
rgb_lagged_sum	14	1000000	100 0.71055425	PASSED						
rgb_lagged_sum	15	1000000	100 0.33713461	PASSED						
rgb_lagged_sum	16	1000000	100 0.18801083	PASSED						
rgb_lagged_sum	17	1000000	100 0.93752798	PASSED						
rgb_lagged_sum	18	1000000	100 0.74780929	PASSED						
rgb_lagged_sum	19	1000000	100 0.77036468	PASSED						
rgb_lagged_sum	20	1000000	100 0.02167148	PASSED						
rgb_lagged_sum	21	1000000	100 0.27125305	PASSED						
rgb_lagged_sum	22	1000000	100 0.25188743	PASSED						
rgb_lagged_sum	23	1000000	100 0.48114446	PASSED						
rgb_lagged_sum	24	1000000	100 0.08743945	PASSED						
rgb_lagged_sum	25	1000000	100 0.65486863	PASSED						
rgb_lagged_sum	26	1000000	100 0.34579238	PASSED						
rgb_lagged_sum	27	1000000	100 0.06829313	PASSED						
rgb_lagged_sum	28	1000000	100 0.86051575	PASSED						

# Simple Cipher

As Implemented By Ruskin and Casper.

# Analysis of Simple Cipher

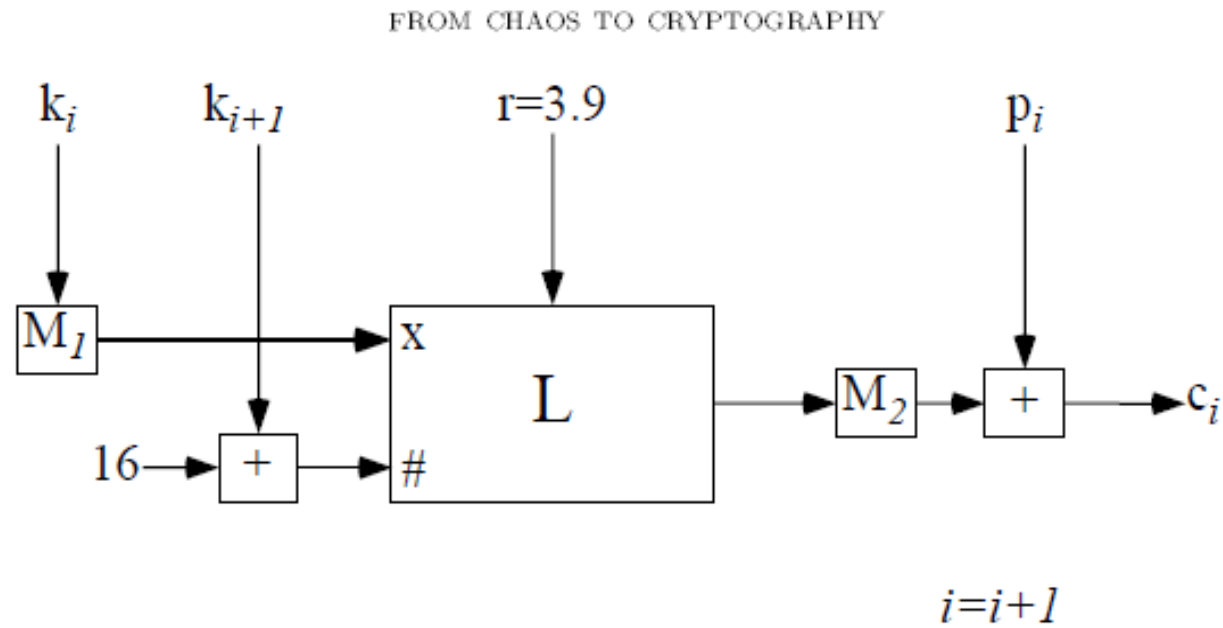


FIGURE 1. Diagram of the Simple Cipher

# Analysis of security:

- ▶ Cipher has some nice statistical properties.
- ▶ Simple cipher was implemented using algorithm by Fridrich Fri(97).

# Fridrich's Implementation

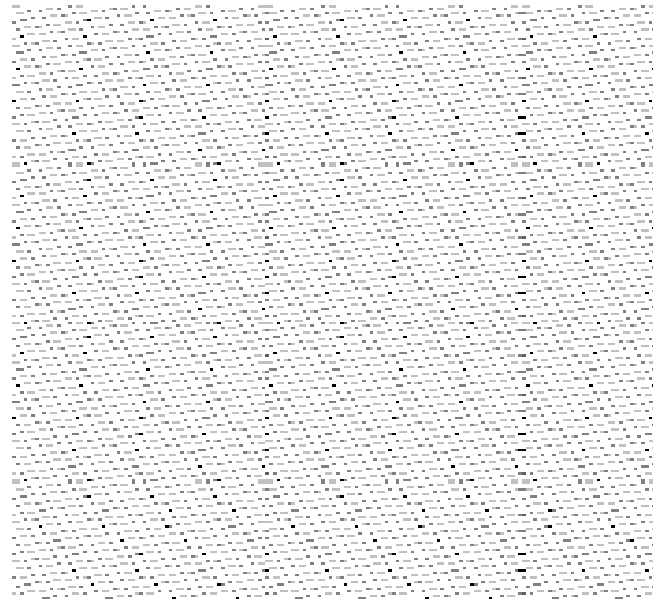


Plain Image for encryption



Encrypted File.

# Using Chaos in Simple Cipher.



Encrypted Image using Simple  
Cipher with Chaos

- Periodic nature observed introduces insecurity in the cipher.
- The pads that are produced by the cipher formed a series that created a pattern of displacement in the cipher text.



# Advanced Cipher

# Schematic Diagram

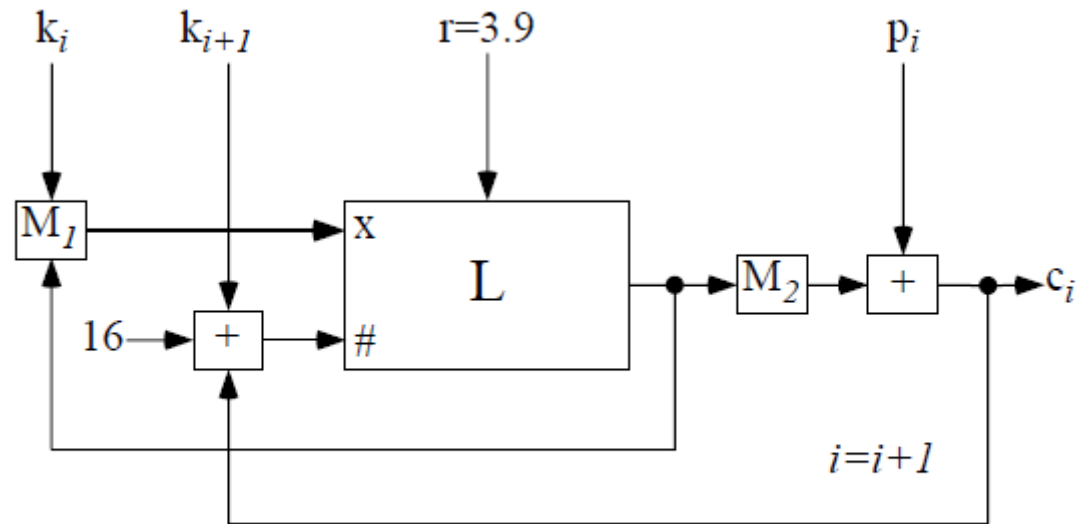


FIGURE 5. Diagram of the Advanced Cipher

# Feedback Mechanism.

- Design so as to vary the pad by both the key values and the output of the previous cipher text byte.

# Implementation of Advanced Cipher



Encrypted Image using  
Advanced Cipher

# Analysis

- The cipher has very good statistical properties.
- Additionally, it was shown that a change in a single bit in the encryption key changed, on average, 49.6% of the bits in the corresponding cipher text.
- There are residual dark spots that appear around high contrast areas of plain image.
- Advanced Cipher does far better job of encryption owing to its lack of periodicity in the cipher text.

# Baptista Algorithm using Logistic Map

---

-Choose a pair  $(r, X_0)$  and determine the interval of interest  $[X_{min}, X_{max}]$ . In this case, the interval of interest is a subset of the phase space over which the chaotic attractor resides.

-Partition the interval  $[X_{min}, X_{max}]$  into  $n$  equal sets and assign each set a unique ASCII character

**for** each character in the plaintext string  $s$  **do**

$X_0 \leftarrow rX_0 (1 - X_0)$

**while**  $X_0$  is not at the site in the interval  $[X_{min}, X_{max}]$  corresponding to the current plaintext character and  $random() > p$  **do**

$X_0 \leftarrow rX_0 (1 - X_0)$

**end while**

**end for**

# Analysis

- The implementation is simple enough.
- The performance of the algorithm depends on the randomness of the PRNG.
- Performance depends on the initial values and properties of logistic map.
- The size of 'p' should be taken care of.
- Merit: For  $r = 3.8$ , besides having the knowledge of S keys the initial conditions play crucial role in decrypting the cipher text.

# Conclusion

- ▶ Its not the most sought after cryptographic primitive. Security properties are not provable and performance is very slow.
- ▶ The relation between the two fields is subtle but very much present.
- ▶ The space of reals also poses a problem.
- ▶ The field is still emerging and the connection between the two seemingly different fields of chaos and cryptography has just been found. As mentioned earlier due to lack of provable security properties the enthusiasm is low.
- ▶ Currently number theory leads the way in cryptography. Chaos theory is a strong contender for leading chaos into the future.



# Future Work

- ▶ Though the security properties and the performance characteristics are not provable and weak, it is an evolving field and future might hold something truly wonderful which might be manifested in the unique intermixing of these two seemingly different fields but with many topological similarities.
- ▶ Future work in this field continues to this day, trying to develop more efficient algorithms. As Computers become more sophisticated and with further advances in information technology, chaos might be the way to future cryptosystems.
- ▶ From information theoretic perspective evolution of quantum cryptography and chaos hold the key to future research.

The background features abstract, overlapping green geometric shapes, primarily triangles and polygons, in various shades of green, creating a modern and dynamic visual effect. The shapes are layered, with some appearing more prominent than others, and they extend from the edges of the frame towards the center.

Thank You