

Q.1 Attribute translation grammar to generate AST.

The grammar used is as follows:

Exp \rightarrow E {print_tree(E.nodept)}

E \rightarrow E₁ + T {E₁.nodept=generate_new('+',E₁.nodept,T.nodept)}

\rightarrow E₁ - T {E₁.nodept=generate_new('-',E₁.nodept,T.nodept)}

\rightarrow T {E₁.nodept=T.nodept}

T₁ \rightarrow T₂ * F {T₁.nodept=generate_new('*',T₂.nodept,F.nodept)}

\rightarrow T₂ / F {T₁.nodept=generate_new('/',T₂.nodept,F.nodept)}

\rightarrow T₂ % F {T₁.nodept=generate_new('%',T₂.nodept,F.nodept)}

\rightarrow F {T₁.nodept=F.nodept}

F \rightarrow (E) {F.nodept=E.nodept}

\rightarrow NUMBER {F.nodept=NUMBER.nodept}

Here

Generate_new(char,node * left ,node* right) : Function that creates new node with content as char , left and right child as second and third argument respectively.

Print_tree(node * root) : Prints inorder of the AST generated.

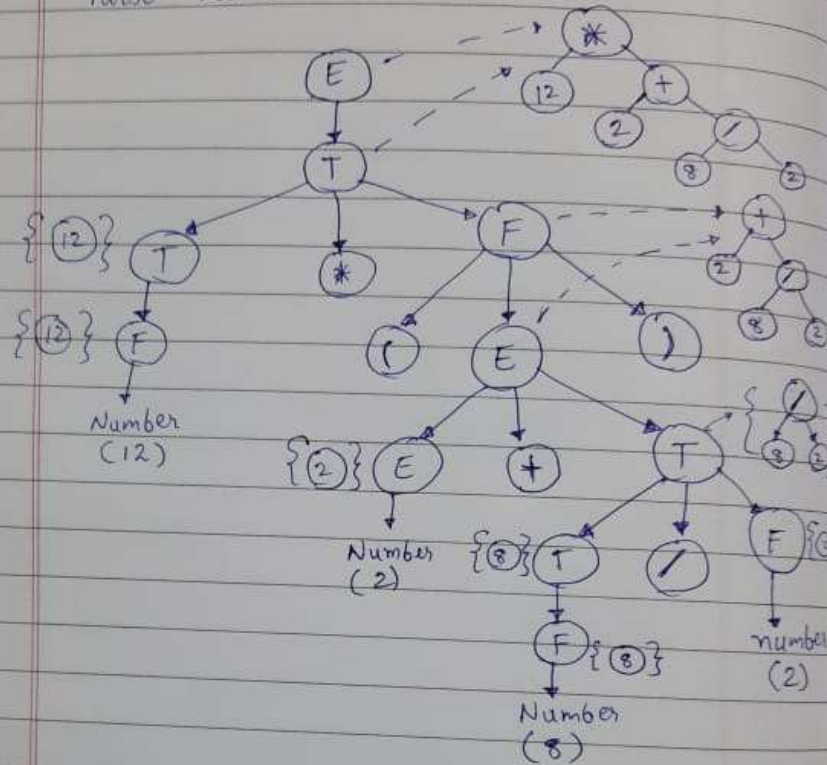
Example:

12*(2+8/2)

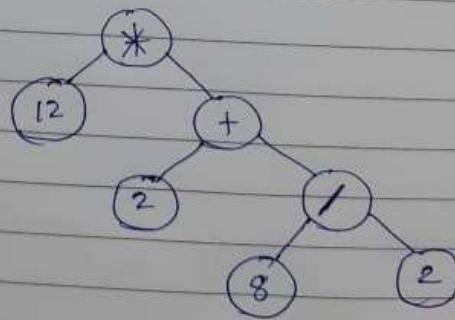
```
aneri@DESKTOP-64G0F59:/mnt/d/semester7/cd_lab/cd_lab9/ast$ ./a.out
12*(2+8/2)

Result
12
*
2
+
8
/
2
```

Parse Tree :-



Final AST:



Q.2. Attribute translation grammar to generate 3 Address Code.

The grammar is as follows:

$\text{Exp} \rightarrow E \{ \text{print}(\text{new_temp}() = E.\text{value}) \}$

$E1 \rightarrow E2 + T \{ t = \text{new_temp}(); \text{print}(t = E2.\text{value} + T.\text{value}); E1.\text{value} = t; \}$

$\rightarrow E2 - T \{ t = \text{new_temp}(); \text{print}(t = E2.\text{value} - T.\text{value}); E1.\text{value} = t; \}$

$\rightarrow T \{ E1.\text{value} = T.\text{value} \}$

$T1 \rightarrow T2 * F \{ t = \text{new_temp}(); \text{print}(t = T2.\text{value} * F.\text{value}); T1.\text{value} = t; \}$

$\rightarrow T2 / F \{ t = \text{new_temp}(); \text{print}(t = T2.\text{value} / F.\text{value}); T1.\text{value} = t; \}$

$\rightarrow T2 \% F \{ t = \text{new_temp}(); \text{print}(t = T2.\text{value} \% F.\text{value}); T1.\text{value} = t; \}$

$\rightarrow F \{ T1.\text{value} = F.\text{value} \}$

$F \rightarrow (E) \{ F.\text{value} = E.\text{value} \}$

$\rightarrow \text{NUMBER} \{ F.\text{value} = \text{lexeme}(\text{NUMBER}) \}$

$\rightarrow \text{ID} \{ F.\text{value} = \text{lexeme}(\text{ID}) \}$

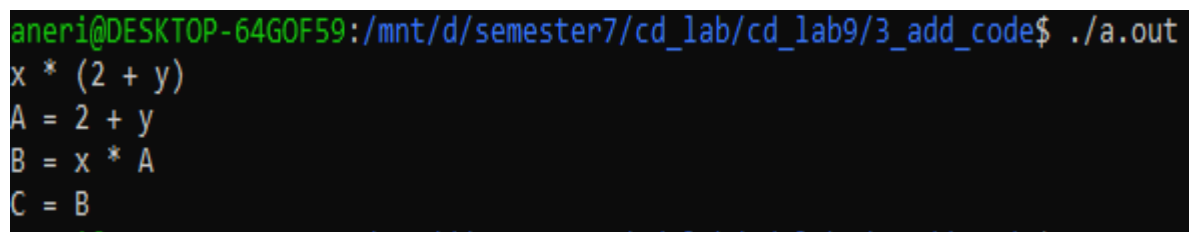
Here

3 Address code are printed as they are generated.

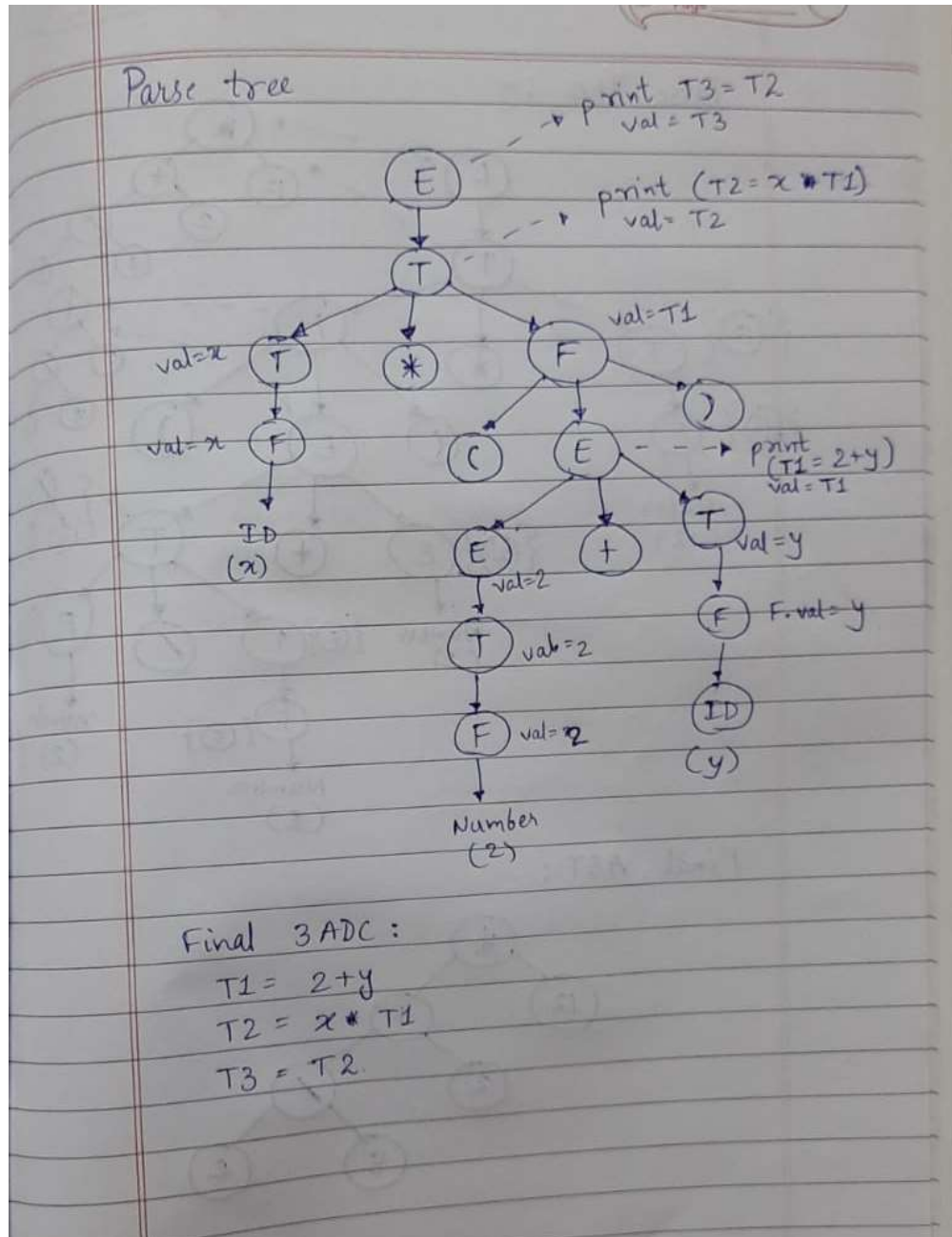
`New_temp()`: Gives a string that represents new temporary variable name.

Example:

$X * (2 + y)$



```
aneri@DESKTOP-64G0F59:/mnt/d/semester7/cd_lab/cd_lab9/3_add_code$ ./a.out
x * (2 + y)
A = 2 + y
B = x * A
C = B
```



Q.3 Attribute translation grammar to generate DAG.

The grammar is as follows:

Exp \rightarrow E {postorder(E.nodept)}

E \rightarrow E₁ + T {E₁.nodept=get_node('+',E₂.nodept,T.nodept)}

\rightarrow E₂ - T {E₁.nodept=get_node('-',E₂.nodept,T.nodept)}

->T {E1.nodept=T.nodept}

T1 ->T2 * F {T1.nodept=get_node('*',T2.nodept,F.nodept)}

->T2/F {T1.nodept=get_node('/',T2.nodept,F.nodept)}

->T2%F {T1.nodept=get_node('%',T2.nodept,F.nodept)}

->F {T1.nodept=F.nodept}

F -> (E) {F.nodept=E.nodept}

->NUMBER {F.nodept=NUMBER.nodept}

Here

Get_node(): If expression is already calculated, returns existing pointer to the node else creates a new node with content as first argument and left and right child as second and third argument.

Postorder(): Prints the nodes in postorder. If the node is already visited and is repeated in post order, (done) is printed beside it.

Eg:

(i) a + b +(a+b)

(ii) a+b+a+b

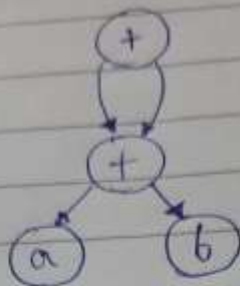
```
aneri@DESKTOP-64GOF59:/mnt/d/semester7/cd_lab/cd_lab9/dag$ ./a.out
a + b +(a+b)

Result
node: a
node: b
node: +
node: a + b (done)
node: +

aneri@DESKTOP-64GOF59:/mnt/d/semester7/cd_lab/cd_lab9/dag$ ./a.out
a+b+a+b

Result
node: a
node: b
node: +
node: a (done)
node: +
node: b (done)
node: +
```

(i)



(ii)

