Recursive Algorithms.

1) Prefix to Postfix Conversion

2) Towers of Hanoi.

Arithmetic exprusion can be represented in 3 A + {B * c} diff formets - Ez. 1) Infix A + B + A *BC 2) Prefix + AB ABC*+ AB+ 3) Postfix Goal: Use recursion to analyze the order in which expressions are to be evaluated. AB *

(with reconsists)

AB *

(with reconsists)

I nested expressions

Inay to present. Prefix: -+ *ABC/EF There are only four] - , +, *, /
possible operands (Assumption) AB * C + EF/ postfix:

Q. What is the base care? finding an operand. General Con: finding an aperator with its left and right operands. + (first openend) (seemed openend) leng tx .= ? post fix.(F) postfix (E) *AB -> AB* EFT Ex. + * AB C + (* AB) (c) protopost(C) prebpist (* AB) * A B prehipest(A) prehipest(B) * AB * C+

Preto Post Algorithm.

(Base (ane) if (expr[o] == openand)

[comination to length(expr)=1]

return expr[o];

(Use f to compute length)

(General Can)

(Use f to compute length)

3. PFI = prepost (operand1)

PF2 = prepost (operand2)

4. return[PFI -> PF2 -> oper];

Ex. XABLen 1 = find un (AB) \rightarrow return 1 AB = AB AB

A * B * C + D

$$\begin{array}{ll}
+ * * ABCD & \longrightarrow AB*C*D+ \\
findlength (** * ABCD) & (5) \\
gcane: * & findlength (* ABCD) & den! \\
expr+1+len! & \longrightarrow len2
\\
\xrightarrow{Ex} & A * B + C \\
+ * ABC & \longrightarrow AB*C+ \\
& length & \rightarrow retent 3 \\
* * AB & \longrightarrow postifix & \longrightarrow pF1 \\
finit opened scened opened = C & \longrightarrow postfix \\
& pF2 = C
\\
& PF1 & \longrightarrow PF2 & \longrightarrow + \\
& AB*C & +
\end{array}$$