

Math 321

Q's 3.7 (47)

$$p = 43 \quad q = 59$$
$$e = 13$$

$$n = 43 \cdot 59 = 2537$$

$$M = 42 \cdot 58 = 2436$$

$$C_1 = 667$$

$$C_2 = 1947$$

$$C_3 = 671$$

$$M_1 = 2_6$$

$$M_2 = 2_6$$

$$M_3 = 2_6$$

$$C = M^e \text{ mod } n$$

$$M = C^d \text{ mod } n$$

d is e's inverse mod M

$$\text{gcd}(M, e) = 1$$

$$\text{gcd}(2436, 13)$$

$$2436 = 187 \cdot 13 + 5$$

$$13 = 2 \cdot 5 + 3$$

$$5 = 1 \cdot 3 + 2$$

$$3 = 1 \cdot 2 + \boxed{1} \text{ gcd}$$

$$2 = 2 \cdot 1 + 0$$

$$1 = 3 - (1) \cdot 2$$

$$1 = 3 - (1)(5 - 3)$$

$$= (-1)(5) + (2)(3)$$

$$1 = (-1)(5) + (2)(13 - 2 \cdot 5)$$

$$= (2)(13) + (-5)(5)$$

$$1 = (2)(13) + (-5)(2436 - 187 \cdot 13)$$

$$1 = (-5)(2436) + (937)(13)$$

Mod 2436

$$P^1 = (-5)(2436) + (937)(13)$$

$$1 \equiv \boxed{937 \cdot 13} \pmod{2436}$$

\uparrow 13's inv.

$$d = 937$$

$$C_1 = 667 \quad C_2 = 1947 \quad C_3 = 671$$

$$M_1 = 2_6 \quad M_2 = 2_6 \quad M_3 = 2_6$$

$$M = C^d \text{ mod } n$$

$$d = 937 \quad n = 2537$$

$$M_1 = 667^{937} \text{ mod } 2537 = \boxed{18} \boxed{08}$$

$$M_2 = 1947^{937} \text{ mod } 2537 = \boxed{11} \boxed{21}$$

$$M_3 = 671^{937} \text{ mod } 2537 = \boxed{04} \boxed{17}$$

to encrypt.

$$\underline{e = 13} \quad \underline{n = 2537}$$

$$3.6(1) \quad (4532)_{10} =$$

n

2 ⁰	= 1	< 512
1	2	10 - 1024
2	4	11 - 2048
3	8	12 - 4096
4	16	13 - 8192
5	32	
6	64	
7	128	
8	256	

$$4532 = 4096 + 256 + 180$$

$$\begin{array}{r} 3 \quad 13 \\ 4096 \\ \underline{256} \\ 180 \end{array}$$

$$\begin{array}{r} 180 + 52 \\ \wedge \\ 32 \quad 20 \\ \wedge \\ 16 \quad 4 \end{array}$$

$$(1, 0, 0, 0, 1, 1, 0, 1, 1, 0, 1, 0, 0)_2$$

4.2 Weak Induction.

$$\left[\underbrace{P(1)}_{\text{Basis Step}} \wedge \underbrace{\forall k (P(k) \rightarrow P(k+1))}_{\text{Weak Inductive Step}} \right] \rightarrow \forall n P(n)$$

Strong Induction

$$\left[\underbrace{P(1)}_{\text{Basis Step}} \wedge \underbrace{\forall k (\{P(1) \wedge P(2) \wedge \dots \wedge P(k)\} \rightarrow P(k+1))}_{\text{Strong Inductive Step}} \right] \rightarrow \forall n P(n)$$

Ex 5 $\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \frac{1}{3 \cdot 4} + \dots + \frac{1}{n(n+1)} = a_n$

$$n=1 \quad a_1 = \frac{1}{2}$$

$$n=2 \quad a_2 = \frac{1}{2} + \frac{1}{6} = \frac{3}{6} + \frac{1}{6} = \frac{4}{6} = \frac{2}{3}$$

$$n=3 \quad a_3 = \frac{1}{2} + \frac{1}{6} + \frac{1}{12} = \frac{2}{3} + \frac{1}{12} = \frac{9}{12} = \frac{3}{4}$$

$$n=4 \quad a_4 = \frac{1}{2} + \frac{1}{6} + \frac{1}{12} + \frac{1}{20} = \frac{3}{4} + \frac{1}{20} = \frac{16}{20} = \frac{4}{5}$$

Conjecture.

$$a_n = \frac{1}{2} + \frac{1}{6} + \dots + \frac{1}{n(n+1)} = \frac{n}{n+1}$$

th⁴: $\forall n \left(\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \frac{1}{3 \cdot 4} + \dots + \frac{1}{n(n+1)} = \frac{n}{n+1} \right)$

pf: Base Step: $P(1)$: let $n=1$

left = $\frac{1}{1 \cdot 2} = \frac{1}{2}$
 right = $\frac{1}{2} = \frac{1}{2}$ } equal! $P(1)$ is true.

Inductive Step: $\forall k (P(k) \rightarrow P(k+1))$

$P(k)$: " $\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \dots + \frac{1}{k(k+1)} = \frac{k}{k+1}$ "

$P(k+1)$: " $\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \dots + \frac{1}{k(k+1)} + \frac{1}{(k+1)(k+2)} = \frac{k+1}{k+2}$ "

assume: $\frac{1}{1 \cdot 2} + \dots + \frac{1}{k(k+1)} = \frac{k}{k+1}$ is true

$$\frac{1}{1 \cdot 2} + \dots + \frac{1}{k(k+1)} + \frac{1}{(k+1)(k+2)} = \frac{k}{k+1} + \frac{1}{(k+1)(k+2)}$$

$$\rightarrow \frac{1}{1 \cdot 2} + \dots + \frac{1}{(k+1)(k+2)} = \frac{k(k+2) + 1}{(k+1)(k+2)} = \frac{k^2 + 2k + 1}{(k+1)(k+2)}$$

$$\rightarrow \frac{1}{1 \cdot 2} + \dots + \frac{1}{(k+1)(k+2)} = \frac{\cancel{(k+1)}(k+1)}{\cancel{(k+1)}(k+2)} = \frac{k+1}{k+2}$$

$$\rightarrow \frac{1}{1 \cdot 2} + \dots + \frac{1}{(k+1)(k+2)} = \frac{k+1}{k+2}$$

So inductive step is true.

i. $\forall n P(n)$ is True.

boards.



2x2 or

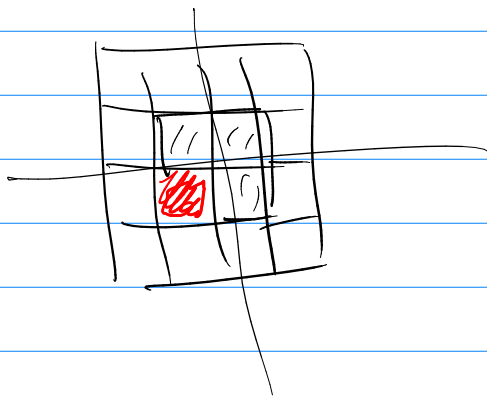
4x4 or


8x8 or

$2^n \times 2^n$

if you remove 1 tile.

then the entire board
can be tiled with



an L-  piece.