

1. Consider a preemptive, priority-based system consisting of three periodic tasks, A, B, and C, with the following properties:

Task	Run Time ( $C_i$ )	Period ( $T_i$ )	Deadline ( $D_i$ )
A	1	4	4
B	5	12	12
C	9	36	36

Suppose that tasks are assigned priorities using a rate monotonic priority assignment.

- (a) Is the task set feasible based on a Utilization Based Test? Yes. Explain briefly.  
Hint:  $3(2^{1/3} - 1) \approx 28/36$  and  $1.0 = 36/36$ .

The task set is simply periodic,  $4 \mid 36$  and  $12 \mid 36$ , so the schedulable utilization is 1.0.  
 $U = 1/4 + 5/12 + 9/36 = 33/36 \leq 1.0$ , so the task set is feasible.

- (b) Compute the worst-case response time for task B using Response Time Analysis: \_\_\_\_\_.  
Show work below.

$W_B^0 = 5$ ,  $W_B^1 = 5 + \text{ceil}(5/4) * 1 = 7$ ,  $W_B^2 = 5 + \text{ceil}(7/4) * 1 = 7$ , so WCRT of B = 7.

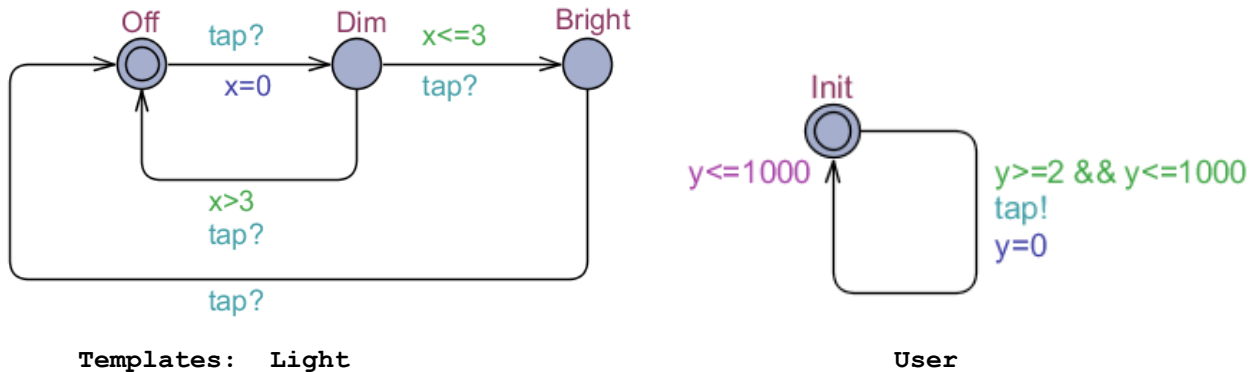
- (c) Suppose that **only** task B needs to access a resource X for **k** time units. What is the maximum value of **k** that can be used and still ensure that the task set remains feasible if access is controlled using NonPreemptable Critical Sections (NPCS)? \_\_\_\_\_. Note that **k** is in the range 0 to 5 because the run-time of task B is 5. Explain briefly.

Using NPCS, task B will block all higher priority tasks when it holds a lock on X for k time units. Since task A can only tolerate up to 3 units of blocking time (it needs one unit to complete its own job), so the maximum value of k is 3.

- (d) Suppose that **only** task B needs to access a resource X for **k** time units. What is the maximum value of **k** that can be used and still ensure that the task set remains feasible if access is controlled using the Original Priority Ceiling Protocol (OPCP)? \_\_\_\_\_. Explain briefly.

With OPCP, task B will only block tasks that also use the shared resource, so it will have no impact on the other tasks. Thus, task A will be able to preempt task B at any time, and k can be as large as B's run-time which is 5.

2. Consider the following timed automata, Light and User. They communicate via a global channel called tap. Light has a local clock x, User has a local clock y and a location invariant of  $y \leq 1000$ .



```
// Global Declarations:
chan tap;

// System Declarations
System Light, User;
```

- (a) Determine which of the following properties are satisfied for the automata shown above; circle all properties that are satisfied, explain **briefly** why the properties that are not circled are not satisfied: **SATISFIED IN RED**

- A[] not deadlock
- Light.Off --> Light.Dim
- Light.Dim --> Light.Bright - if the user never clicks within 3 time units, then the light will just go from Off to Dim to Off, etc.
- Light.Bright --> Light.Dim
- $E \langle \rangle (\text{Light.x} == 1000 \text{ and } \text{User.y} == 1000)$
- $A[] (\text{Light.x} == 2003) \text{ imply } \text{Light.Off}$
- $E \langle \rangle (\text{Light.x} == 2003 \text{ and } \text{User.y} == 1000)$
- $E \langle \rangle (\text{Light.x} == 2004)$  - the max for Light.x is 2003, in Dim for 3, Bright for 1000, and Off for 1000

- (b) As these automata have two clocks, `Light.x` and `User.y`, the reachable state space with respect to the clocks can be viewed as a point in a two-dimensional Cartesian plane, one axis for clock `Light.x` and one axis for clock `User.y`. A point  $(a, b)$ , with non-negative  $a$  and  $b$ , can be used to denote that clock `Light.x` equals  $a$  and clock `User.y` equals  $b$ . Determine the reachable state space for this automaton; e.g., draw a 2-d graph.



3. Consider the following Promela model:

```

int x,y;

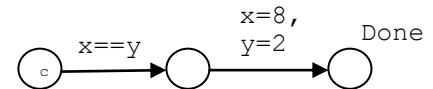
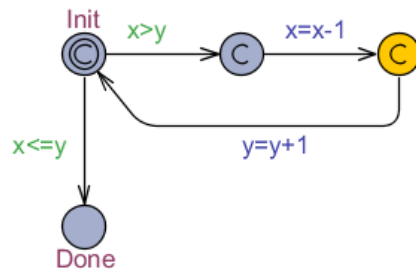
proctype p() {
  do
    :: (x > y) -> x=x-1; y=y+1;
    :: else -> break;
  od;
}

proctype q() {
  (x == y);
  atomic{ x=8; y=2; }
}

init
{
  atomic { x=0; y=0; run p(); run q(); }
}

```

- (a) Proctype **p()** can be modeled in UPAAL using the following template, draw an equivalent UPAAL template for proctype **q()**, mark the initial state **Init**, and the final state **Done**.



- (b) Will the processes terminate; e.g., is the property **A<>(p.Done and q.Done)** satisfied? Explain briefly. **Yes**
- (c) What are the possible final values for **x** and **y**? **x=5,y=5, or x=8,y=2**
- (d) What if the atomic statement is removed around **{ x=8; y=2; }**? Does your model for proctype **q()** need to be changed? If so, draw the new model below. Also, in the following list, circle all of the possible final values for **x** and **y**: **Yes, add another state and transition, the existing transition for x=8, and the new next transition for y=2.**

- $x=8, y=2$
- $x=8, y=0$
- $x=4, y=5$
- $x=5, y=5$
- $x=4, y=2$
- $x=3, y=3$