CIS 721 - Real-Time Systems

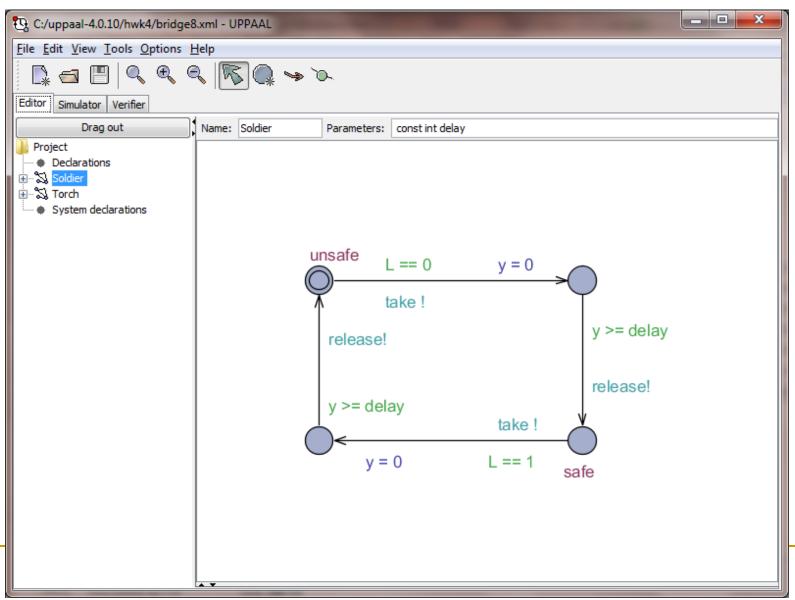
Lecture 31: Advanced UPPAAL Models

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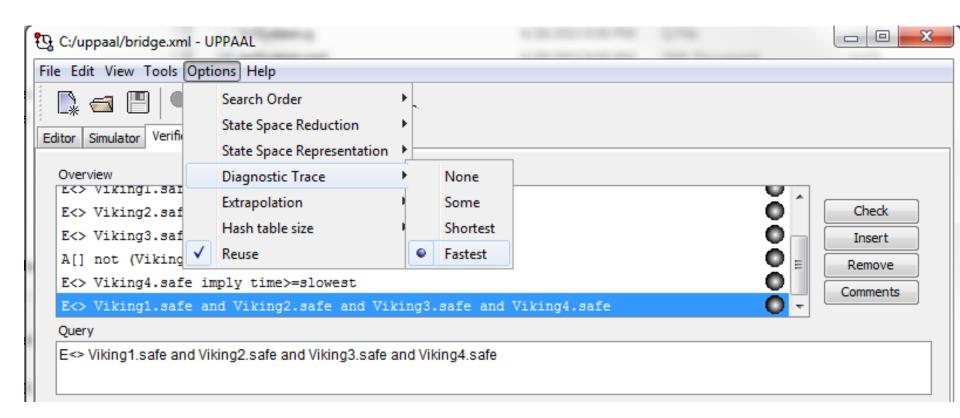
Uppaal Models

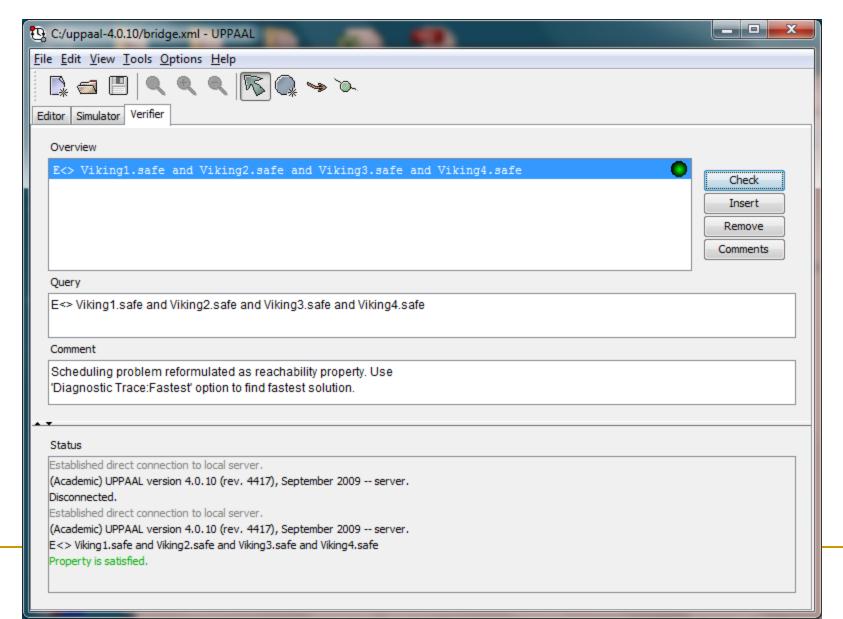
- Vikings' Bridge Problem
- Gossiping Girls' Problem
- Prioritized Token-Based Mutual Exclusion Algorithms

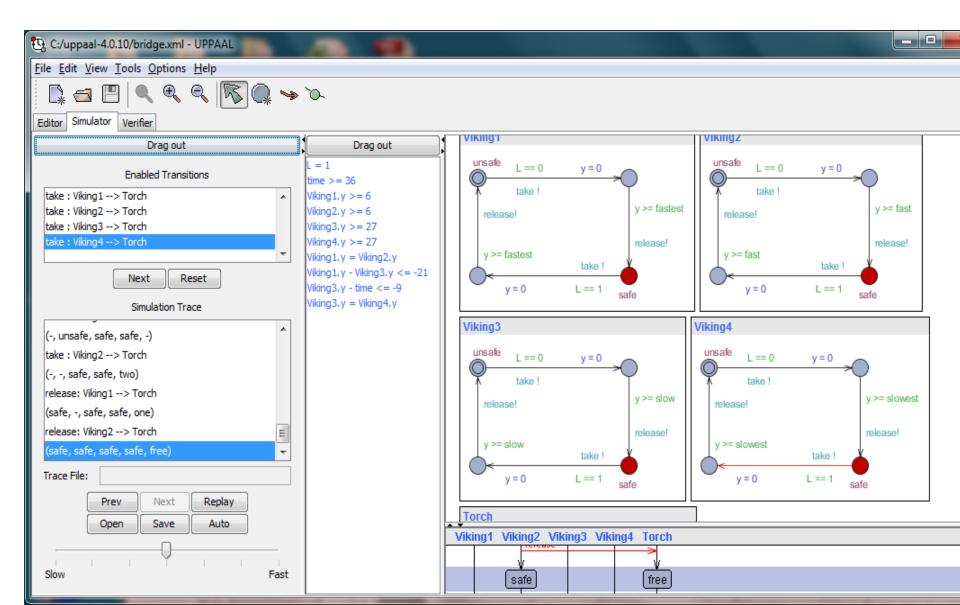
- Adjust the solution to the Vikings' Problem given in bridge.xml to handle four vikings with crossing times of 3, 6, 12, and 15. Determine the minimum time required for all four vikings to cross the bridge.
- Repeat the problem, but vary the time required for the third viking (currently 12); e.g., try 3, 6, 9, and 15, etc.
- What range of values require the same amount of time as the first run with crossing times of 3, 6, 12, and 15?
- Finally, develop a general solution (on paper) for 4 vikings with crossing times of k, 2*k, 4*k, and 5*k, for any fixed k > 0.

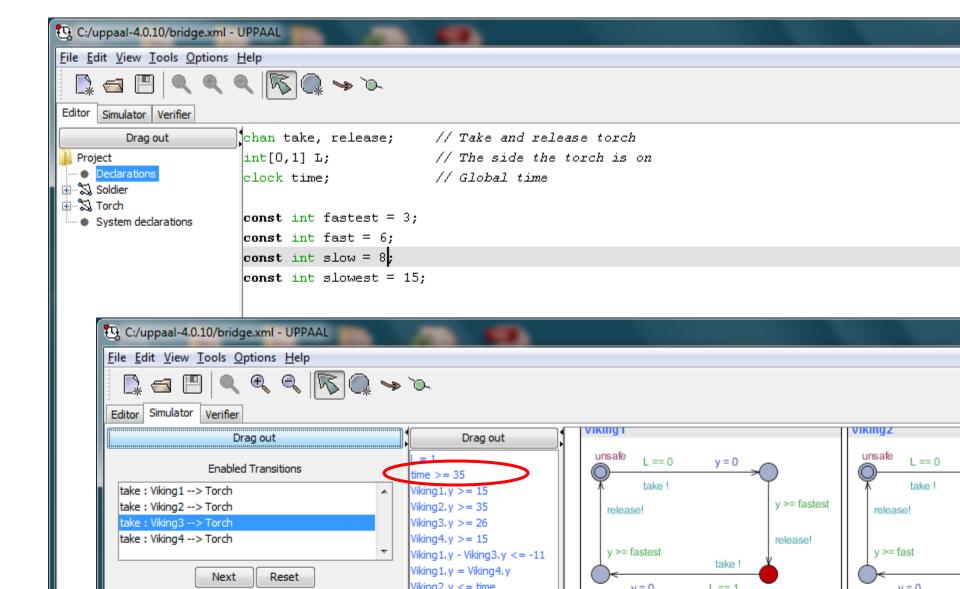


Diagnostic Trace + Fastest





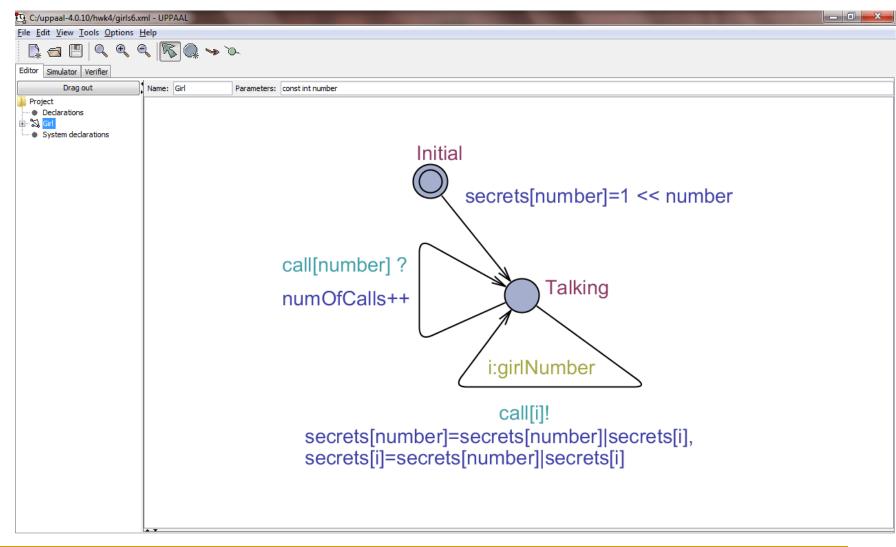


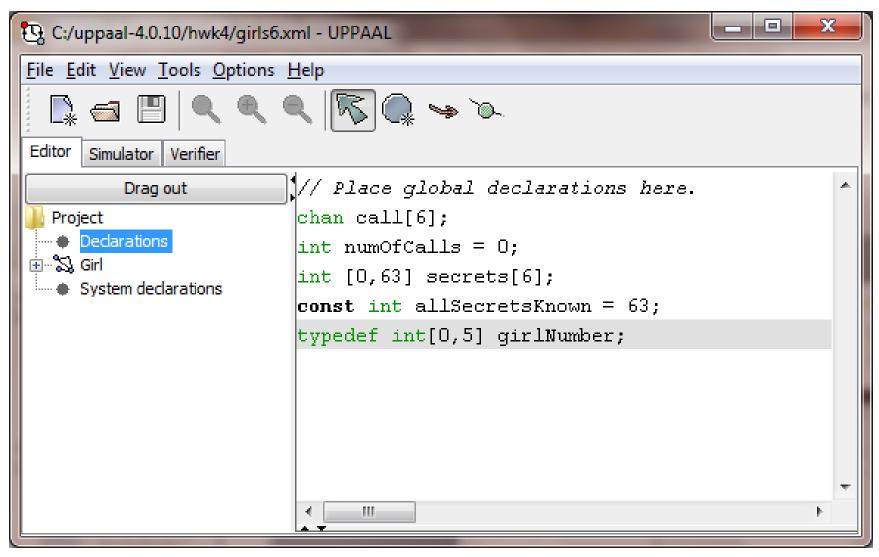


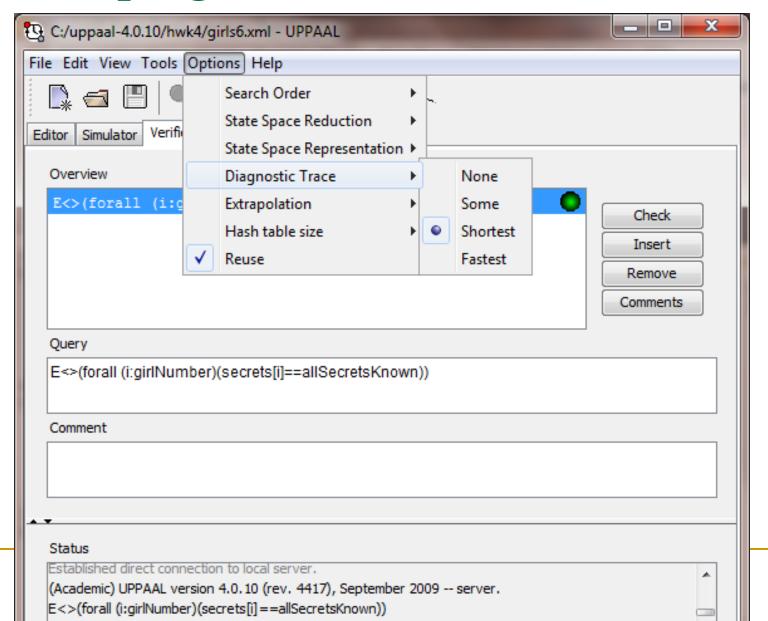
- Values from 9 to 15 for the third viking require the same time, 36, for all four to cross.
- Generalize:
 - □ k=3: 3, 6, 12, 15 => 36 crossing time
 - \sim k=4: 4, 8, 16, 20 => 48 crossing time
 - □ k=5: 5, 10, 20, 25 => 60 crossing time

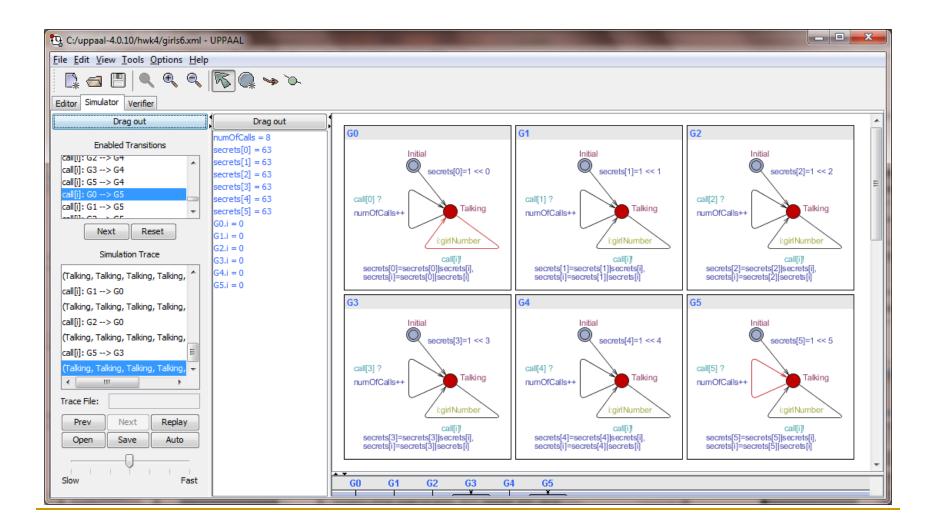
So, total time = 12*k, for crossing times of k, 2*k, 4*k, 5*k

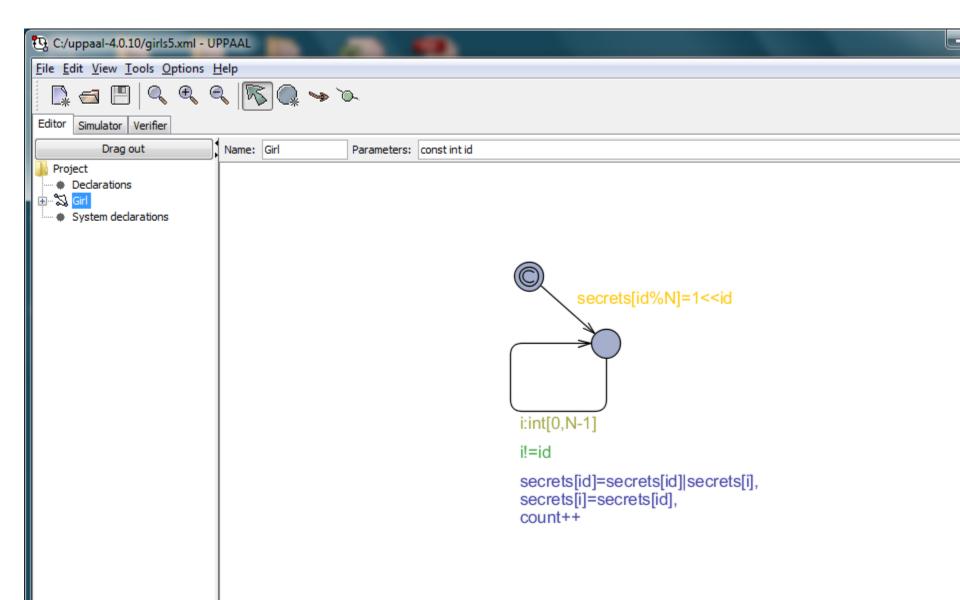
- Model and analyze the following gossiping girls problem in UPPAAL. A number of girls initially know one distinct secret each. Each pair of girls has access to a phone line that can be used to share their secrets. Each time two girls talk to each other they always exchange all secrets that they currently know with each other (thus after the phone call they both know all secrets they knew together before the phone call). Only binary communication (between two girls in each phone call) is possible.
- (a) Use UPPAAL to find the minimal number of phone calls needed for two, three, or four girls to know all secrets.
- (b) Analytically determine how many phone calls are needed to solve the gossiping girls' problem for n girls. Hint: Use a recursive or inductive approach. Consider the solution for n = 2, n = 3, etc. Then, write a proof by induction.











- Special case, n = 4: 4 calls required.
- Other cases:
 - n = 2 = 1 call
 - n = 3 = 3 calls
 - n = 4 = 4 calls
 - = n = 5 = 6 calls
 - n = 6 = 8 calls
 - n = 7 = 10 calls
 - n = 8 = 12 calls
 - □ For (n >= 4), 2*n 4 calls

Distributed Mutual Exclusion Algorithms

Permission based

Lamport

Ricart, et.al.

Maekawa

Sanders

Agrawal, et al.

Singhal

Raynal

Token based

Logical structure

No logical structure

Suzuki, et al.

Ricart, et al.

Singhal

Helary, et al.

Dynamic

(path reversal)

Trehel, et al.

Bernabeu-Auban, et al.

Ginat, et al.

Static

(edge reversal)

Raymond

van de Snepscheut

Neilsen, et al.

Generalized Algorithm



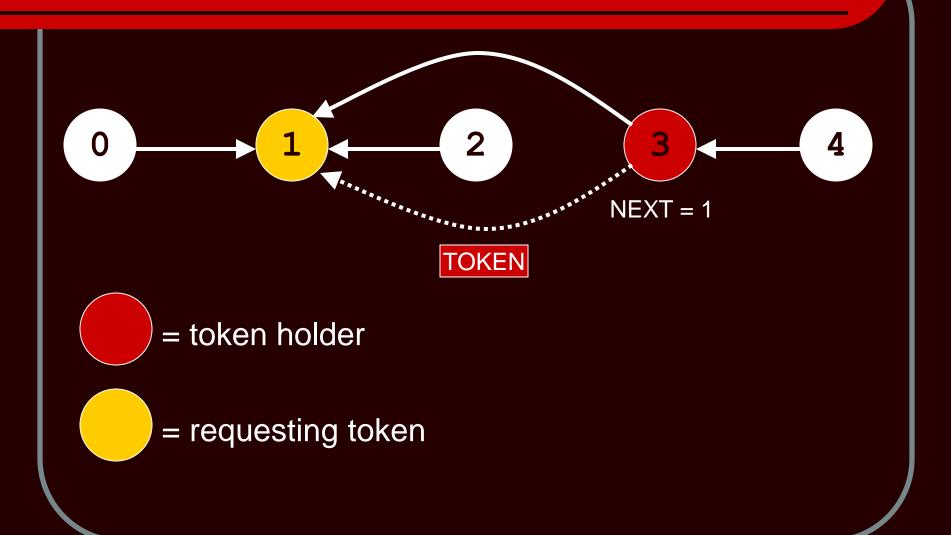
= token holder

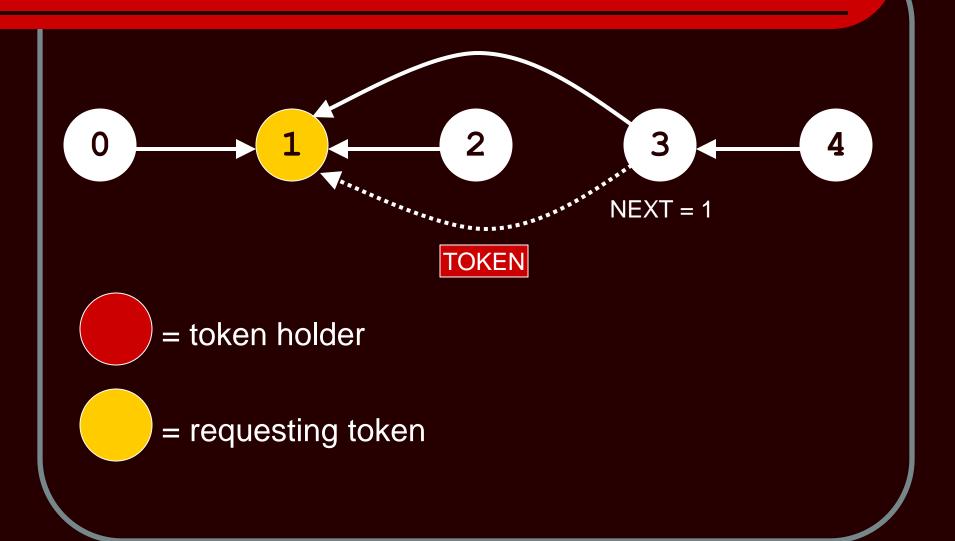


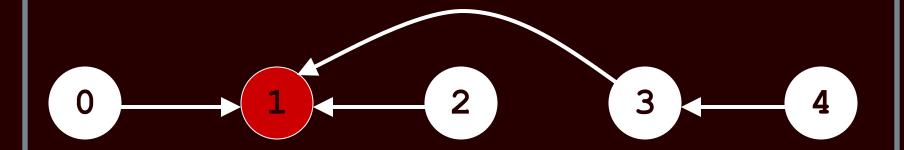
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= token holder





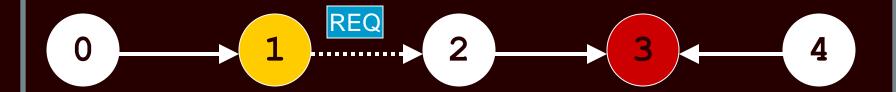








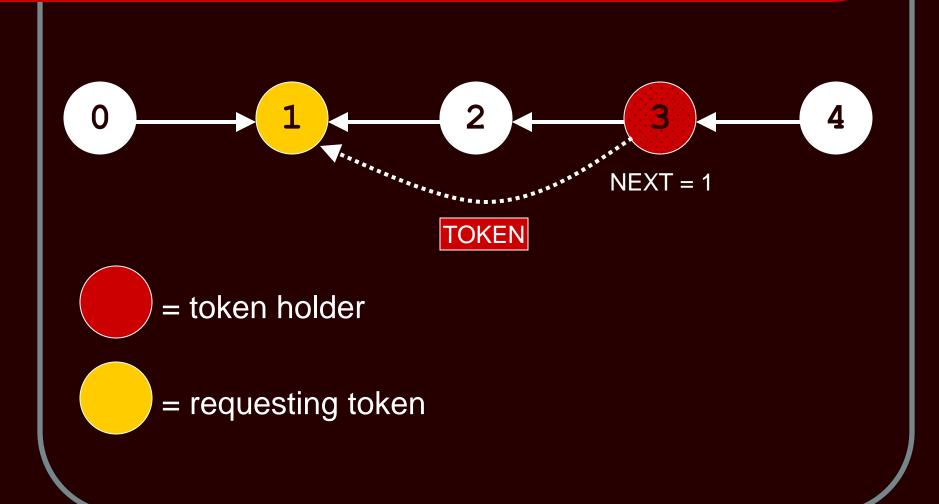
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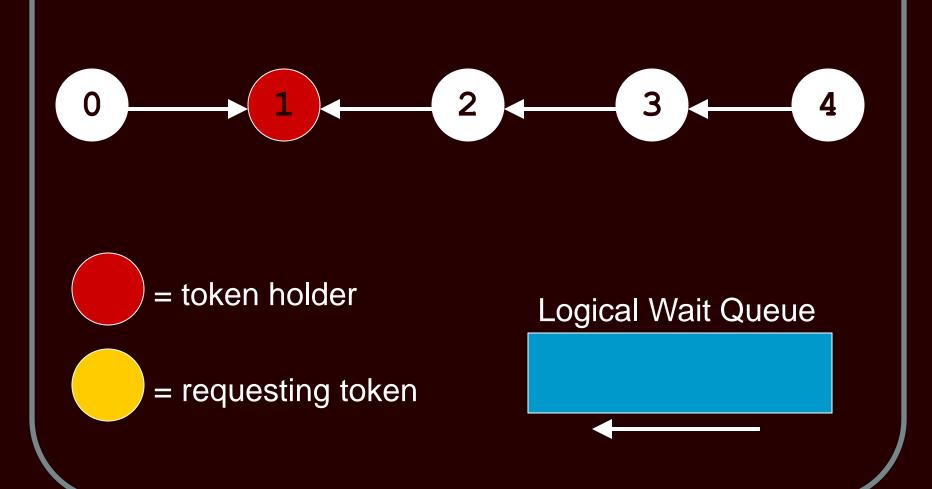


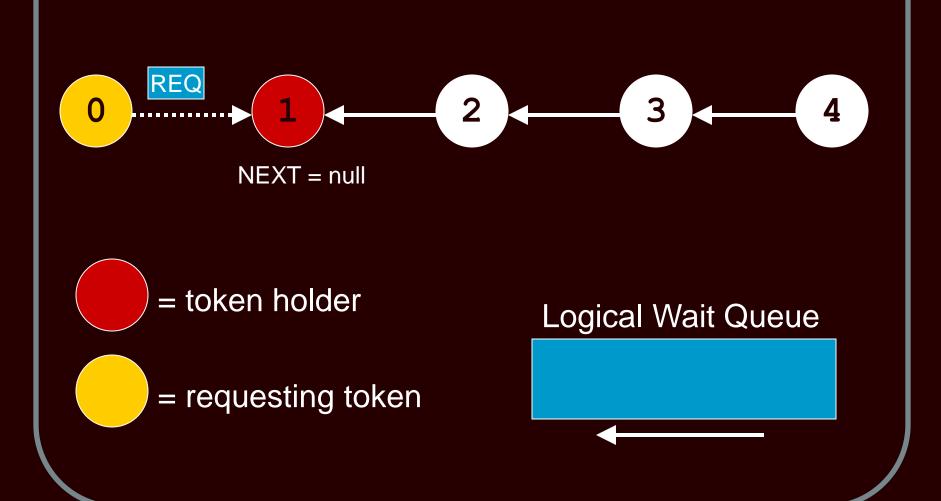
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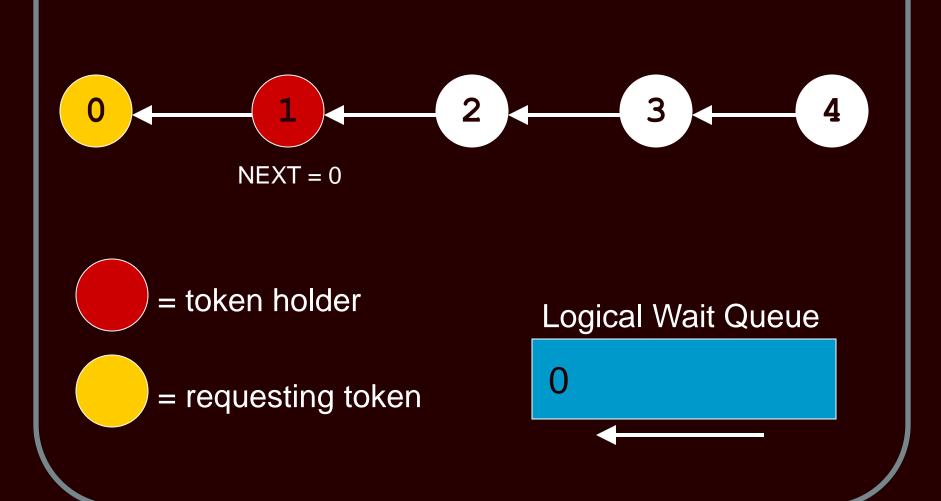


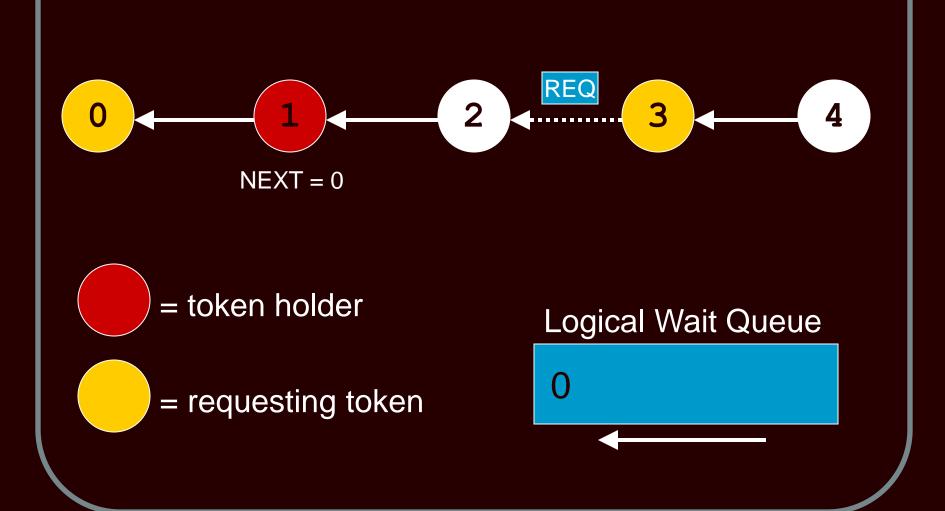


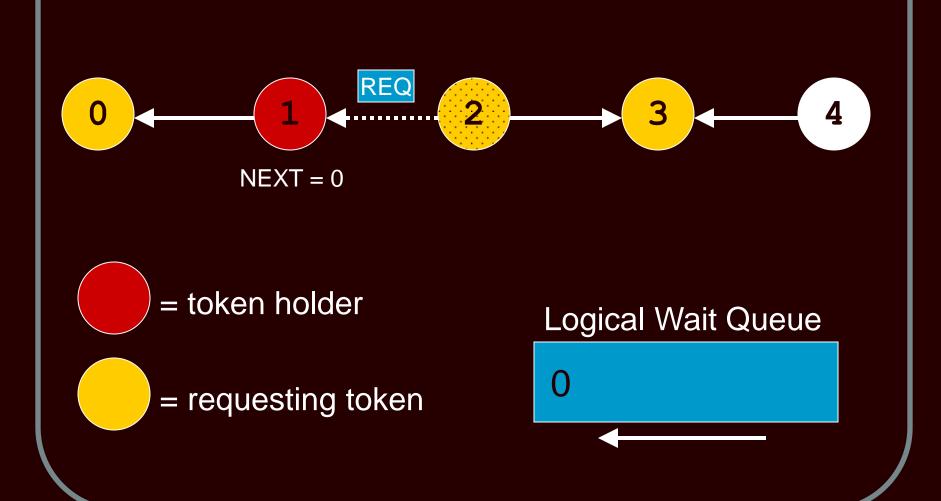
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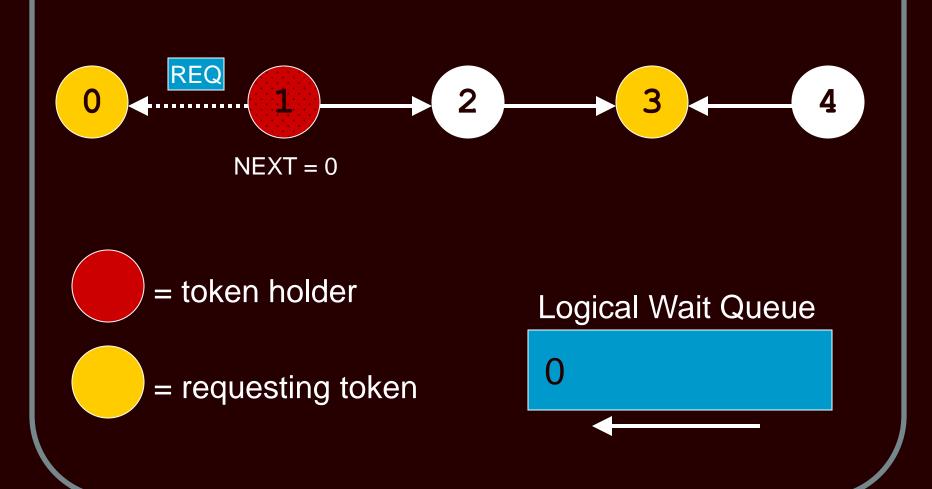


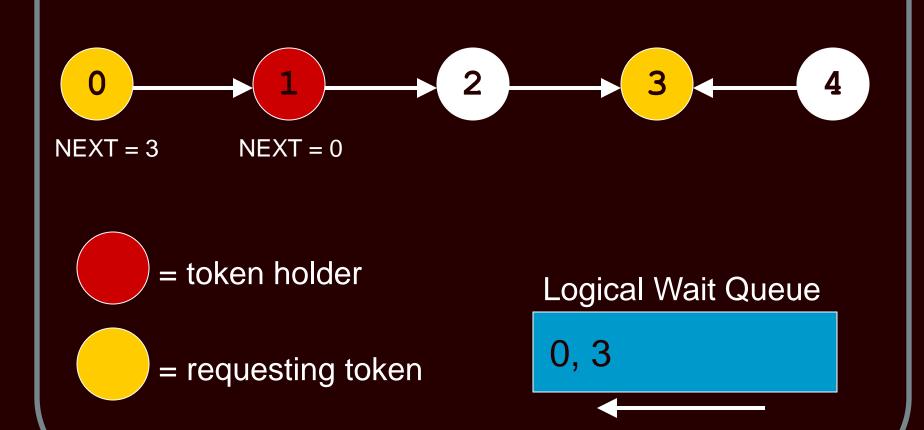


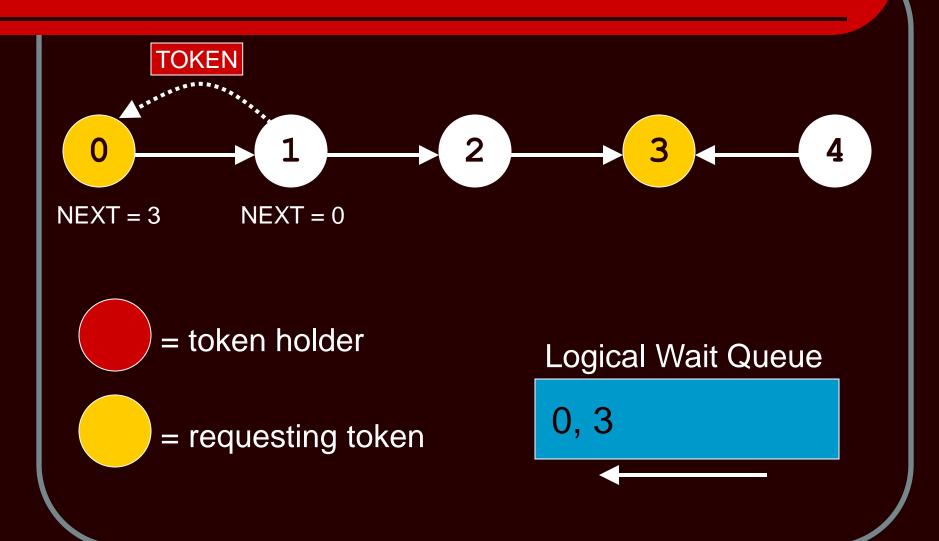




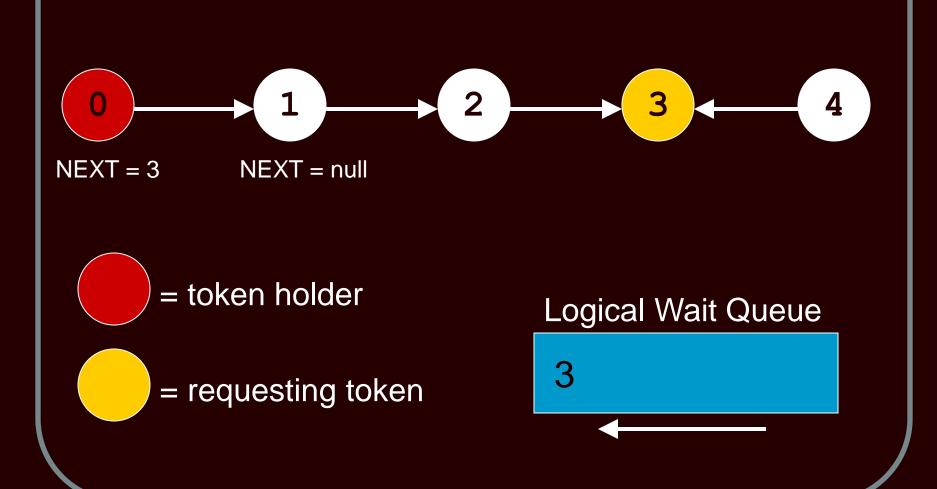




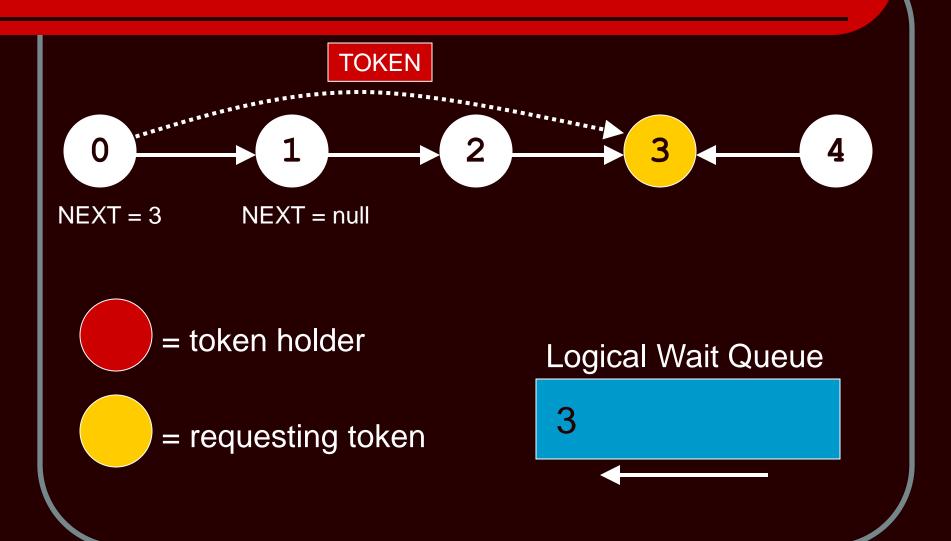




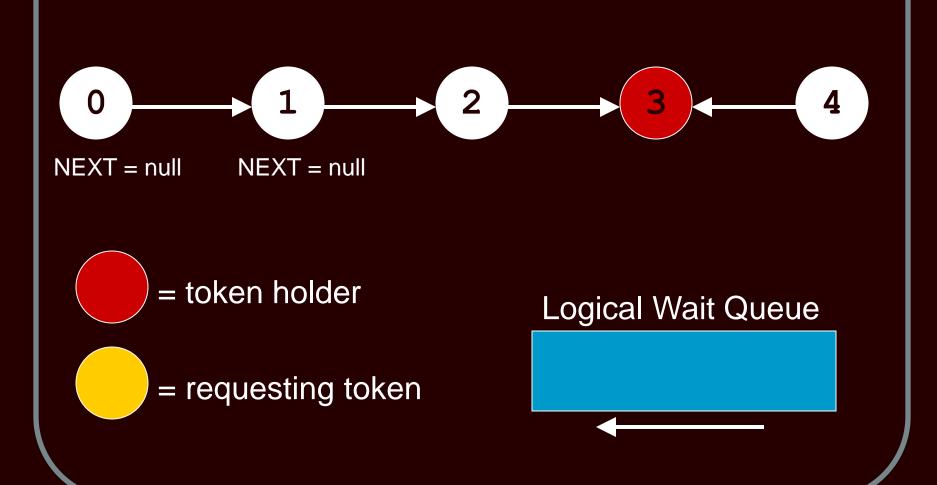
Static Logical Structure – Edge Reversal



Static Logical Structure – Edge Reversal



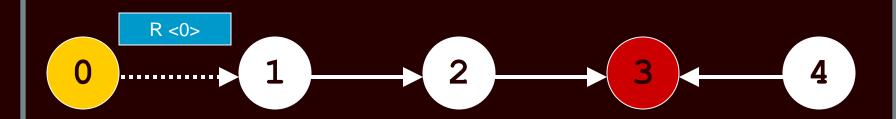
Static Logical Structure – Edge Reversal





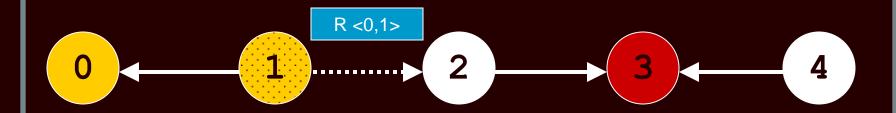
= token holder

= requesting token



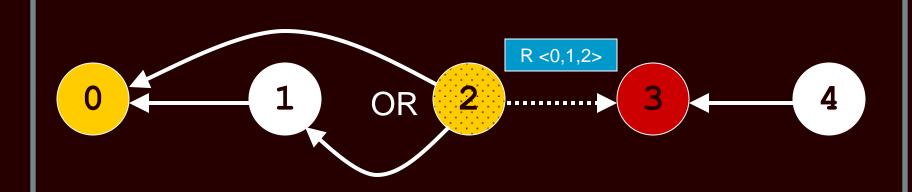






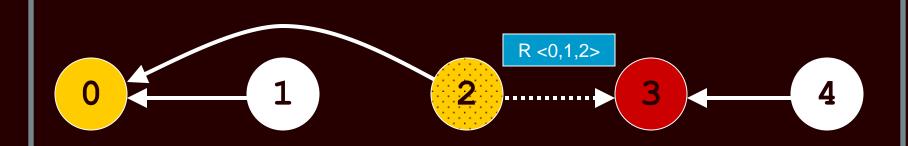


= requesting token



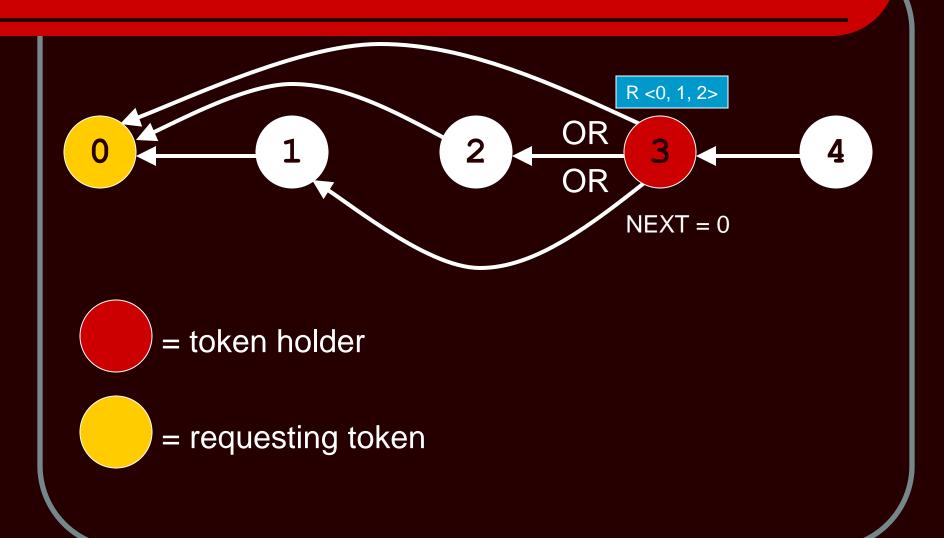


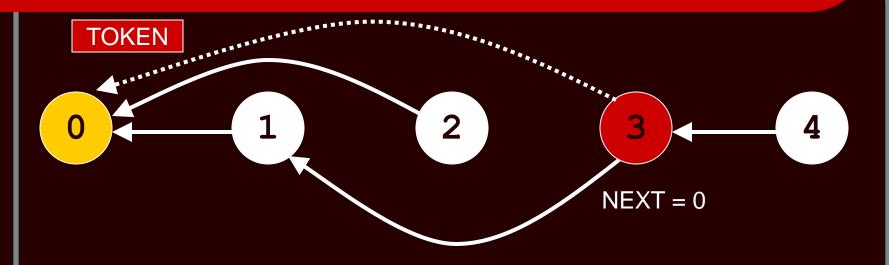






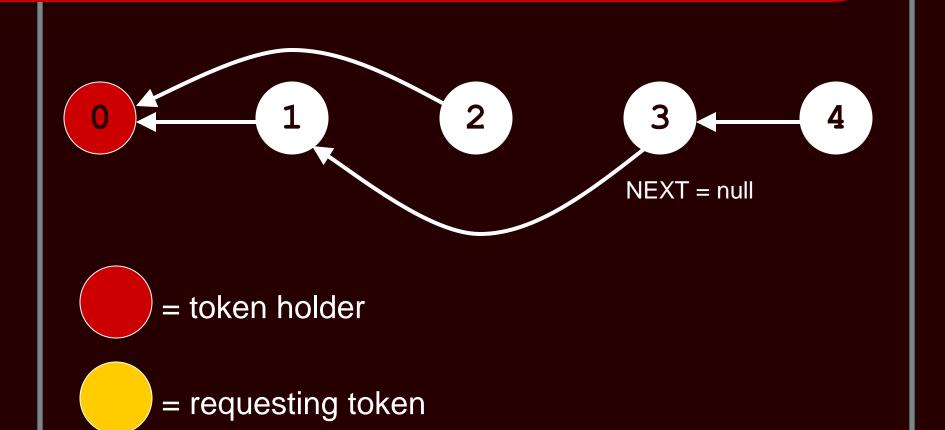
= requesting token





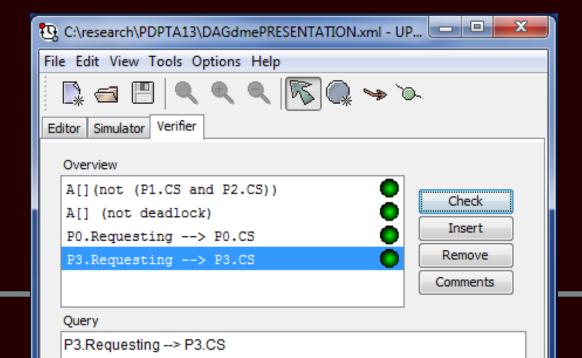






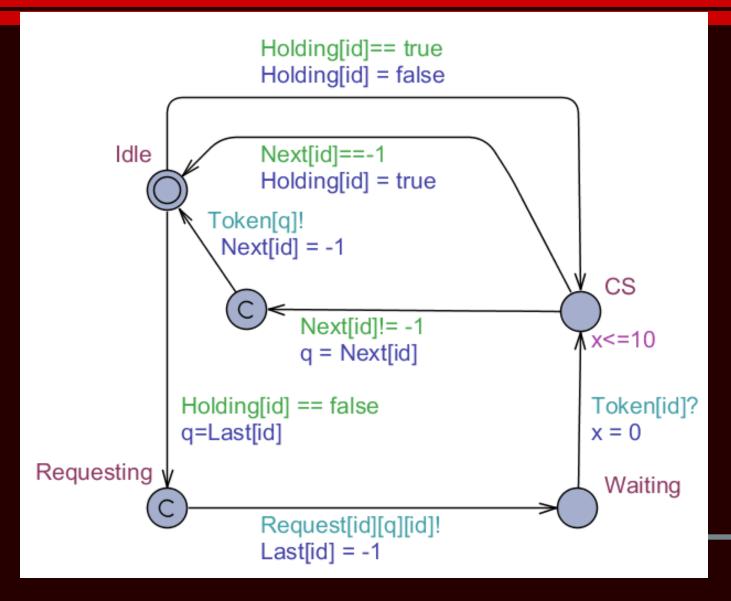
Properties to Verify

- Mutual exclusion is guaranteed.
- The algorithm is deadlock free.
- The algorithm is starvation free.



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UPPAAL Model – ProcessWork(id)

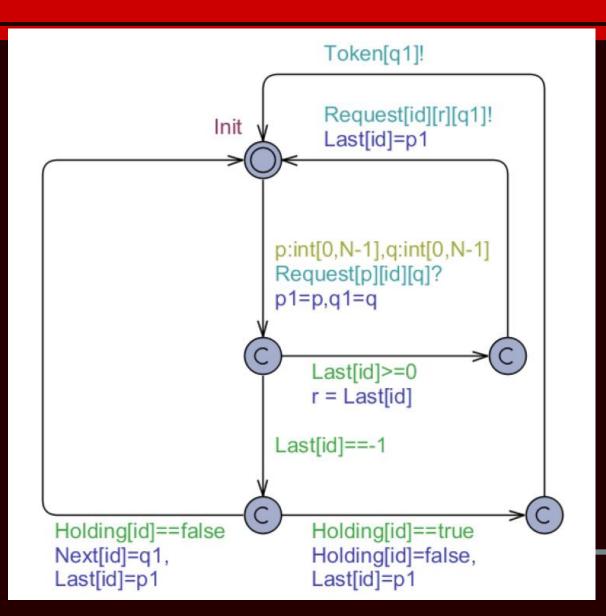


Holding – holding, but not using the token

Next – next node to receive token

Last – next node to send request to (node on path to last request)

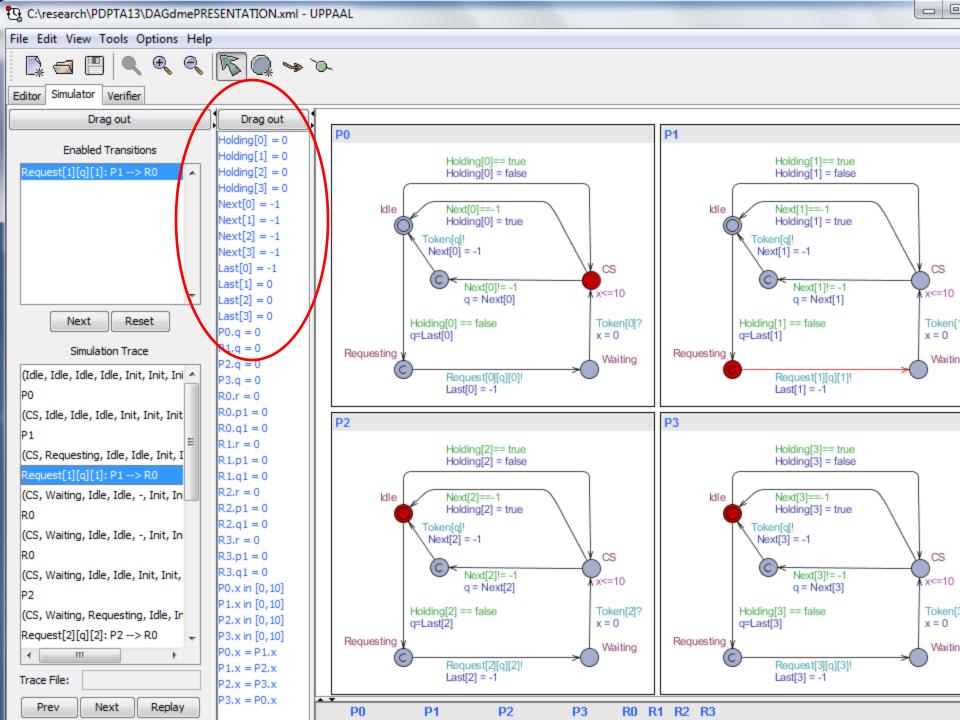
UPPAAL Model – ProcessRequest(id)

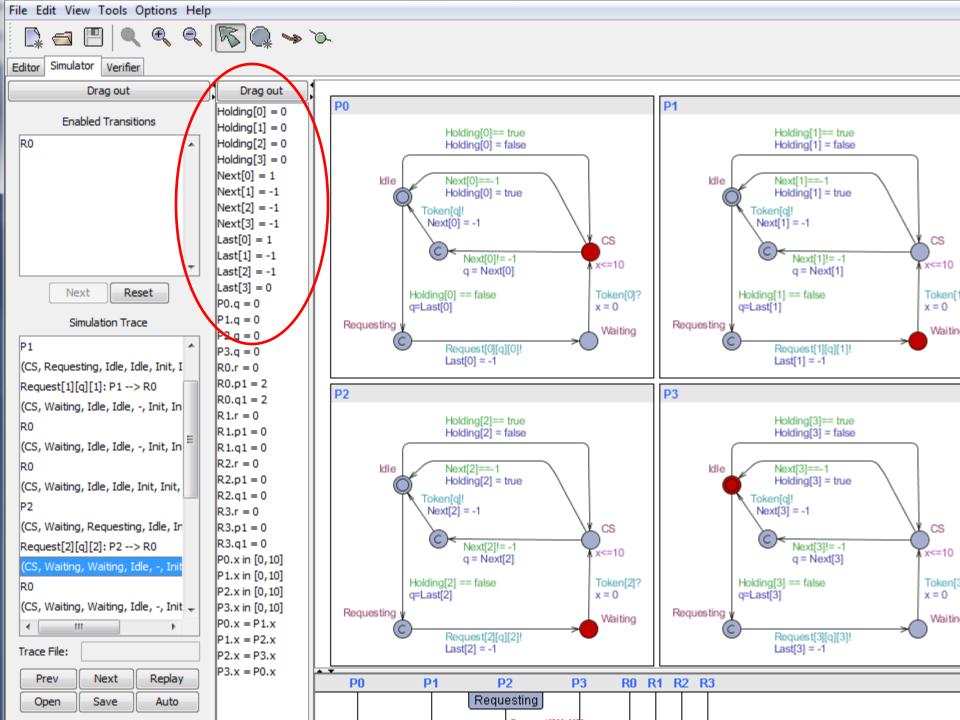


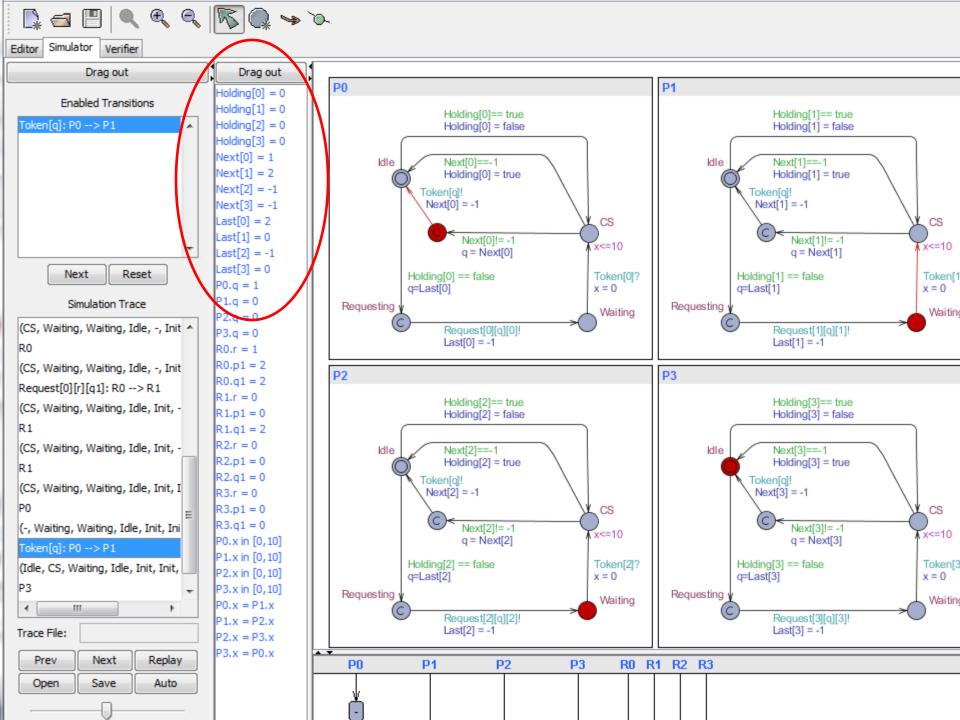
Holding – holding, but not using the token

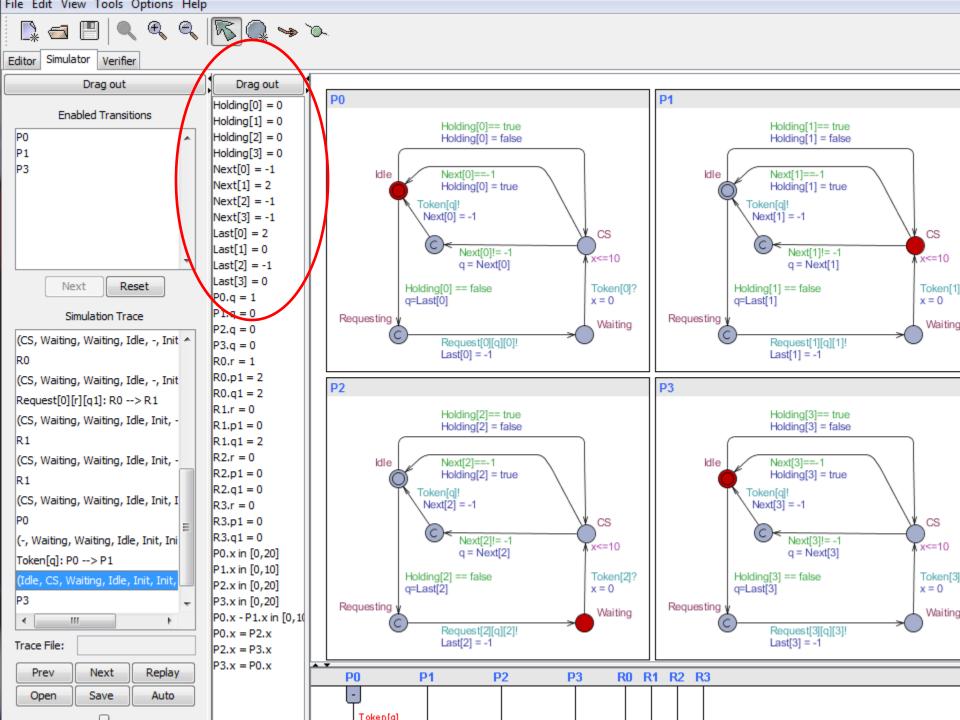
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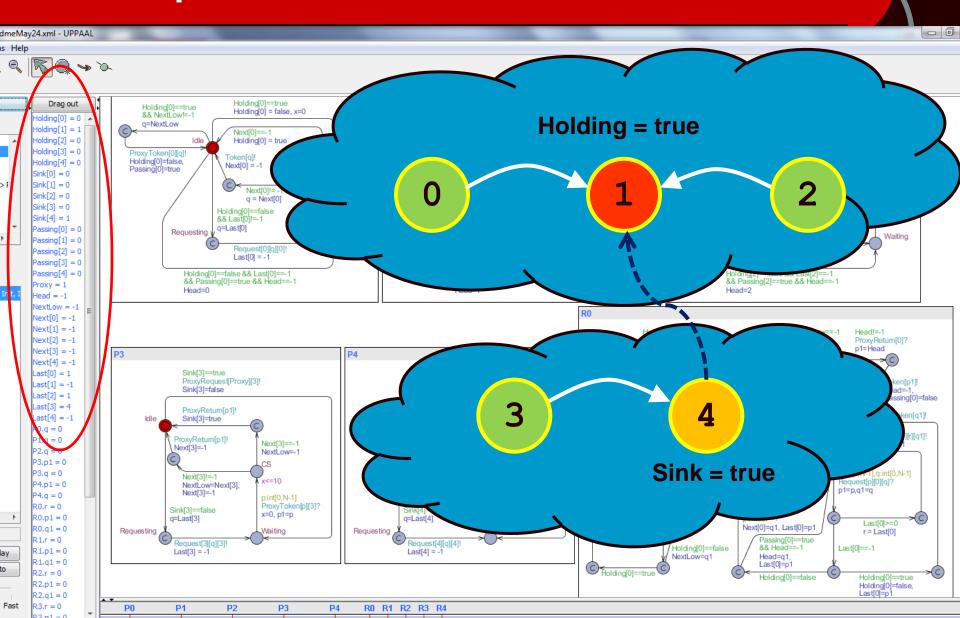




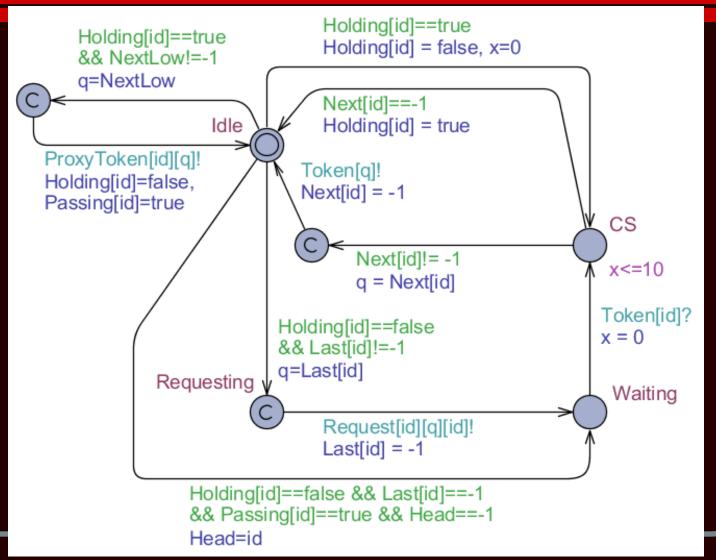
Prioritized Algorithms

- Modeled after real-time scheduling algorithm.
- Each priority level uses distributed algorithm for round-robin processing of requests at the same priority level.
- Proxy requests are passed up to reach a sink node at the highest priority level and to be enqueued in priority order in the token.
- Algorithm correctness is verified using same real-time verification tool, UPPAAL.
- Algorithm performance can also be determined using UPPAAL.

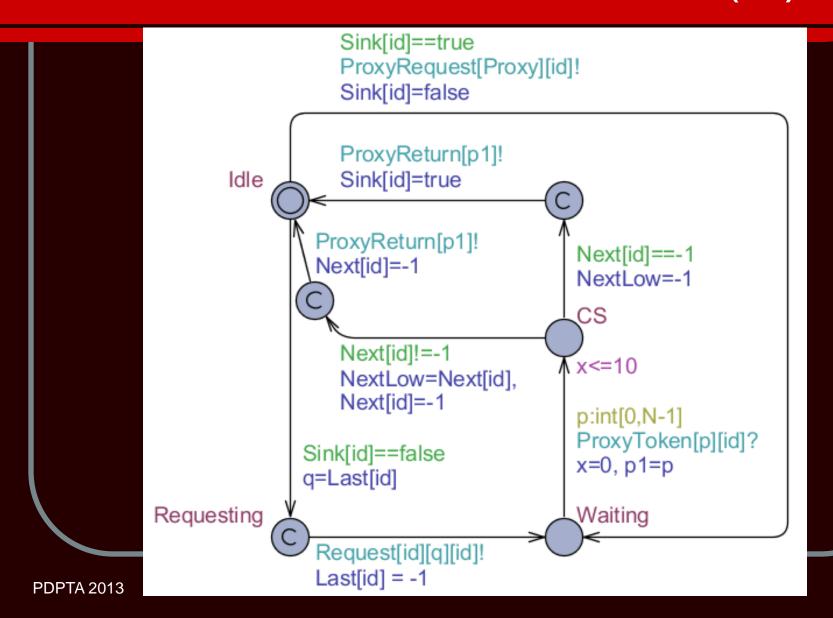
Example



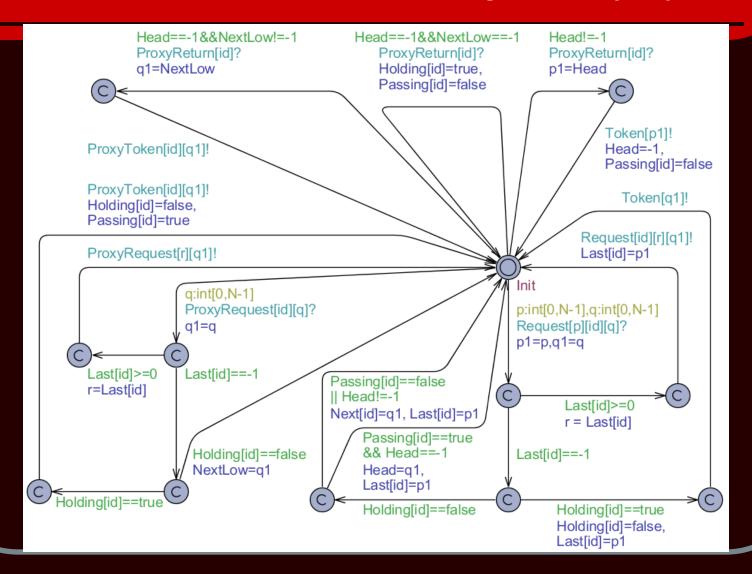
Prioritized ProcessWork(id)

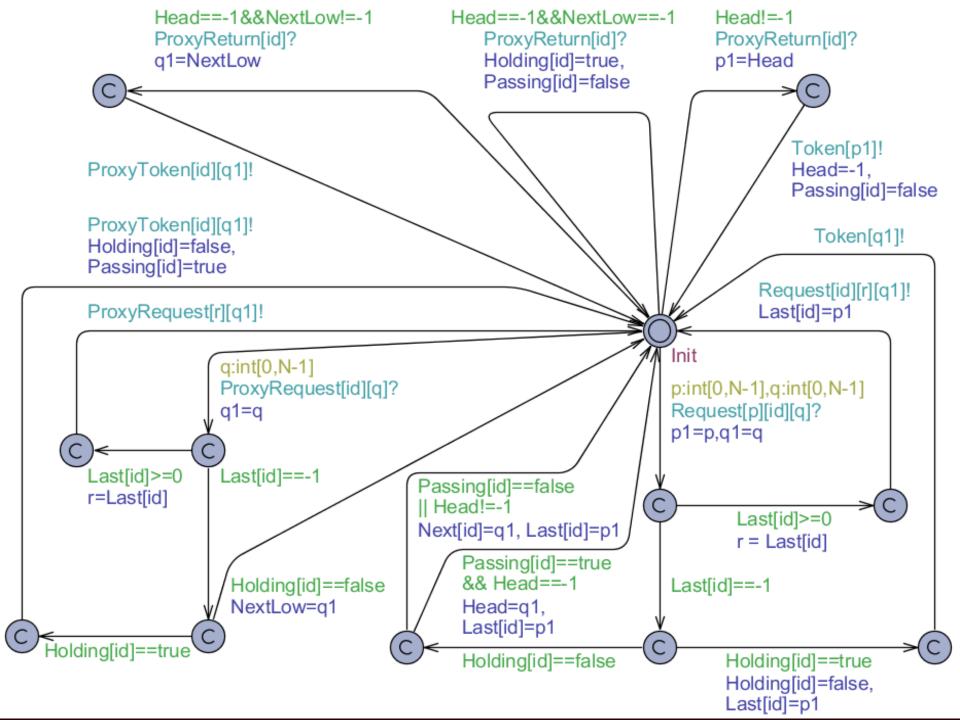


Prioritized ProcessWorkLow(id)

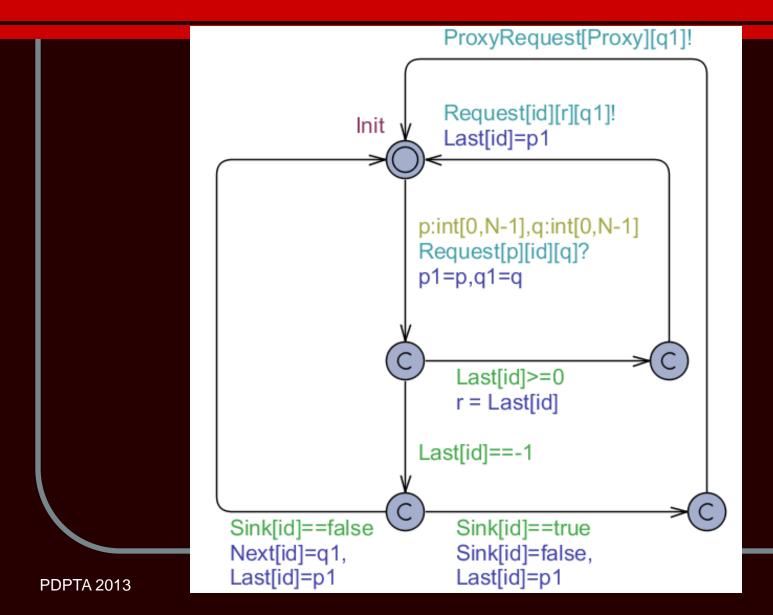


Prioritized ProcessRequest(id)

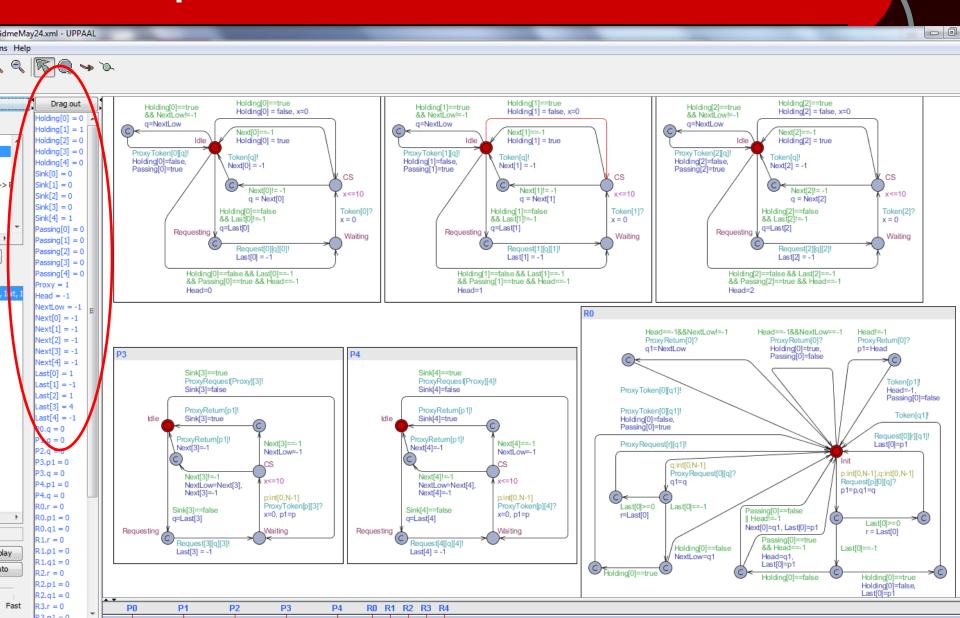




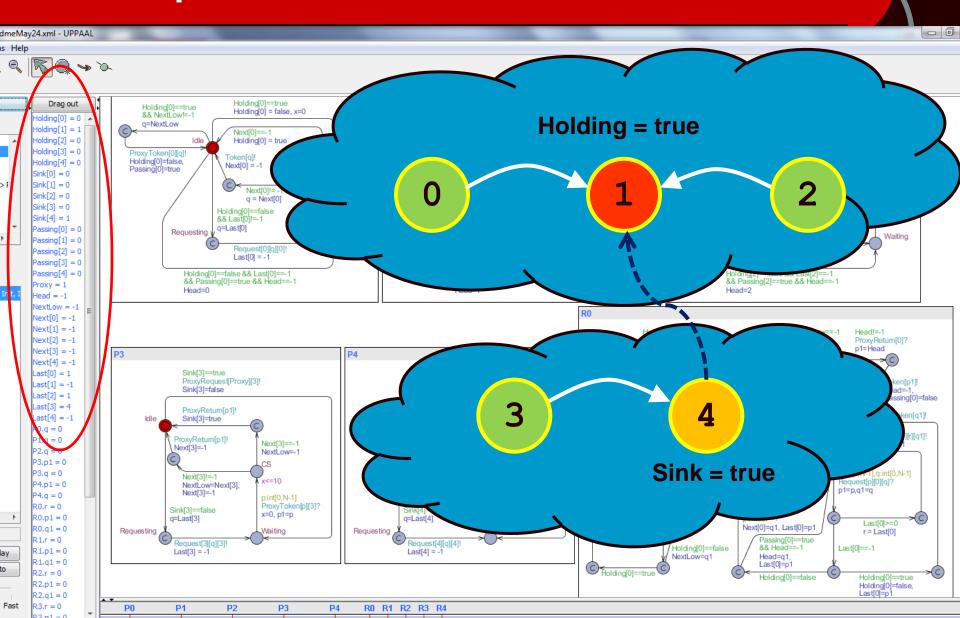
Prioritized ProcessRequestLow(id)



Example

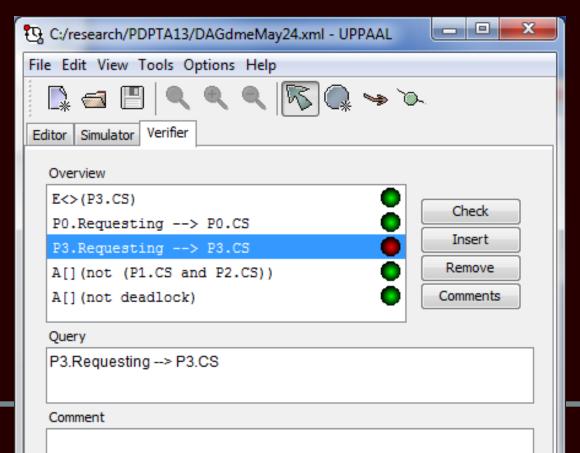


Example



Properties Checked

 Note: Starvation Freedom is not guaranteed for low-priority nodes.



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Summary

- Generalized Algorithm generalizes all algorithms that impose a logical structure on the nodes.
- Using a star topology, the dag-based algorithm achieves performance that exceeds any known algorithm, even the centralized algorithm.
- Generalized Algorithm can be prioritized, and correctness is verified.
- Prioritized Algorithm is realized by passing token between priority levels and enqueuing low-priority requests in the token.

Summary

- Quiz #2 this Friday, 4/25
- Final take-home
- Project Presentations start next week