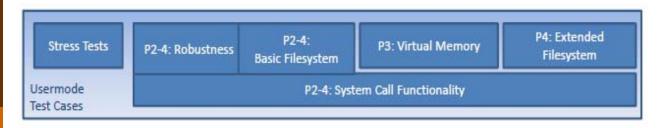
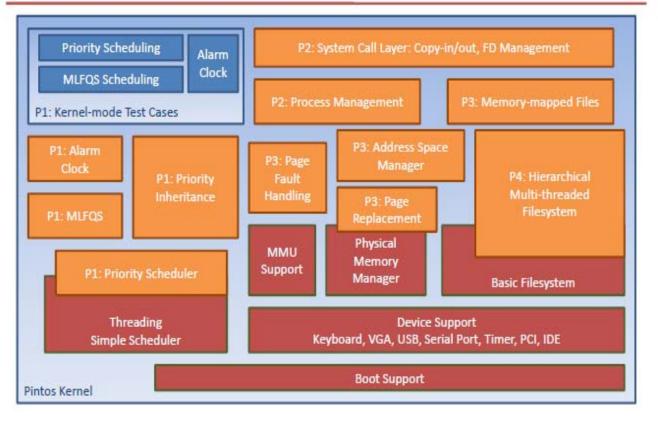
Project 2: User Programs and System Calls

P2: Project 2 – System Calls



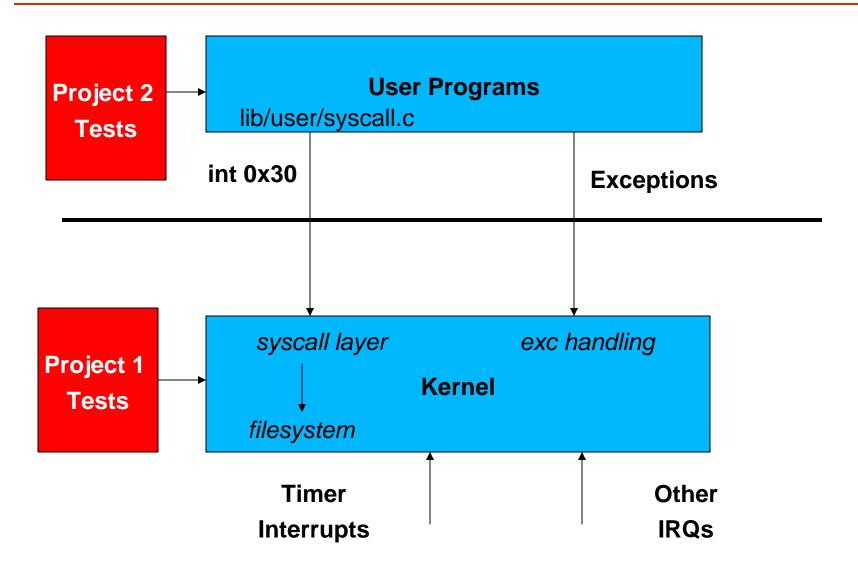


Support Code

Students Create

Public Tests

Project 1 and Project 2



Using the File system

- □ Interfacing with the file system
- No need to modify the file system
- □ Certain limitations
 - No internal synchronization
 - File size fixed
 - File data allocated as a single extent
 - No subdirectories
 - File names limited to 14 chars
 - System crash might corrupt the file system
- □ Files to take a look at: 'filesys.h' & 'file.h'

Some commands

- □ In **userprog/build**, create a new simulated disk
 - pintos-mkdisk fs.dsk --filesys-size=2
- Format the disk
 - pintos -f -q
 - This will only work after your disk is created and kernel is built!
- □ Copy the program "echo" onto the disk
 - pintos -p ../../examples/echo -a echo -- -q
- □ Run the program
 - pintos -q run 'echo x'
- All in a single command:

```
pintos --filesys-size=2 -p ../../examples/echo -a echo -- -f
    -q run 'echo x'
```

```
pintos ... -v -- .. - to from terminal without X11 server
```

- \$ make check or make grade builds the disk automatically
 - You can just copy & paste the commands that make check does!

Various directories

- □ Few user programs:
 - src/examples
- □ Relevant files:
 - userprog/
- □ Other files:
 - threads/

Project 2 Requirements

- □ Process Termination Messages
- Argument Passing
- □ System Calls
- Deny writes to executables

1. Process Termination

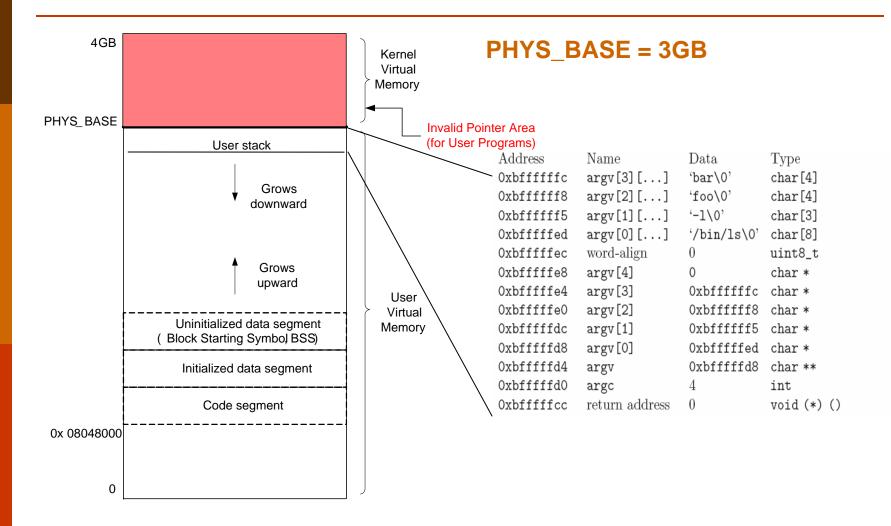
- When a Process Terminates
 - printf ("%s: exit(%d)\n", ...);
- Name: Full name passed to process_execute() in process.c
- Exit Code
- □ Do not print any other message!

2. Argument Passing

- □ No support currently for argument passing
- □ Change *esp = PHYS_BASE to *esp = PHYS_BASE 12 in setup_stack() to get started
- □ Change process_execute() to process multiple arguments
- □ Can limit the arguments to fit in a page(4 kb)
- □ String Parsing: strtok_r() in lib/string.h

```
pgm.c
main(int argc,
        char *argv[]) {
$ pintos run 'pgm alpha beta'
argc = 3
argv[0] = "pgm"
argv[1] = "alpha"
argv[2] = "beta"
```

Memory Layout



Setting up the Stack

How to setup the stack for the program - /bin/ls -I foo bar

```
Address
             Name
                            Data
                                         Type
Oxbffffffc
             argv[3][...]
                            'bar\0'
                                         char[4]
                            'foo\0'
             argv[2][...]
0xbffffff8
                                         char[4]
                            '-1\0'
0xbffffff5
             argv[1][...]
                                         char[3]
             argv[0][...]
                            '/bin/ls\0'
                                         char[8]
Oxbfffffed
0xbfffffec
             word-align
                            0
                                         uint8_t
0xbfffffe8
             argv[4]
                            0
                                         char *
             argv[3]
Oxbfffffe4
                            Oxbffffffc
                                         char *
0xbfffffe0
             argv[2]
                            0xbffffff8
                                         char *
             argv[1]
                            0xbffffff5
0xbfffffdc
                                         char *
0xbfffffd8
             argv[0]
                            0xbfffffed
                                         char *
0xbfffffd4
             argv
                            0xbfffffd8
                                         char **
0xbfffffd0
                            4
             argc
                                         int
             return address
0xbfffffcc
                            0
                                         void (*) ()
```

3. System Calls – already discussed some; e.g. open() system call

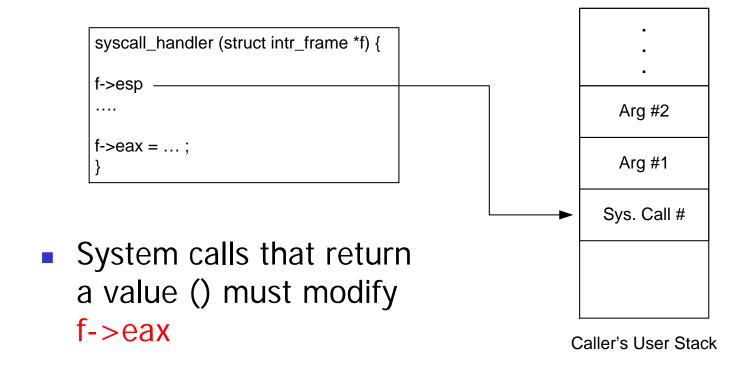
- No Support for system calls currently
- □ Implement the system call handler in userprog/syscall.c
- □ System call numbers defined in lib/syscall-nr.h
- □ Process Control: exit, exec, wait
- □ File system: create, remove, open, filesize, read, write, seek, tell, close
- Others: halt

System Call Details

- □ Types of Interrupts External and Internal
- System calls Internal Interrupts or Software Exceptions
- 80x86 'int' instruction to invoke system calls
- □ Pintos 'int \$0x30' to invoke system call

Continued...

- A system call has:
 - System call number
 - (possibly) arguments
- When syscall_handler() gets control:



System calls – File system

- □ Decide on how to implement the file descriptors keep it simple -- O(n) data structures will entail no deduction
- □ Access granularity is the entire file system add one global lock
- □ write() fd 1 writes to console use putbuf() to write entire buffer to console
- □ read() fd 0 reads from console use input_getc() to get input from keyboard
- Implement the rest of the system calls

System calls – Process Control

- wait(pid) Waits for process pid to die and returns the status pid returned from exit
- □ Returns -1 if
 - pid was terminated by the kernel
 - pid does not refer to child of the calling thread
 - wait() has already been called for the given pid
- exec(cmd) runs the executable whose name is given in command line
 - returns -1 if the program cannot be loaded
- □ exit(status) terminates the current user program, returns status
 - status of 0 indicates success, non zero otherwise

Process Control – continued...

- Parent may or may not wait for its child
- Parent may call wait() after child terminates!
- Implement process_wait() in process.c
- Then, implement wait() in terms of process_wait()
- □ Cond variables and/or semaphores will help
 - Think about what semaphores may be used for and how they must be initialized

```
main() {
int i; pid_t p;
p = exec("pgm a b");
i = wait (p);
... /* i must be 5 */
}
```

```
main() {
int status;
... status = 5;
exit(status);
} pgm.c
```

Memory Access

- Invalid pointers must be rejected. Why?
 - Kernel has access to all of physical memory including that of other processes
 - Kernel like user process would fault when it tries to access unmapped addresses
- □ User process cannot access kernel virtual memory
- User Process after it has entered the kernel can access kernel virtual memory and user virtual memory
- How to handle invalid memory access?

Memory Access – contd...

- Two methods to handle invalid memory access
 - Verify the validity of user provided pointer and then dereference it
 - □ Look at functions in userprog/pagedir.c, threads/vaddr.h
 - Strongly recommended!
 - Check if user pointer is below PHYS_BASE and dereference it
 - Could cause page fault
 - □ Handle the page fault by modifying the page_fault() code in userprog/exception.c
 - Make sure that resources are not leaked

Some Issues to look at...

- □ Check the validity of the system call parameters
- Every single location should be checked for validity before accessing it; e.g. not only for f->esp, but also the locations f->esp+1, f->esp+2 and f->esp+3 should be checked
- Read system call parameters into kernel memory (except for long buffers) write a copy_in function for this purpose.

Denying writes to Executables

- Use file_deny_write() to prevent writes to an open file
- Use file_allow_write() to re-enable write
- □ Closing a file will re-enable writes

Suggested Order of Implementation

- □ Change *esp = PHYS_BASE to *esp = PHYS_BASE 12 to get started
- □ Implement the system call infrastructure
- □ Change process_wait() to a infinite loop to prevent pintos getting powered off before the process gets executed
- Implement exit system call
- Implement write system call
- Start making other changes

Pintos Project 2 Sample Test

- Example Open System Call
- □ Test: tests/userprog/open-normal.c

```
/* Open a file. */
#include <syscall.h>
#include "tests/lib.h"
#include "tests/main.h"
void
test main (void)
  int handle = open ("sample.txt");
  if (handle < 2)
    fail ("open() returned %d", handle);
```

userprog/ - make check

```
gcc -c ../../tests/userprog/open-normal.c -o tests/userprog/open-normal.o
 -g -msoft-float -O -fno-stack-protector -nostdinc -I../.. -I../../lib
 -I../../lib/user -I. -Wall -W -Wstrict-prototypes -Wmissing-prototypes
 -Wsystem-headers -MMD -MF tests/userprog/open-normal.d
    -Wl,--build-id=none -nostdlib -static -Wl,-T,../../lib/user/user.lds
qcc
 tests/userprog/open-normal.o tests/main.o tests/lib.o lib/user/entry.o
 libc.a -o tests/userprog/open-normal
pintos -v -k -T 60 --qemu --filesys-size=2 -p tests/userprog/open-normal
 -a open-normal -p ../../tests/userprog/sample.txt -a sample.txt -- -q
 -f run open-normal < /dev/null 2> tests/userprog/open-normal.errors
 > tests/userprog/open-normal.output
perl -I../.. ../../tests/userprog/open-normal.ck
tests/userprog/open-normal tests/userprog/open-normal.result
pass tests/userprog/open-normal
ar r libc.a lib/debug.o lib/random.o lib/stdio.o lib/stdlib.o lib/string.o
 lib/arithmetic.o lib/ustar.o lib/user/debug.o lib/user/syscall.o
 lib/user/console.o
ranlib libc.a
```

\$ objdump -D open-normal

```
080480a0 <test main>:
 80480a0:
                55
                                         push
                                                %ebp
 80480a1:
                89 e5
                                                %esp,%ebp
                                         mov
 80480a3:
                83 ec 18
                                         sub
                                                $0x18,%esp
 80480a6:
                c7 04 24 6a a7 04 08
                                         movl
                                                $0x804a76a,(%esp)
 80480ad:
                e8 1f 21 00 00
                                         call
                                                804a1d1 <open>
 80480b2:
                83 f8 01
                                                $0x1,%eax
                                         cmp
0804a1d1 <open>:
 804a1d1:
                55
                                         push
                                                %ebp
 804a1d2:
                89 e5
                                                %esp,%ebp
                                         mov
 804a1d4:
                ff 75 08
                                         pushl
                                                0x8(%ebp)
 804a1d7:
                6a 06
                                         push
                                                $0x6
 804a1d9:
                cd 30
                                                $0x30
                                         int
 804a1db:
                83 c4 08
                                         add
                                                $0x8,%esp
 804a1de:
                5d
                                                %ebp
                                         pop
 804a1df:
                c3
                                         ret
0804a76a <.rodata.str1.1>:
 804a76a:
                73 61 6d 70 6c 65 2e 74 78 74 00
                                 е
                 S
```

```
void
syscall_init (void)
{
  intr_register_int (0x30, 3, INTR_ON, syscall_handler,
    "syscall");
  lock_init (&fs_lock);
}
```

```
/* System call handler. */
static void
syscall handler (struct intr frame *f)
 typedef int syscall_function (int, int, int);
 /* A system call. */
 struct syscall
     size t arg cnt; /* Number of arguments. */
     syscall function *func; /* Implementation. */
   };
  /* Table of system calls. */
 static const struct syscall syscall_table[] =
      {0, (syscall_function *) sys_halt},
      {1, (syscall_function *) sys_exit},
      {1, (syscall_function *) sys_exec},
      {1, (syscall_function *) sys_wait},
      {2, (syscall_function *) sys_create},
      {1, (syscall_function *) sys_remove},
      {1, (syscall function *) sys open}, <-- call number 6
```

```
const struct syscall *sc;
unsigned call nr;
int args[3];
/* Get the system call. */
copy in (&call nr, f->esp, sizeof call nr);
if (call nr >= sizeof syscall table / sizeof *syscall table)
 thread exit ();
sc = syscall table + call nr;
/* Get the system call arguments. */
ASSERT (sc->arg cnt <= sizeof args / sizeof *args);
memset (args, 0, sizeof args);
copy in (args, (uint32 t *) f->esp + 1, sizeof *args * sc->arg cnt);
/* Execute the system call,
   and set the return value. */
f\rightarrow eax = sc\rightarrow func (args[0], args[1], args[2]);
```

```
/* Open system call. */
static int
sys open (const char *ufile)
 char *kfile = copy in string (ufile);
  struct file descriptor *fd;
  int handle = -1;
 fd = malloc (sizeof *fd);
  if (fd != NULL)
      lock acquire (&fs lock);
      fd->file = filesys open (kfile);
      if (fd->file != NULL)
```

Copy a byte from user space to kernel space

```
/* Copies a byte from user address USRC to kernel address DST. */
static inline bool
get user (uint8 t *dst, const uint8 t *usrc)
  int eax;
  asm ("movl $1f, %%eax; movb %2, %%al; movb %%al, %0; 1:"
       : "=m" (*dst), "=&a" (eax) : "m" (*usrc));
  return eax != 0;
/* Writes BYTE to user address UDST. */
static inline bool
put user (uint8 t *udst, uint8 t byte)
  int eax;
  asm ("movl $1f, %%eax; movb %b2, %0; 1:"
       : "=m" (*udst), "=&a" (eax) : "q" (byte));
  return eax != 0;
```

Copy a string from user space to kernel space

```
static char *
copy_in_string (const char *us)
  char *ks;
  size_t length;
 ks = palloc_get_page (0);
  if (ks == NULL)
    thread_exit ();
    for (length = 0; length < PGSIZE; length++)</pre>
      if (us >= (char *) PHYS_BASE | | !get_user (ks + length, us++))
          palloc_free_page (ks);
          thread exit ();
      if (ks[length] == '\0')
        return ks;
 ks[PGSIZE - 1] = ' 0';
  return ks;
```

Run a single test

- □ From userprog/build:
 - rm tests/userprog/open-normal.output
 - make test/userprog/open-normal.output

```
pintos -v -k -T 60 --qemu --filesys-size=2 -p tests/userprog/open-normal -a open-normal -p ../../tests/userprog/sample.txt -a sample.txt -- -q -f run open-normal < /dev/null 2> tests/userprog/open-normal.errors > tests/userprog/open-normal.output
```

examples/shell.c

```
int main (void)
 printf ("Shell starting..\n");
  for (;;)
      char command[80];
      /* Read command. */
      printf ("--");
      read line (command, sizeof command);
      /* Execute command. */
      if (!strcmp (command, "exit"))
        break;
      else if (!memcmp (command, "cd ", 3))
          if (!chdir (command + 3))
           printf ("\"%s\": chdir failed\n", command + 3);
      else if (command[0] == '\0')
```

examples/shell.c

```
else if (command[0] == '\0')
         /* Empty command. */
     else
        pid_t pid = exec (command);
         if (pid != PID_ERROR)
           printf ("\"%s\": exit code %d\n", command, wait (pid));
         else
           printf ("exec failed\n");
printf ("Shell exiting.");
return EXIT_SUCCESS;
```

\$ objdump -D examples/shell

080480cc <main>: 80480cc: 55 push %ebp 80480cd: 89 e5 %esp,%ebp mov 80480cf: 83 e4 f0 and \$0xfffffff0, %esp 80480d2: 57 %edi push 80480d3: %esi 56 push 80480d4: 53 %ebx push 80480d5: 83 ec 74 sub \$0x74,%esp 80480d8: c7 04 24 3e al 04 08 movl \$0x804a13e,(%esp) 80480df: e8 23 1d 00 00 call 8049e07 <puts> 80480e4: c7 04 24 5a a1 04 08 movl \$0x804a15a,(%esp) 80480eb: e8 2f 0c 00 00 call 8048d1f <printf> 80480f0: 8d 44 24 18 lea 0x18(%esp),%eax 80480f4: 89 44 24 68 %eax,0x68(%esp) mov 80480f8: 89 c3 %eax,%ebx mov 80480fa: c7 44 24 08 01 00 00 movl \$0x1,0x8(%esp)8048101: 00 8048102: 8d 44 24 6f lea 0x6f(%esp),%eax8048106: 89 44 24 04 %eax,0x4(%esp) mov 804810a: c7 04 24 00 00 00 00 movl \$0x0,(%esp) 8048111: e8 a9 1a 00 00 8049bbf <read> call

\$ objdump -D examples/shell

```
80481f1:
                 80 7c 24 18 00
                                                   $0x0,0x18(%esp)
                                           cmpb
 80481f6:
                 Of 84 e8 fe ff ff
                                                   80480e4 < main + 0x18 >
                                           jе
 80481fc:
                 8d 44 24 18
                                                   0x18(%esp),%eax
                                           lea
 8048200:
                 89 04 24
                                                   %eax,(%esp)
                                           mov
 8048203:
                 e8 50 19 00 00
                                           call
                                                   8049b58 <exec>
08049bbf <read>:
 8049bbf:
                 55
                                                   %ebp
                                           push
 8049bc0:
                 89 e5
                                                   %esp,%ebp
                                           mov
 8049bc2:
                 ff 75 10
                                                   0x10(%ebp)
                                           pushl
 8049bc5:
                 ff 75
                       0c
                                           pushl
                                                   0xc(%ebp)
 8049bc8:
                 ff 75 08
                                           pushl
                                                   0x8(%ebp)
 8049bcb:
                 6a 08
                                           push
                                                   $0x8
 8049bcd:
                 cd 30
                                                   $0x30
                                           int
 8049bcf:
                 83 c4 10
                                           add
                                                   $0x10,%esp
 8049bd2:
                 5d
                                                   %ebp
                                           pop
 8049bd3:
                 c3
                                           ret
```

syscall exec() calls process_execute

```
/* Starts a new thread running a user program loaded from
  FILENAME. The new thread may be scheduled (and may even exit)
  before process execute() returns. Returns the new process's
  thread id, or TID_ERROR if the thread cannot be created. */
tid_t process_execute (const char *file_name)
  struct exec info exec;
  char thread name[15];
  char *save ptr;
  tid_t tid;
  /* Initialize exec_info. */
  exec.file_name = file_name;
  sema_init (&exec.load_done, 0);
  /* Create a new thread to execute FILE_NAME. */
  strlcpy (thread_name, file_name, sizeof thread_name);
  strtok_r (thread_name, " ", &save_ptr);
  tid = thread_create (thread_name, PRI_DEFAULT, start_process, &exec);
```

process_execute

```
/* Starts a new thread running a user program loaded from
   FILENAME. The new thread may be scheduled (and may even exit)
   before process execute() returns. Returns the new process's
   thread id, or TID_ERROR if the thread cannot be created. */
tid t process_execute (const char *file_name)
  if (tid != TID_ERROR)
      sema down (&exec.load done);
      if (exec.success)
        list_push_back (&thread_current ()->children,
                        &exec.wait status->elem);
      else
        tid = TID_ERROR;
  return tid;
```

start_process

```
static void start process (void *exec )
  struct exec_info *exec = exec_;
  struct intr_frame if_;
 bool success;
  /* Initialize interrupt frame and load executable. */
 memset (&if_, 0, sizeof if_);
  if .qs = if .fs = if .es = if .ds = if .ss = SEL UDSEG;
  if .cs = SEL_UCSEG;
  if_.eflags = FLAG_IF | FLAG_MBS;
  success = load (exec->file_name, &if_.eip, &if_.esp);
  /* Allocate wait_status. */
  if (success)
      exec->wait_status = thread_current ()->wait_status
        = malloc (sizeof *exec->wait status);
      success = exec->wait status != NULL;
```

start_process

```
/* Initialize wait status. */
if (success)
     . .
 /* Notify parent thread and clean up. */
 exec->success = success;
sema up (&exec->load done);
if (!success)
  thread_exit ();
 /* Start the user process by simulating a return from an
    interrupt, implemented by intr_exit (in
   threads/intr-stubs.S). Because intr_exit takes all of its
   arguments on the stack in the form of a `struct intr_frame',
   we just point the stack pointer (%esp) to our stack frame
   and jump to it. */
asm volatile ("movl %0, %%esp; jmp intr_exit" : : "g" (&if_) : "memory");
```

Cycle Counters

- Most modern systems have built in registers that are incremented every clock cycle
 - Very fine grained
 - □ Maintained as part of process state
 - In Linux, counts elapsed global time
- Special assembly code instruction to access
- On (recent model) Intel machines:
 - □ 64 bit counter.
 - RDTSC instruction sets %edx to high order 32-bits,
 %eax to low order 32-bits
- Aside: Is this a security issue?

Cycle Counter Period

□Wrap Around Times for 550 MHz machine

- Low order 32 bits wrap around every $2^{32} / (550 * 10^6) = 7.8$ seconds
- High order 64 bits wrap around every 2⁶⁴ / (550 * 10⁶) = 33539534679 seconds
 - □ 1065 years

□For 2 GHz machine

- Low order 32 bits every 2.1 seconds
- High order 64 bits every 293 years

Measuring with Cycle Counter

Idea

- □ Get current value of cycle counter
 - store as pair of unsigned's cyc_hi and cyc_lo
- Compute something
- □ Get new value of cycle counter
- Perform double precision subtraction to get elapsed cycles

```
/* Keep track of most recent reading of cycle counter */
static unsigned cyc_hi = 0;
static unsigned cyc_lo = 0;

void start_counter()
{
   /* Get current value of cycle counter */
   access_counter(&cyc_hi, &cyc_lo);
}
```

Accessing the Cycle Counter

- GCC allows inline assembly code with mechanism for matching registers with program variables
- Code only works on x86 machine compiling with GCC

```
void access_counter(unsigned *hi, unsigned *lo)
{
   /* Get cycle counter */
   asm("rdtsc; movl %%edx,%0; movl %%eax,%1"
        : "=r" (*hi), "=r" (*lo)
        : /* No input */
        : "%edx", "%eax");
}
```

■ Emit assembly with rdtsc and two movl instructions

```
asm("Instruction String"
      : Output List
      : Input List
      : Clobbers List);
                    void access counter
                      (unsigned *hi, unsigned *lo)
                      /* Get cycle counter */
                      asm("rdtsc; movl %%edx,%0; movl %%eax,%1"
                          : "=r" (*hi), "=r" (*lo)
                          : /* No input */
                          : "%edx", "%eax");
```

Instruction String

- Series of assembly commands
 - Separated by ";" or "\n"
 - Use "%%" where normally would use "%"

```
asm("Instruction String"
      : Output List
                      void access counter
        Input List
                         (unsigned *hi, unsigned *lo)
      : Clobbers List
                        /* Get cycle counter */
                        asm("rdtsc; movl %%edx,%0; movl %%eax,%1"
                             : "=r" (*hi), "=r" (*lo)
                             : /* No input */
                             : "%edx", "%eax");
```

- Output List
 - Expressions indicating destinations for values %0, %1, ..., %j
 - Enclosed in parentheses
 - □ Must be *lvalue*
 - Value that can appear on LHS of assignment
 - Tag "=r" indicates that symbolic value (%0, etc.), should be replaced by a register

```
asm("Instruction String"

: Output List
: Input List
: Clobbers List)
}

/* Get cycle counter */
asm("rdtsc; movl %%edx,%0; movl %%eax,%1"
: "=r" (*hi), "=r" (*lo)
: /* No input */
: "%edx", "%eax");
}
Input List
```

- - . . .
 - Enclosed in parentheses
 - Any expression returning value
- Tag "r" indicates that symbolic value (%0, etc.) will come from register

```
asm("Instruction String"
      : Output List
                         void access_counter
      : Input List
                            (unsigned *hi, unsigned *lo)
      : Clobbers List);
                         {
                           /* Get cycle counter */
                           asm("rdtsc; movl %%edx,%0; movl %%eax,%1
                                : "=r" (*hi), "=r" (*lo)
                                : /* No input */
                                : "%edx", "%eax");
```

- □ Clobbers List
 - List of register names that get altered by assembly instruction
 - Compiler will make sure doesn't store something in one of these registers that must be preserved across asm
 - □ Value set before & used after

Completing Measurement

- Get new value of cycle counter
- Perform double precision subtraction to get elapsed cycles
- Express as double to avoid overflow problems

```
double get counter()
  unsigned ncyc_hi, ncyc_lo
  unsigned hi, lo, borrow;
  /* Get cycle counter */
  access_counter(&ncyc_hi, &ncyc_lo);
  /* Do double precision subtraction */
  lo = ncyc_lo - cyc_lo;
 borrow = lo > ncyc_lo;
  hi = ncyc_hi - cyc_hi - borrow;
  return (double) hi * (1 << 30) * 4 + lo;
```

Timing With Cycle Counter

□ Determine Clock Rate of Processor

■ Count number of cycles required for some fixed number of seconds

```
double MHZ;
int sleep_time = 10;
start_counter();
sleep(sleep_time);
MHZ = get_counter()/(sleep_time * 1e6);
```

□ Time Function P()

■ First attempt: Simply count cycles for one execution of P

```
double tsecs;
start_counter();
P();
tsecs = get_counter() / (MHZ * 1e6);
```

Example – testClock.c

```
#include <stdio.h>
#include "clock.h"
               Processor Clock Rate ~= 2673.5 MHz
               cycles = 5343976388.000000, MHz = 2673.526339, cycles/Mhz = 1998849.351153
int main()
               elapsed time = 1.998849 seconds
  double cycles, Mhz;
  Mhz = mhz(1);
  start_counter();
  sleep(2);
  cycles = get counter();
  printf("cycles = %f, MHz = %f, cycles/Mhz = %f\n", cycles, Mhz,
 cycles/Mhz);
  printf("elapsed time = %f seconds \n", cycles/(1.0e6*Mhz));
  return 0;
```

Measurement Pitfalls

□Overhead

- Calling get_counter() incurs small amount of overhead
- Want to measure long enough code sequence to compensate