

Math 322

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(Notebook crashed!) fresh!

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F.S.M with no output

→ Automata

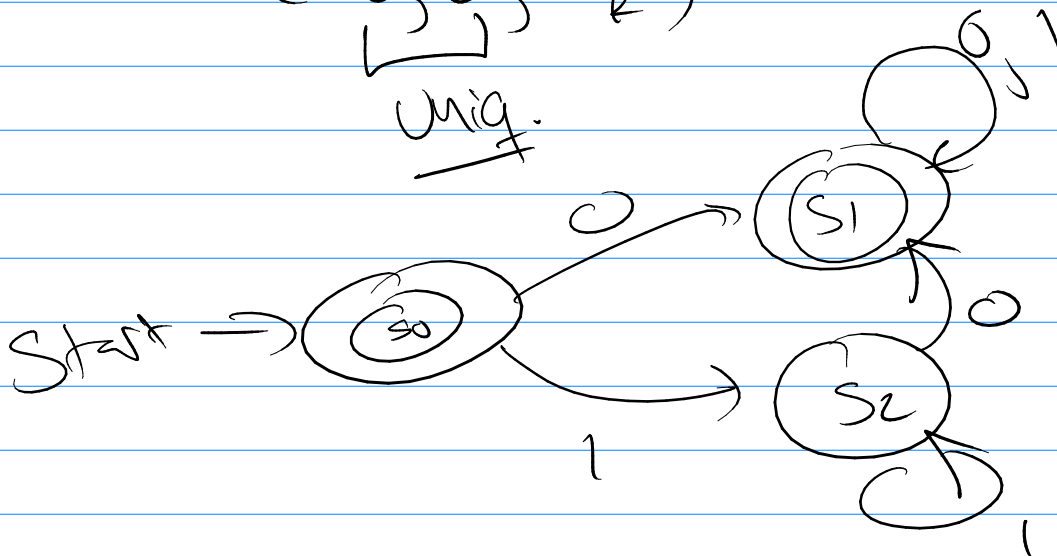
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Deterministic  $M = (S, I, f, s_0, F)$

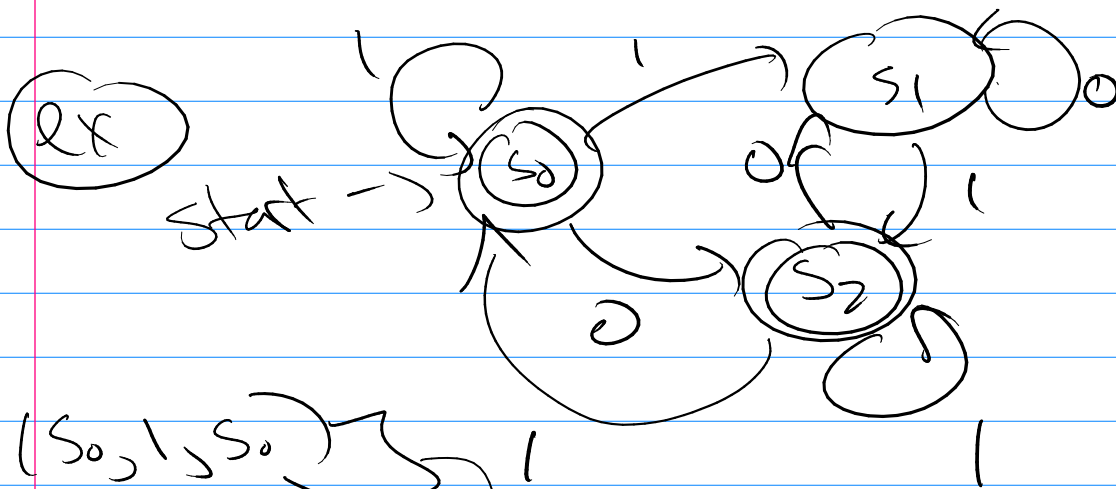
$$f: S \times I \rightarrow S$$

$$(s_i, i, s_k)$$

unig.



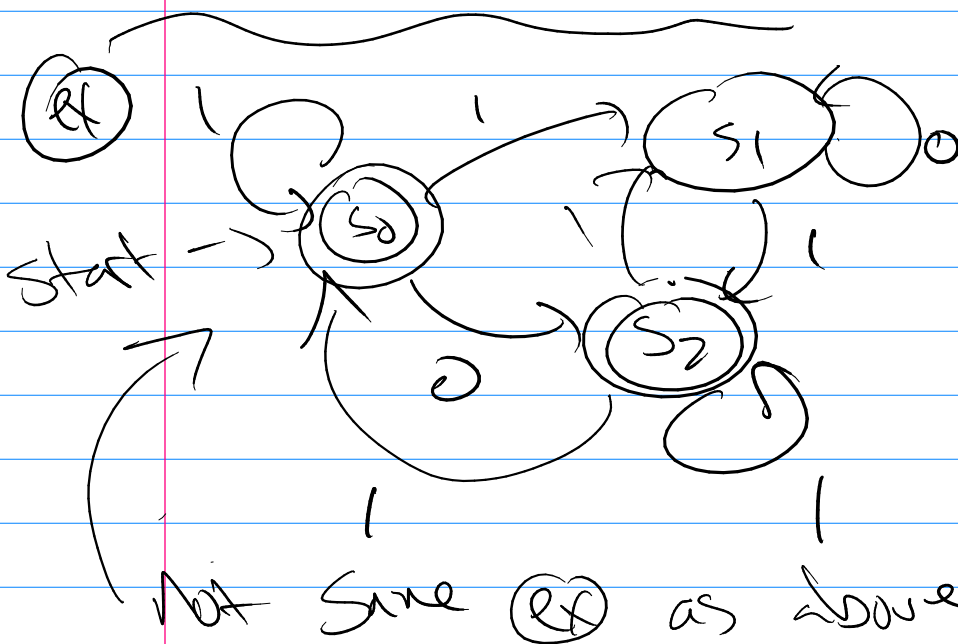
# Non deterministic F.S.A.



$(s_0, 1, s_0)$   
 $(s_0, 1, s_1)$   
 $(s_2, 1, s_0)$   
 $(s_2, 1, s_2)$

Non deterministic

$S \times I \rightarrow S$  Not a function

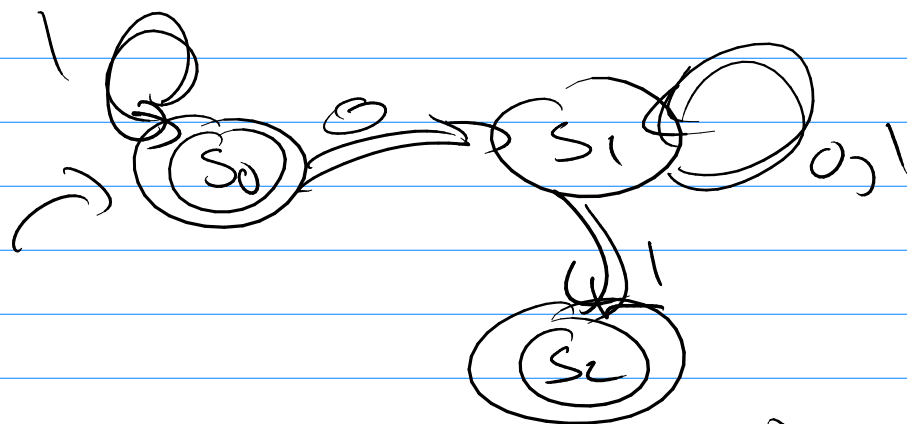


state table

	0	1
s <sub>0</sub>	s <sub>2</sub>	s <sub>0</sub> , s <sub>1</sub>
s <sub>1</sub>		s <sub>2</sub>
s <sub>2</sub>	s <sub>1</sub>	s <sub>1</sub> , s <sub>0</sub> , s <sub>2</sub>

Non-Det. F.S.A  $M = (S, I, f, s_0, F)$   
 $f: S \times I \rightarrow P(S)$

If  $M_1$  is nondeterministic when should  $x \in I^*$  be recognized?

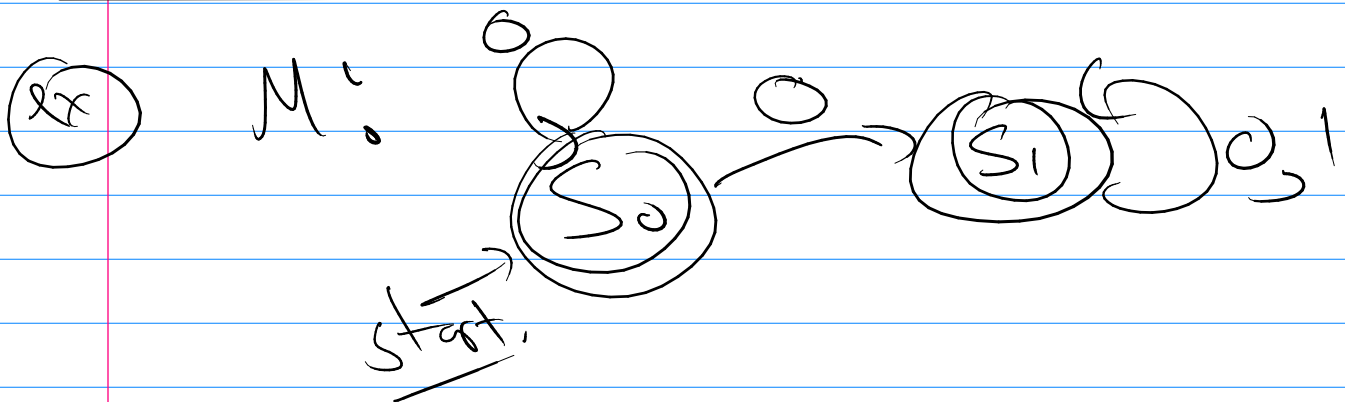
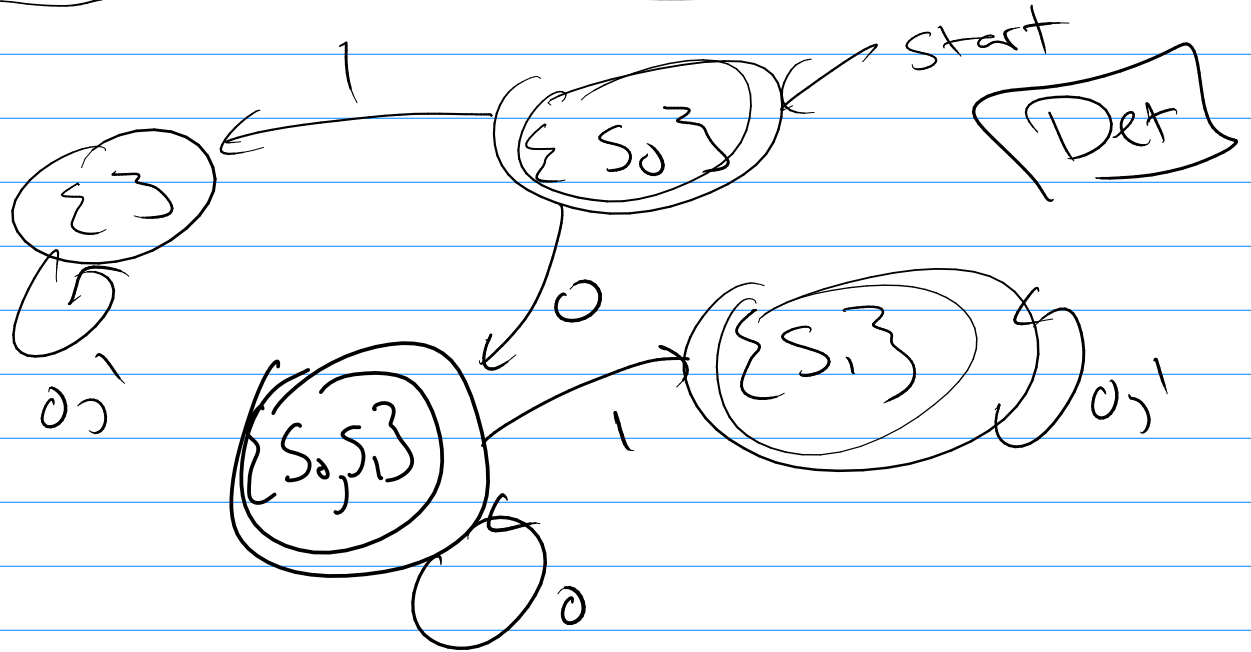


If the machine ends in a final state for any of the possible end states  $\rightarrow M_1$  recog.  $x$

Def.  $L(M_1)$  is all  $x \in I^*$  that the non-det. Machine recognizes.

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Th<sup>m</sup>: if  $M_1$  is non-det. then  $L(M_1)$  has a det. f.s.a.  $(M_2)$  where  $L(M_1) = L(M_2)$



$$L(M) = 0^* \cup 0^n \{0,1\}^* \mid n \geq 1$$

# Language Recognition (12.4)

## F.S.A. w/ Recognition

$th=1$  &  $th=2$

① A set is generated by a regular grammar

iff

② It is a regular set

iff

③ It is recognized by a F.S.A.

$th=1$

#2

#1

Regular Set (Inductively Def.)

a) Basis:  $\{\epsilon\}$ ,  $\{a\}$ ,  $\{x\}$

$x \in I$

are regular

b) Inductive:  $AB$ ,  $A \cup B$ ,  $A^*$   
are regular

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0<sup>n</sup> 1<sup>n</sup>