Distributed Systems, continued

Daniel Andresen

CIS520 – Operating Systems

RPC Failure Semantics

What happens if the server crashes? The client stub should:

- Use UNIX Semantics: Hang forever waiting for a reply that will never come (e.g., put the burden on the application programmer).
- Time out and raise an exception or report failure to the client.
- Time out and retransmit the request only satisfactory
 if the operation is idempotent it does not matter how
 many times it is executed by the server.

Semantics for duplicate requests:

- Exactly once unachievable due to server crash. At most once -the operation has been performed either zero or one times.
- At least once the client stub tries over and over until it gets a proper reply (e.g., this is only ok for idempotent operations). The client stub may use the last-of –many replies by attaching a transaction id to each request.

What happens if the client crashes?

A running server with no waiting client is called an **orphan**.

Orphans may cause problems by:

- using up cpu cycles,
- holding locks on files, or
- sending results that cause confusion when a client machine is rebooted.

Ways of dealing with orphans:

- Extermination: when a machine recovers, it checks if it had any RPC in progress. If so, it asks the server to kill any process running on its behalf. Extermination is performed recursively.
- Expiration: the server is given a fixed amount of time to complete a call. When a client machine recovers, any client must wait until all orphans have expired before submitting new requests.
- Reincarnation: kill all remote activity in the network, use epochs. Reincarnate only if the client cannot be located; that is, only orphans whose clients cannot be found are killed.

Types of Dist. OS support: Network Operating Systems

- users are aware of multiplicity of machines.
- Access to resources of various machines is done explicitly by remote logging (telnet or rlogin) into the appropriate remote machine.
- Networking software creates two-way link.
- Process on remote machine handles interactions acts as proxy for user.
- Access file on remote machine
- Compute just as any local user.
- Transferring data from remote machines to local machines, via the File Transfer Protocol (FTP) mechanism.
- Each computer has its own local file system.
- No real file sharing

Distributed Operating Systems

- Users are not aware of multiplicity of machines.
- Access to remote resources similar to access to local resources.
- Data Migration transfer data by transferring entire file, or transferring only those portions of the file necessary for the immediate task.
- Computation Migration transfers the computation, rather than the data, across the system.
- Process Migration executes an entire process, or parts of it, at different sites.

Distributed Operating Systems Features

- Load balancing distribute processes across network to even the workload.
- Computation speedup subprocesses can run concurrently on different sites.
- Hardware preferences process execution may require specialised processor.
- Software preferences required software might be available at only a particular site.
- Data access run process remotely, rather than transfer all data locally.
- Examples: GLUnix (NOW project), Amoeba, Legion (automated migration/scheduling)

Robustness

- A distributed system may suffer from various types of hardware failure.
- To ensure that a system robust, you need to detect failures, to reconfigure the system to continue and to recover when site or link is repaired

Failure Detection

It is difficult to differentiate between link failure, site failure and message loss.

Handshaking

- Used to detect link/site failure.
- Intermittent sends of "I-am-up" message between sites. If the message is not received, then there is a problem.
- If message is not received, send "Are-you-up" message.

Site failure – no answer; link failure – send by another site.

- message has set time wait for response.
 - If no reply, direct link down/site dawn/alt. Path down/message lost.

Reconfiguration

- If failure has occurred, the discovering site must initiate a procedure to allow the system to reconfigure to continue its normal mode of operation.
- if link is down, every site needs rerouting information.
- If site is down, every site needs to know, so it doesn't try to use services of the failed site.

Recovery from failure

- When a failed link or site is repaired, it must be integrated into the system gracefully and smoothly.
- if link fail, re-establish handshaking between linked sites.
- If site fail, re-establish site, notify others, update local table to current network sites.