



Faculty of Computers and Artificial Intelligence  
Cairo University  
Mid-term Exam



Program: Computer Science / Software Engineering  
Course Name: Concepts of Programming Languages  
Course Code: CS317 / SCS317  
Instructor(s): Dr. Amin Allam

Date: 10/12/2020  
Duration: 1 hour  
Total Marks: 20 marks

ID: ..... Full Name in Arabic: .....

تعليمات هامة:

- حيازة التليفون المحمول مفتوحا داخل لجنة الامتحان يعتبر حالة غش تستوجب العقاب وإذا كان ضروري الدخول بالمحمول فيوضع مغلقا في الحقائب .
- لا يسمح بدخول سماعة الأذن أو البلوتوث .
- لايسمح بدخول أي كتب أو ملازم أو أوراق داخل اللجنة والمخالفة تعتبر حالة غش .

Exam consists of 30 multiple-choice questions in 3 pages

**Qa** ⇒ For questions 1 to 7, consider the following BNF grammar:

```
<expr> -> <expr> * <term> | <term>
<term> -> <var> + <term> | <var>
<var> -> x | y | z
```

**1** The above BNF grammar is:

☐ A unambiguous ☐ B left associative ☐ C right associative ☐ D orthogonal ☐ E semantics

**2** The + operator in the above BNF grammar is:

☐ A unambiguous ☐ B left associative ☐ C right associative ☐ D orthogonal ☐ E semantics

**3** The \* operator in  $x+y*z$  according to the above BNF grammar is applied to:

☐ A y and z ☐ B x and y ☐ C x and z ☐ D  $x+y$  and z ☐ E cannot be determined

**4** The + operator in  $x+y*z$  according to the above BNF grammar is applied to:

☐ A y and z ☐ B x and y ☐ C x and z ☐ D x and  $y*z$  ☐ E cannot be determined

**5** In the above BNF grammar, y is:

☐ A right associative ☐ B terminal ☐ C nonterminal ☐ D metasymbol ☐ E low precedence

**6** The highest precedence operator in the above BNF grammar is:

☐ A  $\rightarrow$  ☐ B \* ☐ C + ☐ D | ☐ E several operators have equal highest precedence

**7** The second line of the above BNF grammar can be replaced in an EBNF grammar with  $\langle \text{term} \rangle \rightarrow \langle \text{var} \rangle \{ + \langle \text{var} \rangle \}$  which is equivalent except for:

☐ A orthogonality ☐ B precedence ☐ C associativity ☐ D efficiency ☐ E terminality

**Qb** **8** In an EBNF grammar having the rule  $\langle \text{expr} \rangle \rightarrow \langle \text{term} \rangle \{ (+|-) [-] (*|+) \langle \text{term} \rangle \}$ ,  $\langle \text{expr} \rangle$  can be expanded to:

☐ A  $\langle \text{term} \rangle + \langle \text{term} \rangle$  ☐ B  $\langle \text{term} \rangle + - \langle \text{term} \rangle$  ☐ C  $\langle \text{term} \rangle -- \langle \text{term} \rangle$   
☐ D  $\langle \text{term} \rangle ++ \langle \text{term} \rangle$  ☐ E  $\langle \text{term} \rangle + \langle \text{term} \rangle * \langle \text{term} \rangle$

**9** In an attribute grammar, the syntax rule  $\langle \text{expr} \rangle[1] \rightarrow \langle \text{expr} \rangle[2] * \langle \text{term} \rangle$  and the predicate  $\langle \text{expr} \rangle[2].\text{type} == \langle \text{term} \rangle.\text{type}$  mean that:

- ☐ A \* operator can be applied to all types    ☐ B types of LHS and RHS of \* operator must match  
☐ C there is an array of two expressions    ☐ D type of LHS expression is assigned to type of RHS term  
☐ E all expression types in the program are the same

**10** In an attribute grammar, the syntax rule  $\langle \text{expr} \rangle[1] \rightarrow \langle \text{expr} \rangle[2] * \langle \text{term} \rangle$  and the attribute computation function  $\langle \text{expr} \rangle[1].\text{type} \leftarrow \langle \text{term} \rangle.\text{type}$  mean that:

- ☐ A \* operator can be applied to all types    ☐ B types of LHS and RHS of \* operator must match  
☐ C there is an array of two expressions    ☐ D type of LHS expression is assigned to type of RHS term  
☐ E all expression types in the program are the same

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**Qc**  $\Rightarrow$  For questions 11 to 19, consider the following C++ program:

```
1 int r=8;
2 void F(int& t){t=5;}
3 int main()
4 {
5     int a=5, b=6; float c=1.2, d=2.9;
6     int x=a+r; float y=c+d; float z=a+c;
7     int* s=new int; F(*s); delete s;
8     return 0;
9 }
```

**11** The following variable is static:

- ☐ A r    ☐ B t    ☐ C x    ☐ D s    ☐ E no static variable exists

**12** The following line is related to orthogonality:

- ☐ A 2    ☐ B 5    ☐ C 6    ☐ D 7    ☐ E no such line exists

**13** The following line is related to dynamic type binding:

- ☐ A 1    ☐ B 2    ☐ C 6    ☐ D 7    ☐ E no such line exists

**14** The following variable is heap-dynamic:

- ☐ A r    ☐ B t    ☐ C s    ☐ D the unnamed variable created by new    ☐ E no such variable exists

**15** The following variable is an alias to another variable:

- ☐ A r    ☐ B t    ☐ C y    ☐ D the unnamed variable created by new    ☐ E no such variable exists

**16** Line 2 checks for types at:

- ☐ A compile time    ☐ B load time    ☐ C run time    ☐ D exception time    ☐ E no check

**17** The following line contains coercion:

- ☐ A 1    ☐ B 2    ☐ C 6    ☐ D 7    ☐ E no such line exists

**18** The following line causes a side effect:

- ☐ A 1    ☐ B 2    ☐ C 5    ☐ D 6    ☐ E no such line exists

**19** The following line contains static value binding:

- ☐ A 1    ☐ B 2    ☐ C 5    ☐ D 1 and 5    ☐ E none of the previous choices

- Qd** **20** One keyword with two different meanings mainly reduces:  
☐ A readability ☐ B writability ☐ C reliability ☐ D efficiency ☐ E generality
- 21** An interpreter interprets a statement inside a loop:  
☐ A once ☐ B twice ☐ C number of times equal to number of loop iterations  
☐ D same as compiler ☐ E zero times
- 22** Abstraction mainly improves:  
☐ A readability ☐ B writability ☐ C reliability ☐ D efficiency ☐ E portability
- 23** A static variable differs from other variables because it must:  
☐ A bind to values at load time ☐ B bind to type at compile time ☐ C bind to storage at load time  
☐ D bind to values at run time ☐ E bind to values at compile time
- 24** If a language supports short-circuit evaluation for `&&` (logical AND), consider `(false && g)`:  
☐ A `g` is evaluated ☐ B `g` is not evaluated ☐ C `g` may be evaluated ☐ D error ☐ E exception
- 25** If a language supports short-circuit evaluation for `&&` (logical AND), consider `(true && g)`:  
☐ A `g` is evaluated ☐ B `g` is not evaluated ☐ C `g` may be evaluated ☐ D error ☐ E exception
- 26** Type checking achieves its maximum reliability if it is done:  
☐ A at compile time ☐ B at load time ☐ C at run time ☐ D when the related function is called  
☐ E immediately before an error occurs
- 27** An `int` variable and an `int*` variable are:  
☐ A compatible for name type equivalence rules ☐ B compatible for structure type equivalence rules  
☐ C all previous choices ☐ D never compatible ☐ E implicitly converted for name type equivalence
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**Qe** ⇒ For questions 28 to 30, consider the following C++ program and assume static scoping:

```
1 int x,y,z,r;  
2 void main()  
3 {  
4     int x,y,z;  
5     while(x<10)  
6     {  
7         int x,y,w;  
8         if(x<5) {int x;}  
9     }  
10 }
```

- 28** Assuming static scoping, the referencing environment of Line 8 does not contain:  
☐ A `x` of line 4 ☐ B `y` of line 7 ☐ C `r` ☐ D `w` ☐ E none of the previous choices
- 29** Assuming static scoping, the referencing environment of Line 5 consists exactly of:  
☐ A `x,y,z,r` of line 1 ☐ B `x,y,z` of line 4 ☐ C `x,y,z` of line 4 and `r` ☐ D `r` ☐ E `x,y,z` of line 1
- 30** Assuming static scoping, the variable `z` of line 4 with respect to the block from line 6 to line 9 is:  
☐ A local ☐ B nonlocal ☐ C global ☐ D invisible ☐ E not related
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