#### Lesson 10

Introduction to algorithms Sorting algorithms

Bubble sort

Selection sort

Counting sort



The goal of an algorithm – to solve a particular problem as efficiently as possible



Algorithm Complexity – a measure for the efficiency of an algorithm



# **Time Complexity**

how fast the algorithm is



# **Space Complexity**

how much memory the algorithm requires



Big O Notation (simplified) O(N) - the algorithm executes approximately N simple operations, where N is the size of the input



Popular complexities for input of size N

Number of operations	Big O	Short name
1, 5, 10, 10 000	O(1)	constant
logN, $logN + 10$	O(logN)	logarithmic
N, 2N, 10N + 5	O( N )	linear
N log N, N log 2N	O(N log N)	linearithmic
$N^2$ , $3N^2 + N + 1$	$O(N^2)$	quadratic
$N^3$ , $N^3 + N^2$	$O(N^3)$	cubic
N <sup>K</sup> , K − constant	$O(N^K)$	polynomial
$2^N$ , $3^N$ , $2^N + N^2$	$O(2^{N})$	exponential
<b>∧</b> !	O( N! )	factorial



# Complexity in the:

best case average case worst case



## Sorting algorithms I

Sorting problem –
rearrange a sequence of items,
so that all the elements in the resulting sequence
are in increasing (decreasing) order



# Sorting algorithms I

Sort an array of integers in increasing order



# Mostly used in practice

Usually using predefined function for sorting



#### **Bubble Sort**

- Simple sorting algorithm
- Worst case scenario of N^2 operations
- Best case scenario of N operations
- Used for small arrays



# **Bubble Sort Philosophy**

- Traverse elements
- Compare adjacent elements
- Swap adjacent elements if the next element is greater



#### **Bubble Sort Pseudo Code**

- Compare array[0] & array[1]
- If array[0] > array [1] swap it.
- Compare array[1] & array[2]
- If array[1] > array[2] swap it.
- •
- Compare array[n-1] & array[n]
- if array[n-1] > array[n] then swap it.
- Repeat the same steps for array[1] to array[n-1]



# Try it! Also try to optimize the algorithm.



#### **Bubble Sort**

```
public static int[] bubbleSort(int[] intArray) {
       int temp = 0;
       for(int i=0; i < intArray.length; i++){</pre>
           for(int j=1; j < (intArray.length - i); j++){</pre>
               if(intArray[j-1] > intArray[j]){
                   //swap the elements!
                   temp = intArray[j-1];
                   intArray[j-1] = intArray[j];
                   intArray[j] = temp;
       return intArray;
```



#### **Selection Sort**

- Used only for small arrays
- Slow for bigger arrays
- Worst case scenario of N^2 operations
- Best case scenario of N^2 operations
- Find the next smallest element and move it to its final position
- Start by the smallest element
- Subset of sorted and unsorted elements



## Selection Sort Algorithm

- Initially, the sorted sublist is empty and the unsorted sublist is the entire input list
- Start with first element of the unsorted subset
- Test with all elements after the assumed min element to find the smallest
- Swap the smallest element to the starting index



# Try it! Also try to optimize the algorithm.



#### **Selection Sort**

```
public static int[] selectionSort(int[] a){
/* a[0] to a[n-1] is the array to sort */
   int temp = 0;
   int iMin = 0:
// advance the position through the entire array
// (could do j < n-1 because single element is also min element)
   for (int j = 0; j < a.length-1; j++) {</pre>
   /* find the min element in the unsorted a[j .. n-1] */
   /* assume the min is the first element */
   iMin = j;
   /* test against elements after j to find the smallest */
   for ( int i = j+1; i < a.length; i++) {</pre>
   /* if this element is less, then it is the new minimum */
       if (a[i] < a[iMin]) {
       /* found new minimum; remember its index */
           iMin = i;
```



#### Selection Sort

```
/* iMin is the index of the minimum element.
    Swap it with the current position */
    if ( iMin != j ) {
        //swap the elements!
        temp = a[j];
        a[j] = a[iMin];
        a[iMin] = temp;
    }
}
return a;
}
```



# **Counting Sort**

Counting sort is an algorithm for sorting small integers.

 Its main idea is to count the number of different objects.

Runs in linear running time in the number of items.

# **Counting Sort Algorithm**

### **Basic implementation**

Find min/max values

 For every element, count in separate array the number of times it was met.

 Then for every value from min to max in the array ouput K times this value, where K is the number calculated in the previous array.

# **Counting Sort Algorithm**

#### Input Data

0	4	2	2	0	0	1	1	0	1	0	2	4	2
_					-	·—-	_	-	_	-		-	

#### Count Array

#### Sorted Data

0	0	0	0	0	1	1	1	2	2	2	2	4	4
---	---	---	---	---	---	---	---	---	---	---	---	---	---

# **Counting Sort Basic Version**

```
public static int[] countingSort(int[] numbers, int max)
   int[] counts = new int[max + 1];
   for (int i = 0; i < numbers.length; i++)</pre>
      counts[numbers[i]]++;
   int[] results = new int[numbers.length];
   int counter = 0;
   for(int i = 0; i <= counts.length-1; i++)</pre>
       while(counts[i]>0)
           results[counter++] = i;
           counts[i]--;
   return results;
```

# **Counting Sort Complexity**

- Best Case O(N+K)
- Average Case O(N+K)
- Worst Case O(N+K)

Where K is the largest element in sequence.

# Sorting Complexity Refresh

		Best	Average	vvorst
•	<b>Bubble Sort</b>	O(n)	$O(n^2)$	$O(n^2)$

• Selection Sort  $O(n^2)$   $O(n^2)$ 

• Counting Sort O(n+k) O(n+k) O(n+k)

• Radix Sort  $O(nlg_b n) O(nlg_b n)$   $O(nlg_b n)$ 

### Resources

Counting and Radix Sort Explained:

http://courses.csail.mit.edu/6.006/spring11/rec/rec11.pdf

# Summary

- Algoritm Complexity
- Most used in practice
- Bubble Sort
- Selection Sort
- Counting Sort

