

Lab2

$$\begin{array}{l}
 \text{2a) (1) } \frac{}{a \in \text{VObjects}} \quad \frac{}{b \in \text{VObjects}} \\
 \\
 \text{(2) } \frac{s1 \in \text{AObjects} \quad s2 \in \text{AObjects} \quad s \in \text{VObjects}}{s1 \& s2 \in \text{AObjects}} \quad \frac{}{s \in \text{AObjects}}
 \end{array}$$

2b) If we want a result of $b \& b \& b$, there is more than one answer:

$$\begin{array}{l}
 A \Rightarrow A \& A \qquad \qquad \qquad A \Rightarrow A \& A \\
 | \quad | \qquad \qquad \qquad | \quad | \\
 A \& A \& A \Rightarrow A \& A \& V \qquad A \& A \& A \Rightarrow V \& A \& A \\
 | \quad | \quad | \qquad \qquad | \quad | \quad | \\
 A \& A \& b \Rightarrow A \& V \& b \qquad b \& A \& A \Rightarrow b \& V \& A \\
 | \quad | \quad | \qquad \qquad | \quad | \quad | \\
 A \& b \& b \Rightarrow V \& b \& b \qquad b \& b \& A \Rightarrow b \& b \& V \\
 | \quad | \quad | \qquad \qquad | \quad | \quad | \\
 b \& b \& b \qquad \qquad \qquad b \& b \& b
 \end{array}$$

Since there is not ONE specific way to find an answer, the grammar is defined as ambiguous.

$$2(c) L(g) ::= \{a^n, b^n, c^n : n \geq 0\}$$

2(d) Only sentences 1 and 4 are in the language generated by this grammar.

$$\text{baab: } S \Rightarrow AaBb \Rightarrow baBb \Rightarrow baab$$

$$\text{bbaab: } S \Rightarrow AaBb \Rightarrow AbaBb \Rightarrow bbaBb \Rightarrow bbaab$$

2(e) Only sentences 1 and 5 are in the language generated by this grammar.

$$\begin{array}{l}
 \text{abcd} \\
 S \\
 / \ / \ \backslash \ \backslash \\
 a \ S \ c \ B \\
 | \quad | \\
 b \quad d \\
 \text{accc} \\
 S \\
 / \ / \ \backslash \ \backslash \\
 a \ S \ c \ B \\
 | \quad | \\
 A \ A \\
 | \quad | \\
 c \quad c
 \end{array}$$

- 3(a)(i) The first grammar is left associative because the recursive non-terminal symbol is on the left of the parse tree.

The 2nd grammar is right associative because the recursive non-terminal symbol is on the right of the parse tree.

- (ii) The two grammars do produce the same expression, one by right associative and the other by left associative. The results will be in a combination of: operand operator operand operator ... operand
The difference is determined by if it adds an operand to the left or right

- (b) Example: $4 - 1 \ll 3$

```
val exp = 4 - 1 << 3
val minus = (4 - 1) << 3
val shift = 4 - (1 << 3)
if (exp == minus)
  println("- has higher precedence over <<")
if (exp == shift)
  println("<< has higher precedence over -")
```

After using the scala interpreter, we discover the - has higher precedence over <<. By using the expression, $4 - 1 \ll 3$, the result was 24. We compare that answer to $(4 - 1) \ll 3$ and $4 - (1 \ll 3)$. $(4 - 1) \ll 3$ is 24 while $4 - (1 \ll 3)$ is -4. Therefore, this shows us that - has higher precedence over <<. This is interesting because of the fact that << is also the same as multiplying by base 2.

- (c) $\text{frac} ::= -n.n|n.n|z.n|-z.n|-n.z|n.z|\text{Explz}.z$
 $\text{Exp} ::= z.n|\text{Enl}-z.n|\text{Enln}.n|\text{Enl}-n.n|\text{Enln}.z|\text{Enl}-n.z|\text{En}$
 $z ::= 0$
 $n ::= 1|2|3|4|5|6|\dots|\text{linfinity}$