

Ch3: Minimum-Time Trajectory Generation

Time-Optimal Time Scaling

Path, Time Scaling, and Trajectory

Path $\mathcal{C}(s)$ is a purely geometric description of the sequence of configurations achieved by the robot:

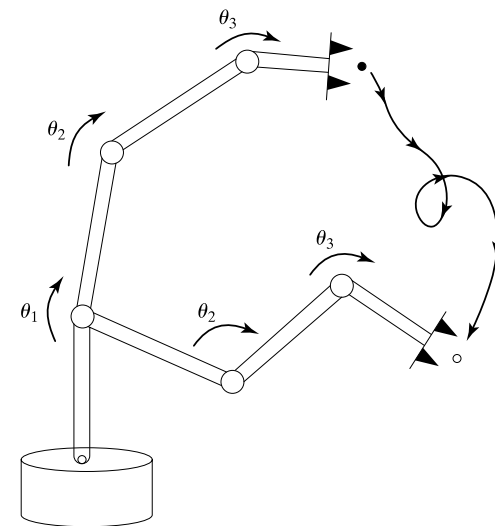
$$\mathcal{C}: \underbrace{[0,1]}_{s \in [0,1]: \text{scalar path parameter (0 at the start and 1 at the end of the path)}} \rightarrow \mathbb{C} \quad \text{Robot's C-space}$$

- As s increases from 0 to 1, the robot moves along the path.

Time Scaling $s(t)$ specifies the times when those robot configurations are reached:

$$s: [0, t_f] \rightarrow [0,1]$$

Trajectory $\mathcal{C}(s(t))$ or $\mathcal{C}(t)$ specifies the robot configuration as a function of time, i.e., the combination of a path and a time scaling.



Some Examples of Path Planning $\mathcal{C}(s)$, $s \in [0,1]$

- Point-to-Point Straight-Line Path in Joint Space: $\boldsymbol{\theta}(s) = \boldsymbol{\theta}_{\text{start}} + s(\boldsymbol{\theta}_{\text{end}} - \boldsymbol{\theta}_{\text{start}})$
 $\boldsymbol{\theta} \in \mathbb{R}^n$

- Point-to-Point Straight-Line Path in Task Space (in Cartesian Space \mathbb{R}^3):

$$\mathbf{x}(s) = \mathbf{x}_{\text{start}} + s(\mathbf{x}_{\text{end}} - \mathbf{x}_{\text{start}}) \quad \mathbf{x} \in \mathbb{R}^m: \text{minimum set of coordinates}$$

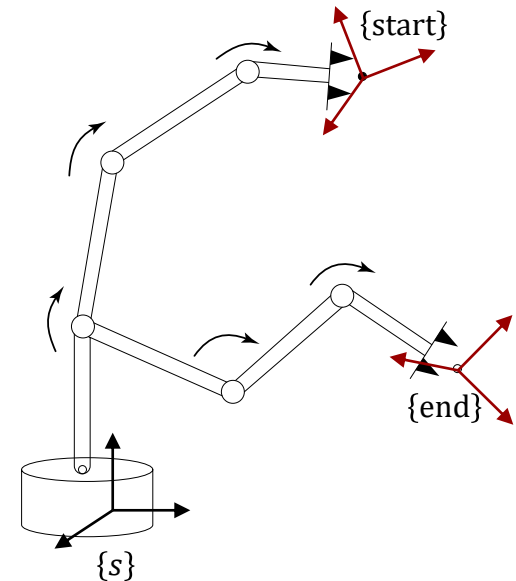
or

$$\mathbf{p}(s) = \mathbf{p}_{\text{start}} + s(\mathbf{p}_{\text{end}} - \mathbf{p}_{\text{start}}) \quad \mathbf{p} \in \mathbb{R}^3, \mathbf{R} \in SO(3)$$

$$\mathbf{R}(s) = \mathbf{R}_{\text{start}} \exp(\underbrace{\log(\mathbf{R}_{\text{start}}^T \mathbf{R}_{\text{end}})}_{\mathbf{R}_{\text{start,end}}} s)$$

- Point-to-Point Straight-Line Path in Task Space (in $SE(3)$):

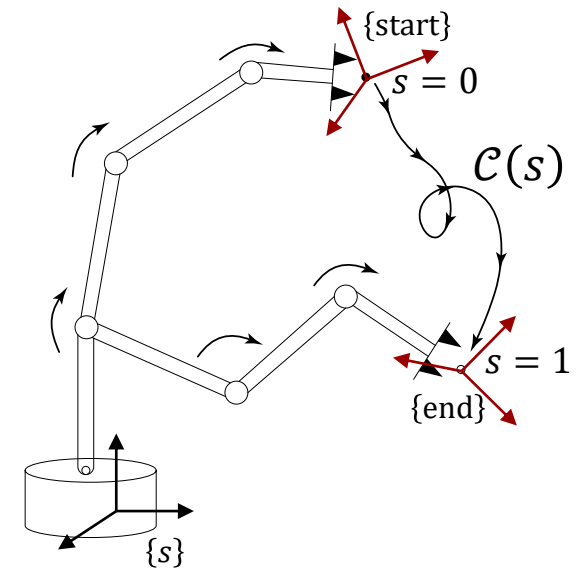
$$\mathbf{T}(s) = \mathbf{T}_{\text{start}} \exp(\underbrace{\log(\mathbf{T}_{\text{start}}^{-1} \mathbf{T}_{\text{end}})}_{\mathbf{T}_{\text{start,end}}} s) \quad \mathbf{T} = (\mathbf{R}, \mathbf{p}) \in SE(3)$$



Time-Optimal Time Scaling

Consider a case where the **path** $\mathcal{C}(s)$, $s \in [0,1]$, is fully specified by the task or an obstacle-avoiding path planner. The **time-optimal time scaling** is finding a **time scaling** $s(t)$ that minimizes the time of motion along the path, subject to the robot's actuator limits.

A time-optimal trajectory maximizes the robot's productivity.



Actuation Constraints as a Function of s

In practice, the robot dynamics and joint actuator limits dependent on $(\boldsymbol{\theta}, \dot{\boldsymbol{\theta}})$, thus, the maximum available velocities and accelerations change along the path.

$$\boldsymbol{\tau} = \mathbf{M}(\boldsymbol{\theta})\ddot{\boldsymbol{\theta}} + \dot{\boldsymbol{\theta}}^T \boldsymbol{\Gamma}(\boldsymbol{\theta})\dot{\boldsymbol{\theta}} + \mathbf{g}(\boldsymbol{\theta}) \quad (1)$$

$$\tau_i^{\min}(\boldsymbol{\theta}, \dot{\boldsymbol{\theta}}) \leq \tau_i \leq \tau_i^{\max}(\boldsymbol{\theta}, \dot{\boldsymbol{\theta}}) \quad i = 1, \dots, n \quad (\text{actuation constraints}) \quad (2)$$

A path $\mathcal{C}(s)$ can be always expressed in joint space $\boldsymbol{\theta}(s) \in \mathbb{R}^n$ using inverse kinematics. Thus,

$$\dot{\boldsymbol{\theta}} = \frac{d\boldsymbol{\theta}}{ds} \dot{s}, \quad \ddot{\boldsymbol{\theta}} = \frac{d\boldsymbol{\theta}}{ds} \ddot{s} + \frac{d^2\boldsymbol{\theta}}{ds^2} \dot{s}^2$$

Dynamics along the path:

$$\begin{aligned} (1) \quad \rightarrow \quad \boldsymbol{\tau} &= \underbrace{\left(\mathbf{M}(\boldsymbol{\theta}(s)) \frac{d\boldsymbol{\theta}}{ds} \right)}_{\mathbf{m}(s) \in \mathbb{R}^n} \ddot{s} + \underbrace{\left(\mathbf{M}(\boldsymbol{\theta}(s)) \frac{d^2\boldsymbol{\theta}}{ds^2} + \left(\frac{d\boldsymbol{\theta}}{ds} \right)^T \boldsymbol{\Gamma}(\boldsymbol{\theta}(s)) \frac{d\boldsymbol{\theta}}{ds} \right)}_{\mathbf{c}(s) \in \mathbb{R}^n} \dot{s}^2 + \underbrace{\mathbf{g}(\boldsymbol{\theta}(s))}_{\mathbf{g}(s) \in \mathbb{R}^n} \\ &= \mathbf{m}(s)\ddot{s} + \mathbf{c}(s)\dot{s}^2 + \mathbf{g}(s) \quad (3) \end{aligned}$$

Actuation Constraints as a Function of s

$$(2) \rightarrow \tau_i^{\min}(s, \dot{s}) \leq \tau_i \leq \tau_i^{\max}(s, \dot{s}) \quad (4)$$

$$(3), (4) \rightarrow \tau_i^{\min}(s, \dot{s}) \leq m_i(s)\ddot{s} + c_i(s)\dot{s}^2 + g_i(s) \leq \tau_i^{\max}(s, \dot{s}) \quad (5)$$

Let define $L_i(s, \dot{s})$ and $U_i(s, \dot{s})$ be the minimum and maximum \ddot{s} satisfying (5):

$$\left\{ \begin{array}{ll} \text{- If } m_i(s) > 0 : & \begin{aligned} L_i(s, \dot{s}) &= \frac{\tau_i^{\min}(s, \dot{s}) - c_i(s)\dot{s}^2 - g_i(s)}{m_i(s)} \\ U_i(s, \dot{s}) &= \frac{\tau_i^{\max}(s, \dot{s}) - c_i(s)\dot{s}^2 - g_i(s)}{m_i(s)} \end{aligned} \\ \text{- If } m_i(s) < 0 : & \begin{aligned} L_i(s, \dot{s}) &= \frac{\tau_i^{\max}(s, \dot{s}) - c_i(s)\dot{s}^2 - g_i(s)}{m_i(s)} \\ U_i(s, \dot{s}) &= \frac{\tau_i^{\min}(s, \dot{s}) - c_i(s)\dot{s}^2 - g_i(s)}{m_i(s)} \end{aligned} \end{array} \right.$$

By defining $L(s, \dot{s}) = \max_i L_i(s, \dot{s})$ and $U(s, \dot{s}) = \min_i U_i(s, \dot{s})$ as the lower and upper bounds on \ddot{s} at (s, \dot{s}) , (5) can be written as $L(s, \dot{s}) \leq \ddot{s} \leq U(s, \dot{s})$

Time-optimal Time-scaling Problem

Given a path $\theta(s)$, $s \in [0,1]$, an initial state $(s_0, \dot{s}_0) = (0,0)$, and a final state $(s_f, \dot{s}_f) = (1,0)$, find a monotonically increasing twice-differentiable time-scaling $s(t)$, $s: [0, t_f] \rightarrow [0,1]$ that

(a) satisfies: $s(0) = \dot{s}(0) = \dot{s}(t_f) = 0$ and $s(t_f) = 1$,

(b) minimizes the total travel time t_f along the path while respecting the actuator constraints:

$$L(s, \dot{s}) \leq \ddot{s} \leq U(s, \dot{s})$$

$\equiv \dot{s}(t) \geq 0$ (robot moves forward along the path)

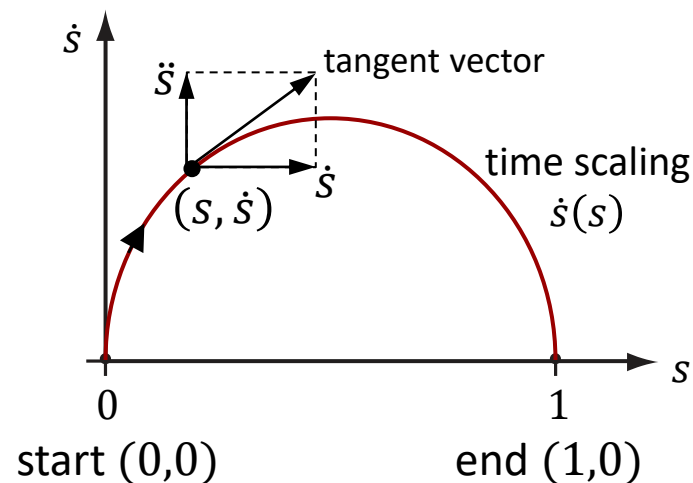
This problem is easily visualized in the (s, \dot{s}) **phase plane**.

(s, \dot{s}) Phase Plane

(s, \dot{s}) Phase Plane

(s, \dot{s}) **phase plane** is defined as a plane with s running from 0 to 1 on a horizontal axis and \dot{s} on a vertical axis.

A time scaling $s(t)$ of the path is any curve $\dot{s}(s)$ in the phase plane that moves monotonically to the right from $(0,0)$ to $(1,0)$ in the top-right quadrant.

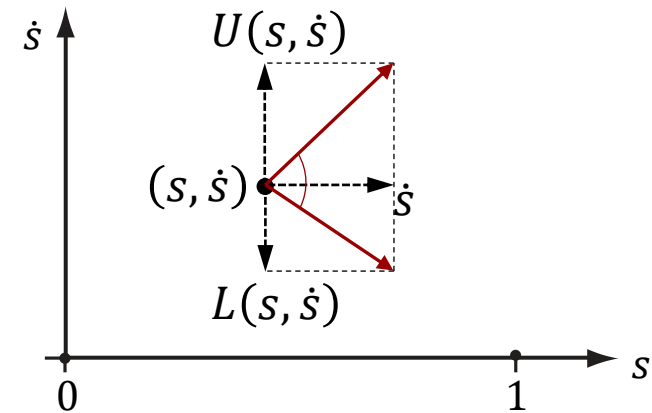


Among all these curves, we are looking for a time-optimal curve that satisfy the actuator/acceleration constraints $L(s, \dot{s}) \leq \ddot{s} \leq U(s, \dot{s})$.

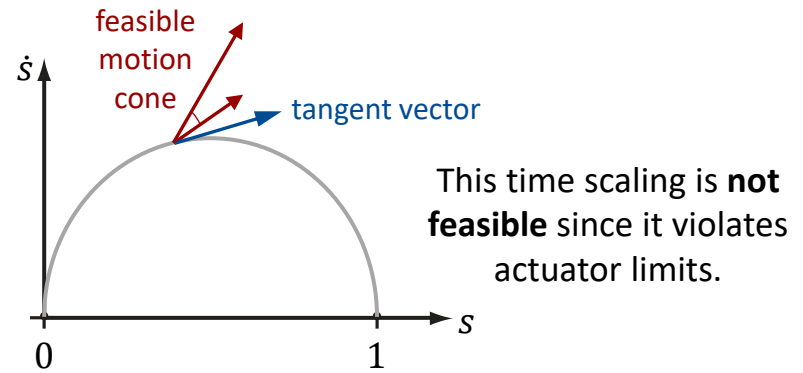
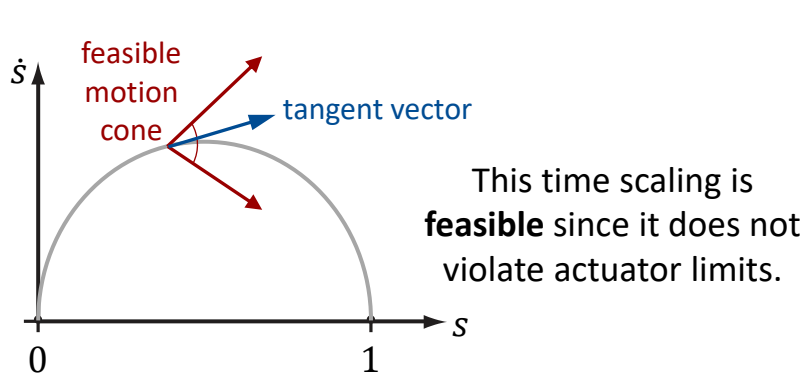
Feasible Motion Cone

By drawing the range of feasible accelerations $L(s, \dot{s}) \leq \ddot{s} \leq U(s, \dot{s})$ according to the dynamics at any state (s, \dot{s}) , we find a cone called the **feasible motion cone**.

Note: Vector \dot{s} is proportional to the height of the point along the \dot{s} axis.

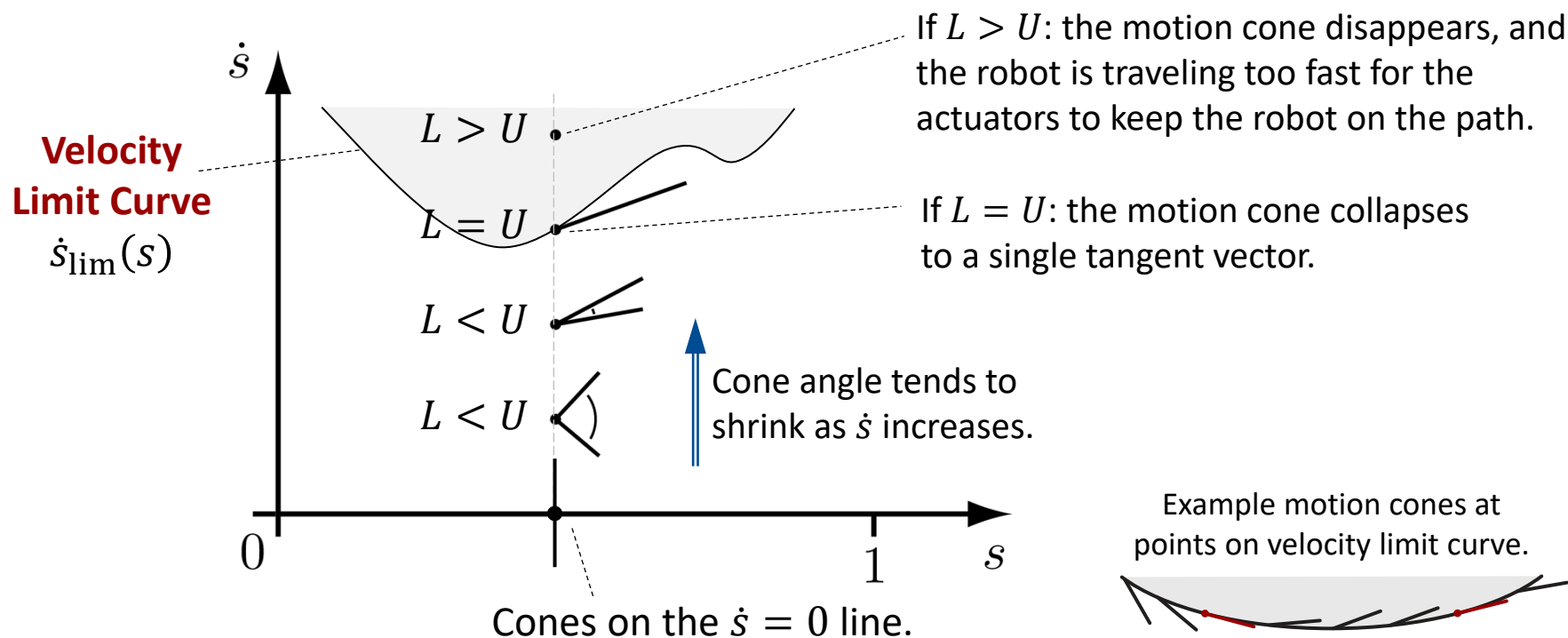


At each state (s, \dot{s}) , the **tangent vector** to the time scaling must lie **inside feasible motion cone** to satisfy the actuator limits (or the acceleration constraints).



Velocity Limit Curve

Let keep s constant but increase \dot{s} :



At states on velocity limit curve, only a single acceleration is possible; at states above this curve (inadmissible states), the robot leaves the path immediately; and at states below the curve, there is a cone of possible tangent vectors.

Bang-Bang Time Scaling

The total time of motion t_f can be written as

$$t_f = \int_0^{t_f} 1 dt = \int_0^{t_f} \frac{ds}{\dot{s}} dt = \int_0^1 \frac{dt}{ds} ds = \int_0^1 \dot{s}^{-1}(s) ds$$

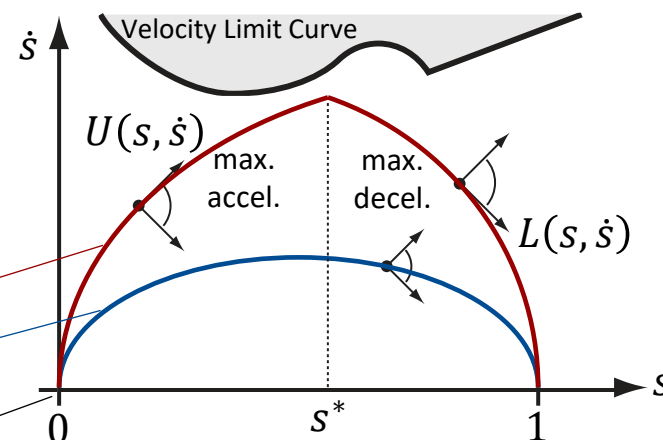
For a minimum-time motion, \dot{s}^{-1} should be as small as possible, and therefore, \dot{s} must be as large as possible, at all s , while still satisfying the acceleration constraints $L(s, \dot{s}) \leq \ddot{s} \leq U(s, \dot{s})$ and the boundary constraints $s(0) = \dot{s}(0) = \dot{s}(t_f) = 0, s(t_f) = 1$.

This implies that the time scaling must always operate either at the limit $U(s, \dot{s})$ (the upper edge of the motion cones) or at the limit $L(s, \dot{s})$ (the lower edge of the motion cones), and we should determine switching point s^* between these limits.

Time-optimal “bang-bang” time scaling

A non-optimal time scaling

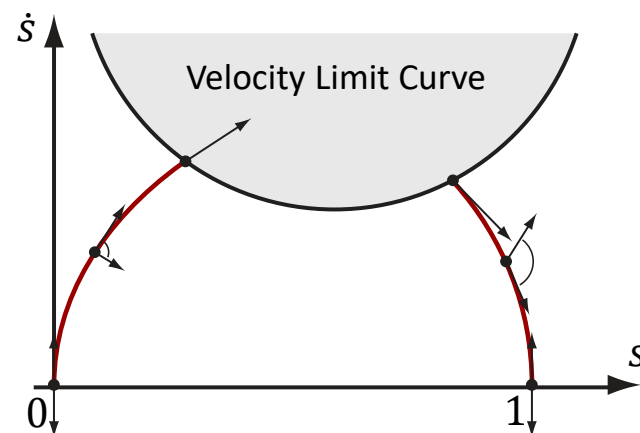
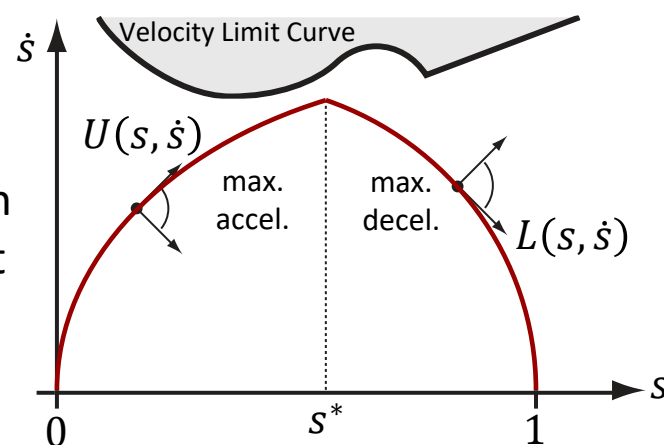
The curve must be normal to the s -axis when $\dot{s} = 0$.



Bang-Bang Time Scaling

In general, the time scaling is calculated by numerically integrating $\ddot{s} = U(s, \dot{s})$ (the maximum possible accelerations) forward in s from $(0,0)$, integrating $\ddot{s} = L(s, \dot{s})$ (the maximum possible decelerations) backward in s from $(1,0)$, and finding the intersection (switching point s^*) of these curves.

However, in some cases, the existence of a velocity limit curve prevents a single-switch solution (two curves do not intersect and run into the velocity limit curve). In these cases, bang-bang control is not possible, and it requires an algorithm to find multiple switching points.



Time-Scaling Algorithm

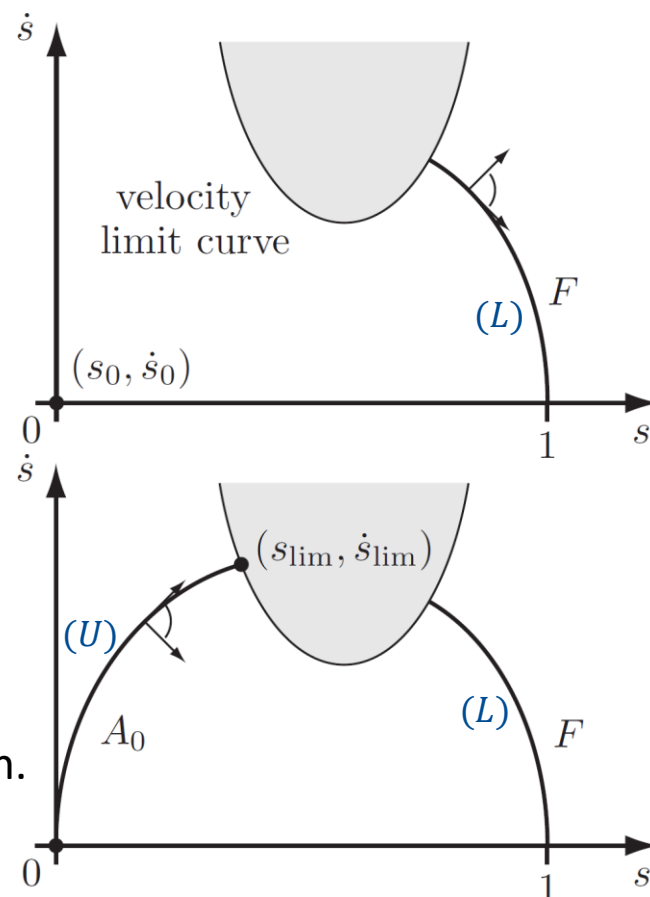
Time-Scaling Algorithm

Since time-optimal trajectories consist of only maximum acceleration $U(s, \dot{s})$ and minimum acceleration $L(s, \dot{s})$, we need to find the switches between U and L :

Step 1: Integrate $\ddot{s} = L(s, \dot{s})$ backward in time from $(1,0)$ until (i) the velocity limit curve is penetrated ($L(s, \dot{s}) > U(s, \dot{s})$) or (ii) $s = 0$. Call this curve F .

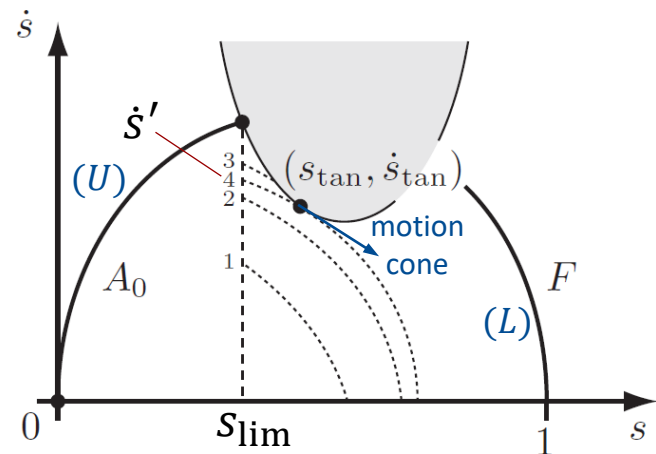
Step 2: Integrate $\ddot{s} = U(s, \dot{s})$ forward in time from $(0,0)$ until (i) it intersects F or (ii) until the velocity limit curve is penetrated ($L(s, \dot{s}) > U(s, \dot{s})$). Call this curve A_0 .

- If (i) happens, the problem is solved.
- If (ii) happens, let $(s_{\text{lim}}, \dot{s}_{\text{lim}})$ be the point of penetration.

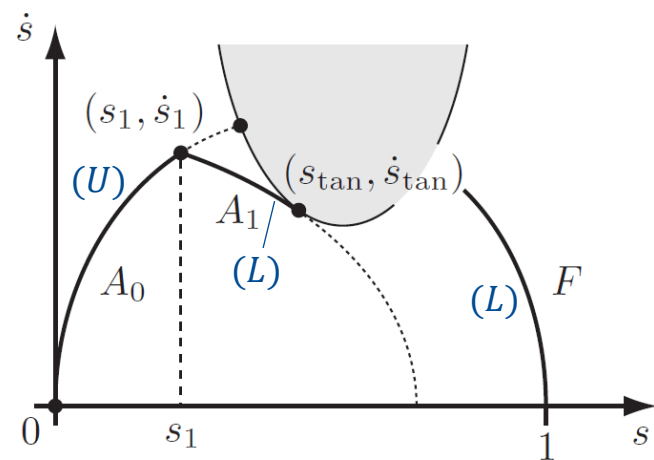


Time-Scaling Algorithm

Step 3: Perform a binary search (or half-interval search) on the velocity in the range $[0, \dot{s}_{\text{lim}}]$ at s_{lim} to find the velocity \dot{s}' such that the curve integrating $\ddot{s} = L(s, \dot{s})$ forward in time from $(s_{\text{lim}}, \dot{s}')$ touches the velocity limit curve tangentially (or comes closest to the curve within a specified tolerance without hitting it) at $(s_{\text{tan}}, \dot{s}_{\text{tan}})$.



Step 4: Integrate $\ddot{s} = L(s, \dot{s})$ backward in time from $(s_{\text{tan}}, \dot{s}_{\text{tan}})$ until it intersects A_0 at (s_1, \dot{s}_1) . Call this curve A_1 . (s_1, \dot{s}_1) is the first switch point from maximum acceleration to maximum deceleration.



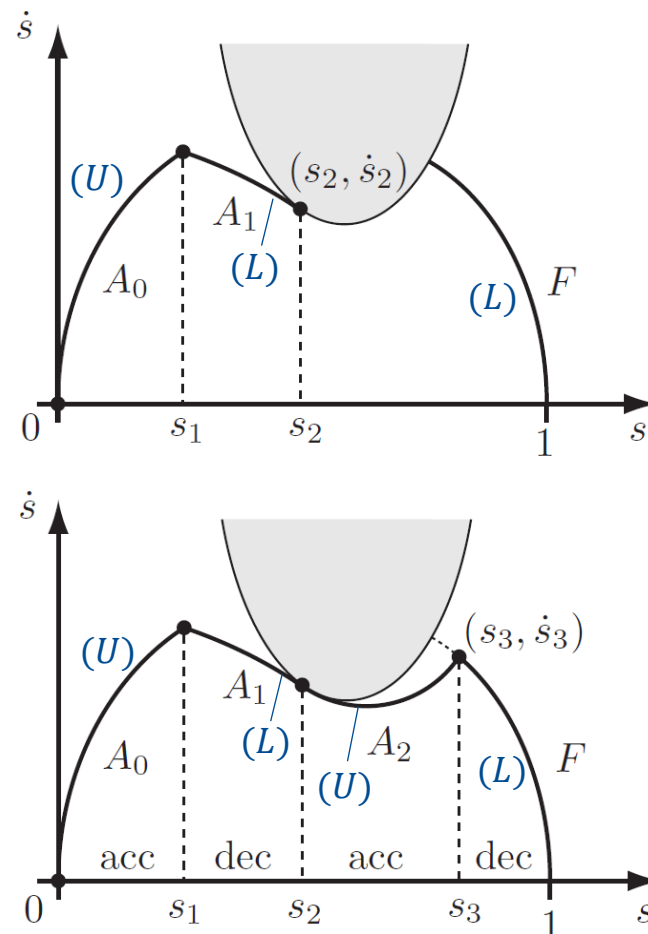
Time-Scaling Algorithm

Step 5: Mark the tangent point $(s_{\text{tan}}, \dot{s}_{\text{tan}})$ as the switch point (s_2, \dot{s}_2) from maximum deceleration to maximum acceleration.

Step 6: Go back to **Step 2**, i.e., integrate $\ddot{s} = U(s, \dot{s})$ forward in time from (s_2, \dot{s}_2) until (i) it intersects F or (ii) until the velocity limit curve is penetrated again ($L(s, \dot{s}) > U(s, \dot{s})$). Call this curve A_2 .

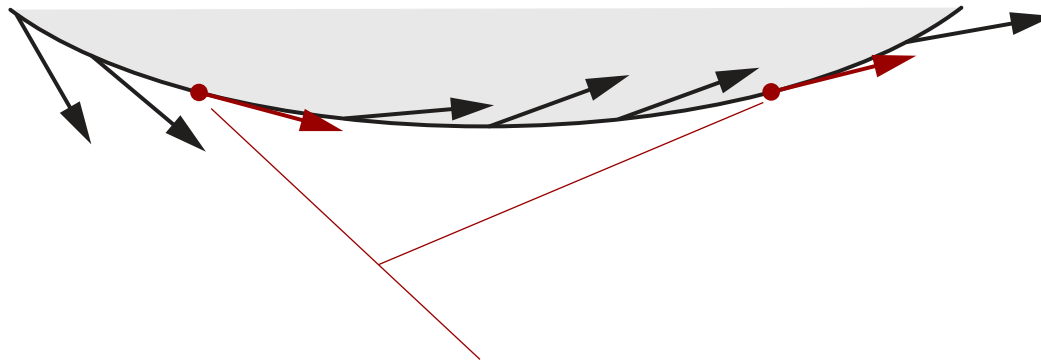
- If (i) happens, the intersection point (s_3, \dot{s}_3) is the final switch point from maximum acceleration to maximum deceleration and the algorithm is complete.
- If (ii) happens, let $(s_{\text{lim}}, \dot{s}_{\text{lim}})$ be the new point of penetration and repeat the process from **Step 3**.

Note: Now, by having $\dot{s}(s)$ and integrating $\dot{\theta} = \frac{d\theta}{ds} \dot{s}$, we can find $\theta(t)$.



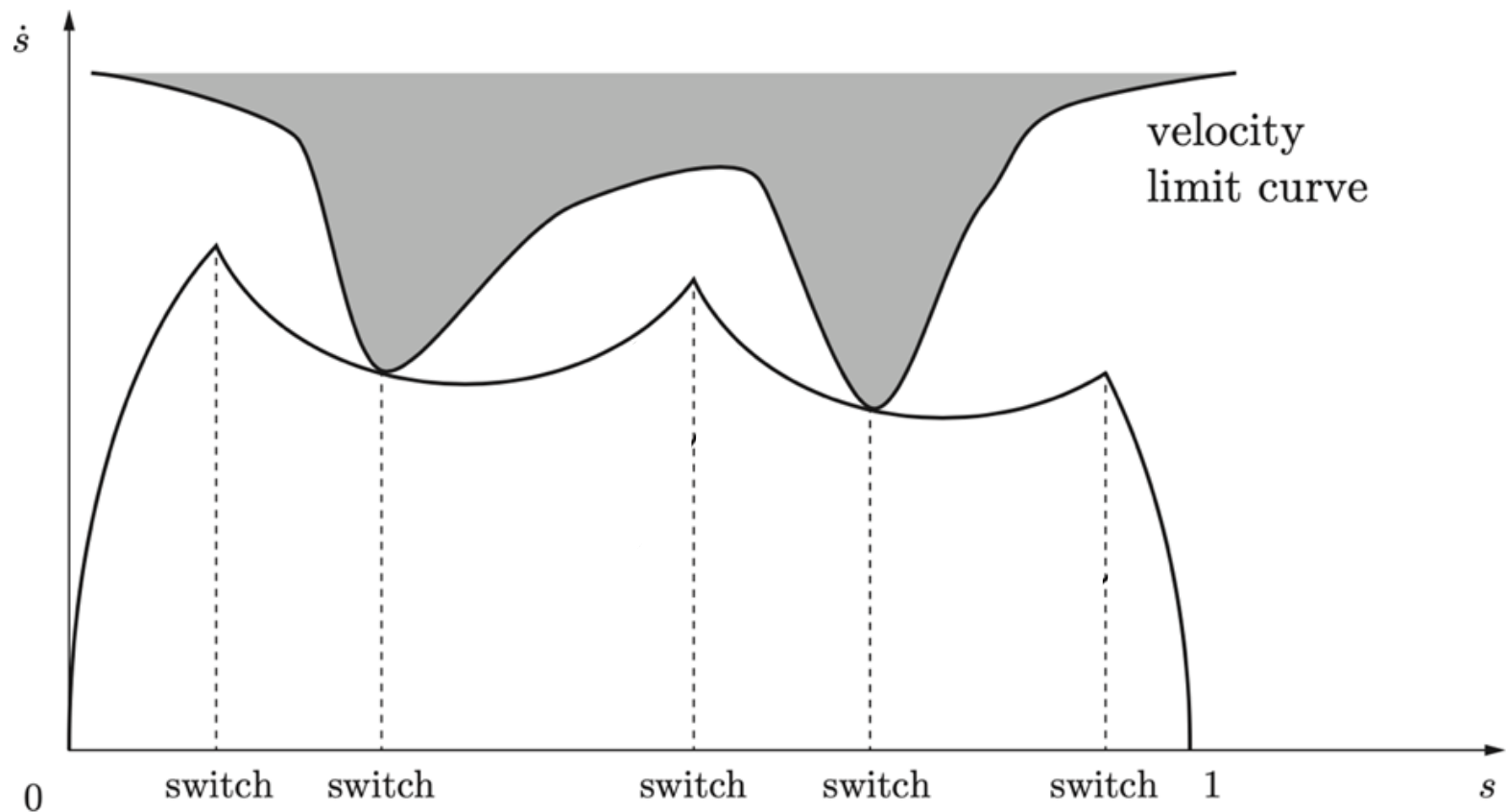
A Quick Method to Find Switch Points on Velocity Limit Curve

The only points on the velocity limit curve ($L(s, \dot{s}) = U(s, \dot{s})$) that can be part of an optimal solution are those where the feasible motion vector is tangent to the velocity limit curve ($s_{\text{tan}}, \dot{s}_{\text{tan}}$). Therefore, the binary search in the time-scaling algorithm can be replaced by a more computationally efficient approach of numerical construction of the velocity limit curve and a searching on this curve for points where the motion vector is tangent to the curve.



The points that can belong to a time-optimal time scaling.

An Example of Multi-Switch Time-Optimal Time Scaling



Example

Draw the feasible time-optimal time-scaling for a driver rushing home with the max braking and max acceleration integral curves shown.

