Natural Language Processing CSE 325/425

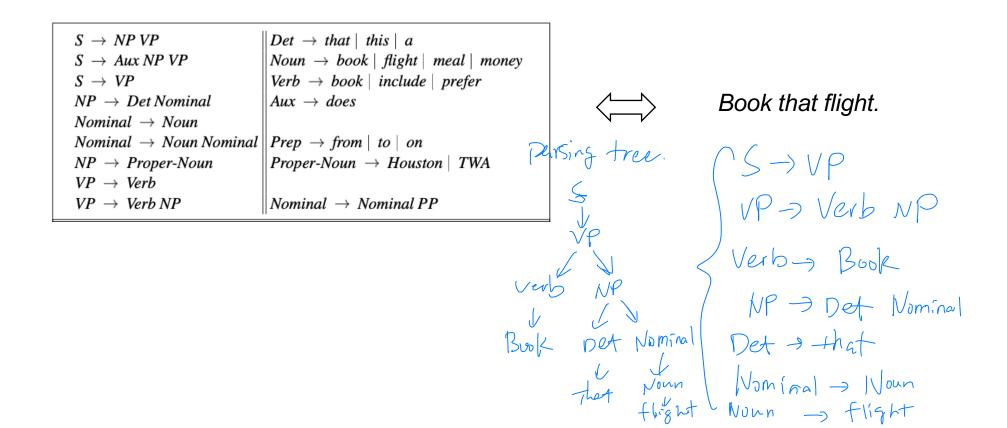


Lecture 16:

Syntactic parsing

Syntactic parsing: assign valid parsing trees to a sentene, using a CFG.

- derivations of a sentence starting from S can be represented as a tree.
- generation of sentences using a CFG is the reverse process.



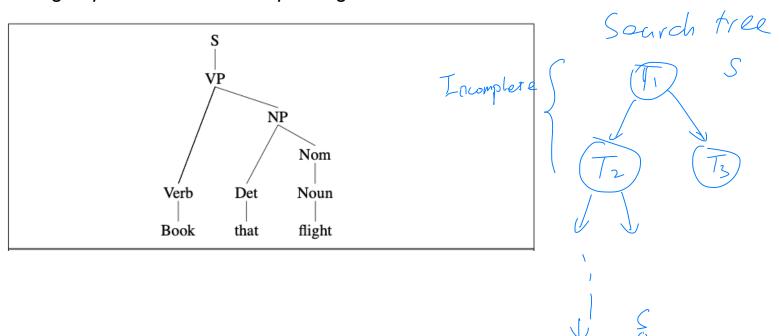
Syntactic parsing as search

Search algorithms are commonly used in computer science and Al.

Input: sentence "Book that flight"

Output: parsing trees with root = S and leaves = the sentence

- Top-down and buttom-up searches.
 - Note: not to go up and down a fixed parsing tree but in a search tree.

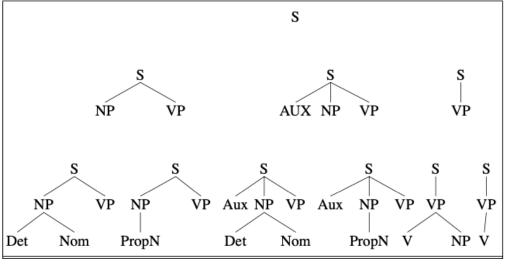


Syntactic parsing as search

Top-down search

- start from the root = S
- expand a non-terminal on an unfinished tree, using a rule in the given CFG.
- Pruning: when there is no hope of matching
- Stop with a match of the input sentence

```
S 	oup NP VP S 	oup Aux NP VP S 	oup VP S 	oup
```

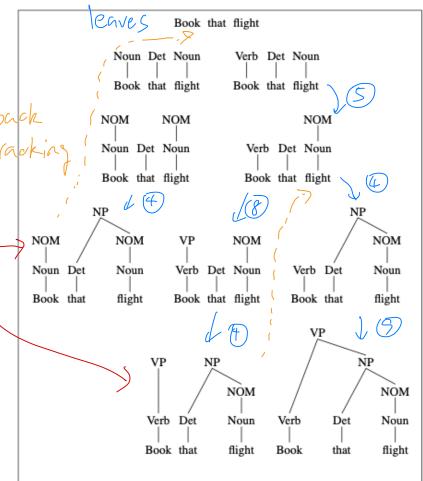


Syntactic parsing as search

Bottom-up search

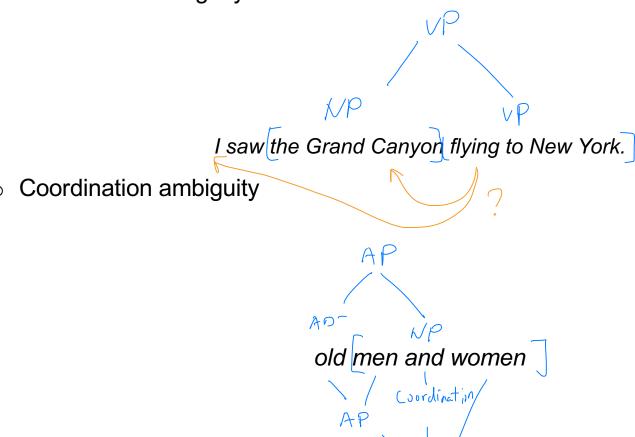
- start from the leaves (bottom) = input sentence,
- create a non-terminal from existing symbols, using a rule in the given CFG.
- Pruning: when can't find a match of any right-hand-side match
- Stop when arriving at root = S.

```
S 	oup NP VP
S 	oup Aux NP VP
S 	oup VP
S 	oup
```



Discussion of search-based methods

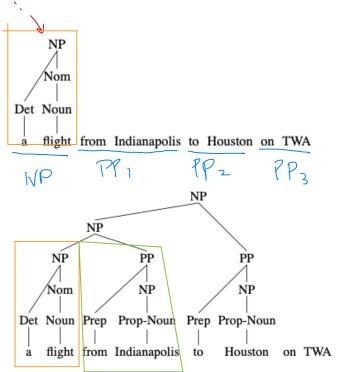
- Ambiguity and multiple "valid" parses.
 - Attachment ambiguity

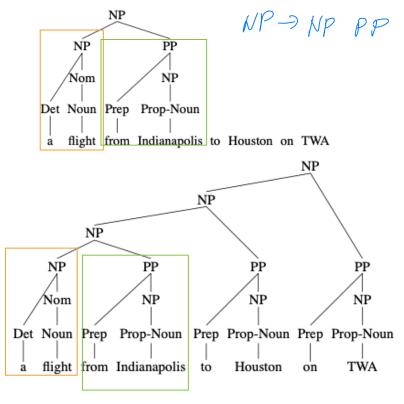


Discussion of search-based methods

Costly: repeatly parsing the same constituent.

o the PP and the attachment ambiguity lead to exponential time complexity $C(n) = \frac{1}{n+1} {2n \choose n}$





- It is a dynamic programming (DP) algorithm.
- Properties of DP:
 - Overlapping sub-problems: can re-use solutions to sub-problems;
 - Subproblems = parsing a sub-part of a sentence.
 - Optimal structures: optimal solutions => optimal sub-problem solutions.
 - successful parse of a constituent => successful parse of its parts.
- Reduce the cost of search by caching the solution to subproblems.
 - don't repeatly parsing "a flight".

Chomsky normal form (CNF)

- Allow simple a data structure for the DP algorithm called "CKY"
- Restrict CFG to the following form
 - There can be at most two symbols on the right hand side of any rule;
 - If there are two, both should be non-terminals (e.g. A→BC);
 - If there is one, it should be a terminal (e.g. $A\rightarrow a$).
- There is no loss of expressiveness by the following conversions:
 - o mixed production $A \rightarrow Bc \Rightarrow (A \rightarrow BC \text{ and } C \rightarrow c)$
- mit Production $A \rightarrow B \Rightarrow$ eliminate B by changing any $B \rightarrow \gamma$ to $A \rightarrow \gamma$.
 - long right hand side $A \rightarrow BCD \Rightarrow (A \rightarrow XD \text{ and } X \rightarrow BC)$;
 - can handle more than three symbols by recursively doing this conversion.

Chomsky normal form (CNF)

Conversion algorithm

- 1. handle mixed RHS
- 2. handle unit productions
- 3. handle long RHS

\mathcal{L}_1 Grammar	\mathcal{L}_1 in CNF
$S \rightarrow NP VP$	$S \rightarrow NP VP$
$S \rightarrow Aux NP VP$	$S \rightarrow XI VP$
	$XI \rightarrow Aux NP$
$S \rightarrow VP$	$S o book \mid include \mid prefer$
	$S \rightarrow Verb NP$
	$S \rightarrow X2 PP$
	$S \rightarrow Verb PP$
	$S \rightarrow VPPP$
$NP \rightarrow Pronoun$	$NP \rightarrow I \mid she \mid me$
$NP \rightarrow Proper-Noun$	$NP \rightarrow TWA \mid Houston$
$NP \rightarrow Det\ Nominal$	$NP \rightarrow Det Nominal$
$Nominal \rightarrow Noun$	$Nominal \rightarrow book \mid flight \mid meal \mid money$
$Nominal \rightarrow Nominal Noun$	$Nominal \rightarrow Nominal Noun$
$Nominal \rightarrow Nominal PP$	$Nominal \rightarrow Nominal PP$
$VP \rightarrow Verb$	$\mathit{VP} o \mathit{book} \mathit{include} \mathit{prefer}$
$VP \rightarrow Verb NP$	$VP \rightarrow Verb NP$
$VP \rightarrow Verb NP PP$	$VP \rightarrow X2 PP$
	$X2 \rightarrow Verb NP$
$VP \rightarrow Verb PP$	$VP \rightarrow Verb PP$
$VP \rightarrow VP PP$	$VP \rightarrow VP PP$
$PP \rightarrow Preposition NP$	$PP \rightarrow Preposition NP$

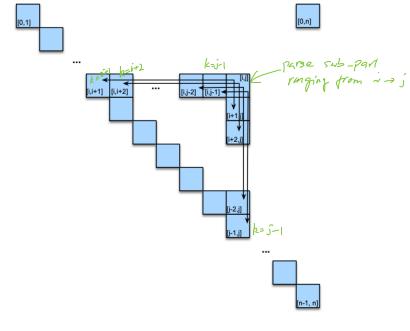
Use a matrix of size (n+1) x (n+1) if the sentence has n words.

- The cell (*i*, *j*) represents all possible non-terminals that can generate the constituent ranging *i* to *j*.
- Due to CNF, each constituent can be broken down into two sub-constituents:

$$\circ$$
 $(i, j) = (i, k) + (k), j)$

$$0 j > i + 1$$
 and $0 \le i < k < j \le n$

$$23 = 0$$
 $5 = 2,3,4,5$



```
function CKY-PARSE(words, grammar) returns table
 for i \leftarrow from j-2 down to 0 do
                                                                                     \vec{v} \vec{s} \vec{k} (0,2) = (0,1) + (1,2)
       for k \leftarrow i+1 to j-1 do
          for all \{A \mid A \rightarrow BC \in grammar \text{ and } B \in table[i,k] \text{ and } C \in table[k,j]\}
                table[i,j] \leftarrow table[i,j] \cup A
                                   verb > book }
                                                                             S,VP,X2
                                                                                        S,VP,X2
                                                 Det -> the.
```

