Projection and its Importance in Scientific Computing

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Additional reference materials:

[1] R.Barrett, M.Berry, T.F.Chan, J.Demmel, J.Donato, J. Dongarra, V. Eijkhout, R.Pozo, C.Romine, and H.Van der Vorst, *Templates for the Solution of Linear Systems: Building Blocks for Iterative Methods (2nd edition)* http://netlib2.cs.utk.edu/linalg/html templates/Templates.html

[2] Yousef Saad, Iterative methods for sparse linear systems (1st edition)

http://www-users.cs.umn.edu/~saad/books.html

Topics

as related to high-performance scientific computing

Projection in Scientific Computing

Sparse matrices, parallel implementations

PDEs, Numerical solution, Tools, etc.

Iterative Methods

Topics

on new architectures - multicore, GPUs (CUDA & OpenCL), MIC

Projection in Scientific Computing

Sparse matrices, parallel implementations

PDEs, Numerical solution, Tools, etc.

Iterative Methods

Outline

- Part I
 - Fundamentals
- Part II
 - Projection in Linear Algebra
- Part III
 - Projection in Functional Analysis (e.g. PDEs)
- HPC with Multicore and GPUs

Part I Fundamentals



Projection in Scientific Computing

[an example – in solvers for PDE discretizations]

Electronic structure calculations

Density functional theory

Many-body Schrödinger equation (exact but exponential scaling)

$$\{-\sum_{i}\frac{1}{2}\nabla_{i}^{2}+\sum_{i,j}\frac{1}{\mid r_{i}-r_{j}\mid}+\sum_{i,j}\frac{Z}{\mid r_{i}-R_{j}\mid}\}\Psi(r_{1},...r_{N})=E\Psi(r_{1},...r_{N})$$

- Nuclei fixed, generating external potential (system dependent, non-trivial)
- · N is number of electrons



Kohn Sham Equation: The many body problem of interacting electrons is reduced to non-interacting electrons (single particle problem) with the same electron density and a different effective potential (cubic scaling).

$$\begin{split} \{-\frac{1}{2}\nabla^{2} + \int & \frac{\rho(r')}{|r-r'|} dr' + \sum_{l} \frac{Z}{|r-R_{l}|} + V_{xc} \} \psi_{i}(r) = E_{i} \psi_{i}(r) \\ \rho(r) = & \sum |\psi_{i}(r)|^{2} = |\Psi(r_{1},...r_{N})|^{2} \end{split}$$

- V_{xc} represents effects of the Coulomb interactions between electrons
- · ρ is the density (of the original many-body system)

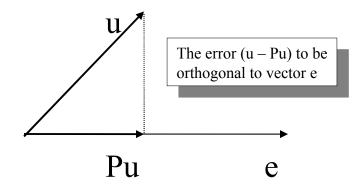
 V_{xc} is not known except special cases \Rightarrow use approximation, e.g. Local Density Approximation (LDA) where V_{xc} depends only on ρ

What is Projection?

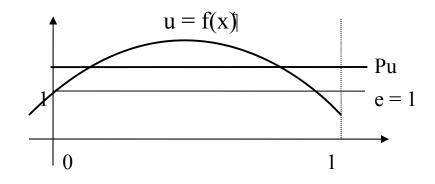
Here are two examples

(from linear algebra)

(from functional analysis)



P: orthogonal projection of vector u on e



P: best approximation (projection) of f(x)in span{ e } \subset C[0,1]

Definition

• **Projection** is a linear transformation P from a linear space V to itself such that

$$P^2 = P$$

equivalently

Let V is direct sum of subspaces V₁ and V₂

$$V = V_1 \oplus V_2$$

i.e. for $\forall u \in V$ there are unique $u_1 \in V_1$ and $u_2 \in V_2$ s.t.

$$u = u_1 + u_2$$

Then P: $V \rightarrow V_1$ is defined for $\forall u \in V$ as $Pu \equiv u_1$



Importance in Scientific Computing

• To compute approximations $Pu \approx u$ where $\dim V_1 \ll \dim V$

$$V = V_1 \oplus V_2$$

When computation directly in V is not feasible or even possible.

A few examples:

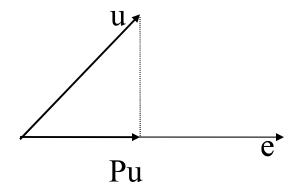
- Interpolation (via polynomial interpolation/projection)
- Image compression
- Sparse iterative linear solvers and eigensolvers
- Finite Element/Volume Method approximations
- Least-squares approximations, etc.

Projection in R²

• In R^2 with Euclidean inner-product, i.e. for $x, y \in R^2$

$$(\mathbf{x}, \mathbf{y}) = \mathbf{x}_1 \mathbf{y}_1 + \mathbf{x}_2 \mathbf{y}_2 \quad (= \mathbf{y}^T \mathbf{x} = \mathbf{x} \cdot \mathbf{y})$$

and $\|\mathbf{x}\| = (\mathbf{x}, \mathbf{x})^{1/2}$



P : orthogonal projection of vector u on e

$$Pu = \frac{(u, e)}{\parallel e \parallel^2} e \qquad (Exercise)$$

i.e. for
$$||e|| = 1$$

Pu = (u, e) e

Projection in Rⁿ/Cⁿ

Similarly to R²

P : Orthogonal projection of u into span $\{e_1, ..., e_m\}$, $m \le n$.

Let e_i , $i = 1 \dots m$ is **orthonormal** basis, i.e.

$$(e_i, e_j) = 0$$
 for $i \neq j$ and
 $(e_i, e_j) = 1$ for $i = j$
 $\mathbf{P} \mathbf{u} = (\mathbf{u}, \mathbf{e}_1) \mathbf{e}_1 + \dots + (\mathbf{u}, \mathbf{e}_m) \mathbf{e}_m$ (Exercise)

Orthogonal projection of u on e₁



How to get an orthonormal basis?

Can get one from every subspace by **Gram-Schmidt** orthogonalization:

```
<u>Input</u>: m linearly independent vectors x_1, ..., x_m
```

 \underline{Output} : m orthonormal vectors $x_1, ..., x_m$

1.
$$x_1 = x_1 / ||x_1||$$

2. do
$$i = 2, m$$

3.
$$x_i = x_i - (x_i, x_1) x_1 - \dots - (x_i, x_{i-1}) x_{i-1}$$
 (Exercise: $x_i \perp x_1, \dots, x_{i-1}$)
4. $x_i = x_i / ||x_i||$ Orthogonal projection of x_i on x_1

4.
$$X_i = X_i / ||X_i||$$
 Orthogonal projection of X_i on X_i

5. enddo

Known as **Classical Gram-Schmidt** (CGM) orthogonalization



How to get an orthonormal basis?

• What if we replace line 3 with the following (3')?

3.
$$x_{i} = (x_{i} - (x_{i}, x_{1}) x_{1}) \dots - (x_{i}, x_{i-1}) x_{i-1}$$

3'. $do j = 1, i-1$
 $x_{i} = x_{i} - (x_{i}, x_{j}) x_{j}$

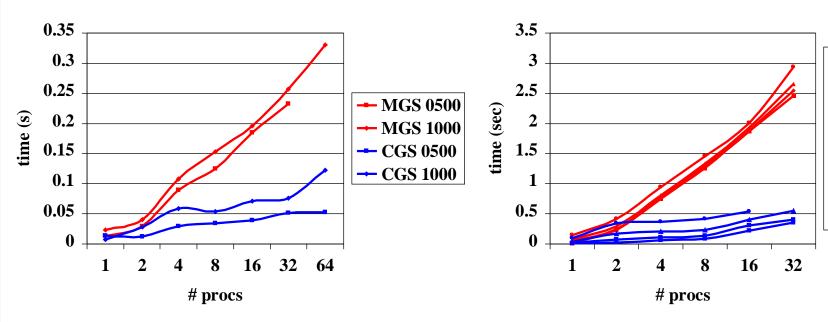
enddo

- Equivalent in exact arithmetic (Exercise) but not with round-off errors (next)!
- Known as Modified Gram-Schmidt (MGS) orthogonalization

CGS vs MGS

[Results from Julien Langou:]

Scalability of MGS and CGS on two different clusters for matrices of various size $m=[500\ 1000\ 2000\ 4000]$ per processor, n=100





- MGS 0500

- MGS 1000

MGS 2000

-- CGS 0500

-- CGS 1000

-- CGS 2000

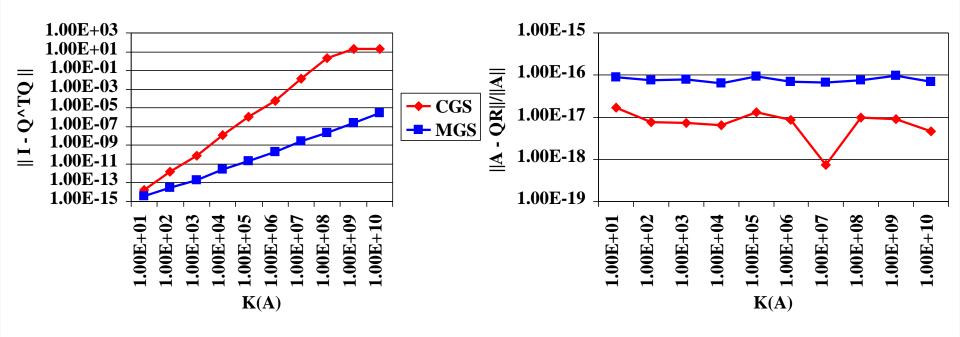
-- CGS 4000

MGS 4000

CGS vs MGS

[Results from Julien Langou:]

Accuracy of MGS vs CGS on matrices of increasing condition number



QR factorization

• Let $A = [x_1, ..., x_m]$ be the input for CGS/MGS and

$$Q = [q_1, ..., q_m]$$
 the output;

R: an upper $m \times m$ triangular matrix defined from the CGR/MGS. Then

$$A = Q R$$



Other QR factorizations

What about the following?

[known as Cholesky QR]

- 1. $G = A^T A$
- 2. $G = L L^T$ (Cholesky factorization)
- 3. $Q = A(L^T)^{-1}$
- Does Q have orthonormal columns (i.e. $Q^TQ=I$), i.e. $A = QL^T$ to be a QR factorization (Exercise)
- When is this feasible and how compares to CGS and MGS?

Other QR factorizations

Feasible when n >> m

P1

Allows efficient parallel implementation:

blocking both computation and communication

P2 ...

Investigate numerically accuracy and scalability (compare to CGS and MGS)

Exercise

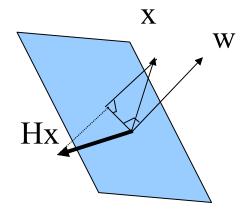
How is done in LAPACK?

Using Householder reflectors

$$H = I - 2 w w^T$$

• w = ? so that

$$H x_1 = \alpha e_1$$



 \Rightarrow w = ... (compute or look at the reference books)

Allows us to construct

LAPACK implementation: "delayed update" of the trailing matrix +

"accumulate transformation" to apply it as BLAS 3

Part II Projection in Linear Algebra



Projection into general basis

- How to define projection without orthogonalization of a basis?
 - Sometimes is not feasible to orthogonalize
 - Often the case in functional analysis
 (e.g. Finite Element Method, Finite Volume Method, etc.)
 where the basis is "linearly independent" functions (more later, and Lecture 2)
- We saw if $X = [x_1, ..., x_m]$ is an orthonormal basis

(*)
$$P u = (u, x_1) x_1 + ... + (u, x_m) x_m$$

• How does (*) change if X are just linearly independent?

$$Pu = ?$$



Projection into a general basis

• The problem:

Find the coefficients $C = (c_1 \dots c_m)^T$ in

$$Pu = c_1 x_1 + c_2 x_2 + \dots + c_m x_m = X C$$

so that

$$u - Pu \perp span\{x_1, ..., x_m\}$$

or \Leftrightarrow so that the error e in

$$\mathbf{u} = \mathbf{P}\mathbf{u} + \mathbf{e}$$

is \perp span $\{x_1, ..., x_m\}$

Projection into a general basis

• (1)
$$\mathbf{u} = \mathbf{P} \mathbf{u} + \mathbf{e} = \mathbf{c}_1 \mathbf{x}_1 + \mathbf{c}_2 \mathbf{x}_2 + \dots + \mathbf{c}_m \mathbf{x}_m + \mathbf{e}$$

• Multiply (1) on both sides by "*test*" vector/function \mathbf{x}_{j} (terminology from functional analysis) for j = 1,..., m

$$(\mathbf{u}, \mathbf{x}_{j}) = \mathbf{c}_{1}(\mathbf{x}_{1}, \mathbf{x}_{j}) + \mathbf{c}_{2}(\mathbf{x}_{2}, \mathbf{x}_{j}) + \dots + \mathbf{c}_{m}(\mathbf{x}_{m}, \mathbf{x}_{j}) + (\mathbf{e}, \mathbf{x}_{j})$$

- i.e., **m** equations for **m** unknowns
- In matrix notations \Leftrightarrow (X^T X) C = X^T u (Exercise) \Leftrightarrow X^T Pu = X^T u
- X^TX is the so called Gram matrix (nonsingular; why?) => there exists a unique solution C





Normal equations

System

$$(X^T X) C = X^T u$$

is known also as Normal Equations

• The Method of Normal Equations:

Finding the projection (approximation) Pu XC (approximation) of u in X by solving the Normal Equations system

Least Squares (LS)

Equivalently, system

Pu

$$(X^T X)^{\bullet}C = X^T u$$

gives also the solution of the LS problem

$$\min_{C \in R^m} \parallel X \; C - u \parallel$$

since
$$\|\mathbf{v}_1 - \mathbf{u}\|^2 = \|(\mathbf{v}_1 - \mathbf{P}\mathbf{u}) - \mathbf{e}\|^2 = \|\mathbf{v}_1 - \mathbf{P}\mathbf{u}\|^2 + \|\mathbf{e}\|^2$$

 $\geq \|\mathbf{e}\|^2 = \|\mathbf{P}\mathbf{u} - \mathbf{u}\|^2 \text{ for } \forall \mathbf{v}_1 \in \mathbf{V}_1$



LS

• Note that the usual notations for LS is: For $A \in \mathbb{R}^{n \times m}$, $b \in \mathbb{R}^n$ find

$$\min_{\mathbf{x} \in \mathbb{R}^m} \| \mathbf{A} \mathbf{x} - \mathbf{b} \|$$

Solving LS with QR factorization

Let
$$A = Q R$$
, $Q^{T}A = R = \begin{pmatrix} R_{1} \\ 0 \end{pmatrix}$, $Q^{T}b = \begin{pmatrix} c \\ d \end{pmatrix}$ m - m

Then

$$||Ax - b||^2 = ||Q^TAx - Q^Tb||^2 = ||R_1x - c||^2 + ||d||^2$$

i.e. we get minimum if x is such that

$$R_1 x = c$$

Projection and iterative solvers

• The problem : Solve $Ax = b in R^n$

Iterative solution: at iteration **i** extract an approximate \mathbf{x}_i from just a subspace $V = \text{span}[v_1, ..., v_m]$ of R^n

• How? As on slide 22, impose constraints: $b - Ax \perp subspace W = span[w_1,...,w_m] \text{ of } R^n, \text{ i.e.}$

(*)
$$(Ax, w_i) = (b, w_i)$$
 for $\forall w_i \in W = \text{span}[w_1, ..., w_m]$

- Conditions (*) known also as Petrov-Galerkin conditions
- Projection is *orthogonal*: V and W are the same (Galerkin conditions) or *oblique*: V and W are different

Matrix representation

• Let
$$V = [v_1, ..., v_m], W = [w_1, ..., w_m]$$

Find $y \in \mathbb{R}^m$ s.t. $\mathbf{x} = \mathbf{x_0} + \mathbf{V} \mathbf{y}$ solves $A\mathbf{x} = \mathbf{b}$, i.e. $\mathbf{A} \mathbf{V} \mathbf{y} = \mathbf{b} - A\mathbf{x_0} = \mathbf{r_0}$

subject to the orthogonality constraints:

$$\mathbf{W}^{\mathsf{T}} \mathbf{A} \mathbf{V} \mathbf{y} = \mathbf{W}^{\mathsf{T}} \mathbf{r}_{0}$$

• The choice for V and W is crucial and determines various methods (more in Lectures 13 and 14)



A General Projection Algorithm

Prototype from Y.Saad's book

- Until convergence, Do:
- 2. Select a pair of subspaces K and L
- 3. Choose bases $V = [v_1, \dots, v_m]$ and $W = [w_1, \dots, w_m]$ for K and L
- 4. r := b Ax
- 5. $y := (W^T A V)^{-1} W^T r$
- $6. \quad x := x + Vy$
- EndDo



Projection and Eigen-Solvers

• The problem : Solve $Ax = \lambda x \qquad \text{in } R^n$

- As in linear solvers: at iteration i extract an approximate
 x_i from a subspace V = span[v₁, ..., v_m] of Rⁿ
- How? As on slides 22 and 26, impose constraints: $\lambda x Ax \perp \text{subspace } W = \text{span}[w_1, ..., w_m] \text{ of } R^n, \text{ i.e.}$ (*) $(Ax, w_i) = (\lambda x, w_i) \text{ for } \forall w_i \in W = \text{span}[w_1, ..., w_m]$
- This procedure is known as **Rayleigh-Ritz**
- Again projection can be *orthogonal* or *oblique*

Matrix representation

• Let
$$V = [v_1, ..., v_m], W = [w_1, ..., w_m]$$

Find $y \in \mathbb{R}^m$ s.t. $\mathbf{x} = \mathbf{V} \mathbf{y}$ solves $A\mathbf{x} = \lambda \mathbf{x}$, i.e. $\mathbf{A} \mathbf{V} \mathbf{y} = \lambda \mathbf{V} \mathbf{y}$

subject to the orthogonality constraints:

$$\mathbf{W}^{\mathsf{T}} \mathbf{A} \mathbf{V} \mathbf{y} = \lambda \mathbf{W}^{\mathsf{T}} \mathbf{V} \mathbf{y}$$

• The choice for V and W is crucial and determines various methods (more in Lectures 4 and 5)

Part III Projection in PDEs

Projection in Functional Spaces

- The discussion so far can be applied to any functional innerproduct space (examples to follow)
- An important space is **C[a, b]**, the space of continuous functions on [a, b], with inner-product

$$(f, g) = \int_{a}^{b} f(x) g(x) dx$$

and induced norm

$$|| f || = (f, f)^{1/2}$$



Projection in Functional Spaces

- Projection P: $V \rightarrow V_1$ where $V = V_1 \oplus V_2$
- In functional analysis and scientific computing V_1 is usually taken as
 - Piecewise polynomials
 - In PDE approximation (FEM/FVM), Numerical integration, etc.
 - Trigonometric functions

$$\{ \sin(n x), \cos(n x) \}_{n=0}, x \in [0, 2\pi]$$

Orthogonal relative to

$$2\pi$$

$$(f, g) = \int_{0}^{\infty} f(x) g(x) dx$$
 (Exercise)

Normal equations / LS

• Exercise:

$$f(x) = \sin(x)$$

Find the projection in $V_1 = \text{span}\{x, x^3, x^5\}$ on interval [-1, 1] using inner-product

$$(f, g) = \int f(x) g(x) dx$$

and norm $|| f || = (f,f)^{1/2}$



Normal equations / LS

- Leads to Gram matrix that is very ill-conditioned (called Hilbert matrix: Gram matrix for polynomials 1, x, x², x³, ...)
- For numerical stability is better to orthogonalize the polynomials
- There are numerous examples of orthonormal polynomial sets
 - * Legendre, Chebyshev, Hermite, etc
 - * Check the literature for more if interested

Integration via Polynomial Interpolation

Take

$$\int f(x) dx \approx \int p(x) dx$$

where p is a polynomial approximation to f

• Taking p a polynomial interpolating f at n+1 fixed nodes x_i leads to quadrature formulas

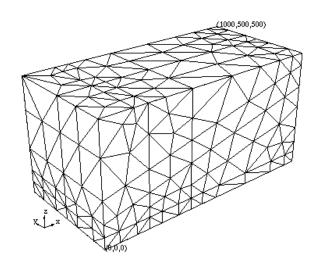
$$\int f(x) dx \approx A_0 f(x_0) + ... + A_n f(x_n)$$

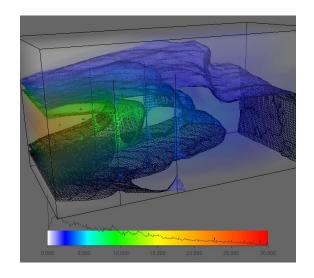
that are exact exact for polynomials of degree $\leq n$

• Smart choice of the nodes x_i (Gaussian quadrature) leads to formulas that are exact for polynomials of degree $\leq 2n+1$

Galerkin Projection

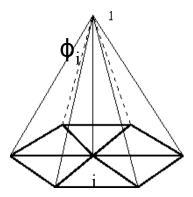
- Numerical PDE discretizations have a common concept:
 - Represent computational domain with mesh
 - Approximate functions and operators over the mesh





Galerkin Projection

• Finite dimensional spaces (e.g. V₁) can can be piecewise polynomials defined over the mesh, e.g.



- Numerical solution of PDE (e.g. FEM)
 - Boundary value problem: Au = f, subject to boundary conditions
 - Get a "weak" formulation: (Au, ϕ) = (f, ϕ) multiply by test function ϕ and integrate over the domain

a(u,
$$\phi$$
) = \phi> for ϕ S

- Galerkin (FEM) problem: Find u_h S_h S s.t. $a(u_h, \phi_h) = \langle f, \phi_h \rangle \text{ for } \phi_h S_h$



Learning Goals

- To refresh some linear algebra essentials that are of fundamental importance for scientific computing
- The idea and application of **Petrov-Galerkin conditions** as a way of defining computationally feasible formulations (approximations)
- Some generic examples demonstrating the ideas in
 - Linear algebra
 - Functional analysis
 (to get more specific in the following lectures)