

```
#####
#          cs315 Week 8
#
#  -> Floating point representation & addition & subtraction
#
#####
```

Floating-point numbers

- \* used to represent real values
- \* have finite precision

--> single-precision - 32 bits (1 word)

- \* 1 bit sign
- \* 8 bit exponent
- \* 23 bit magnitude

--> double-precision - 64 bits (2 words)

- \* 1 bit sign
- \* 11 bit exponent
- \* 52 bit magnitude

- ~ doubles require two registers or two memory locations
- ~ doubles MUST start in an even numbered register
- ~ doubles occupy both the even register and the following odd register

Example:

li.d \$f4, 3.14 stores a double in both \$f4 and \$f5

registers may be written with the 'concatenation' bar to denote they are part of a double

Example:

li.d \$f4, 3.14      # \$f4|\$f5 <-- 3.14

Floating-point standard (FPS) will be discussed more in lecture. See appendix F in MIPS book for floating-point instructions

--> NOTE: appendix F may have typos

FPS consist of 3 parts:

- 1) sign
- 2) exponent
- 3) magnitude

--> sign, exponent, and magnitude (which includes hidden bit)

FPS format (this is the format we use to represent floating-point numbers):

```
|-- sign --|-- magnitude --|-- (hidden bit) --|-- fraction--|
   ^         ^         ^         ^
   1 bit    unsigned binary  1 bit    unsigned binary
               (8 bits)                (7 bits)
```

We should note that:

```
|-- sign --|-- (hidden bit) --|-- fraction --|    <<<----<<< sign magnitude (9 bits)
   ^
   1 bit
```

Definition of mantissa:

Mantissa = (hidden bit) + fraction      <<<----<<< 8 bits mantissa in FPS representation (or simply class representation of floating-point numbers)

```

           ^         ^
           1 bit    7 bits
```

Sign bit:

- \* is only 1 bit
- \* 0 for positive, 1 for negative (similar to sign magnitude number system)

#### Exponent:

- \* 8 bits for 32 bit IEEE 754 standard
- \* 8 bits for the FPS we will use in class
- \* bias 127
  - > bias exponent = actual exponent + 127
  - > actual exponent = bias exponent - 127

##### Example:

if FPS exponent (bias exponent): 131, then actual exponent is: 4

- \* bias is used to allow negative exponents

##### Example:

if FPS exponent (bias exponent): 125, then actual exponent is: -2

#### Magnitude:

- \* 23 bits for 32 bit IEEE 754 standard
- \* 7 bits for the FPS we will use in class (which includes hidden bit)
- \* shortened to 7 for brevity, but the process used for these 7 is the same processed used for all 23
- \* hidden bit will always be a 1
- \* not represented in hardware ('hidden')
  - > less circuitry, therefore less material cost and less energy use
  - > we represent numbers in normal format:  
i.e. 1.001

^

this is hidden bit (always 1), we show it in parenthesis

- \* used in the arithmetic we will be doing
- \* show in very clear parenthesis

#### Normalizing:

- \* normal form: 1.-----

##### Example:

1.23456  
3.14151

#####

#### \* Precision

- > precision will be lost when using FPS
  - unavoidable trade-off for using a finite number of bits
- > lost when converting, lost when shifting values, etc.
  - converting -7.4 to 8-bit binary then back to decimal demonstrates loss of precision
  - >>> -7.325 is the closest we get to -7.4

#### \* FPS addition of same sign

- 1) convert both values to binary
- 2) normalize both values
- 3) match lower exponent to higher exponent (raise exponent n, right-shift n bits)
- 4) add magnitudes, keep overflow carry bit
- 5) normalize result if necessary
- 6) set result sign to same sign as operands

#### \* FPS addition of different signs

- 1) convert both values to binary
- 2) normalize both values
- 3) match lower exponent to higher exponent
- 4) convert sign to 10th bit (+ = 0, - = 1, is used to handle signed overflow)
- 5) convert both values to 2's complement
- 6) add magnitudes (as 2's complement)
- 7) convert result to sign-magnitude
- 8) convert 10th bit to sign
- 9) normalize if necessary

Example 1:

$$\begin{array}{r} 3.125 \\ + 2.500 \\ \hline 5.625 \text{ (expected)} \end{array}$$

3 => 11

|       |   |
|-------|---|
| 0.125 |   |
| 0.250 | 0 |
| 0.500 | 0 |
| 0.000 | 1 |

001  $\Rightarrow$  .625

$$3.125 \Rightarrow 11.001 \underset{\substack{\uparrow \\ \text{normalize}}}{\Rightarrow} 1.1001 * 2^1$$

$0 - 100000000$  (1)  $1001000 \leq 3.125$   
 $\uparrow \quad \quad \quad \uparrow \quad \quad \quad \uparrow \quad \quad \quad \uparrow$   
 sign exponent hidden bit (we show it in parenthesis) fraction

$$2 \Rightarrow 10$$
$$\begin{array}{r} 0.500| \\ 0.000| \quad 1 \\ \hline 1 \Rightarrow 0.50 \end{array}$$
$$2.500 \Rightarrow 10.1 \Rightarrow 1.01 * 2^1$$

^  
normalize

$$\begin{array}{ccccccc} 0 & - & 10000000 & (1) & 0100000 & \leq & 2.500 \\ \wedge & & \wedge & \wedge & \wedge & & \\ \text{sign} & & \text{exponent} & & \text{fraction} & & \\ & & & \wedge & & & \\ & & & \text{hidden bit (we show it in parenthesis)} & & & \end{array}$$

Essentially:

$$1.1001 * 2^1 + 1.01 * 2^1 = 2^1 * (1.1001 + 1.01)$$

exponents are the same, so there is no need to match exponents

```

      3.125 => 0 - 100000000 (1) 1001000
+     2.500 => 0 - 100000000 (1) 0100000
-----

```

0 - 10000001 (1) 0110100  $\Rightarrow 1.0110100 * 2^2 = 101.10100 * 2^0 = 101.10100 \Rightarrow + 5.625$  (correct result)

sign extend to 10 bits when two numbers are the same sign (because magnitude is will get larger when we add two numbers with the same sign):  
--> so more bits might be needed to represent the result of addition

$$\begin{array}{r}
 00 \ (1) \ 1001000 \\
 + \ 00 \ (1) \ 0100000 \\
 \hline
 \end{array}$$

0 1 (0) 1101000 -> (1) 0110100 <<< normalize by shift right 1 time which is essentially divide by 2 <<< thus we need to compensate division by 2, with adding exponent by 1  
 ^  
 sign of the result is positive

#####

Example 2:

```

  3.125
-  2.500
-----
  0.625 (expected)

```

```

  3.125 => 0 - 10000000 (1) 1001000
-  2.500 => 0 - 10000000 (1) 0100000
-----

```

sign extend to 10 bits when two numbers are different sign (we could sign extend it to 9 bits because we are subtracting two numbers and magnitude is getting smaller, but lets be consistent):

Essentially:

$$1.1001 * 2^1 - 1.01 * 2^1 = 2^1 * (1.1001 - 1.01)$$

exponents are the same, so there is no need to match exponents

```

  00 (1) 1001000  rewrite the number  -> 00 (1) 1001000
-  00 (1) 0100000  find additive inverse -> 11 (0) 1100000
-----

```

```

  00 (1) 0100000  <-----|
  11 (0) 1011111  |
+                  |
+                  1
-----
  11 (0) 1100000

```

sign of the result is positive which was expected (however, it is not in normal form)

0 - 01111110 (1) 0100000 => 1.0100000 \* 2<sup>-1</sup> = 0.1010000 => + 0.625 (correct result)

1 (or biased exponent of: 128) - 2 = -1 <-- 01111110 <<< biased exponent of 126

#####

Example 3:

```

  (-3.125)
+  2.500
-----
 -0.625 (expected)

```

```

  (-3.125) => 1 - 10000000 (1) 1001000
+  2.500    => 0 - 10000000 (1) 0100000
-----

```

sign extend to 10 bits when two numbers are different sign (we could sign extend it to 9 bits because we are subtracting two numbers and magnitude is getting smaller, but lets be consistent):

Essentially:

$$-1.1001 * 2^1 + 1.01 * 2^1 = 2^1 * (-1.1001 + 1.01)$$

exponents are the same, so there is no need to match exponents

```

|-----|
|  1 (1) 1001000  find additive inverse  -> 11 (0) 0111000
|+  0 (1) 0100000  rewrite the number    -> 00 (1) 0100000
|-----|
|  ^
|  sign

```

11 (1) 1011000 -> number is negative here, convert it back to sign magnitude (|- sign |- magnitu  
sign of the result is negative

```

|
|----> 10 (1) 1001000 <-- sign magnitude number
        11 (0) 0110111
        +
        -----
        1
        11 (0) 0111000

```

11 (1) 1011000 >-- convert result back to sign magnitude (positive) --> 00 (0) 0101000 -> (1) 0100000 <<< normalized by shift left 2 times, which is essentially multiply by 2, thus we need to compensate multiplication by 2 with subtract exponent with 2

1 - 01111110 (1) 0100000 => 1.0100000 \* 2<sup>-1</sup> = 0.1010000 = - 0.625 (correct result)

1 (or biased exponent of: 128) - 2 = -1 <-- 01111110 <<< biased exponent of 126

#####  
Example 4:

27 => 11011

```

0.125|
0.250| 0
0.500| 0
0.000| 1
-----
001 => .625

```

27.125 => 11011.001 => 1.1011001 \* 2<sup>4</sup> (biased exponent = 4 + 127 = 131 <<< 10000011)

$\uparrow$                        $\uparrow$   
 normalize              fraction

7 => 111

```

0.500|
0.000| 1
-----
1 => 0.500

```

7.5 => 111.1 => 1.111 \* 2<sup>2</sup> (biased exponent = 2 + 127 = 129 <<< 10000001)

$\uparrow$                        $\uparrow$   
 normalize               $\uparrow$   
                             fraction

```

+ 27.125  FPS = 0 - 10000011 (1) 1011001
   7.50   FPS = 0 - 10000001 (1) 1110000
-----
34.625 (expected)

```

```

+ 0 - 10000011 (1) 1011001  rewrite FPS  --> 0 - 10000011 (1) 1011001
+ 0 - 10000001 (1) 1110000  match exponents --> 0 - 10000011 (0) 0111100
-----
                                0 - 10000100 (1) 0001010 => 1.0001010 * 2^5
                                                                = 100010.10 * 2^0
                                                                = 100010.10
                                                                = +34.5 (correct result)

```

10000001 => 129 + 2 (to match) <----|

smaller exponent, we need to match it with larger exponent. We matched the exponent, but we also need to compensate for adding exponent with 2 thus, shift right 2 times (divide by 2)

add two positive numbers, without exponents

```
00 (1) 1011001
+ 00 (0) 0111100
-----
```

```
01 (0) 0010101 --> (1) 0001010
^
```

sign of result is positive

#####

Example 5:

exponents are matched <-|

```
29.877 FPS = 0 - 10000011 (1) 1101111
- 23.62 FPS = 0 - 10000011 (1) 0111100
-----
```

6.257 (expected)

```
-----|
| 0 - 10000011 (1) 1101111
- 0 - 10000011 (1) 0111100
-----
| 0 - 10000001 (1) 1001100 => 1.1001100 * 2^2
|                               = 110.01100 * 2^0
|                               = 110.01100
--> 131 (biased)                = 110.01100
so true exponent:                = +6.375 (correct result)
* 131 - 127 = 4
```

|---- add a positive numbers with additive inverse of second number {a - b = a + (-b)}, without exponents

```
|-> (1) 1101111 rewrite the number (convert from sign magnitude to 2's complement) --> 00 (1) 1101111
+ (1) 0111100 calculate additive inverse --> 11 (0) 1000100
-----
```

```
00 (0) 0110011 --> (1) 1001100
^
```

sign of result is positive

we shifted left 2 times to  
normalize the result, so we  
need to compensate it by  
subtracting exponent with 2  
thus, result exponent:  
131 - 2 = 129  
^  
which is: 10000001

#####

Example 6:

```
(- 63.874) FPS = 1 - 10000100 (1) 1111111
- (152.69) FPS = 0 - 10000110 (1) 0011000
-----
-216.564 (expected)
```

remember, to match exponent, we <-----|  
match smaller one with larger one |  
and don't forget shifting to |  
compensate for adjusting exponent |

```
1 - 10000100 (1) 1111111 match exponents --> 1 - 10000110 (0) 0111111
- 0 - 10000110 (1) 0011000 rewrite FPS --> 0 - 10000110 (1) 0011000
-----
```

-> then find additive inverse of both numbers, add them up two number:

```

+   (0) 0111111 -->   11 (1) 1000001
+   (1) 0011000 --> +  11 (0) 1101000
-----
(1) 1010111           11 (0) 0101001 --> -215
                        ^

```

sign of the result should be negative

now find the unsigned representation of the result:

```

11 (0) 0101001 <-- 2's complement
00 (1) 1010111 <-- unsigned binary <<< normal form already, so we don't do anything with it

1 - 10000110 (1) 1010111 => - 1.1010111 * 2^7 = - 11010111 * 2^0 = - 11010111 = -215 (correct result)

```

#####  
Example 7:

```

-   5153.475   FPS = 0 - 10001011 (1) 0100001
-   49.875     FPS = 0 - 10000100 (1) 1000111
-----
5103.6 (expected)

```

```

-   0 - 10001011 (1) 0100001   rewrite FPS   --> 0 - 10001011 (1) 0100001
-   0 - 10000100 (1) 1000111   match exponents --> 0 - 10001011 (0) 0000001
-----

```

--> find additive inverse of second number

```

-   (1) 0100001 -->   00 (1) 0100001
-   (0) 0000001 --> +  11 (1) 1111111
-----
                        00 (1) 0100000 --> (1) 0100000, therefore:--> 0 - 10001011 (1) 0100000
                        ^

```

sign of the result should be positive, also number is already in normal form, so we don't do anything with it

```

0 - 10001011 (1) 0100000 => + 1.0100000 * 2^12 = 1010000000000 * 2^0 = + 1010000000000 = +5120 (correct result)

```

#####