MergeSort + Heap

MergeSort definition

Merge Sort is a Divide and Conquer algorithm. It divides input array in two halves, calls itself for the two halves and then merges the two sorted halves. The merge() function is used for merging two halves. The merge(arr, l, m, r) is key process that assumes that arr[l.m] and arr[m+1.r] are sorted and merges the two sorted sub-arrays into one.

- MergeSort is a "stable" sort.
- $O(n imes \log n)$ worst-case, best-case and average-case

Stable sort

Stable sorting algorithms maintain the relative order of records with equal keys (i.e. values). That is, a sorting algorithm is stable if whenever there are two records R and S with the same key and with R appearing before S in the original list, R will appear before S in the sorted list

Basically sorting if two object are equal in a sense of compareTo() method, they remain in the same order as original list even after the sort is done.

Pseudocode

```
def mergeSort(arr):
   if len(arr) > 1:
        mid = len(arr) // 2 # Finding the mid of the array (integer division)
       L = arr[:mid]
                      # Dividing the array elements
        R = arr[mid:]
                             # into 2 halves
        mergeSort(L)
                      # Sorting the first half
                             # Sorting the second half
        mergeSort(R)
        i = j = k = 0
        while i < len(L) and j < len(R):
            if L[i] < R[j]: # If left array is less than right array then use left array</pre>
                arr[k] = L[i]
                             # Don't forget to increment index of left array index
                i+=1
            else:
                arr[k] = R[j] # If right array is less than right array then use right array
                             # Don't forget to increment index of right array index
               j += 1
                             # Regardless, increment main array's index
            k+=1
        while i < len(L):</pre>
                             # Checking if any element is remaining in left array
            arr[k] = L[i]
            i += 1
                             # Increment left array's index
            k+=1
                             # Increment main array's index
        while j < len(R):</pre>
                             # Checking if any element is remaining in right array
            arr[k] = R[j]
                              # Increment right array's index
            j+=1
            k+=1
                             # Increment main array
```

Exercise

Let's try to sort [9, 8, 7, 6, 5, 4, 3, 2, 1] using MergeSort

Solution

```
5, 4, 3, 2, 1]
          7, 6,
                                              // initial state
                   [5, 4,
                                              // split +1
                           3, 2, 1]
         7, 6]
                  [5, 4]
                            [3, 2, 1]
         [7, 6]
                                              // split +2
         [7] [6]
                   [5] [4]
                            [3] [2,
                                              // split +3
                    [5]
                                              // split +4
          [7] [6]
                             [3]
    [8]
    [8]
             [6]
                   [5] [4]
                             [3]
                                [1,
                                              // merge -4
[9]
                                2,
                   [4,
[8,
         [6,
                             [1,
                                     3]
                                              // merge -3
             7]
                        5]
                             3, 4,
                                     5]
[6,
                    [1, 2,
                                              // merge -2
[1,
                        6,
                                 8,
                                               // merge -1
```

Heap

A binary heap is a binary tree where:

- "MinHeap" smallest value is always at the top
- "MaxHeap" smallest value is always at the top

Heap Construction

Most efficient implementation of heap uses an array where:

- given index i, to get:
 - \circ left child index: 2 imes i+1
 - \circ right child index: 2 imes i+2
- parent index: $\lfloor \frac{i-1}{2} \rfloor$

Exercise

Represent this tree "MaxHeap" into an array of size $10\,$



Solution

[20, 13, 9, 8, 5, 3, 7, 6, 2, 1]

Pseudocode for "MaxHeap"

```
# To heapify subtree rooted at index i.
# n is size of heap
def heapify(arr, n, i):
    largest = i  # Initialize largest as root
    l = 2 * i + 1 # left = 2*i + 1
    r = 2 * i + 2 # right = 2*i + 2
    # See if left child of root exists and is
    # greater than root
    if l < n and arr[i] < arr[l]:</pre>
        largest = l
    # See if right child of root exists and is
    # greater than root
    if r < n and arr[largest] < arr[r]:</pre>
        largest = r
    # Change root, if needed
    if largest != i:
        # Swap
        swap(arr, i, largest)
        # Heapify the root.
        heapify(arr, n, largest)
```

Pseudocode for "HeapSort"

```
# The main function to sort an array of given size
def heapSort(arr):
    n = len(arr)

# Build a maxheap
    for i in range(n, -1, -1):  # start from n, go until -1 and increment by -1
        heapify(arr, n, i)

# One by one extract elements
    for i in range(n-1, 0, -1):  # start from n-1, go until 0 and increment by -1
        # Swap
        swap(arr, i, 0)
        heapify(arr, i, 0)
```

Lab Exercise

MergeSort