

Midterm - Program Analysis

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1.1)

Forward analysis:

$$E = \text{init}(S_*) = \{1\}$$

Backward analysis:

$$E = \text{final}(S_*) = \{7\}$$

1.2)

Very busy expressions

$$\iota = \emptyset$$

Reaching definitions

$$\iota = \{(x, ?) | x \in FV(S_*)\} = \{(a, ?), (b, ?), (x, ?), (y, ?)\}$$

Or in Scala:

```
val tuples: [(String, Long)] = Utils.vars(stmt).map((_, ?))
```

1.3)

Very busy expressions

$$\perp = AExp = \{x > 0, a * a, y + b, x - 1\}$$

Or in Scala:

```
val exprs: Set[Expression] = Utils.aexpr(stmt)
```

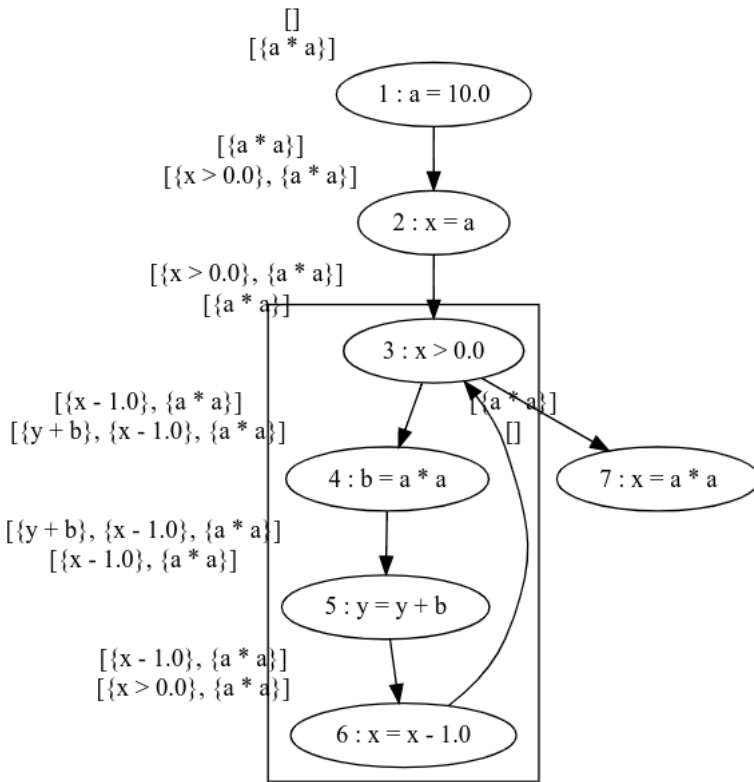
Reaching definitions

$$\perp = \emptyset$$

1.4)

Very busy expression is a **must** analysis and **backward**. So we need to use a backward CFG and **intersection** lattice.

- We need to start from the last
- Go backwards and if there are two paths going up, then we need to use intersection to join the paths
- We can change the direction of arrows to be backward
 - To get to label 3 we need to **intersect** the results 6 and 7



1.5)
1.6)
Definitions:

$$Kill(B) = \{ X \text{ op } Y \mid \text{either } X \text{ or } Y \text{ defined before use of } x \text{ op } Y \text{ in } B \}$$

$$Gen(B) = \{ x \text{ op } Y \mid X \text{ op } Y \text{ used in } B \text{ before any definition of } X \text{ or } Y \}$$

Transfer Equation

$$In(B) = (Out(B) - Kill(B)) \cup Gen(B)$$

Confluence Equation

$$Out(B) = \cap In(S) \text{ for all } S \in Succ(B)$$

Formula for GEN:

$[x := a] \text{ then } AExp(a)$
 $[b] = AExp(b)$

$GEN(1) = \{ \}$
 $GEN(2) = \{ a * a \}$
 $GEN(3) = \{ x - 1, x > 0 \}$
 $GEN(4) = \{ a * a \}$
 $GEN(5) = \{ y + b \}$
 $GEN(6) = \{ x - 1, x > 0 \}$
 $GEN(7) = \{ a * a \}$

Formula for KILL:

$[x := a] \text{ then } \{ a' \setminus in AExp \mid x \setminus in FV(a') \}$
 $[b] = \emptyset$

```

KILL(1) = { a*a }
KILL(2) = { x-1, x>0 }
KILL(3) = { }
KILL(4) = { y+b }
KILL(5) = { y+b }
KILL(6) = { x-1, x>0 }
KILL(7) = { x-1, x>0 }

```

```

OUT(1) = { a*a }
OUT(2) = { a*a, x>0 }
OUT(3) = { a*a }
OUT(4) = { a*a, x-1, y+b }
OUT(5) = { a*a, x-1 }
OUT(6) = { a*a, x>0 }
OUT(7) = { }

```

```

IN(1) = OUT(1) - KILL(1) + GEN(1) = { }
IN(2) = OUT(2) - KILL(2) + GEN(2) = { a*a }
IN(3) = OUT(3) - KILL(3) + GEN(3) = { a*a, x>0 }
IN(4) = OUT(4) - KILL(4) + GEN(4) = { a*a, x-1 }
IN(5) = OUT(5) - KILL(5) + GEN(5) = { a*a, x-1, y+b }
IN(6) = OUT(6) - KILL(6) + GEN(6) = { a*a, x-1 }
IN(7) = OUT(7) - KILL(7) + GEN(7) = { a*a }

```

Basically `a*a` is always available and it could be replaced everywhere it is used.

Implementation

```

case class AExp(exps: Set[Expression]) extends Lattice[AExp] {
  override def lub(that: AExp): AExp = AExp(exps intersect (that.exps))
}

case class VB(stmt: Statement) extends Analysis[AExp] {
  override val cfg: CFG = BackwardCFG(stmt)
  override val extremalValue: AExp = AExp(Set())
  override val bottom: AExp = AExp(Set() ++ Util.aexp(stmt))
  override val entry: mutable.Map[Node, AExp] = real_exit
  override val exit: mutable.Map[Node, AExp] = real_entry

  override def transfer(stmt: Statement, l: AExp): AExp = {
    def kill_gen(y: String, e: Expression) = {
      AExp((l.exps).filter(!Util.fv(_).contains(y)) ++ Util.aexp(e))
    }

    def gen(e: Expression) = AExp(l.exps ++ Util.aexp(e))

    stmt match {
      case ExprStmt(AssignExpr(_, LVarRef(name), e)) => kill_gen(name, e)
      case ExprStmt(expr) => gen(expr)
      case VarDeclStmt(IntroduceVar(y), expr) => expr match {
        case EmptyExpr() => l
        case _ => kill_gen(y, expr)
      }
      case IfStmt(cond, _, _) => gen(cond)
      case WhileStmt(cond, _) => gen(cond)
      case _ => l // no change!
    }
  }
}

```

2.1)

Constant Propagation is a **forward analysis**.

Forward analysis:

$$E = \text{init}(S_*) = \{1\}$$

Backward analysis:

$$E = final(S_*) = \{10\}$$

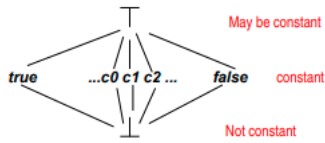
2.2)

For constant propagation extremal value

$$\iota_{CP} = \lambda x. \top$$

The least element (opposite of extremal value)

$$\perp$$



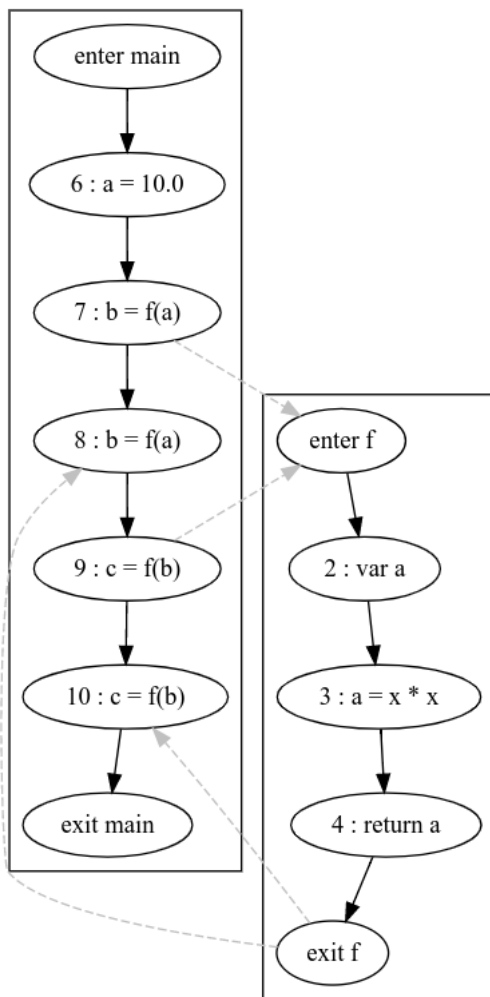
Basically the idea is the worst conclusion of running constant propagation analysis could be *everything may be a constant*.

2.3)

Each call site's `stmt.id` can be a context: $\{7, 9\}$

2.4)

P.S. I copied this from my solution of Homework 5 (context sensitive uninitialized variable). The IDs may not completely match the IDs in the question.



2.5)

$$CP_o(1) = \{[7] \mapsto \{a \mapsto 10, x \mapsto 10\}, [9] \mapsto \{a \mapsto 10, x \mapsto 100\}\}$$

$$CP_{\bullet}(1) = \{[7] \mapsto \{a \mapsto 10, x \mapsto 10\}, [9] \mapsto \{a \mapsto 10, x \mapsto 100\}\}$$

$$CP_{\circ}(2) = CP_{\bullet}(1)$$

$$CP_{\bullet}(2) = \hat{f}_2(CP_{\circ}(2)) = \{[7] \mapsto \{a \mapsto \top, x \mapsto 10\}, [9] \mapsto \{a \mapsto \top, x \mapsto 100\}\}$$

$$CP_{\circ}(3) = CP_{\bullet}(2)$$

$$CP_{\bullet}(3) = \hat{f}_3(CP_{\circ}(3)) = \{[7] \mapsto \{a \mapsto 10 \hat{*} 10, x \mapsto 10\}, [9] \mapsto \{a \mapsto 100 \hat{*} 100, x \mapsto 100\}\} = \{[7] \mapsto \{a \mapsto 100, x \mapsto 10\}, [9] \mapsto \{a \mapsto 100, x \mapsto 100\}\}$$

$$CP_{\circ}(4) = CP_{\bullet}(3)$$

$$CP_{\bullet}(4) = \hat{f}_4(CP_{\circ}(4)) = \{[7] \mapsto \{a \mapsto 100, x \mapsto 10, _r \mapsto 100\}, [9] \mapsto \{a \mapsto 10000, x \mapsto 100, _r \mapsto 10000\}\}$$

$$CP_{\circ}(5) = CP_{\bullet}(4)$$

$$CP_{\bullet}(5) = \hat{f}_5(CP_{\circ}(5)) = CP_{\circ}(5)$$

$$CP_{\circ}(6) = CP_{\bullet}(5)$$

$$CP_{\bullet}(6) = \hat{f}_6(CP_{\circ}(6)) = \{[7] \mapsto \{a \mapsto 10, x \mapsto 10, _r \mapsto 100\}, [9] \mapsto \{a \mapsto 10, x \mapsto 100, _r \mapsto 10000\}\}$$

$$CP_{\circ}(7) = CP_{\bullet}(6)$$

$$CP_{\bullet}(7) = \hat{f}_7(CP_{\circ}(7)) = \{[7] \mapsto \{a \mapsto 10, x \mapsto 10, _r \mapsto 100\}\}$$

$$CP_{\circ}(8) = CP_{\bullet}(7)$$

$$CP_{\bullet}(8) = \hat{f}_8(CP_{\bullet}(4), CP_{\circ}(8)) = \{[7] \mapsto \{a \mapsto 100, b \mapsto 100, x \mapsto 10, _r \mapsto 100\}\}$$

$$CP_{\circ}(9) = CP_{\bullet}(8)$$

$$CP_{\bullet}(9) = \hat{f}_9(CP_{\circ}(9)) = \{[7] \mapsto \{a \mapsto 100, b \mapsto 100, x \mapsto 100, _r \mapsto 100\}\}$$

$$CP_{\circ}(10) = CP_{\bullet}(9)$$

$$CP_{\bullet}(10) = \hat{f}_{10}(CP_{\bullet}(4), CP_{\circ}(9)) = \{[9] \mapsto \{a \mapsto 10000, c \mapsto 100, x \mapsto 100, _r \mapsto 10000\}\}$$

The final result or $_r$ is $10000 = 10 * 10 * 10 * 10 = 100 * 100$ which makes sense. Therefore our constant propagation run successfully.