Amirkabir University of Technology
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Processes (فرآیندها)

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Definition

* مهمترین کار OS ← اجرای برنامه‌های کاربر

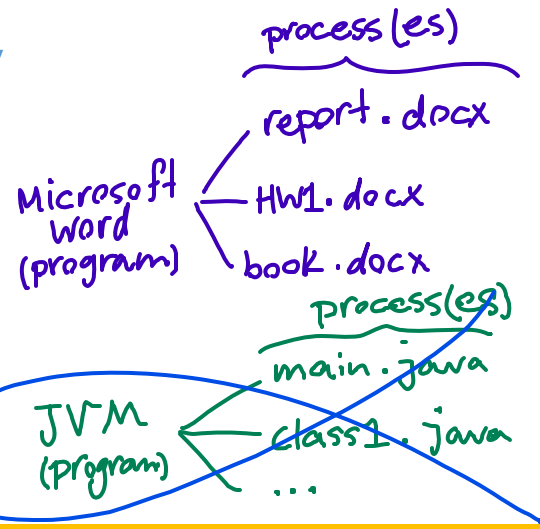
- **Process**
  - A program in execution; process execution must progress in sequential fashion
    - In time-sharing sys: unit of work → الزاما Process عدم نشده
  - All processes are executed concurrently

Parallel → واقعا موازی

concurrently → یکم از یکی یکم از یکی دیگه

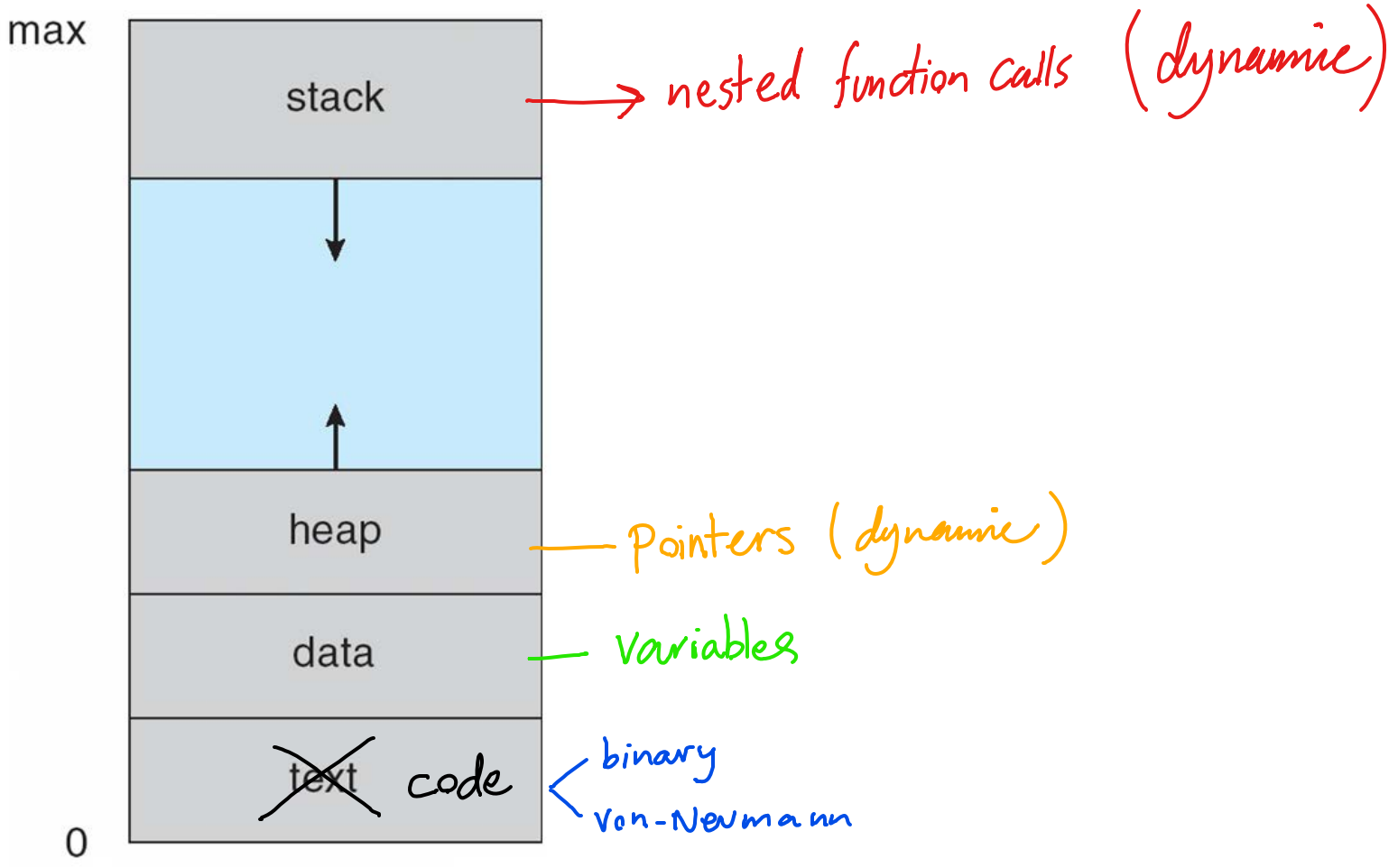
- **Process vs. Job?**
  - Passive: program
  - Active: process
    - Program becomes process when executable file loaded into memory
    - One program can be several processes
  - Question?
    - java program

program , process رابطه : one-to-many



کد بایت خروجی اجرا شده ← runnable است

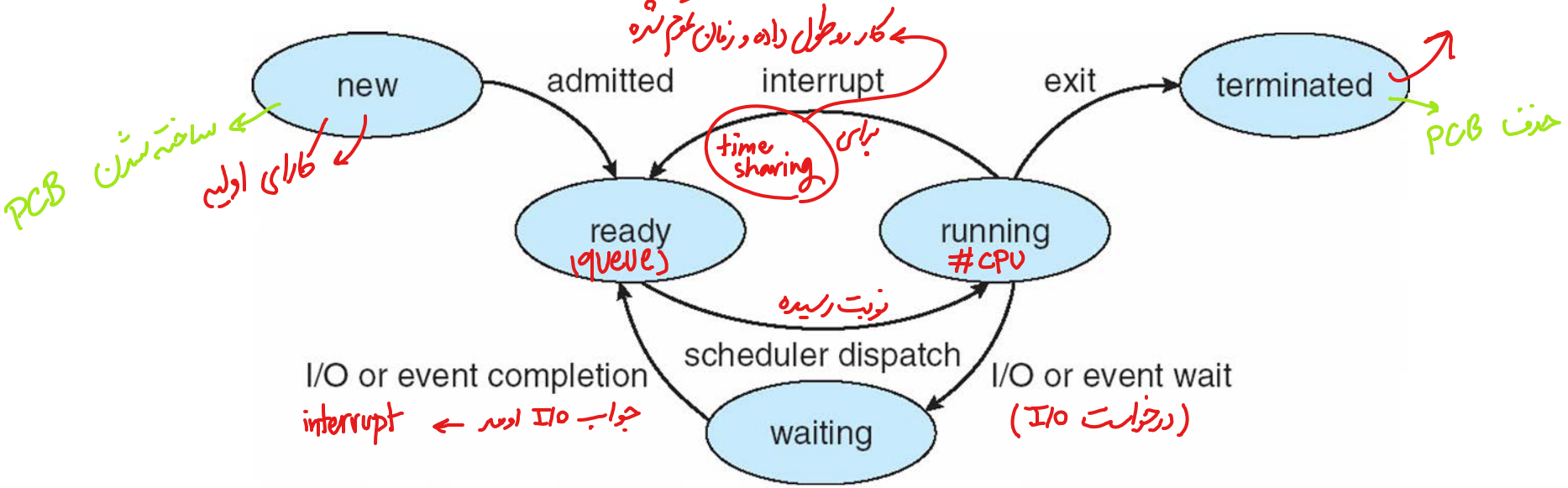
Process in memory



توی Process های مختلف از یک program قسمت text و data و heap با هم مساوی هستند و stack با هم متفاوت هستند.

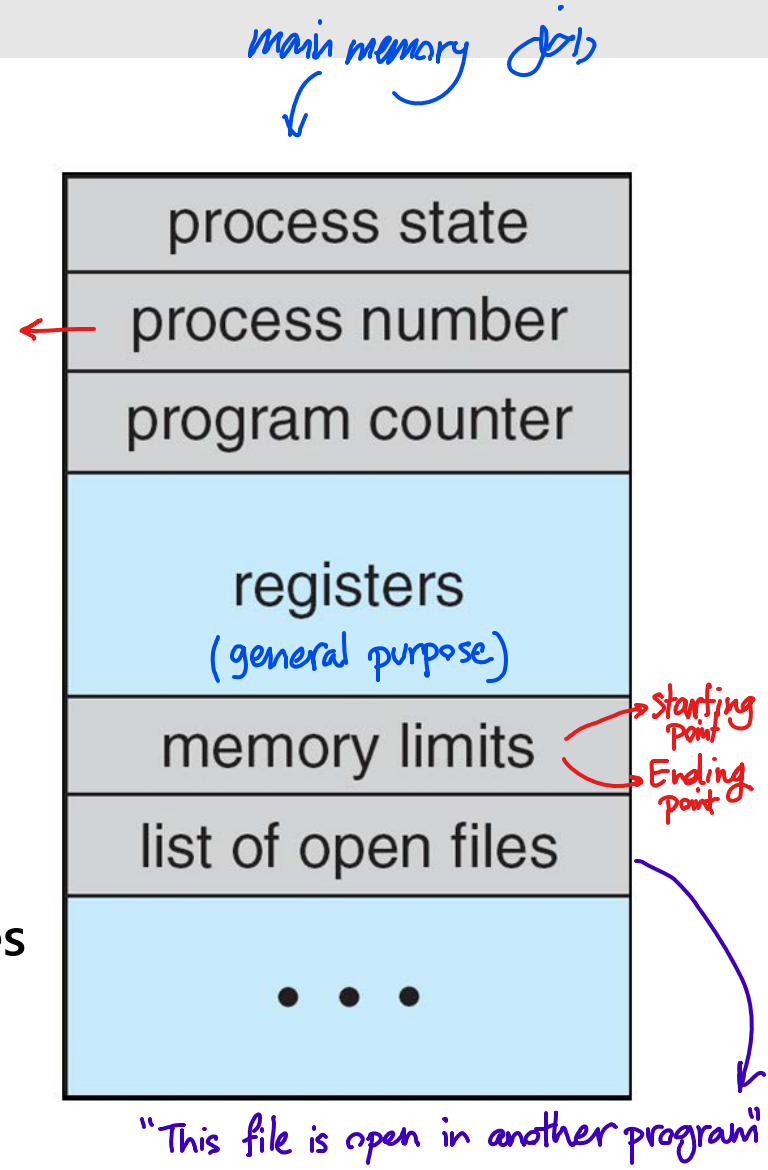
Process state

- As a process executes, it changes **state**
  - **new**: The process is being created
  - **running**: Instructions are being executed
  - **waiting**: The process is waiting for some event to occur
  - **ready**: The process is waiting to be assigned to a processor
  - **terminated**: The process has finished execution

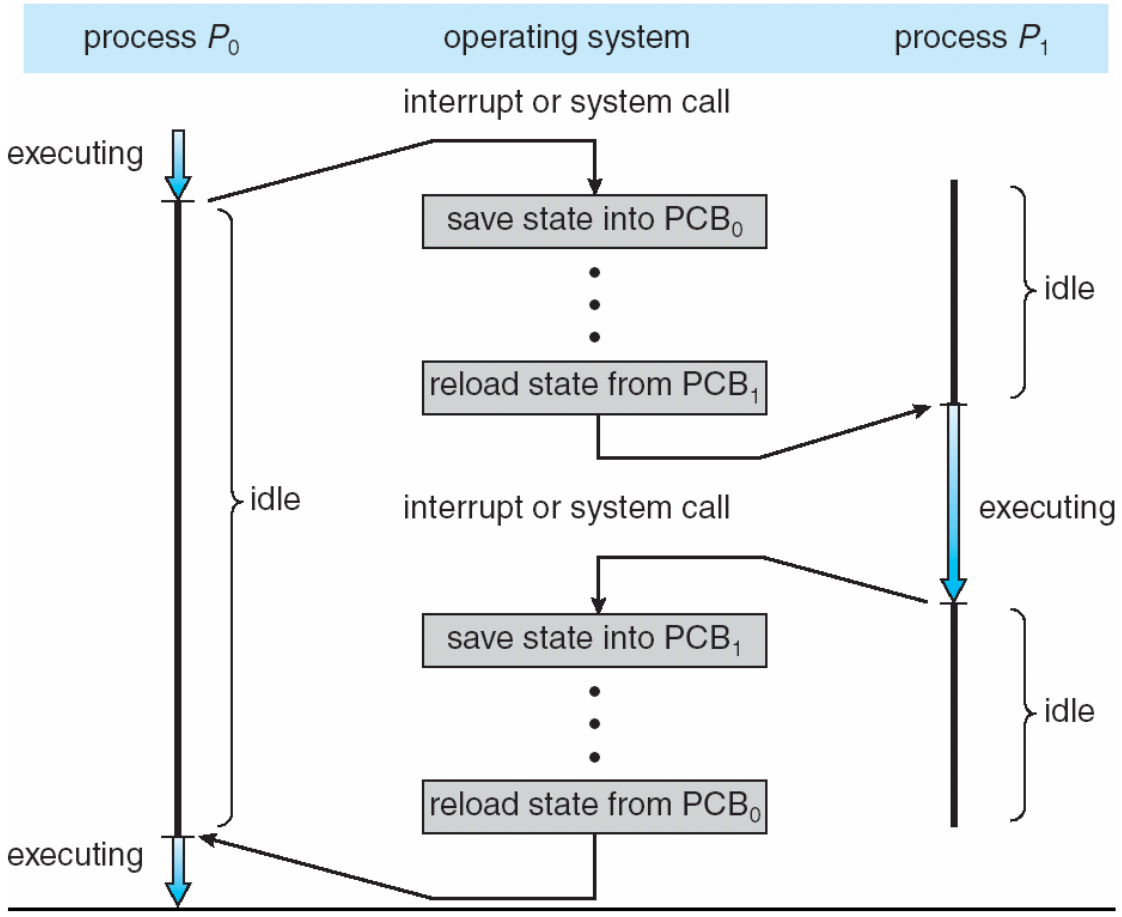


Process Control Block (PCB)

- How to manage processes?
- Information associated with each process (also task control block) → Linux
  - Process state
  - Program counter
  - CPU registers – contents of all process-centric registers
  - CPU scheduling info. – priorities, scheduling queue pointers
  - Memory-management info. – memory allocated to the process
  - Accounting info. – CPU used, clock time elapsed since start, time limits
  - I/O status info. – I/O devices allocated to process, list of open files



CPU switch from process to process



عامل رفتن از process به OS
- interrupt : عامل خارجی
- system call : " داخلی

* هر PCB با کمک ID به فرایند مربوط به خودش link میشه.

⚠ توی idle ها از HLT استفاده نمیشه چون OS دارن اجرا میشه.

* برای switch کردن بین برنامه‌ها نیاز به OS داریم.

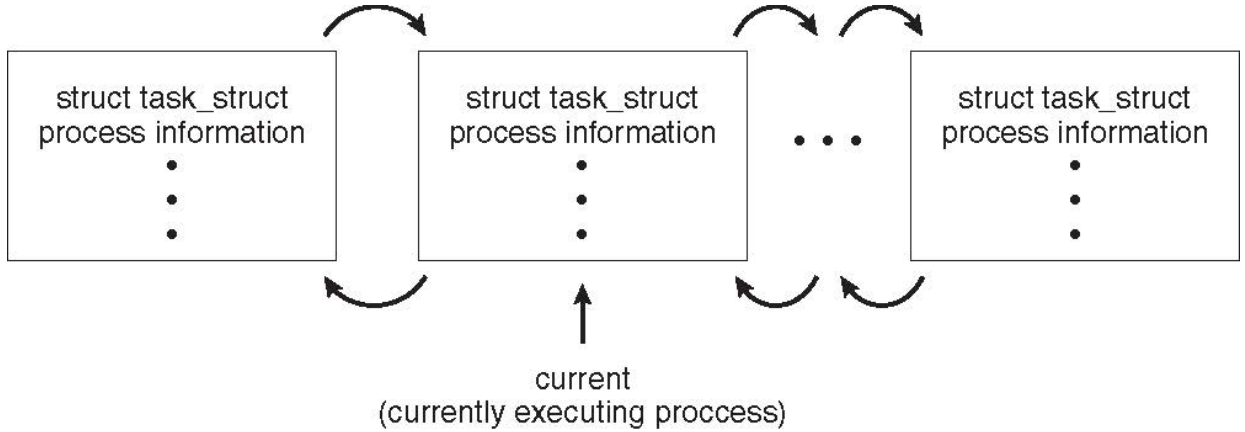
Process representation in Linux

Represented by the C structure `task_struct`

گشتن میان process ها

O(n) ←

```
pid t_pid; /* process identifier */
long state; /* state of the process */
unsigned int time_slice /* scheduling information */
struct task_struct *parent; /* this process's parent */
struct list_head children; /* this process's children */
struct files_struct *files; /* list of open files */
struct mm_struct *mm; /* address space of this process */
```



Process scheduling

- **Process scheduler selects among available processes for next execution on CPU**

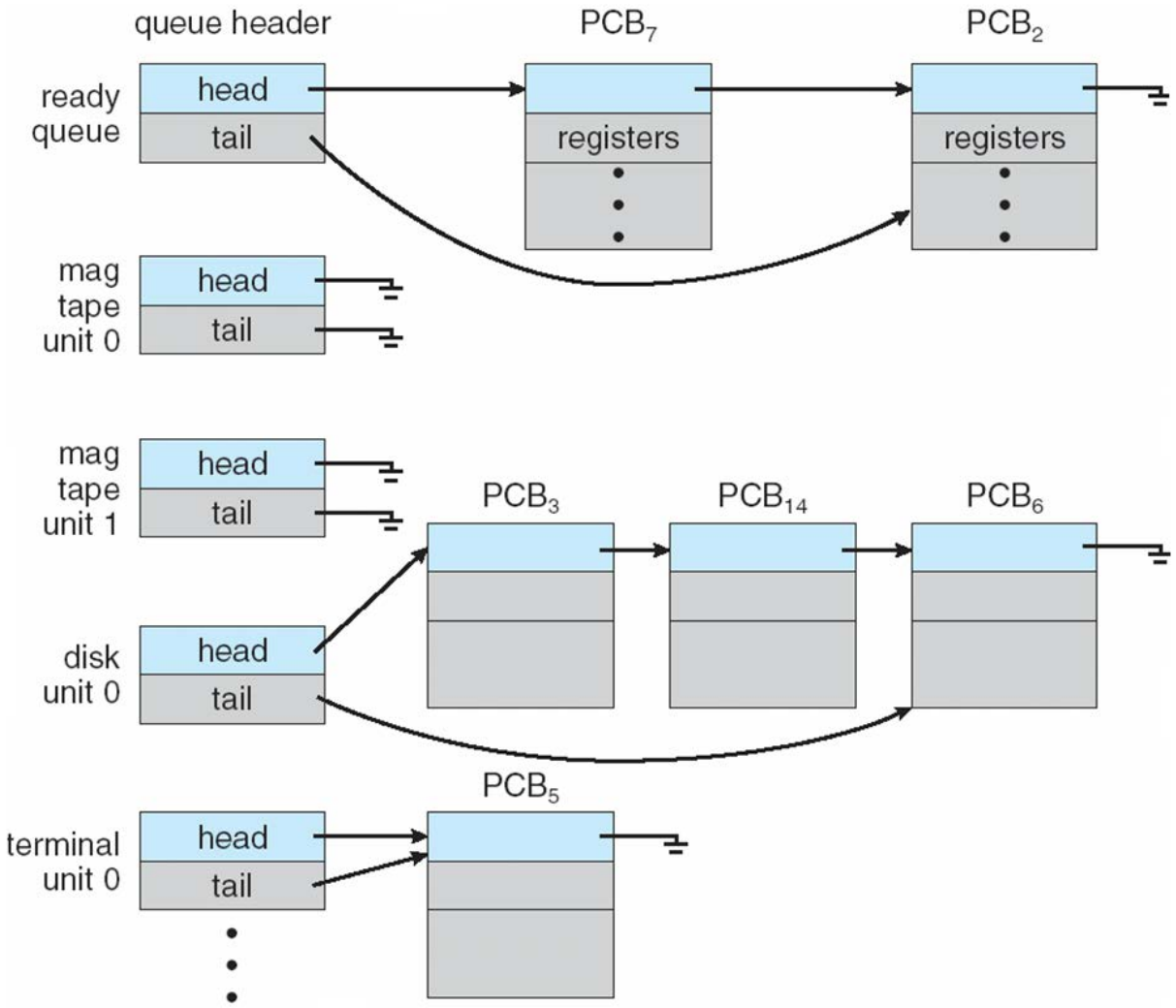
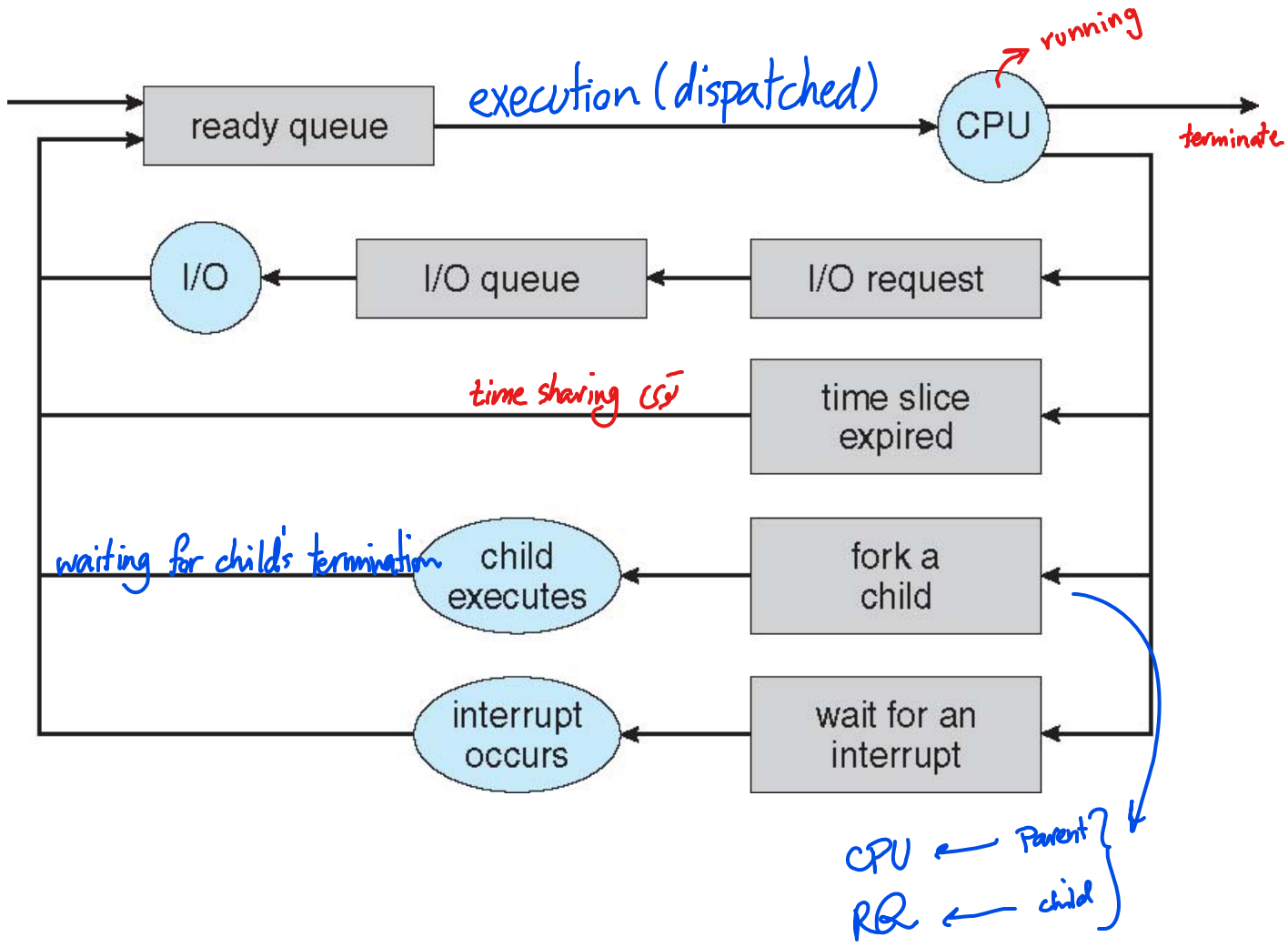


Diagram representation of process scheduling

- **Queueing diagram** represents queues, resources, flows



Schedulers

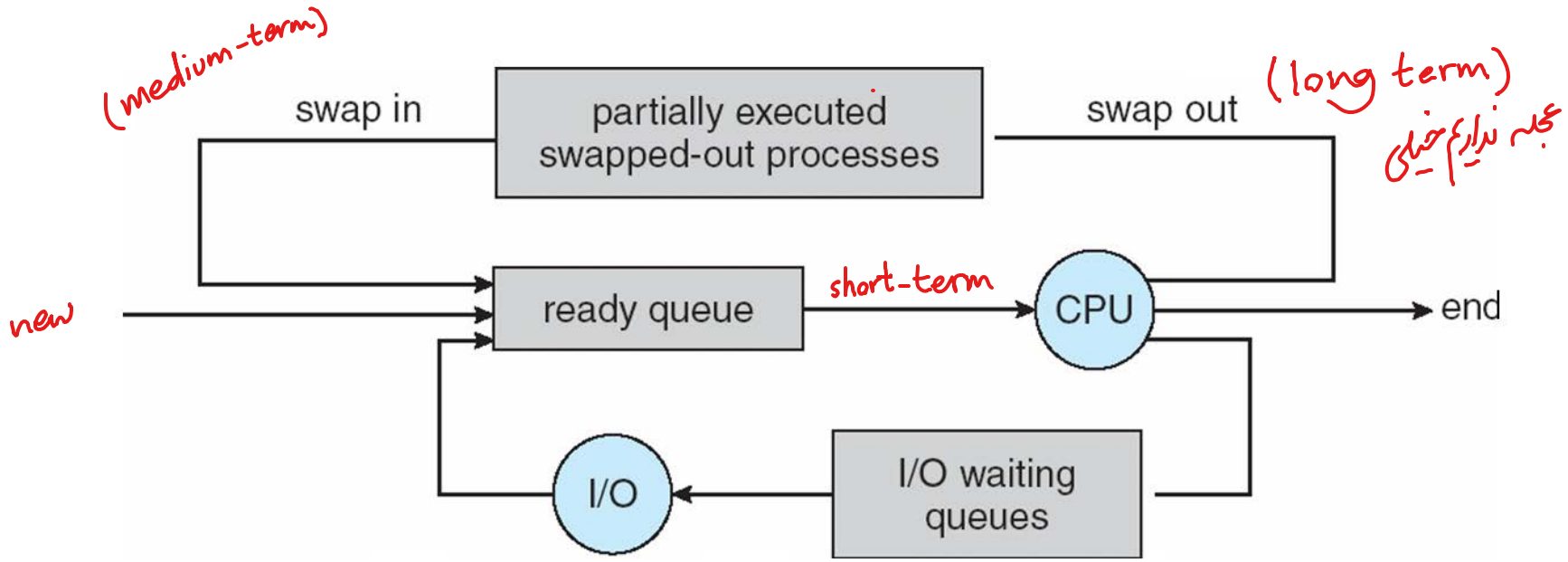
- **Short-term scheduler** (or **CPU scheduler**)
  - selects which process should be executed next and allocates CPU
    - Sometimes the only scheduler in a system
    - Short-term scheduler is invoked frequently (milliseconds) ⇒ (must be fast)
- **Long-term scheduler** (or **job scheduler**)
  - selects which processes should be brought into the ready queue
    - Long-term scheduler is invoked infrequently (seconds, minutes) ⇒ (may be slow)
    - The long-term scheduler controls the **degree of multiprogramming**
- **Processes:**
  - I/O-bound
    - spends more time doing I/O than computations, many short CPU bursts
  - CPU-bound
    - spends more time doing computations; few very long CPU bursts
- **Long-term scheduler strives for good *process mix***

time sharing systems (UNIX - Microsoft windows) have no long-term scheduler but put every new process in Mem. for short-term.

Example of standard API

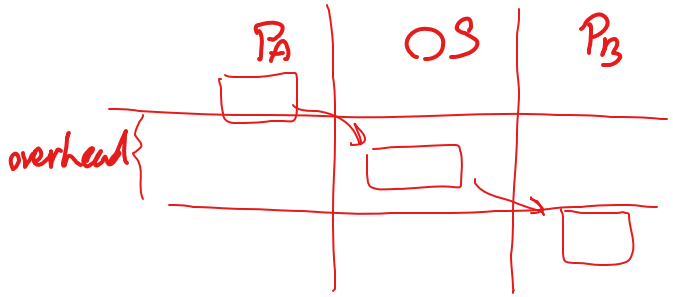
برعکس long-term میاد Ready queue رو گسترش میده.

- **Medium-term scheduler**
  - Can be added if degree of multiple programming needs to decrease
  - Remove process from memory, store on disk, bring back in from disk to continue execution: **swapping**



Context switch (تعویض متن)

- When CPU switches to another process, the system must save the state of the old process and load the saved state for the new process via a context switch



- Context of a process represented in the PCB
- Context-switch time is overhead; the system does no useful work while switching
  - The more complex the OS and the PCB ➔ the longer the context switch
- Time dependent on hardware support
  - Some hardware provides multiple sets of registers per CPU ➔ multiple contexts loaded at once

memory speed
# of regs must be copied
instructions

↓overhead ⇐ میزان ↑ پیچیدگی نرم افزار

Process creation

- **Parent** vs. **Child**

(int)

- Generally, process identified and managed via a process identifier (pid)

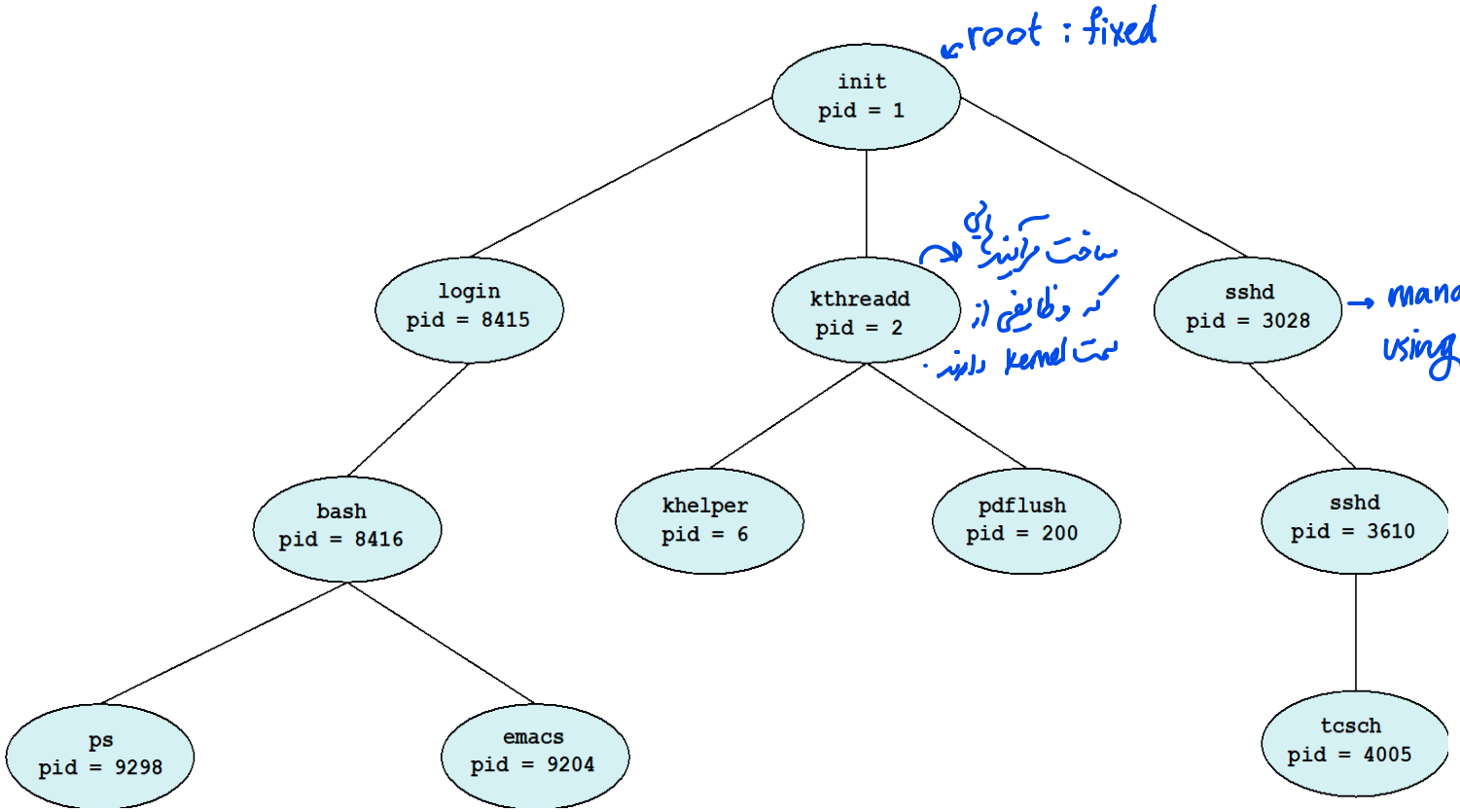
→ variables , memory allocations , ...

- **Resource** sharing options
  - Parent and children share all resources
  - Children share subset of parent's resources
  - Parent and child share no resources

- Execution options
  - Parent and children execute concurrently
  - Parent waits until children terminate

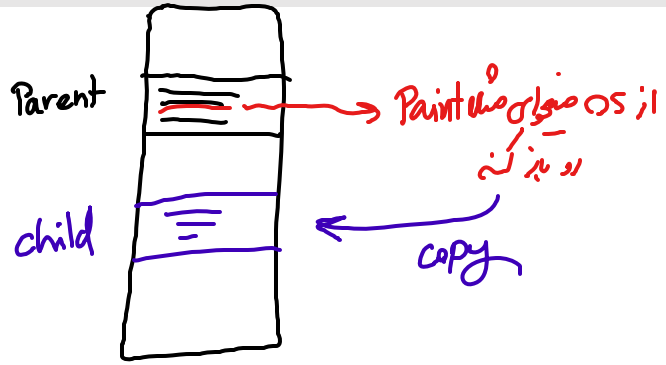
A tree of processes in Linux

لود اتوماتیک نه دستی (همیاع خودمون چیزی Run کنیم)



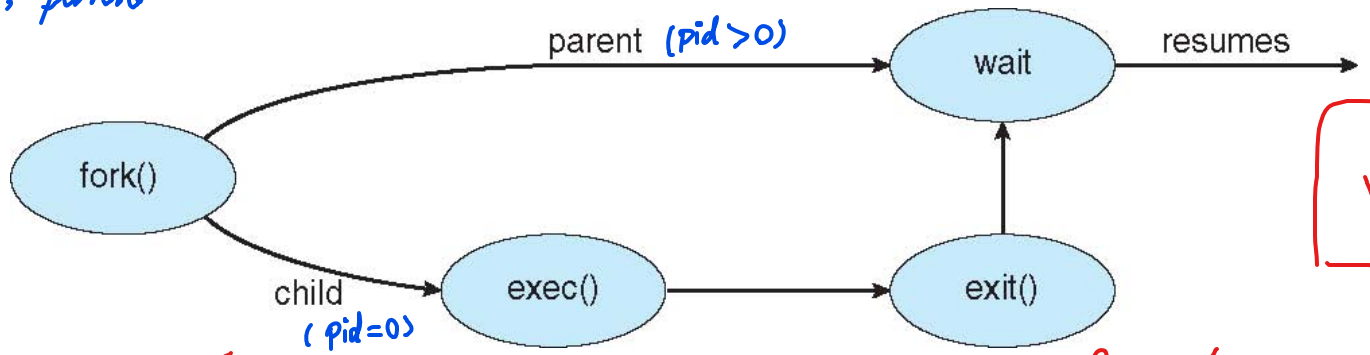
Process creation

- **Address space**
  - Child duplicate of parent
  - Child has a program loaded into it



- **UNIX examples**
  - `fork()` system call creates new process
  - `exec()` system call used after a `fork()` to replace the process' memory space with a new program

return code for fork = { 0 : child ; pid : parent }



∀ fork() ∃ wait()

fork() → copy → Ready queue → update parent → overwrite child

Process creation with C

جلسه ۲۲ / ۷ / ۹۹ روبین

POSIX

```
#include <sys/types.h>
#include <stdio.h>
#include <unistd.h>

int main()
{
pid_t pid;

    /* fork a child process */
    pid = fork();

    if (pid < 0) { /* error occurred */
       fprintf(stderr, "Fork Failed");
       return 1;
    }
    else if (pid == 0) { /* child process */
       execlp("/bin/ls","ls",NULL);
    }
    else { /* parent process */
       /* parent will wait for the child to complete */
       wait(NULL);
       printf("Child Complete");
    }

    return 0;
}
```

Windows

```
#include <stdio.h>
#include <windows.h>

int main(VOID)
{
STARTUPINFO si;
PROCESS_INFORMATION pi;

    /* allocate memory */
    ZeroMemory(&si, sizeof(si));
    si.cb = sizeof(si);
    ZeroMemory(&pi, sizeof(pi));

    /* create child process */
    if (!CreateProcess(NULL, /* use command line */
     "C:\\WINDOWS\\system32\\mspaint.exe", /* command */
     NULL, /* don't inherit process handle */
     NULL, /* don't inherit thread handle */
     FALSE, /* disable handle inheritance */
     0, /* no creation flags */
     NULL, /* use parent's environment block */
     NULL, /* use parent's existing directory */
     &si,
     &pi))
    {
       fprintf(stderr, "Create Process Failed");
       return -1;
    }
    /* parent will wait for the child to complete */
    WaitForSingleObject(pi.hProcess, INFINITE);
    printf("Child Complete");

    /* close handles */
    CloseHandle(pi.hProcess);
    CloseHandle(pi.hThread);
}
```

Process termination

- **Child → Parent**
  - Process' resources are deallocated when:
    - *exit(n)*
    - *return()* in *main()*
  - Catch exit status → *wait()*
    - `pid = wait(&status);`
- **Parent → Child**
  - *abort()*
  - Why?
    - Child has exceeded allocated resources
    - Task assigned to child is no longer required
    - The parent is exiting and the operating systems does not allow a child to continue if its parent terminates

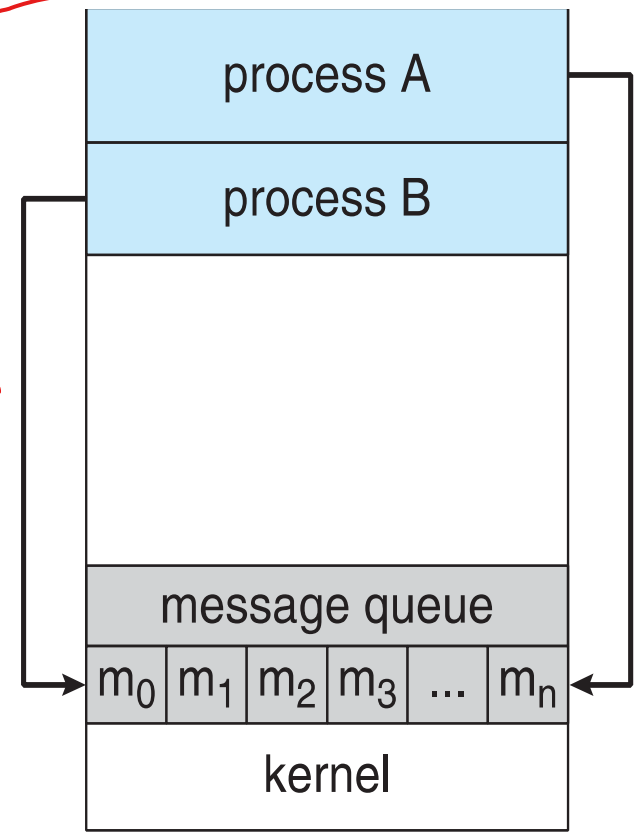
Problems of process termination

- **zombie process**
  - No parent waiting → child : terminated منتظر نمیمونه Parent هیچ وقت
- **orphan process**
  - Parent termination without wait

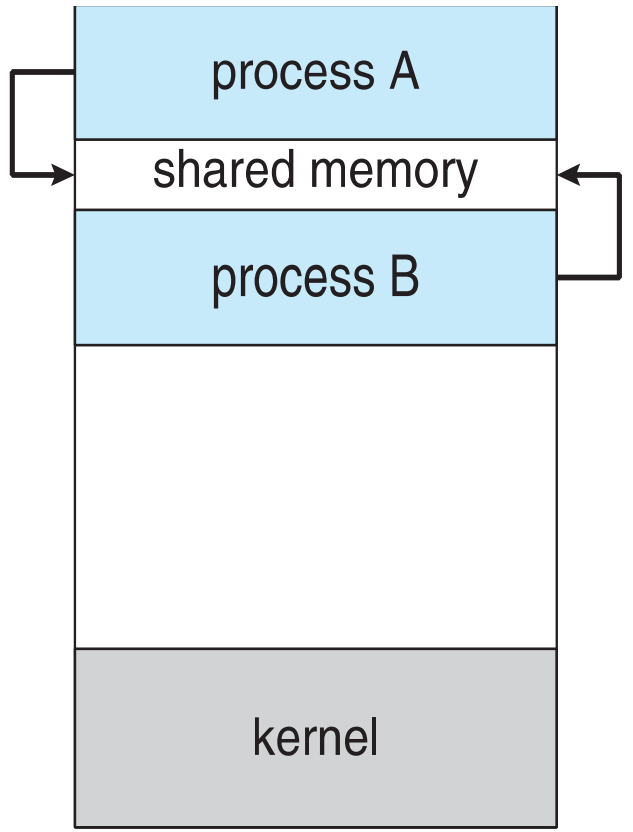
- **Multi process example: Chrome Browser**
  - Browser, Renderer, Plugins, etc

Interprocess communication (*IPC*)

- **Process:**
  - independent vs. cooperating
- **Cooperating process:**
  - Shared memory
  - Message passing



(a)



(b)

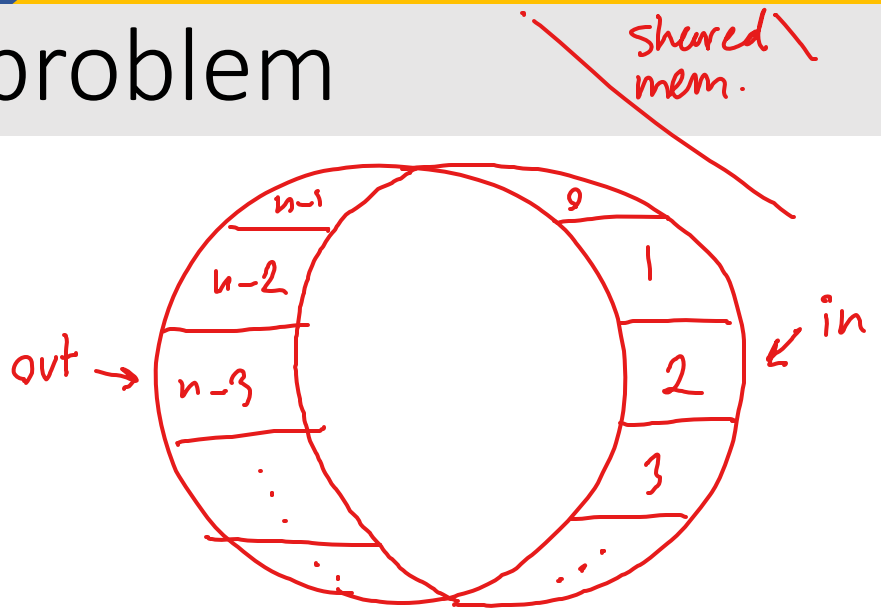
Circular buffer & producer-consumer problem

in = out → empty buffer

bounded buffer

```
#define BUFFER_SIZE 10
typedef struct {
  . . .
} item;

item buffer[BUFFER_SIZE];
int in = 0;
int out = 0;
```




```
item next_produced;
while (true) {
        /* produce an item in next produced */
        while (((in + 1) % BUFFER_SIZE) == out)
                ; /* do nothing */
        buffer[in] = next_produced;
        in = (in + 1) % BUFFER_SIZE;
}
```

```
item next_consumed;
while (true) {
        while (in == out) ; /* do nothing */

        next_consumed = buffer[out];
        out = (out + 1) % BUFFER_SIZE;

        /* consume the item in next consumed */
}
```

Message passing

- **Direct communication (unidirectional)**
  - **send** (*P, message*) – send a message to process P
  - **receive**(*Q, message*) – receive a message from process Q

- **Indirect communication (uni & bidirectional)**
  - Messages are directed and received from mailboxes (or ports)
  - Can be used by multiple processes
  - Primitives are defined as:
    - **send**(*A, message*) – send a message to mailbox A
    - **receive**(*A, message*) – receive a message from mailbox A

Synchronization

- **Blocking vs. non-blocking**
- **Blocking is considered synchronous**
  - Blocking send
  - Blocking receive
- **Non-blocking is considered asynchronous**
  - Non-blocking send
  - Non-blocking receive
    - The receiver receives
      - A valid message
      - Null message
- **Different combinations possible**
  - If both send and receive are blocking, we have a rendezvous

POSIX examples of shared memory: (sender->receiver)

```c
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <fcntl.h>
#include <sys/shm.h>
#include <sys/stat.h>

int main()
{
/* the size (in bytes) of shared memory object */
const int SIZE = 4096;
/* name of the shared memory object */
const char *name = "OS";
/* strings written to shared memory */
const char *message_0 = "Hello";
const char *message_1 = "World!";

/* shared memory file descriptor */
int shm_fd;
/* pointer to shared memory obect */
void *ptr;

    /* create the shared memory object */
    shm_fd = shm_open(name, O_CREAT | O_RDWR, 0666);

    /* configure the size of the shared memory object */
    ftruncate(shm_fd, SIZE);

    /* memory map the shared memory object */
    ptr = mmap(0, SIZE, PROT_WRITE, MAP_SHARED, shm_fd, 0);

    /* write to the shared memory object */
    sprintf(ptr,"%s",message_0);
    ptr += strlen(message_0);
    sprintf(ptr,"%s",message_1);
    ptr += strlen(message_1);

    return 0;
}
```

```c
#include <stdio.h>
#include <stdlib.h>
#include <fcntl.h>
#include <sys/shm.h>
#include <sys/stat.h>

int main()
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const int SIZE = 4096;
/* name of the shared memory object */
const char *name = "OS";
/* shared memory file descriptor */
int shm_fd;
/* pointer to shared memory obect */
void *ptr;

    /* open the shared memory object */
    shm_fd = shm_open(name, O_RDONLY, 0666);

    /* memory map the shared memory object */
    ptr = mmap(0, SIZE, PROT_READ, MAP_SHARED, shm_fd, 0);

    /* read from the shared memory object */
    printf("%s",(char *)ptr);

    /* remove the shared memory object */
    shm_unlink(name);

    return 0;
}
```
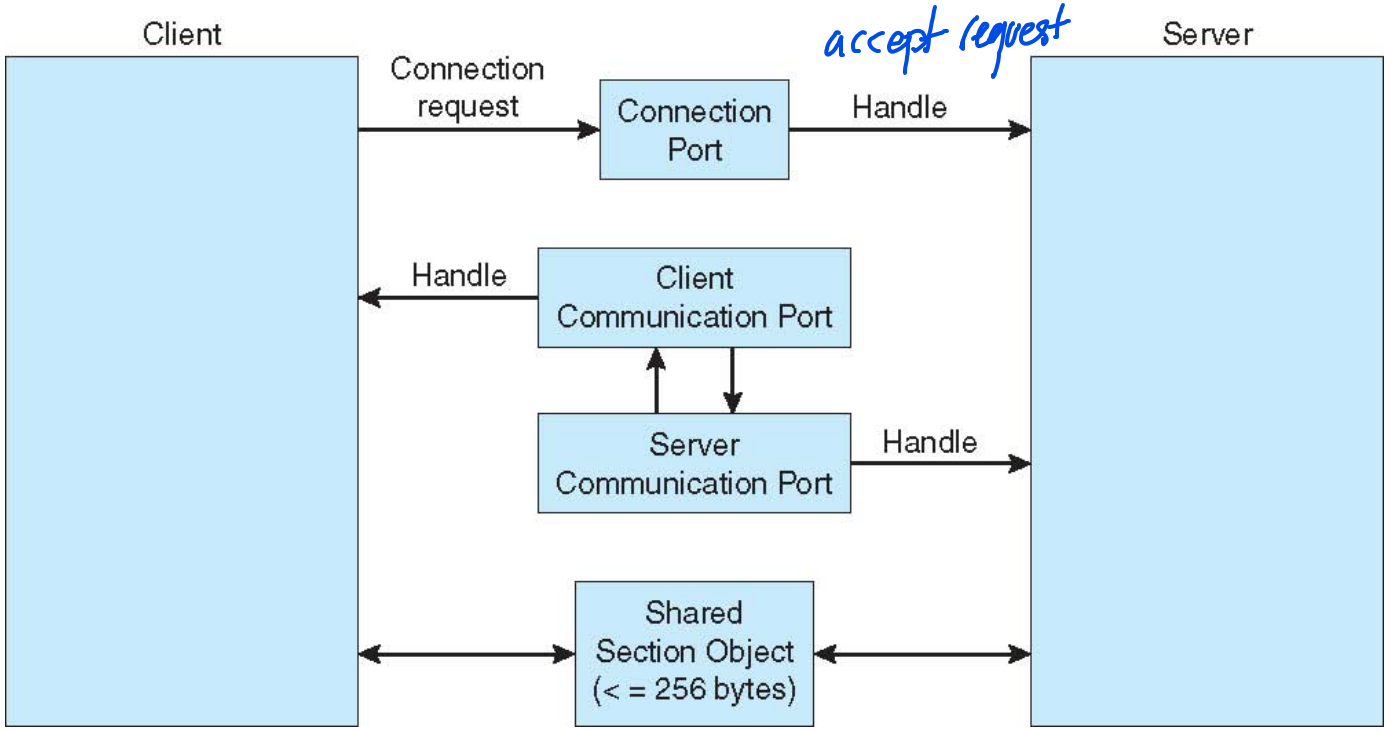

Local procedure calls in Windows (ALPC)

Advanced



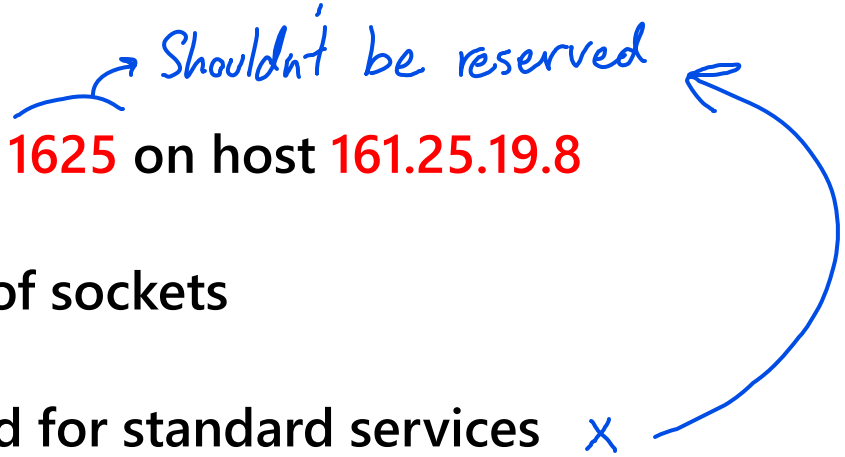
windows محدودیت برای

Communications in client-server systems

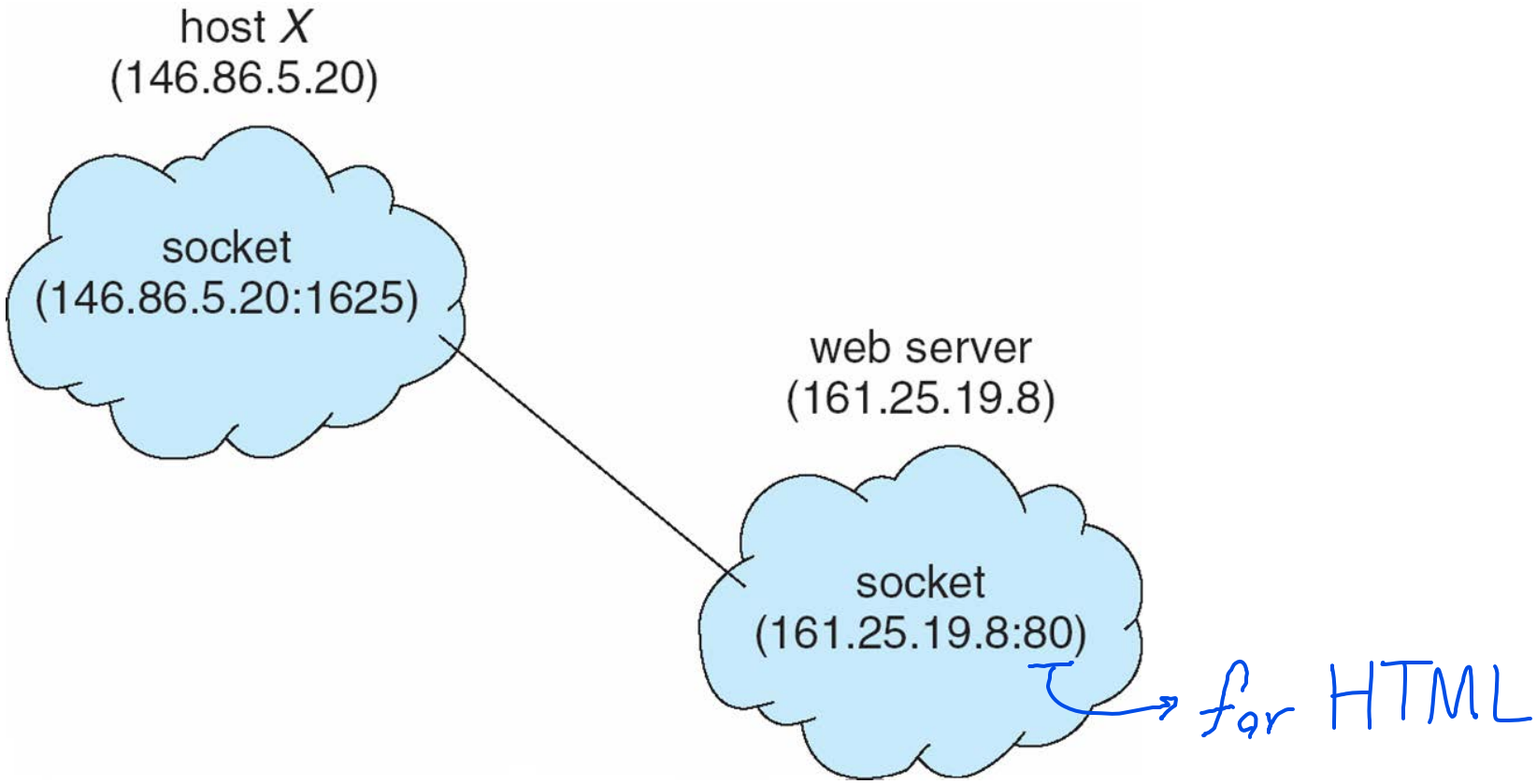
- Sockets
- Remote Procedure Calls (windows)
- Pipes
- Remote Method Invocation (Java)

Sockets

- A socket is defined as an endpoint for communication
- Concatenation of IP address and port – a number included at start of message packet to differentiate network services on a host
- The socket 161.25.19.8:1625 refers to port 1625 on host 161.25.19.8
- Communication consists between a pair of sockets
- All ports below 1024 are *well known*, used for standard services
- Special IP address 127.0.0.1 (loopback) to refer to system on which process is running



Socket communication



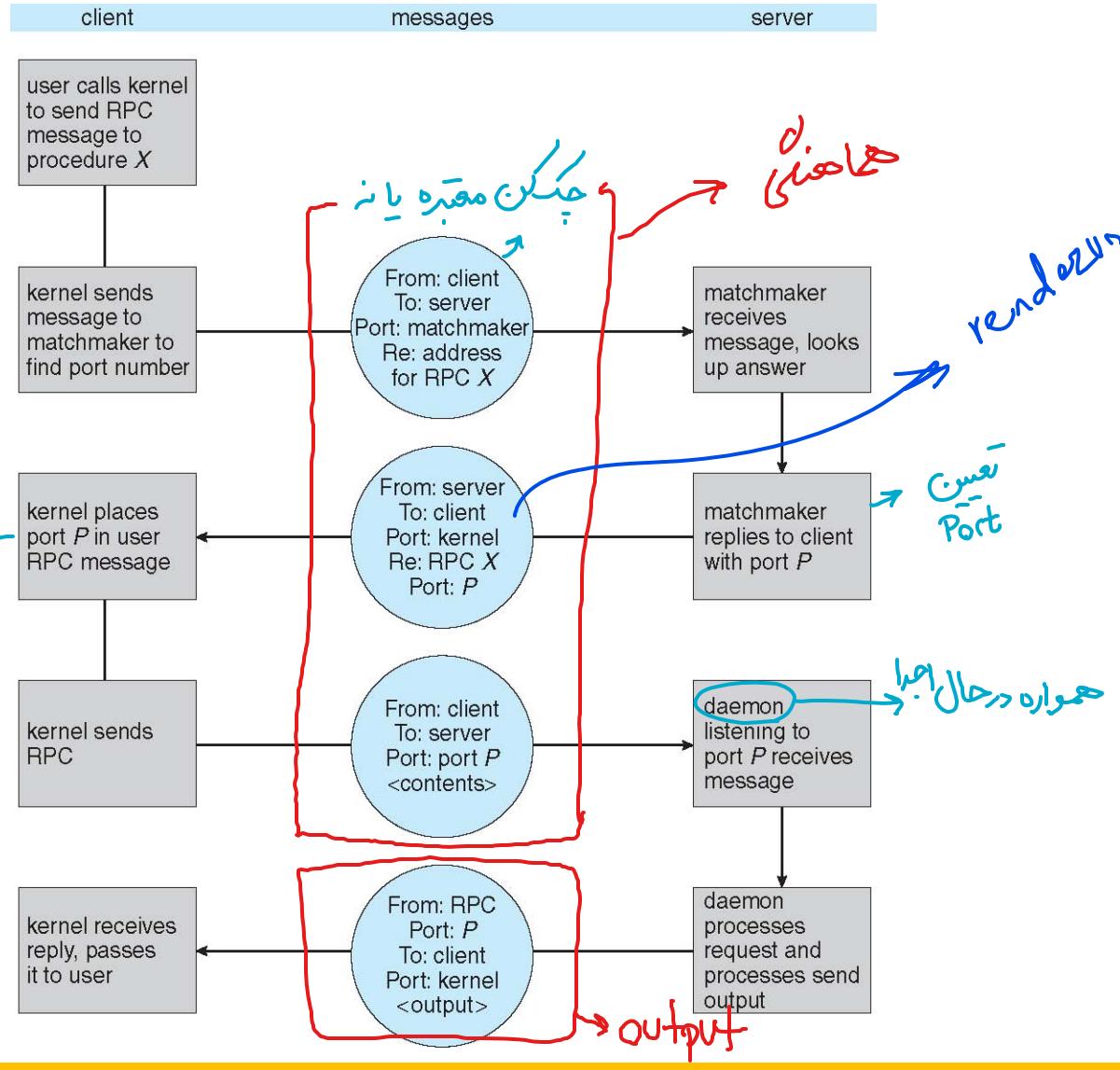
Sockets in Java

- Three types of sockets
  - Connection-oriented (TCP)
  - Connectionless (UDP)
  - `MulticastSocket` class– data can be sent to multiple recipients

- Consider this "Date" server:

```
import java.net.*;
import java.io.*;

public class DateServer
{
   public static void main(String[] args) {
      try {
        ServerSocket sock = new ServerSocket(6013);

        /* now listen for connections */
        while (true) {
          Socket client = sock.accept();

          PrintWriter pout = new
           PrintWriter(client.getOutputStream(), true);

          /* write the Date to the socket */
          pout.println(new java.util.Date().toString());

          /* close the socket and resume */
          /* listening for connections */
          client.close();
        }
      }
      catch (IOException ioe) {
        System.err.println(ioe);
      }
   }
}
```

Execution of RPC (Remote Procedure Call)



Pipes (#Parent, #Child)

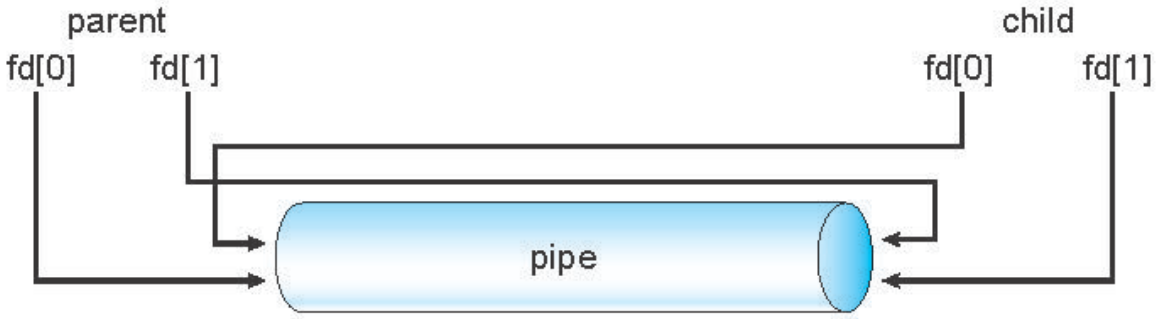
- Acts as a conduit allowing two processes to communicate
- Issues:
  - Is communication unidirectional or bidirectional?
  - In the case of two-way communication, is it half or full-duplex? (تبادل همزمان دوطرفه)
  - Must there exist a relationship (i.e., *parent-child*) between the communicating processes?
  - Can the pipes be used over a network?

- **Ordinary pipes** (unidirectional)
  - cannot be accessed from outside the process that created it. Typically, a parent process creates a pipe and uses it to communicate with a child process that it created.

- **Named pipes** (bidirectional)
  - can be accessed without a parent-child relationship.

Ordinary Pipes

- Ordinary Pipes allow communication in standard producer-consumer style
- Producer writes to one end (the write-end of the pipe)
- Consumer reads from the other end (the read-end of the pipe)
- Ordinary pipes are therefore unidirectional
- Require parent-child relationship between communicating processes



- Windows calls these anonymous pipes
- See Unix and Windows code samples in textbook

Ordinary pipe (POSIX), parent-child

```
#include <sys/types.h>
#include <stdio.h>
#include <string.h>
#include <unistd.h>

#define BUFFER_SIZE 25
#define READ_END 0
#define WRITE_END 1

int main(void)
{
char write_msg[BUFFER_SIZE] = "Greetings";
char read_msg[BUFFER_SIZE];
int fd[2];
pid_t pid;
    /* create the pipe */
    if (pipe(fd) == -1) {
       fprintf(stderr,"Pipe failed");
       return 1;
    }

    /* fork a child process */
    pid = fork();

    if (pid < 0) { /* error occurred */
       fprintf(stderr, "Fork Failed");
       return 1;
    }
```

```
    if (pid > 0) { /* parent process */
       /* close the unused end of the pipe */
       close(fd[READ_END]);

       /* write to the pipe */
       write(fd[WRITE_END], write_msg, strlen(write_msg)+1);

       /* close the write end of the pipe */
       close(fd[WRITE_END]);
    }
    else { /* child process */
       /* close the unused end of the pipe */
       close(fd[WRITE_END]);

       /* read from the pipe */
       read(fd[READ_END], read_msg, BUFFER_SIZE);
       printf("read %s",read_msg);

       /* close the write end of the pipe */
       close(fd[READ_END]);
    }

    return 0;
}
```

Ordinary pipe (windows), parent

```
#include <stdio.h>
#include <stdlib.h>
#include <windows.h>

#define BUFFER_SIZE 25

int main(VOID)
{
HANDLE ReadHandle, WriteHandle;
STARTUPINFO si;
PROCESS_INFORMATION pi;
char message[BUFFER_SIZE] = "Greetings";
DWORD written;
   /* set up security attributes allowing pipes to be inherited */
   SECURITY_ATTRIBUTES sa = {sizeof(SECURITY_ATTRIBUTES),NULL,TRUE};
   /* allocate memory */
   ZeroMemory(&pi, sizeof(pi));

   /* create the pipe */
   if (!CreatePipe(&ReadHandle, &WriteHandle, &sa, 0)) {
      fprintf(stderr, "Create Pipe Failed");
      return 1;
   }

   /* establish the START_INFO structure for the child process */
   GetStartupInfo(&si);
   si.hStdOutput = GetStdHandle(STD_OUTPUT_HANDLE);

   /* redirect standard input to the read end of the pipe */
   si.hStdInput = ReadHandle;
   si.dwFlags = STARTF_USESTDHANDLES;

   /* don't allow the child to inherit the write end of pipe */
   SetHandleInformation(WriteHandle, HANDLE_FLAG_INHERIT, 0);

   /* create the child process */
   CreateProcess(NULL, "child.exe", NULL, NULL,
    TRUE, /* inherit handles */
    0, NULL, NULL, &si, &pi);

   /* close the unused end of the pipe */
   CloseHandle(ReadHandle);

   /* the parent writes to the pipe */
   if (!WriteFile(WriteHandle, message,BUFFER_SIZE,&written,NULL))
      fprintf(stderr, "Error writing to pipe.");

   /* close the write end of the pipe */
   CloseHandle(WriteHandle);

   /* wait for the child to exit */
   WaitForSingleObject(pi.hProcess, INFINITE);
   CloseHandle(pi.hProcess);
   CloseHandle(pi.hThread);
   return 0;
}
```

Ordinary pipe (windows), child

```
#include <stdio.h>
#include <windows.h>

#define BUFFER_SIZE 25

int main(VOID)
{
HANDLE Readhandle;
CHAR buffer[BUFFER_SIZE];
DWORD read;

   /* get the read handle of the pipe */
   ReadHandle = GetStdHandle(STD_INPUT_HANDLE);

   /* the child reads from the pipe */
   if (ReadFile(ReadHandle, buffer, BUFFER_SIZE, &read, NULL))
     printf("child read %s",buffer);
   else
     fprintf(stderr, "Error reading from pipe");

   return 0;
}
```

Named pipes

- **Named Pipes** are more **powerful** than ordinary pipes (?) → بعد از بین رفتن Parent باز هم قابل استفاده اند.
- **Communication is bidirectional**
- **No** parent-child relationship is necessary between the communicating processes
- **Several** processes can use the named pipe for communication
- **Provided on both UNIX and Windows systems**

Questions?