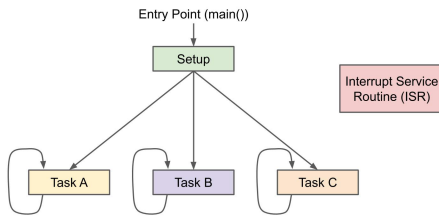


## TASK SCHEDULING

- Block diagram of a multi-task program:

What our code looks like



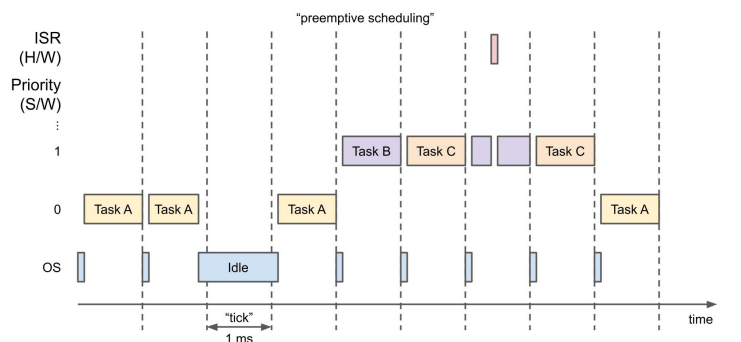
each task runs concurrently in its own while loop.

- can also set up ISRs (interrupt service routines) that can preempt any of the tasks, to execute some code.

- ISR →
- handles hardware timer overflows
  - pin state changes
  - new communication on a bus

## TASK SCHEDULING IN SINGLE CORE SYSTEMS

What actually happens\*



→ CPU divides the tasks into time slices

↳ tasks appear to run concurrently

→ schedulers determine which task to run in each time slice

→ time slice = 1ms = 1 tick

↳ hardware timer creates interrupt every 1ms

↳ ISR for timer runs the scheduler

↳ chooses the task to run next

→ at each tick interrupt → task w/ highest priority is chosen to run

→ task w/ the same highest priority → uses round-robin execution

→ if a task w/ higher priority than running task is available → immediately runs

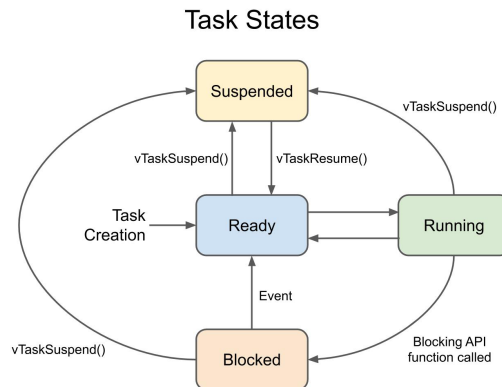
↓  
doesn't wait  
for next tick

→ hardware interrupt <sup>higher priority</sup> > running software

↳ ∴ hardware ISR can interrupt any task

→ can assign task priority using `vTaskPrioritySet()`

## TASK STATES



**Ready state** → task enters this state when it is created

↳ this means the task is ready to run

→ during each tick → 1 task in ready state is chosen by scheduler to run

**Running state** → when tasks are running

→ can go to ready state again by the scheduler

**Blocked state** → when functions cause the task to wait eg. `vTaskDelay()`

→ occurs when tasks wait for an event to occur eg. timer on `vTaskDelay()` to expire

→ allows other tasks to run instead of the blocked task

**Suspended state** → like putting task to sleep eg. use `vTaskSuspend()`

→ any task can put any other task in this mode (including itself)

→ tasks go back to Ready state by calling `vTaskResume()`