

# Heuristic Search

- Best First Search
  - A\*
  - IDA\*
  - Beam Search
- Generate and Test
- Local Searches
  - Hill Climbing
    - Simple Hill Climbing
    - Steepest Ascend Hill Climbing
    - Stochastic Hill Climbing
  - Simulated Annealing
- Local beam search
- Genetic Algorithms

# State spaces

$\langle S, P, I, G, W \rangle$

- state space

$S$

- set of states

$P = \{(x, y) | x, y \in S\}$

- set of production rules

$I \in S$

- initial state

$G \subseteq S$

- set of goal states

$W: P \rightarrow \mathbb{R}^+$

- weight function

# Best first Search

*BestFirstSearch(state space  $\Sigma = \langle S, P, I, G, W \rangle, f$ )*

Open  $\leftarrow \{ \langle f(l), l \rangle \}$

Closed  $\leftarrow \emptyset$

**while** Open  $\neq \emptyset$  **do**

$\langle f_x, x \rangle \leftarrow ExtractMin(\text{Open})$

**if** Goal( $x, \Sigma$ ) **then return**  $x$

*Insert( $x$ , Closed)*

**for**  $y \in Child(x, \Sigma)$  **do**

**if**  $y \notin \text{Closed}$  **then**

*Insert( $f(y), y$ ), Open)*

**return** fail

Open is implemented as a *priority queue*

# Best First Search

- **Example:** The Eight-puzzle
- The task is to transform the initial puzzle

2	8	3
1	6	4
7		5

- into the goal state

1	2	3
8		4
7	6	5

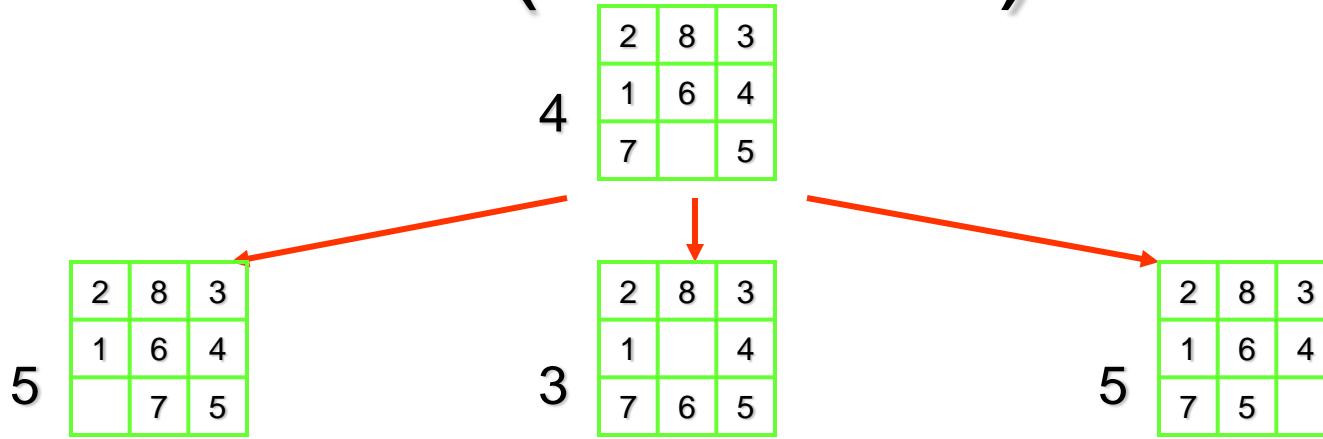
- by shifting numbers up, down, left or right.

# Best-First (Heuristic) Search

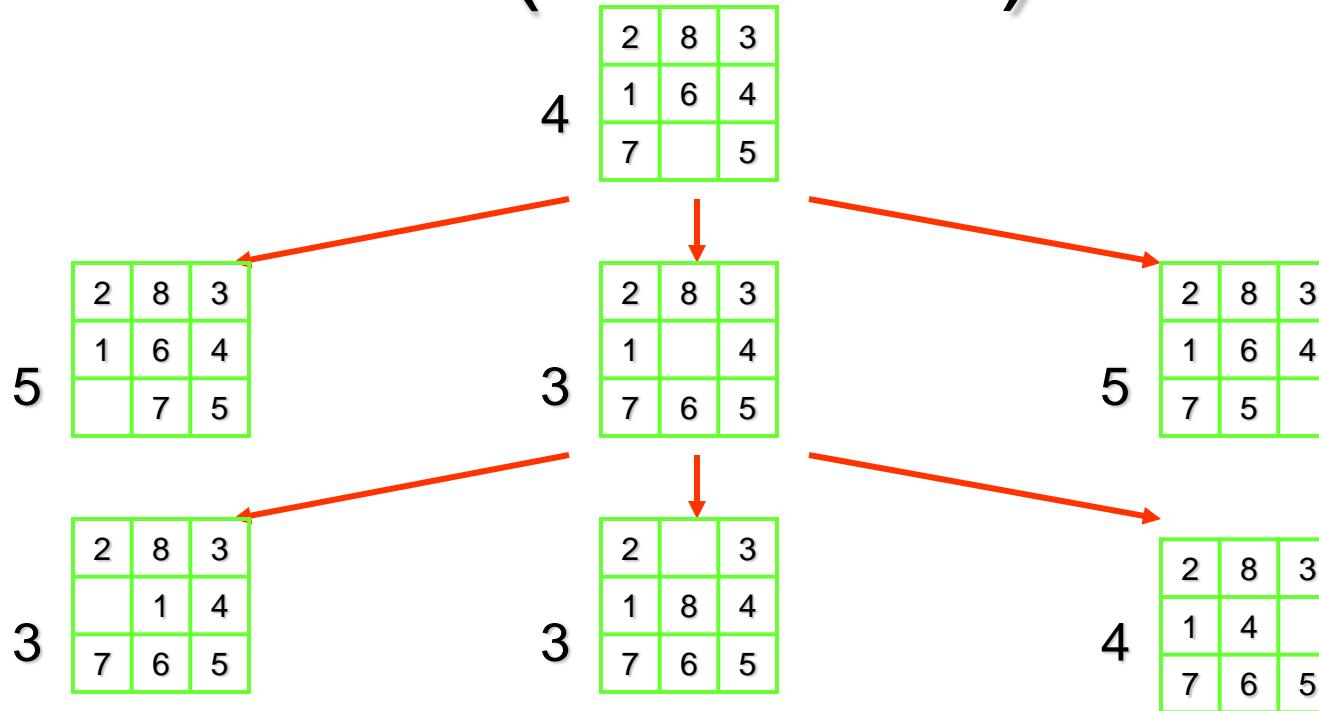
2	8	3
1	6	4
7		5

4

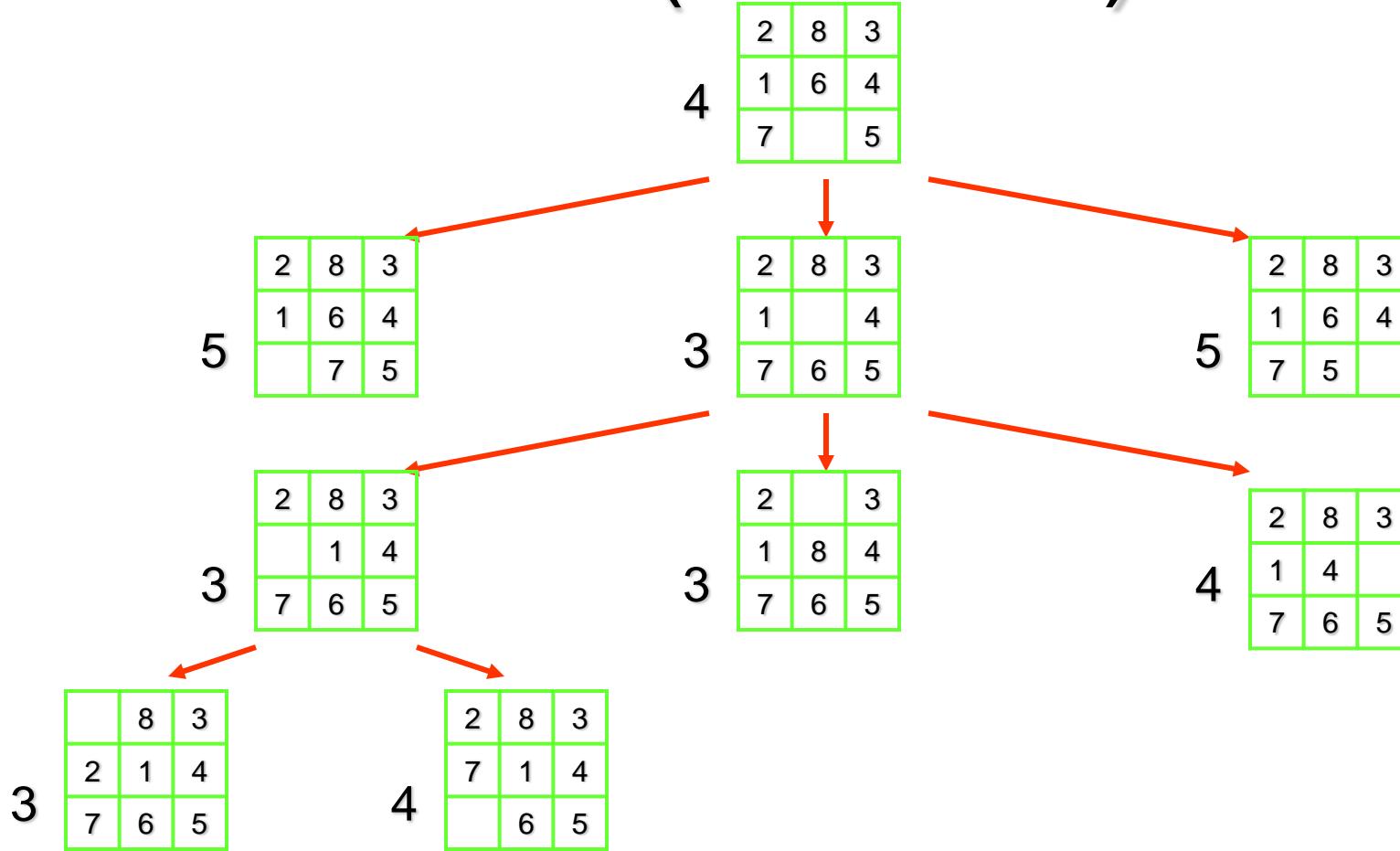
# Best-First (Heuristic) Search



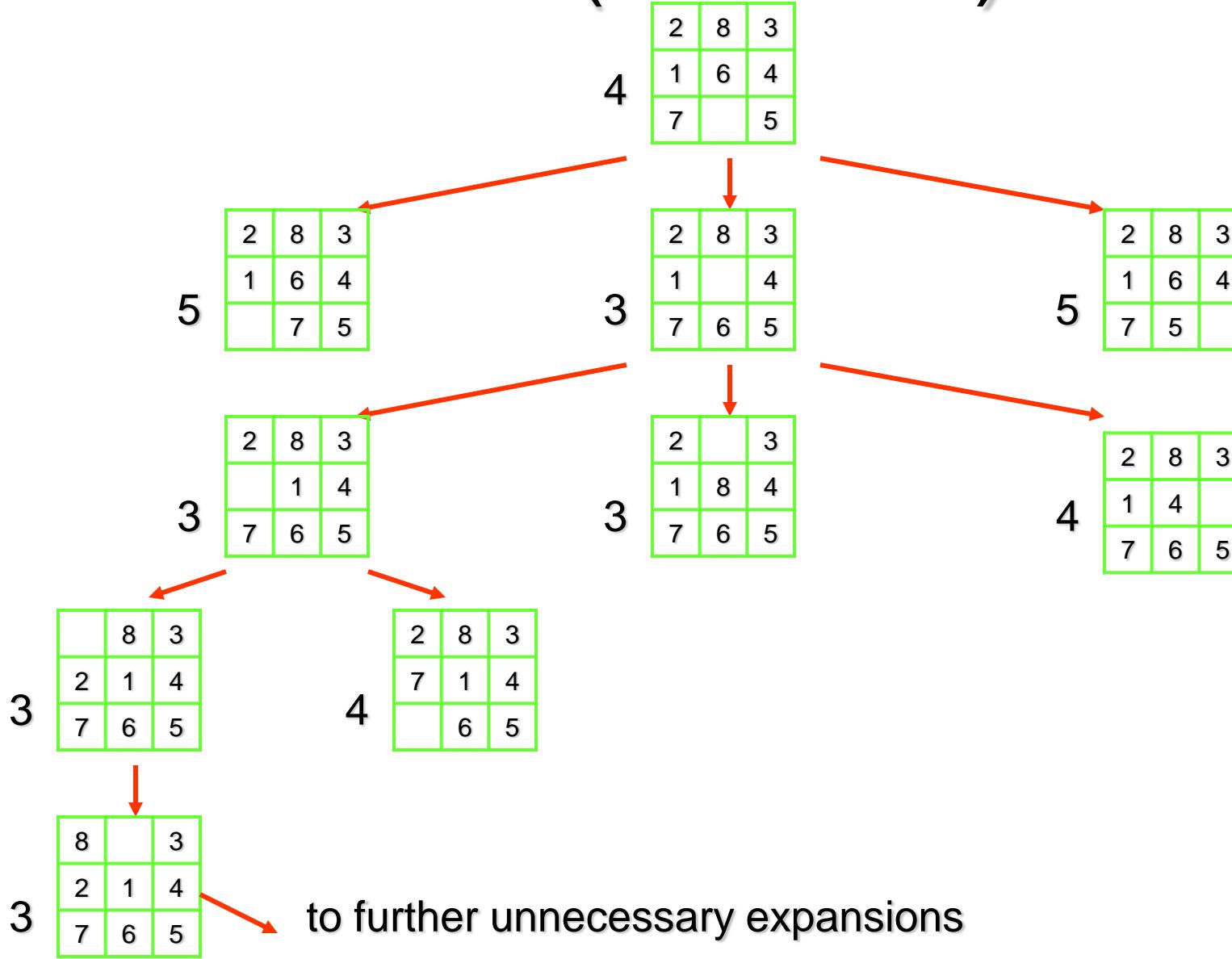
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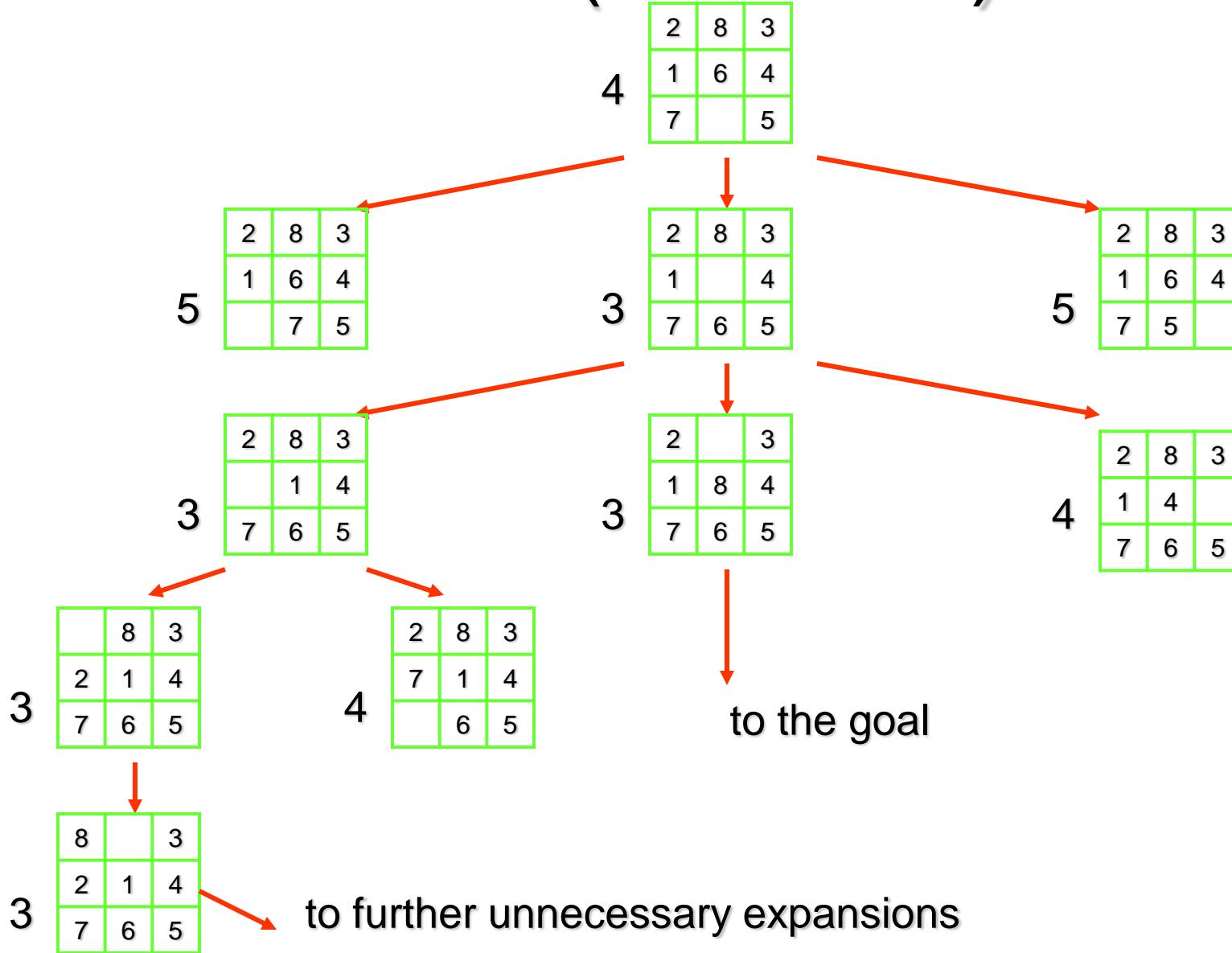
# Best-First (Heuristic) Search



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# Best First Search - A\*

$\langle S, P, I, G, W \rangle$

- state space

$h^*(x)$

- a minimum path weight from  $x$  to the goal state

$h(x)$

- heuristic estimate of  $h^*(x)$

$g(x)$

- a minimum path weight from  $I$  to  $x$

$$f(x) = g(x) + h(x)$$

# Search strategies - A\*

*A\*Search(state space  $\Sigma = \langle S, P, I, G, W \rangle, h$ )*

Open  $\leftarrow \{ \langle h(l), 0, l \rangle \}$

Closed  $\leftarrow \emptyset$

**while** Open  $\neq \emptyset$  **do**

$\langle f_x, g_x, x \rangle \leftarrow ExtractMin(Open)$  [minimum for  $f_x$ ]

**if** Goal( $x, \Sigma$ ) **then return**  $x$

*Insert*( $\langle f_x, g_x, x \rangle$ , Closed)

**for**  $y \in Child(x, \Sigma)$  **do**

$g_y = g_x + W(x, y)$

$f_y = g_y + h(y)$

**if** there is no  $\langle f, g, y \rangle \notin$  Closed with  $f \leq f_y$  **then**

*Insert*( $\langle f_y, g_y, y \rangle$ , Open) [replace existing  
 $\langle f, g, y \rangle$ , if present]

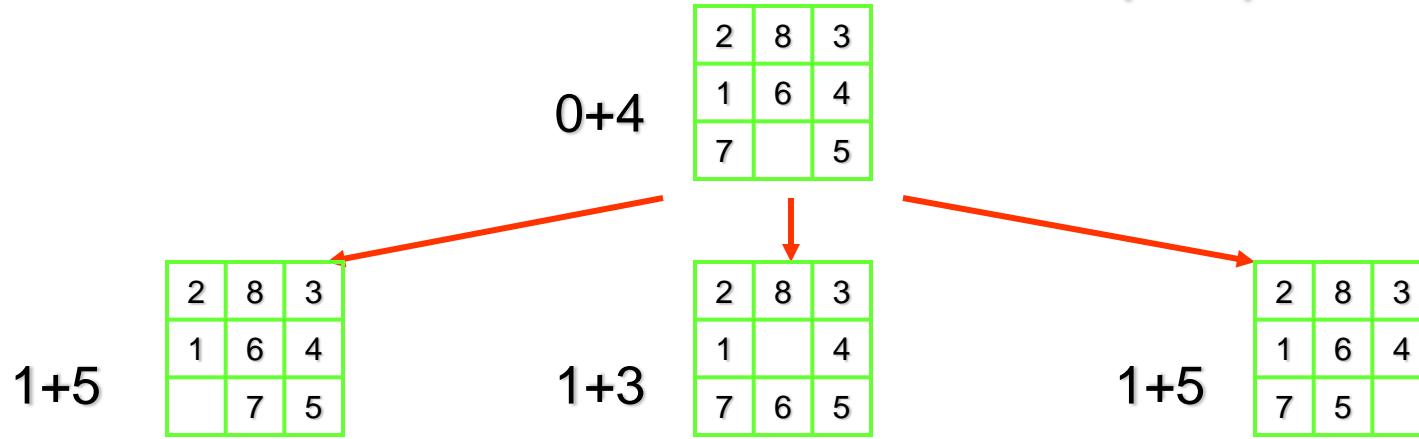
**return** fail

# Best-First Search (A<sup>\*</sup>)

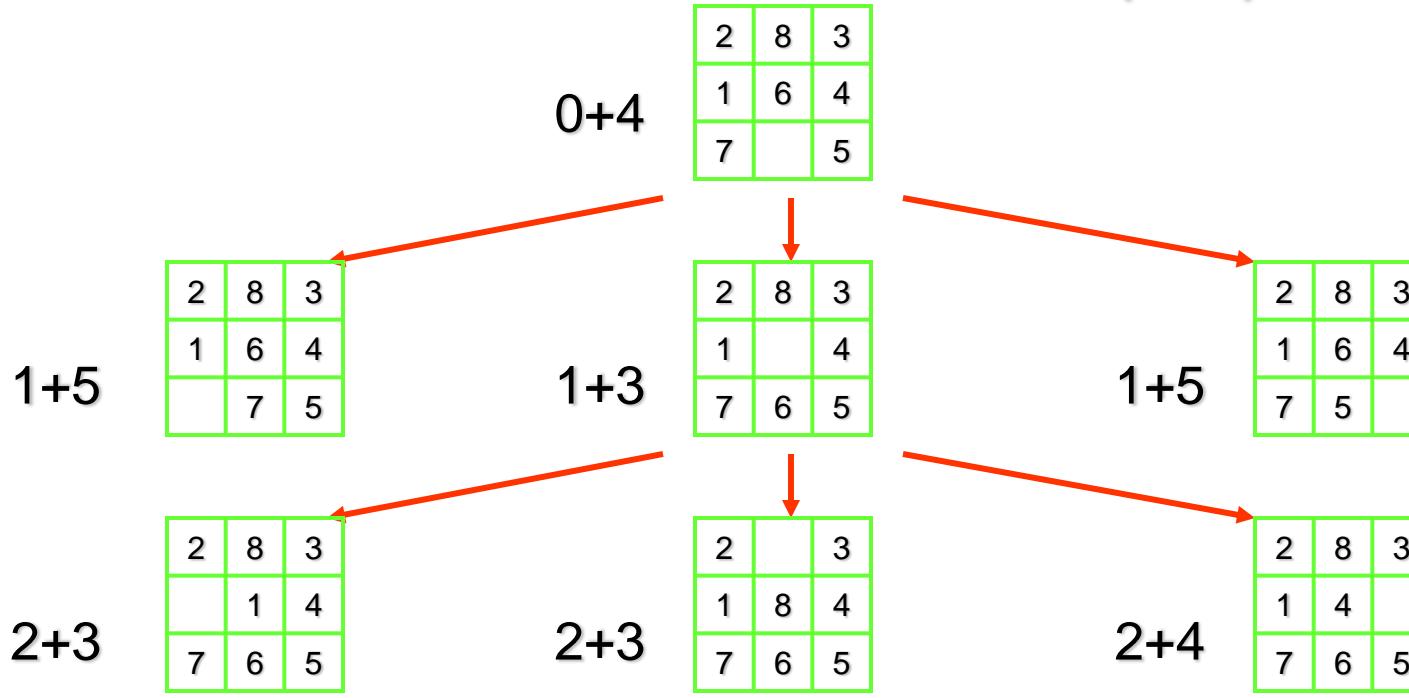
0+4

2	8	3
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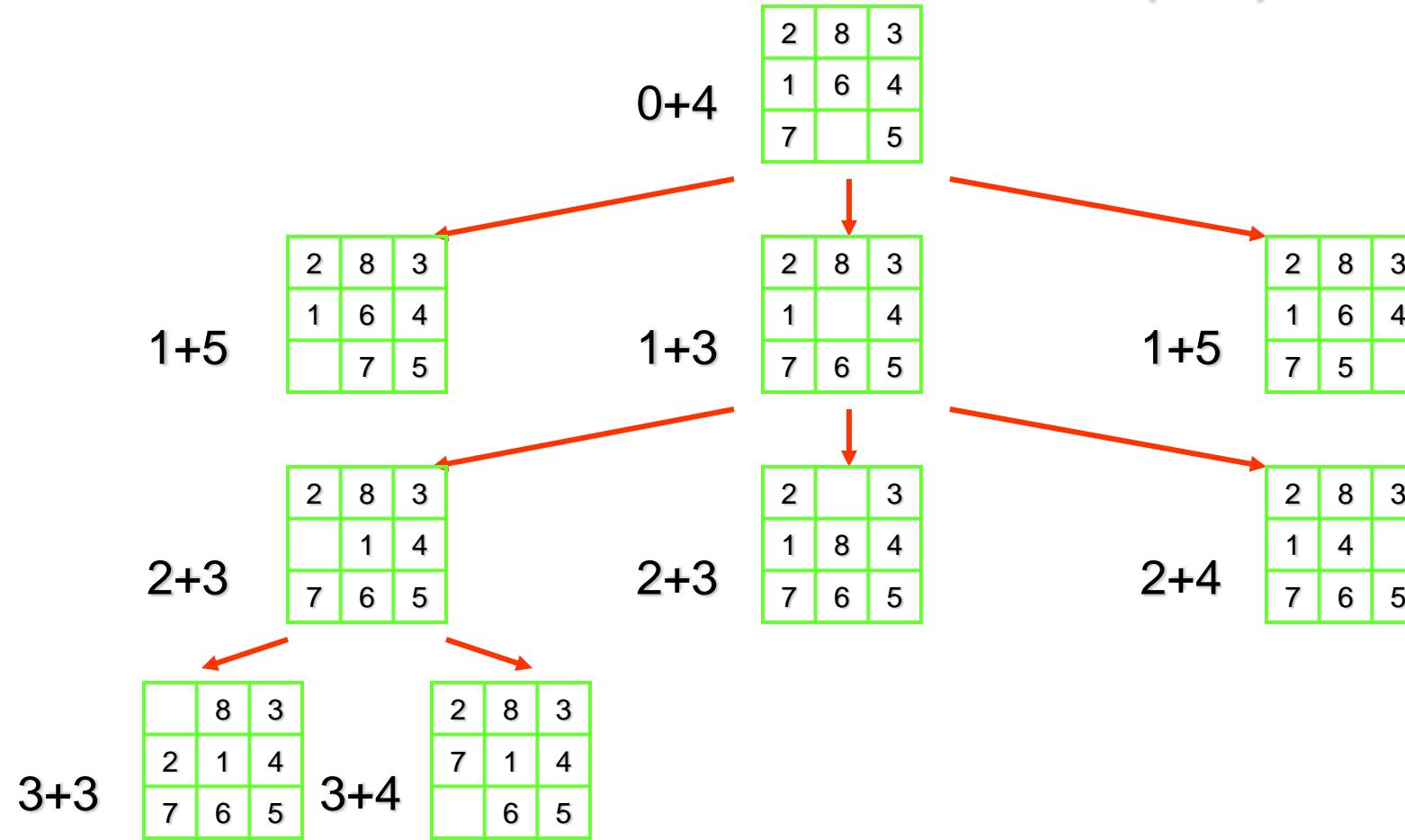
# Best-First Search ( $A^*$ )



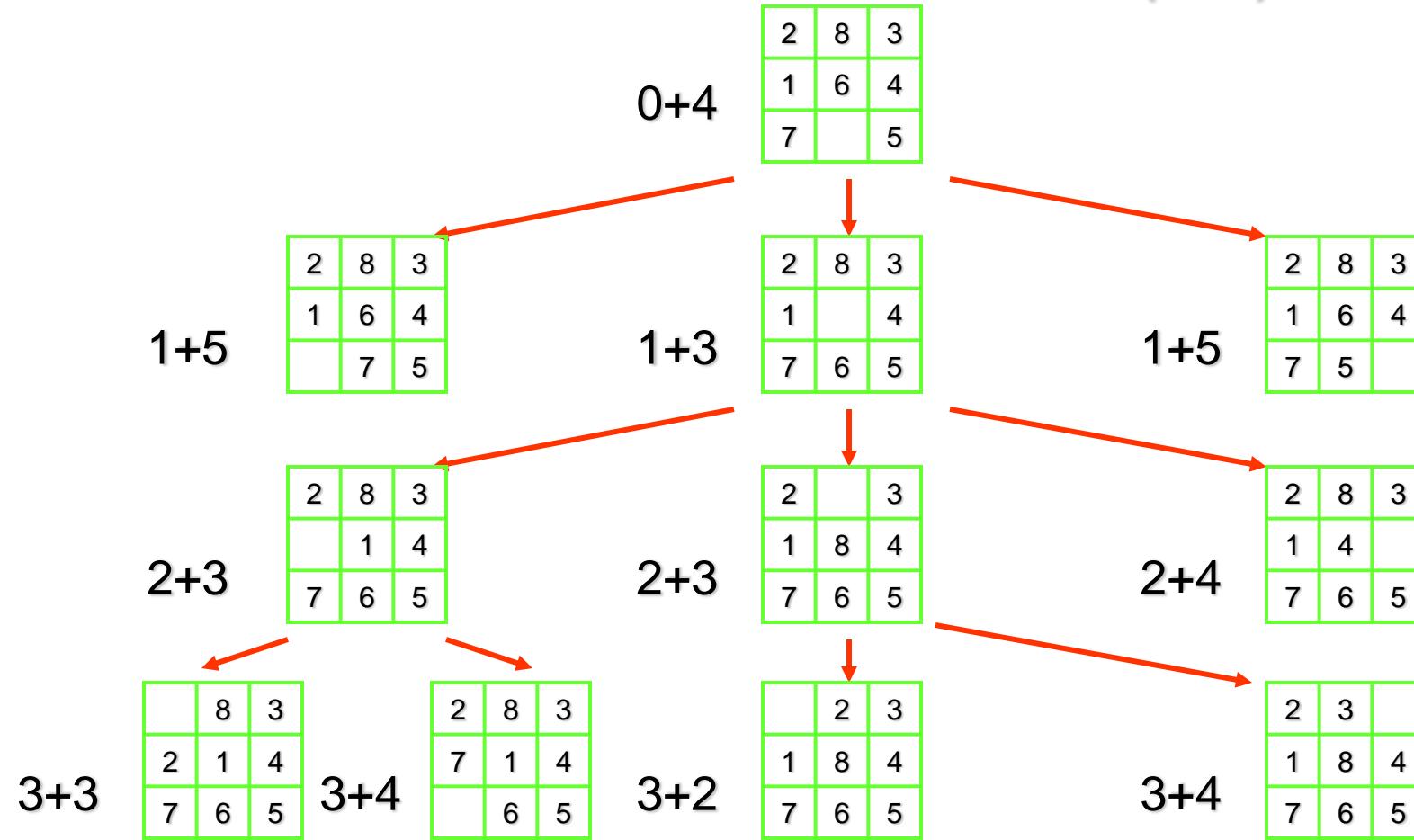
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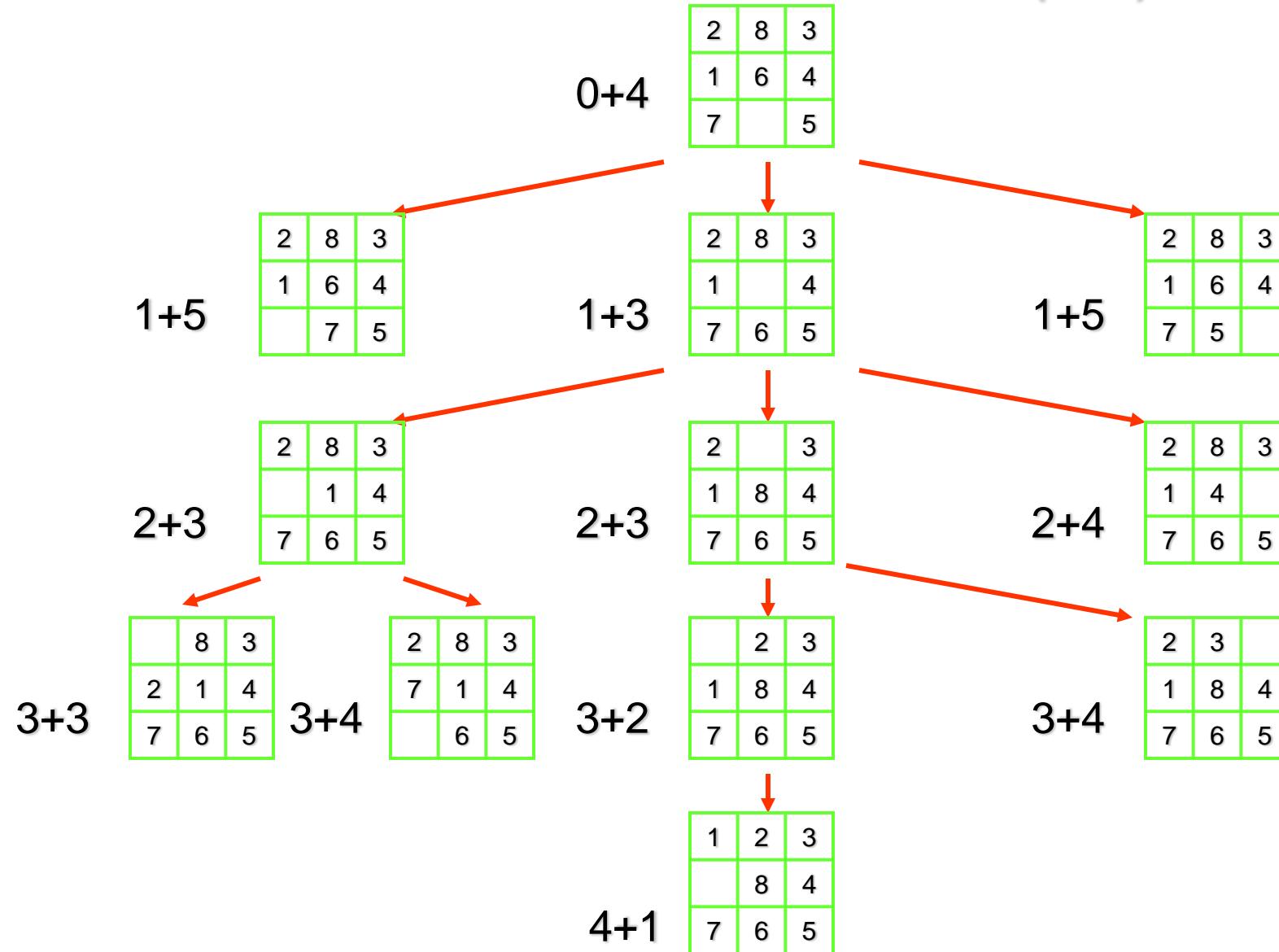
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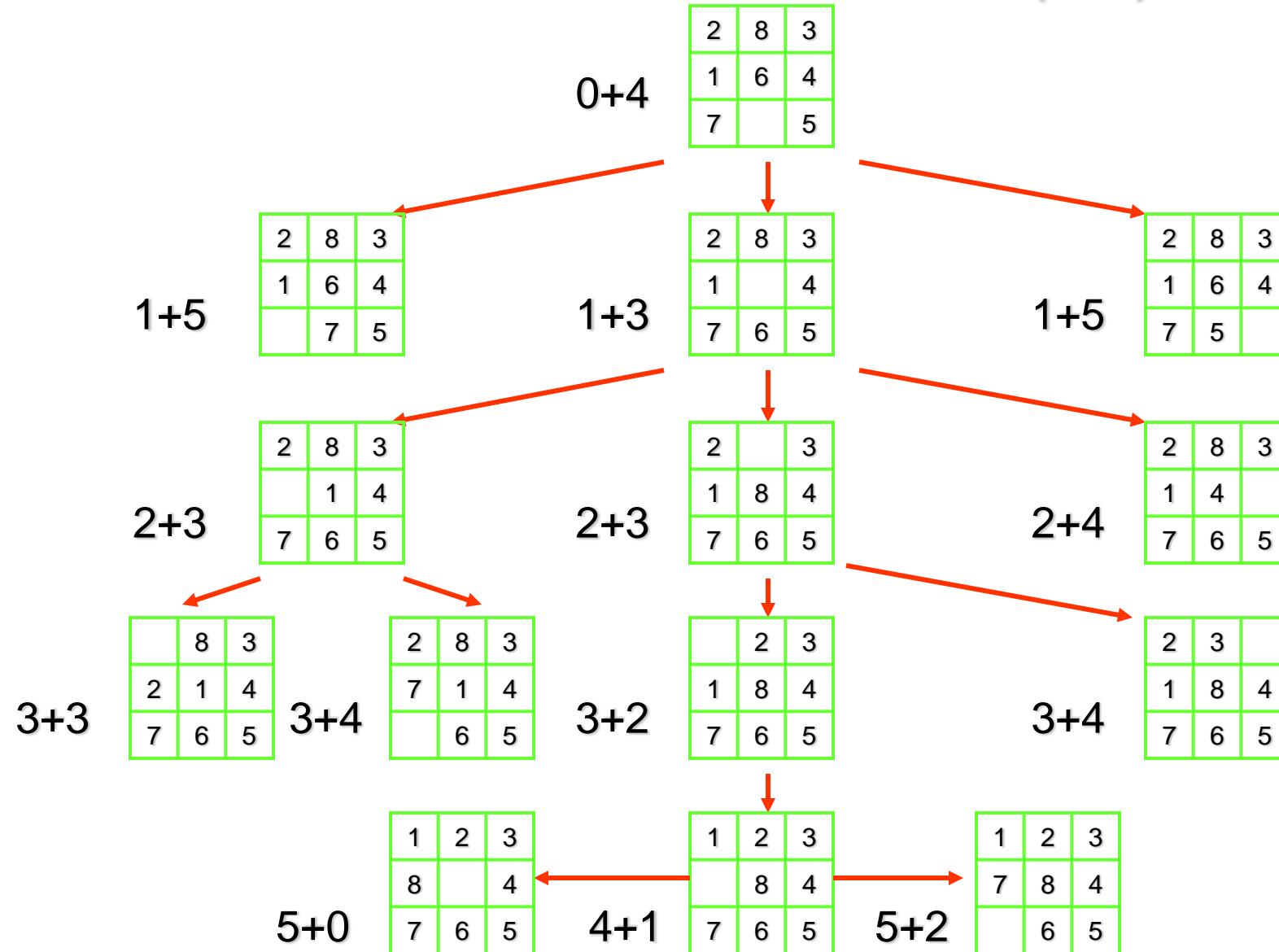
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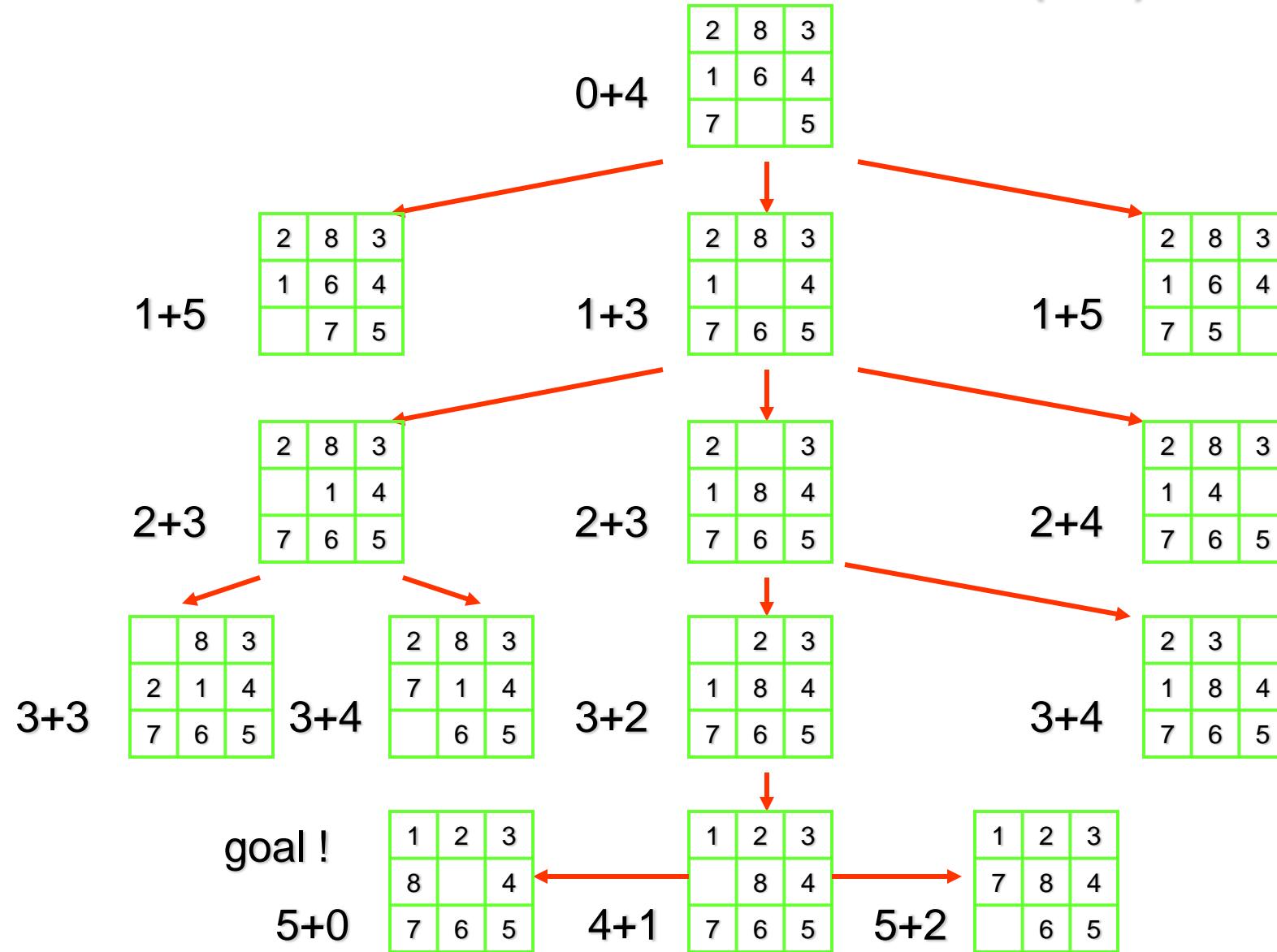
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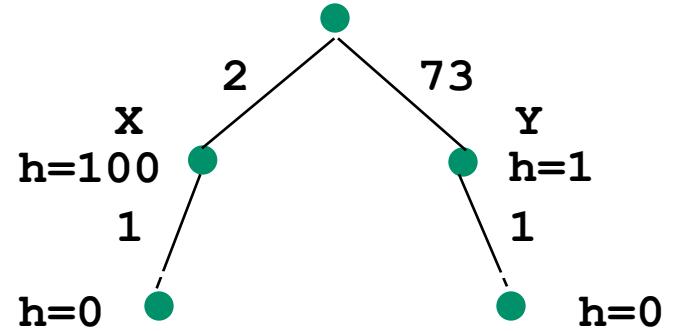
# Admissible search

## Definition

An algorithm is **admissible** if it is guaranteed to return an optimal solution (with minimal possible path weight from the start state to a goal state) whenever a solution exists.

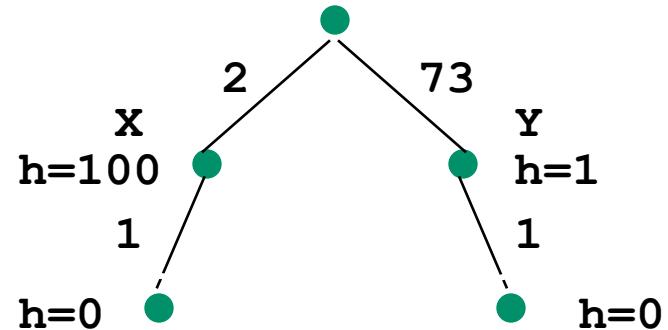
# Admissibility

- What must be true about  $h$  for  $A^*$  to find optimal path?



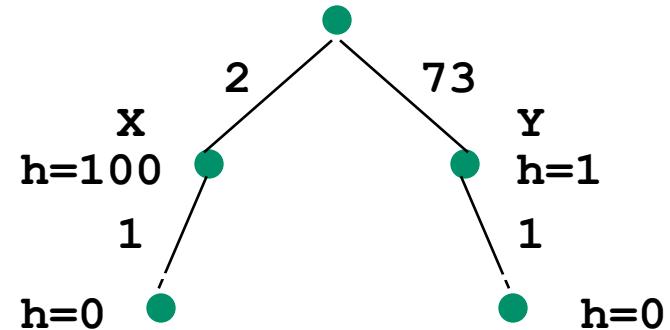
# Admissibility

- What must be true about  $h$  for  $A^*$  to find optimal path?
- $A^*$  finds optimal path if  $h$  is admissible;  $h$  is admissible when it never overestimates.



# Admissibility

- What must be true about  $h$  for  $A^*$  to find optimal path?
- $A^*$  finds optimal path if  $h$  is admissible;  $h$  is admissible when it never overestimates.
- In this example,  $h$  is not admissible.



$$g(X) + h(X) = 102$$

$$g(Y) + h(Y) = 74$$

Optimal path is not found!

# Admissibility of a Heuristic Function

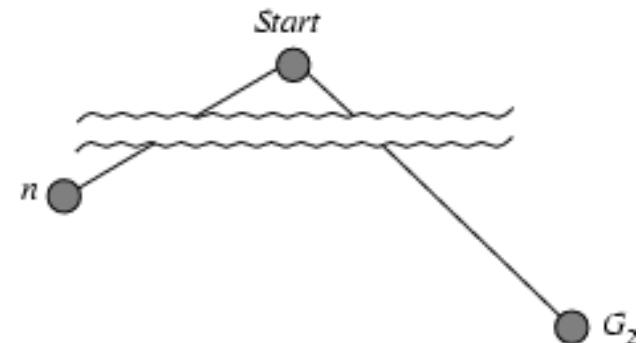
## Definition

A heuristic function  $h$  is said to be **admissible** if

$$0 \leq h(n) \leq h^*(n) \text{ for all } n \in S.$$

# Optimality of A\* (proof)

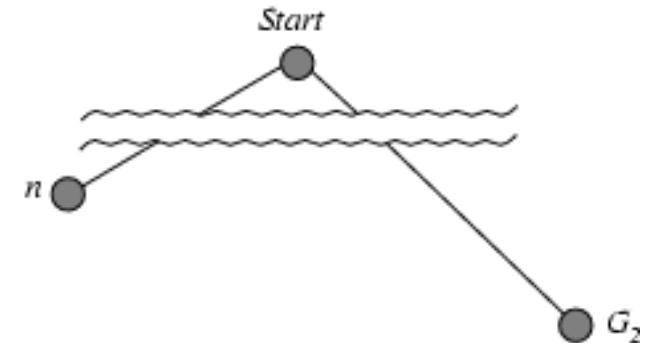
- Suppose some suboptimal goal  $G_2$  has been generated and is in the fringe. Let  $n$  be an unexpanded node in the fringe such that  $n$  is on a shortest path to an optimal goal  $G$ .



- $f(G_2) = g(G_2)$  since  $h(G_2) = 0$
- $f(G) = g(G)$  since  $h(G) = 0$
- $g(G_2) > g(G)$  since  $G_2$  is suboptimal
- $f(G_2) > f(G)$  from above

# Optimality of A\* (proof)

- Suppose some suboptimal goal  $G_2$  has been generated and is in the fringe. Let  $n$  be an unexpanded node in the fringe such that  $n$  is on a shortest path to an optimal goal  $G$ .



- $f(G_2) > f(G)$  from above
- $h(n) \leq h^*(n)$  since  $h$  is admissible
- $g(n) + h(n) \leq g(n) + h^*(n)$
- $f(n) \leq f(G)$ , Hence  $f(G_2) > f(n)$ ,
- and  $A^*$  will never select  $G_2$  for expansion

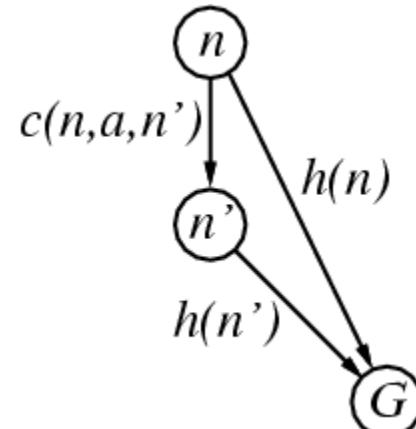
# Consistent heuristics

- A heuristic is consistent if for every node  $n$ , every successor  $n'$  of  $n$  generated by any action  $a$ ,
- 

$$h(n) \leq c(n,a,n') + h(n')$$

If  $h$  is consistent, we have

$$\begin{aligned} f(n') &= g(n') + h(n') \\ &= g(n) + c(n,a,n') + h(n') \\ &\geq g(n) + h(n) \\ &\geq f(n) \end{aligned}$$

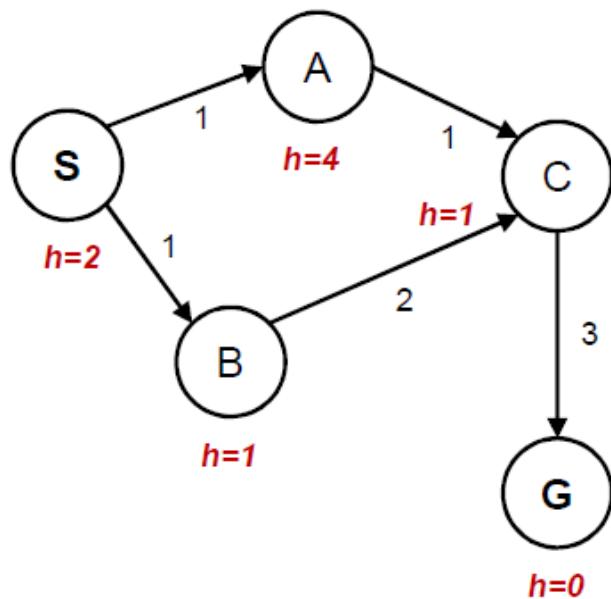


- i.e.,  $f(n)$  is non-decreasing along any path.

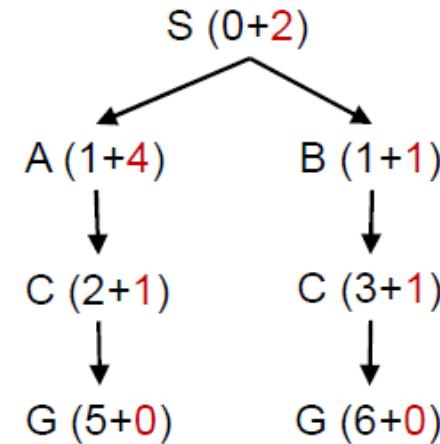
# Consistent heuristics

## A\* Graph Search Gone Wrong?

State space graph



Search tree

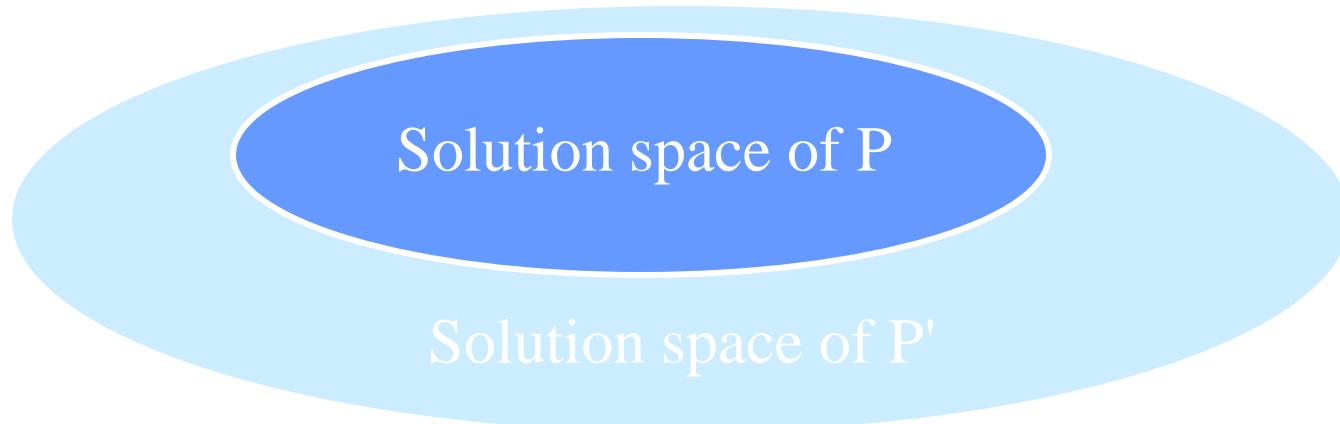


# Extreme Cases

- If we set  $h=0$ , then  $A^*$  is Breadth First search;  $h=0$  is admissible heuristic (when all costs are non-negative).
- If  $g=0$  for all nodes,  $A^*$  is similar to Steepest Ascend Hill Climbing
- If both  $g$  and  $h =0$ , it can be as dfs or bfs depending on the structure of Open

# How to chose a heuristic?

- ❖ Original problem P
  - A set of constraints
  - P is complex
- ❖ Relaxed problem P'
  - removing one or more constraints
  - P' becomes simpler
- ❖ Use cost of a best solution path from n in P' as h(n) for P
- ❖ **Admissibility:**
$$h^* \leq h$$
$$\text{cost of best solution in } P \geq \text{cost of best solution in } P'$$

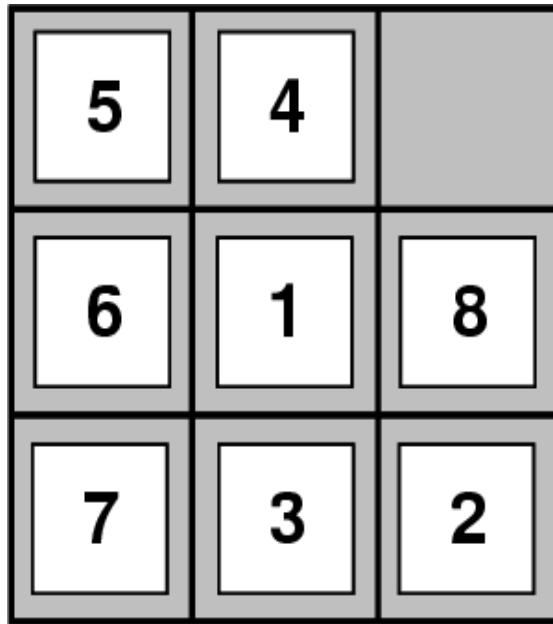


# How to chose a heuristic - 8-puzzle

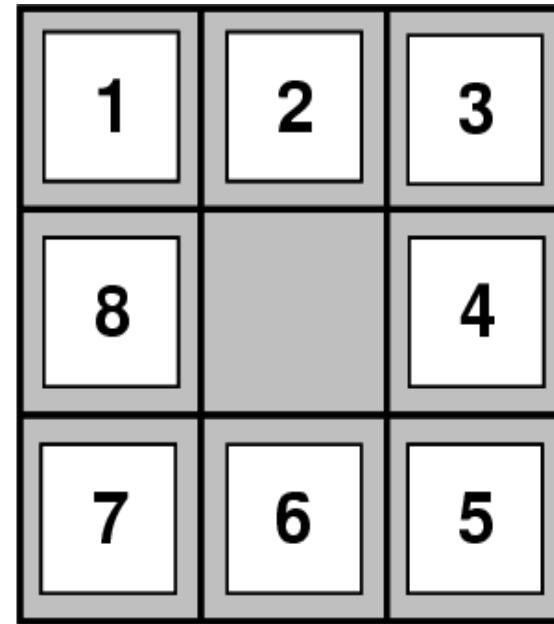
- ❖ Example: 8-puzzle
  - Constraints: to move from cell A to cell B
    - cond1: there is a tile on A
    - cond2: cell B is empty
    - cond3: A and B are adjacent (horizontally or vertically)
  - Removing cond2:
    - h2 (sum of Manhattan distances of all misplaced tiles)
  - Removing both cond2 and cond3:
    - h1 (# of misplaced tiles)

Here  $h_2 \geq h_1$

# How to chose a heuristic - 8-puzzle



Start State



Goal State

$$h_1(\text{start}) = 7$$

$$h_2(\text{start}) = 18$$

# Performance comparison

d	Search Cost (nodes)			Effective Branching Factor		
	IDS	A* ( $h_1$ )	A* ( $h_2$ )	IDS	A* ( $h_1$ )	A* ( $h_2$ )
2	10	6	6	2.45	1.79	1.79
4	112	13	12	2.87	1.48	1.45
6	680	20	18	2.73	1.34	1.30
8	6384	39	25	2.80	1.33	1.24
10	47127	93	39	2.79	1.38	1.22
12	3644035	227	73	2.78	1.42	1.24
14	-	539	113	-	1.44	1.23
16	-	1301	211	-	1.45	1.25
18	-	3056	363	-	1.46	1.26
20	-	7276	676	-	1.47	1.27
22	-	18094	1219	-	1.48	1.28
24	-	39135	1641	-	1.48	1.26

Data are averaged over 100 instances of 8-puzzle, for various solution lengths. Note: there are still better heuristics for the 8-puzzle.

# Comparing heuristic functions

- Bad estimates of the remaining distance can cause extra work!
- Given two algorithms  $A_1$  and  $A_2$  with admissible heuristics  $h_1$  and  $h_2 < h^*(n)$ ,
- if  $h_1(n) < h_2(n)$  for all non-goal nodes  $n$ , then  $A_1$  expands at least as many nodes as  $A_2$
- $h_2$  is said to be more **informed** than  $h_1$  or,  $A_2$  **dominates**  $A_1$

# Weighted A\* (WA\*)

$$f_w(n) = (1-w)g(n) + w h(n)$$

w = 0 - Breadth First

w = 1/2 - A\*

w = 1 - Best First

# IDA\*: Iterative deepening A\*

- To reduce the memory requirements at the expense of some additional computation time, combine uninformed iterative deepening search with A\*
- Use an *f*-cost limit instead of a depth limit

# Iterative Deepening A\*

- In the first iteration, we determine a “**f-cost limit**”  $f'(n_0) = g'(n_0) + h'(n_0) = h'(n_0)$ , where  $n_0$  is the start node.
- We expand nodes using the **depth-first algorithm** and backtrack whenever  $f'(n)$  for an expanded node  $n$  exceeds the cut-off value.
- If this search does not succeed, determine the **lowest f'-value** among the nodes that were visited but not expanded.
- Use this f'-value as the **new limit value** and do another depth-first search.
- Repeat this procedure until a goal node is found.

# IDA\* Algorithm - Top level

```
function IDA*(problem) returns solution
    root := Make-Node(Initial-State[problem])
    f-limit := f-Cost(root)
loop do
    solution, f-limit := DFS-Contour(root, f-limit)
    if solution is not null, then return solution
    if f-limit = infinity, then return failure
end
```

# IDA\* contour expansion

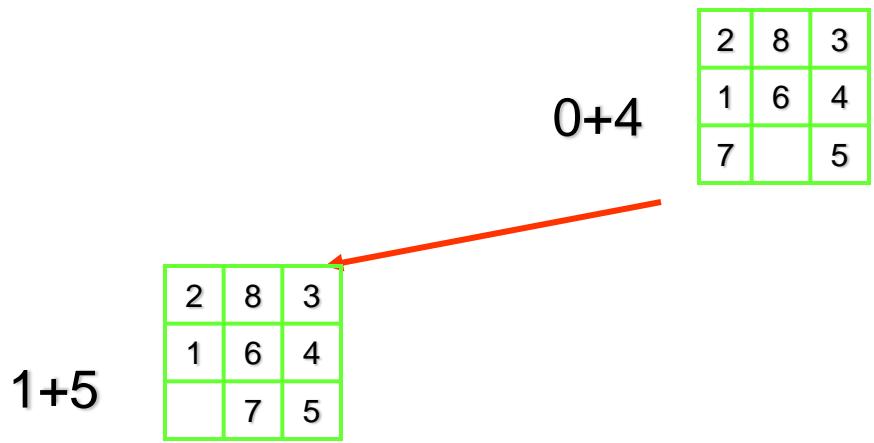
```
function DFS-Countour(node,f-limit) returns
    solution sequence and new f-limit
if f-Cost[node] > f-limit then return (null, f-Cost[node])
if Goal-Test[problem](State[node]) then return (node,f-limit)
for each node s in Successor(node) do
    solution, new-f := DFS-Contour(s, f-limit)
    if solution is not null, then return (solution, f-limit)
    next-f := Min(next-f, new-f);
end
return (null, next-f)
```

# IDA\*: Iteration 1, Limit = 4

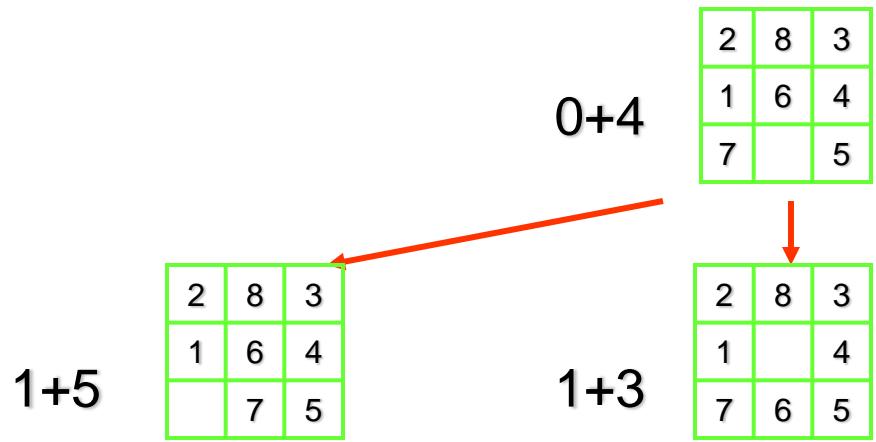
0+4

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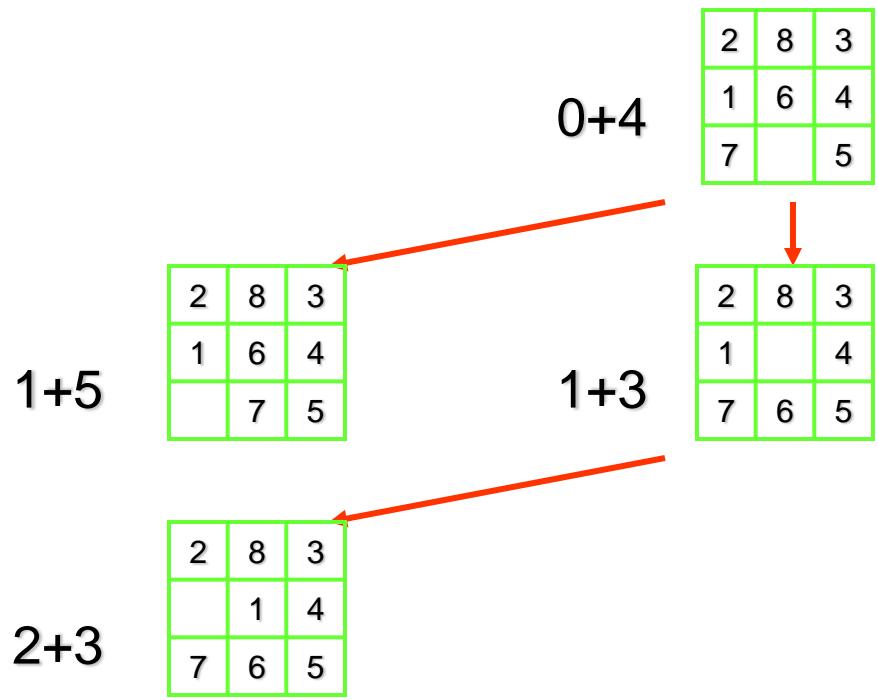
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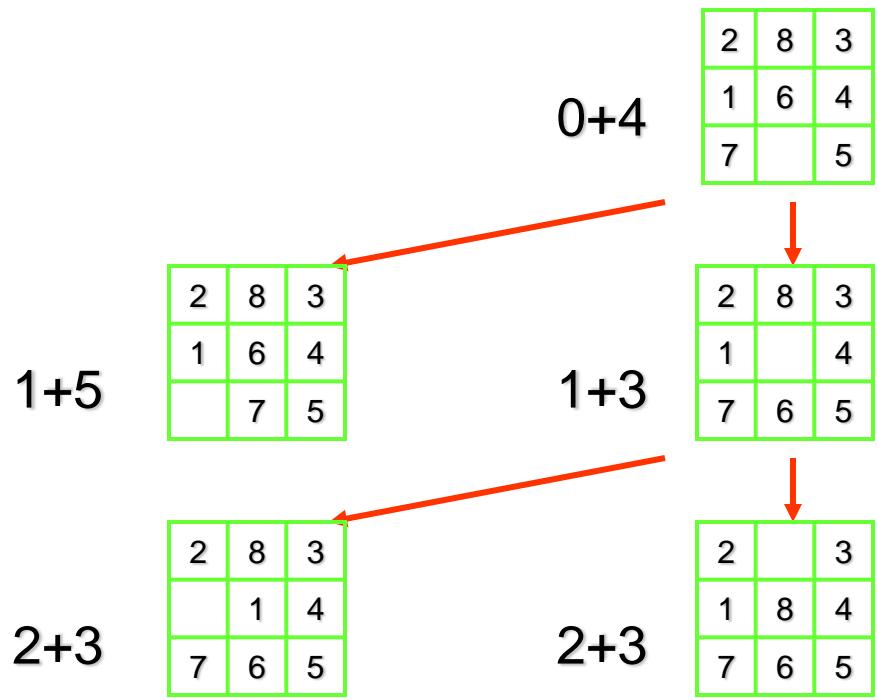
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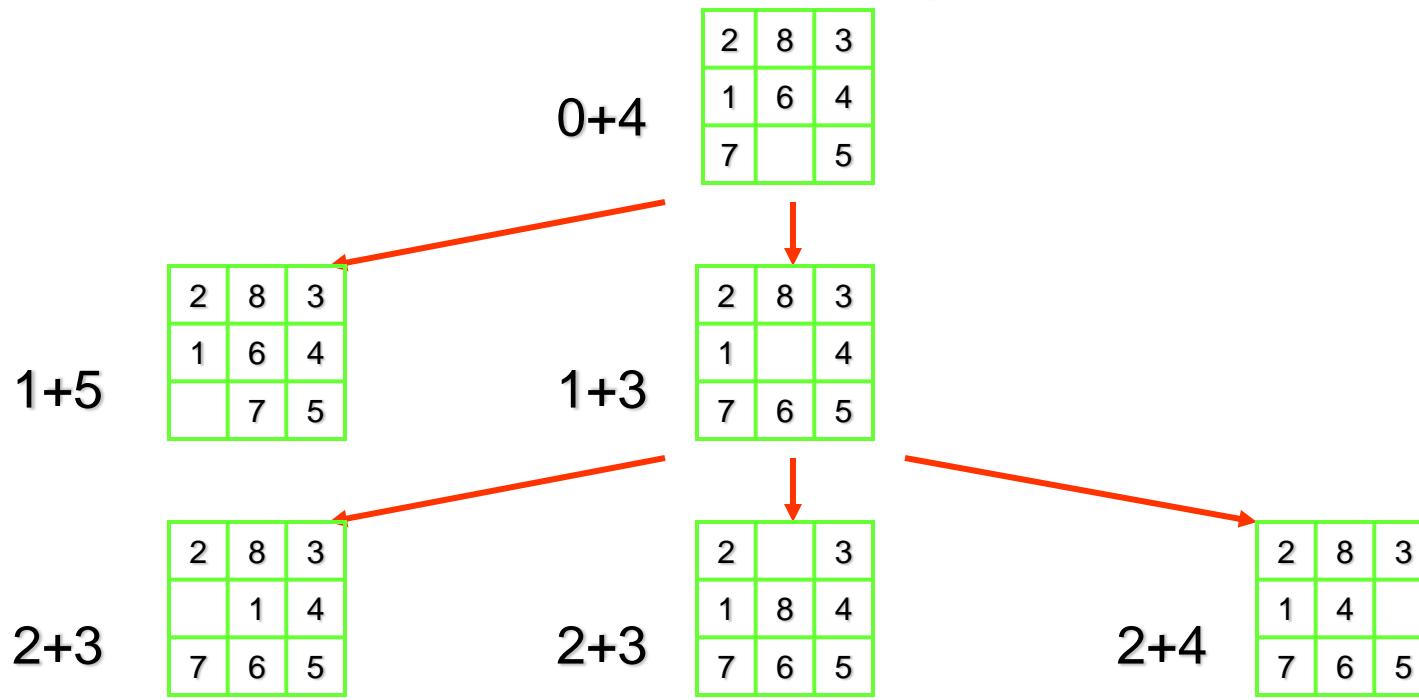
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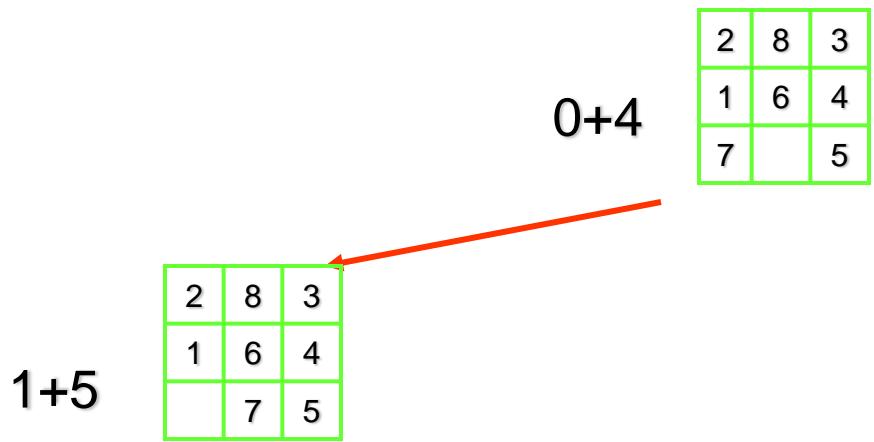


# IDA\*: Iteration 2, Limit = 5

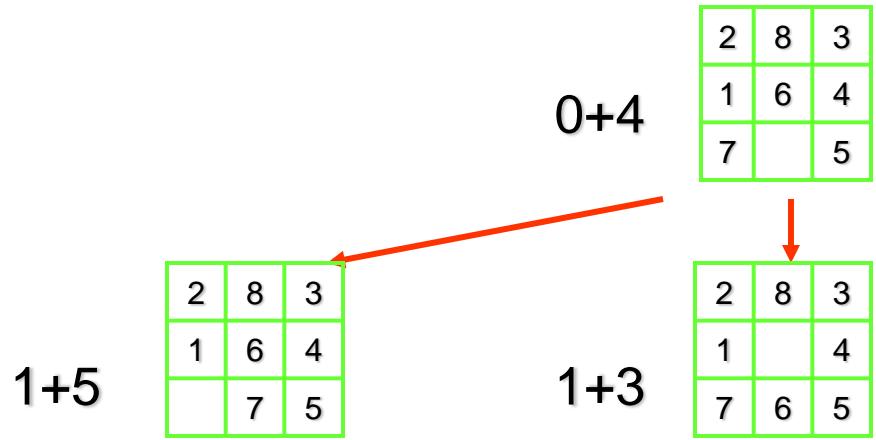
0+4

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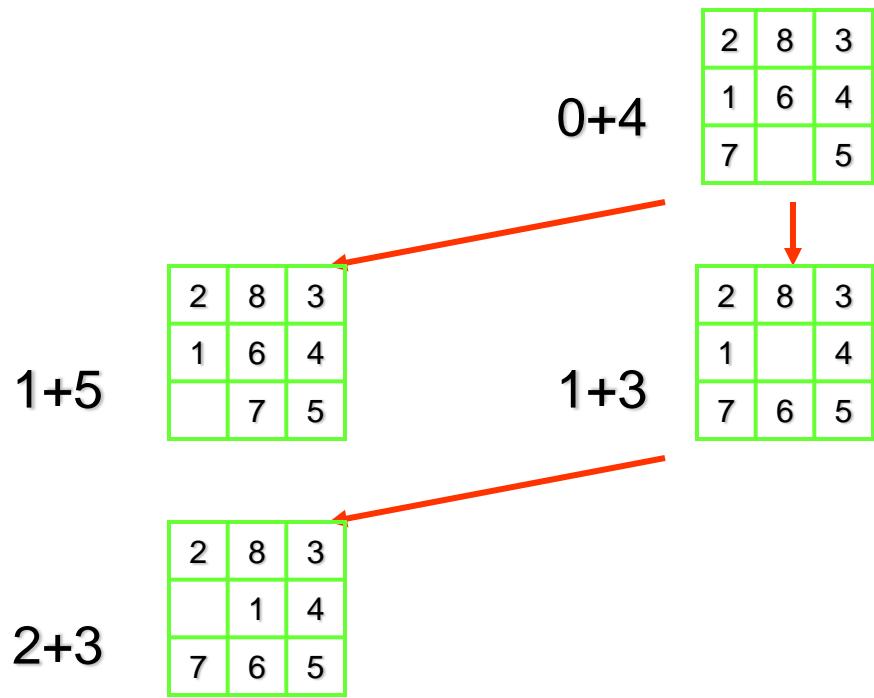
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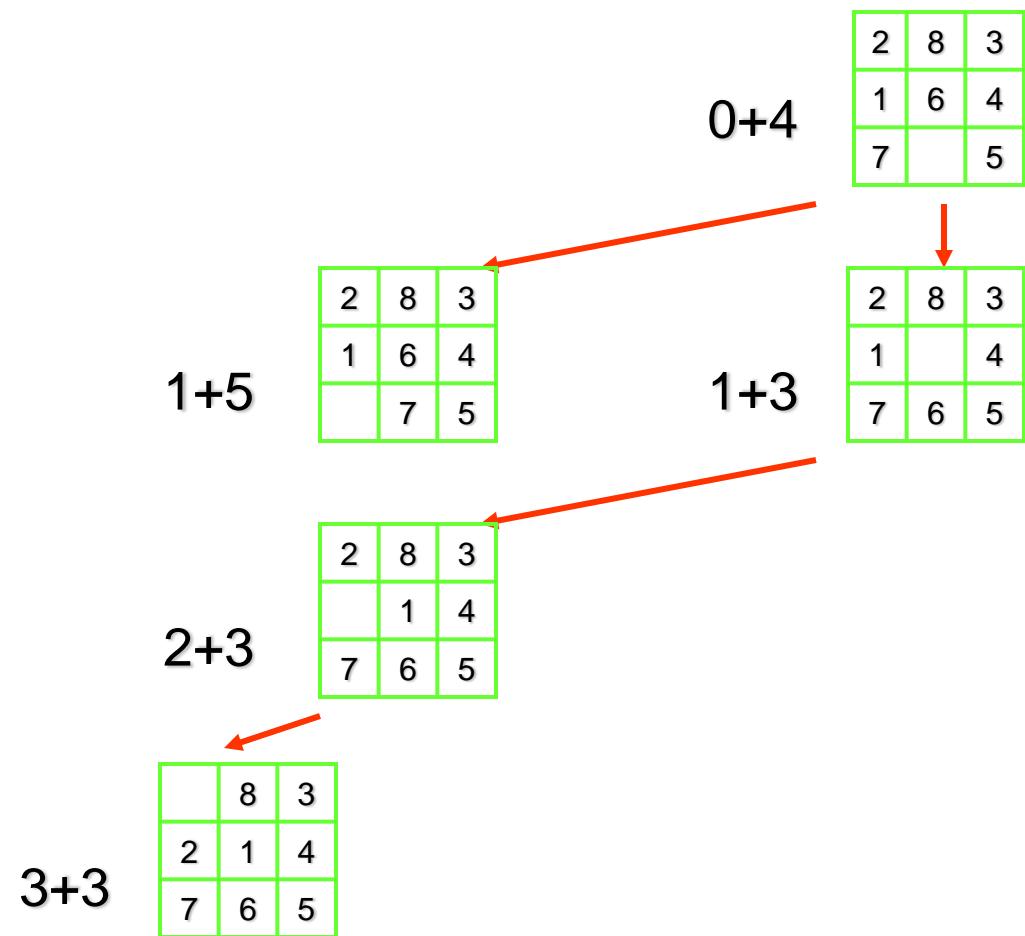
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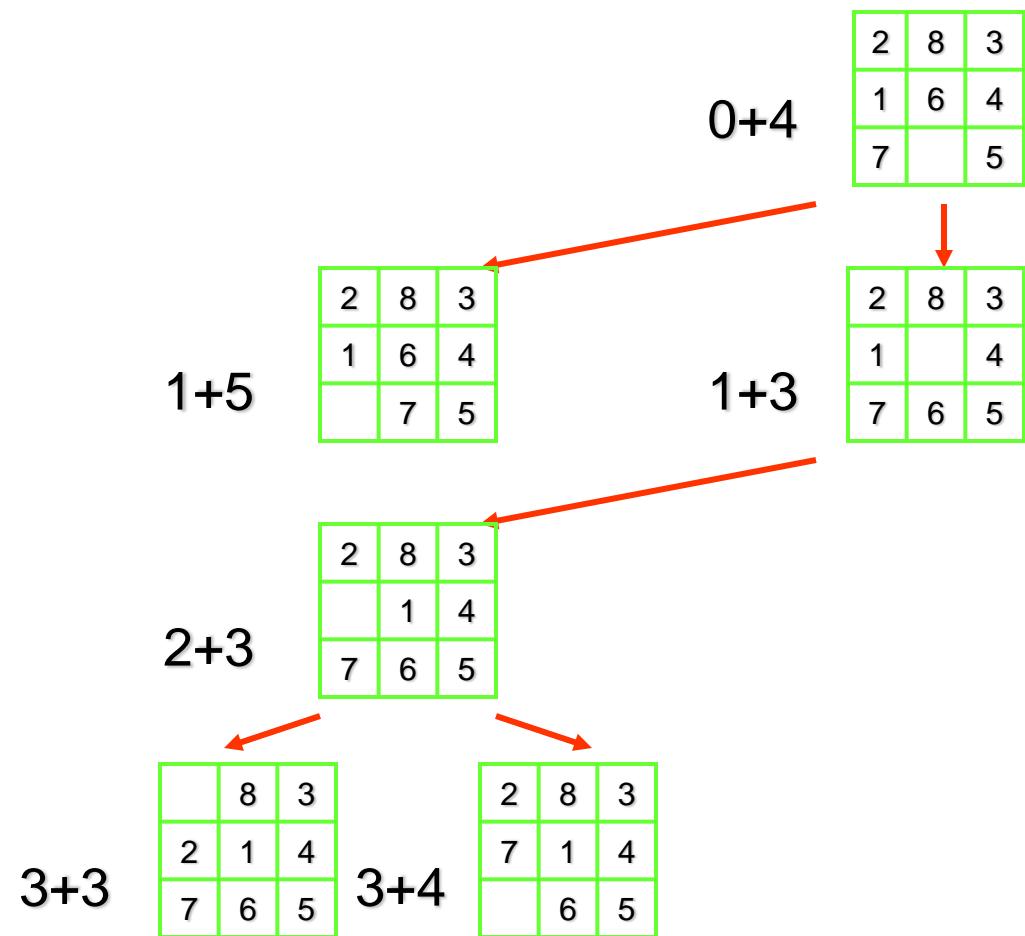
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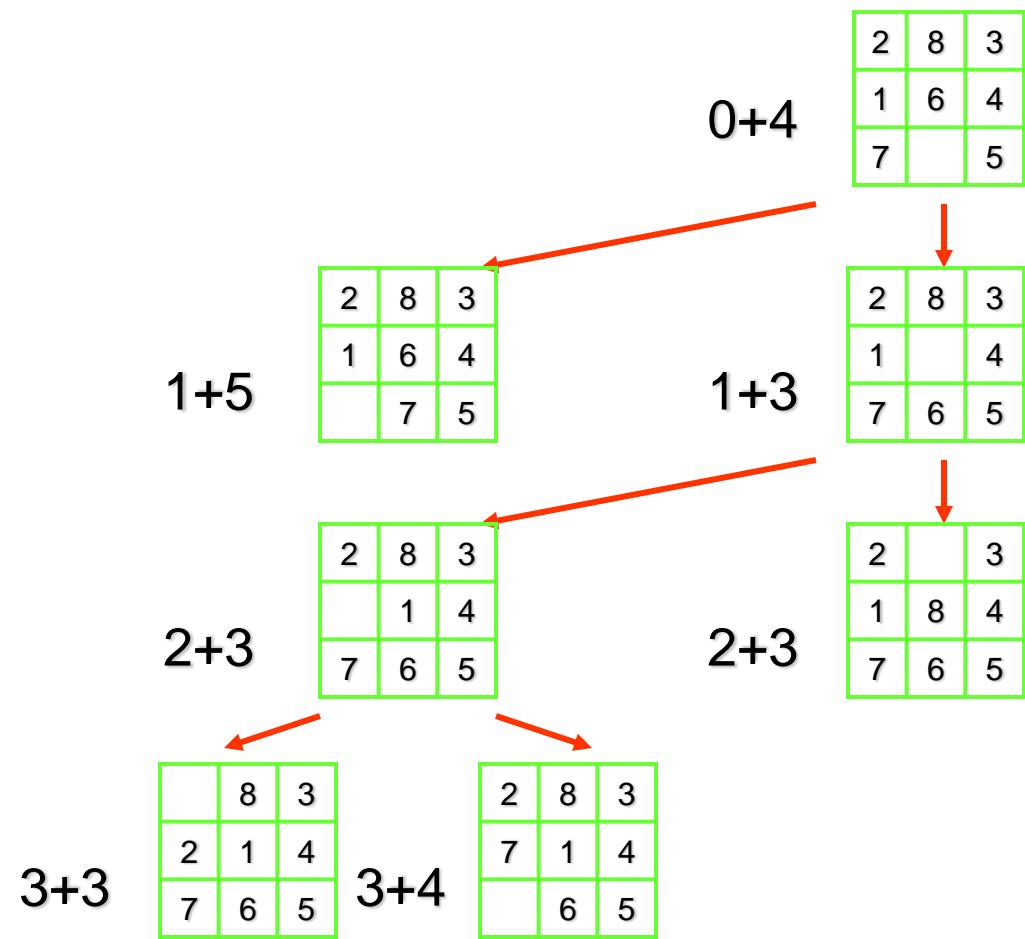
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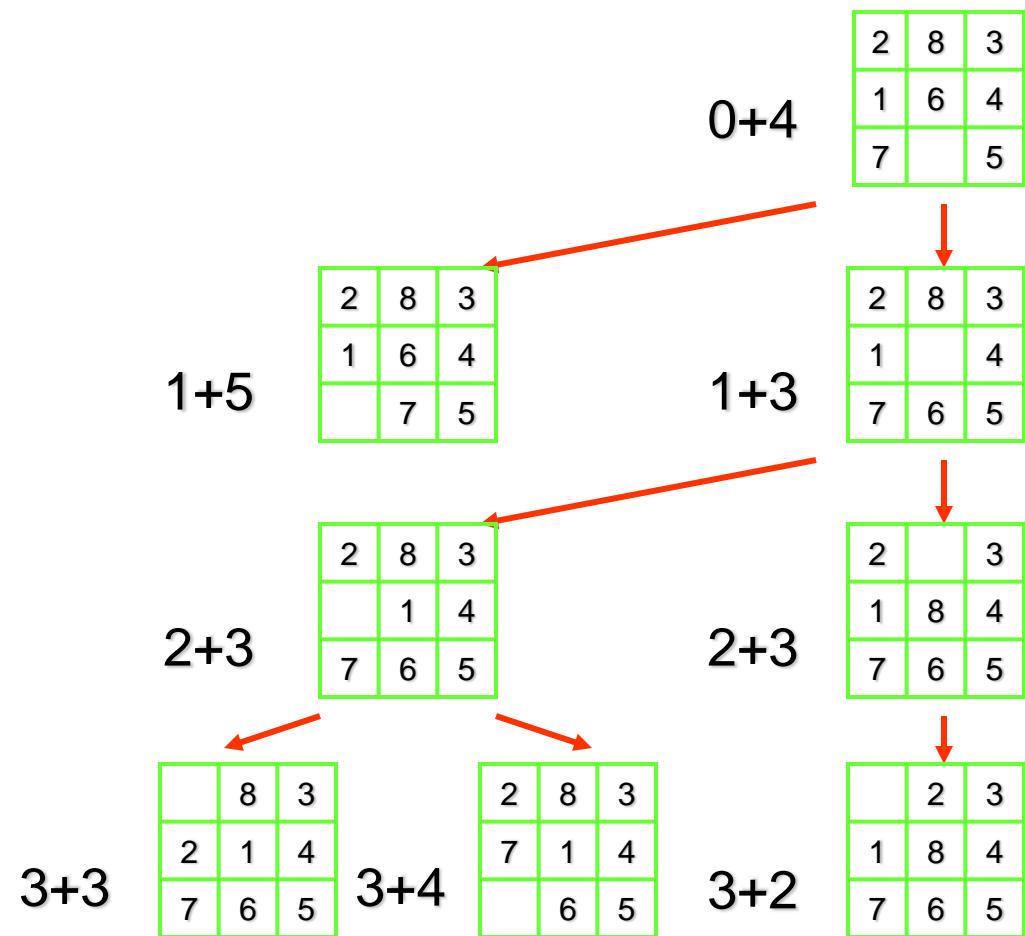
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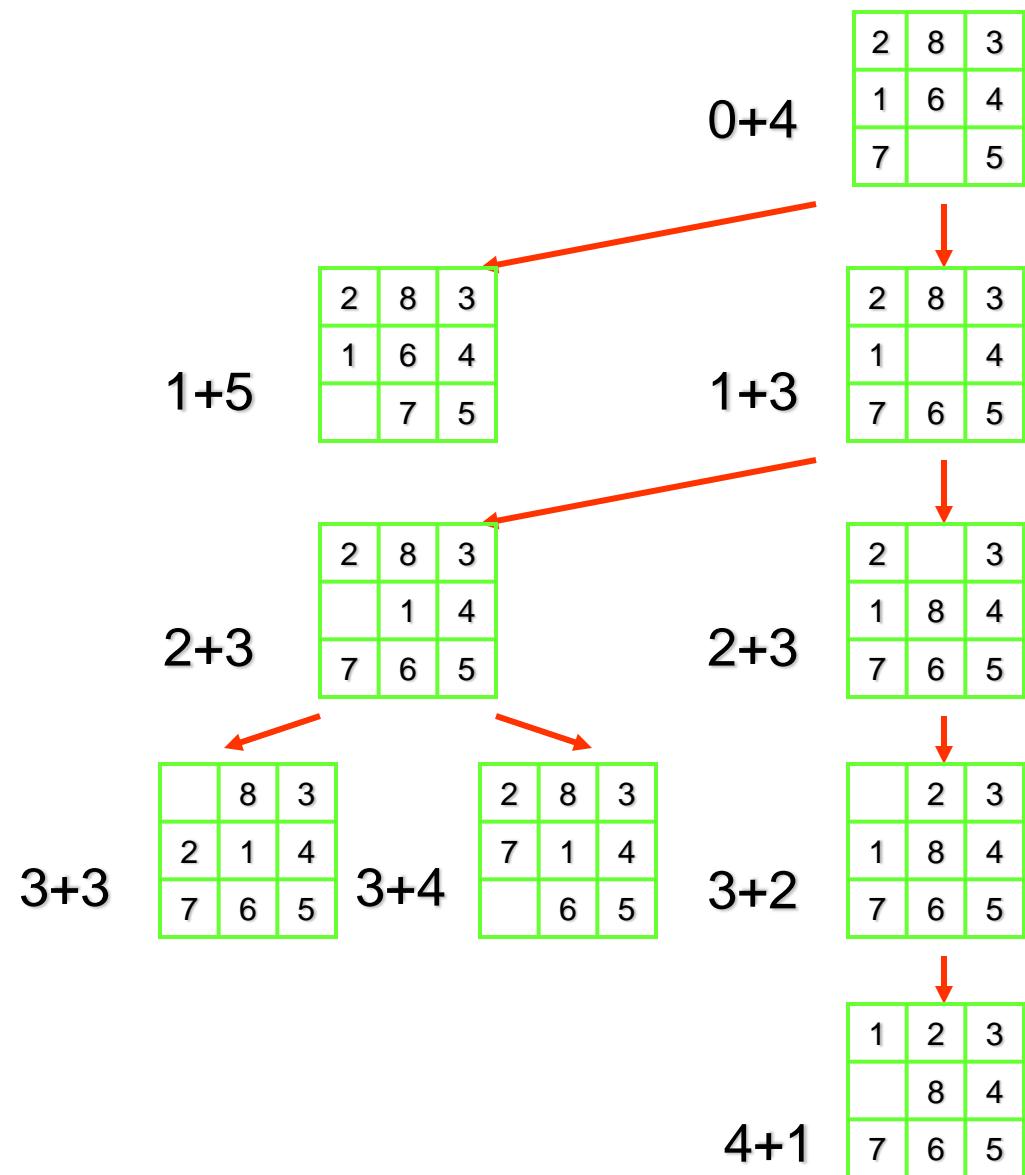
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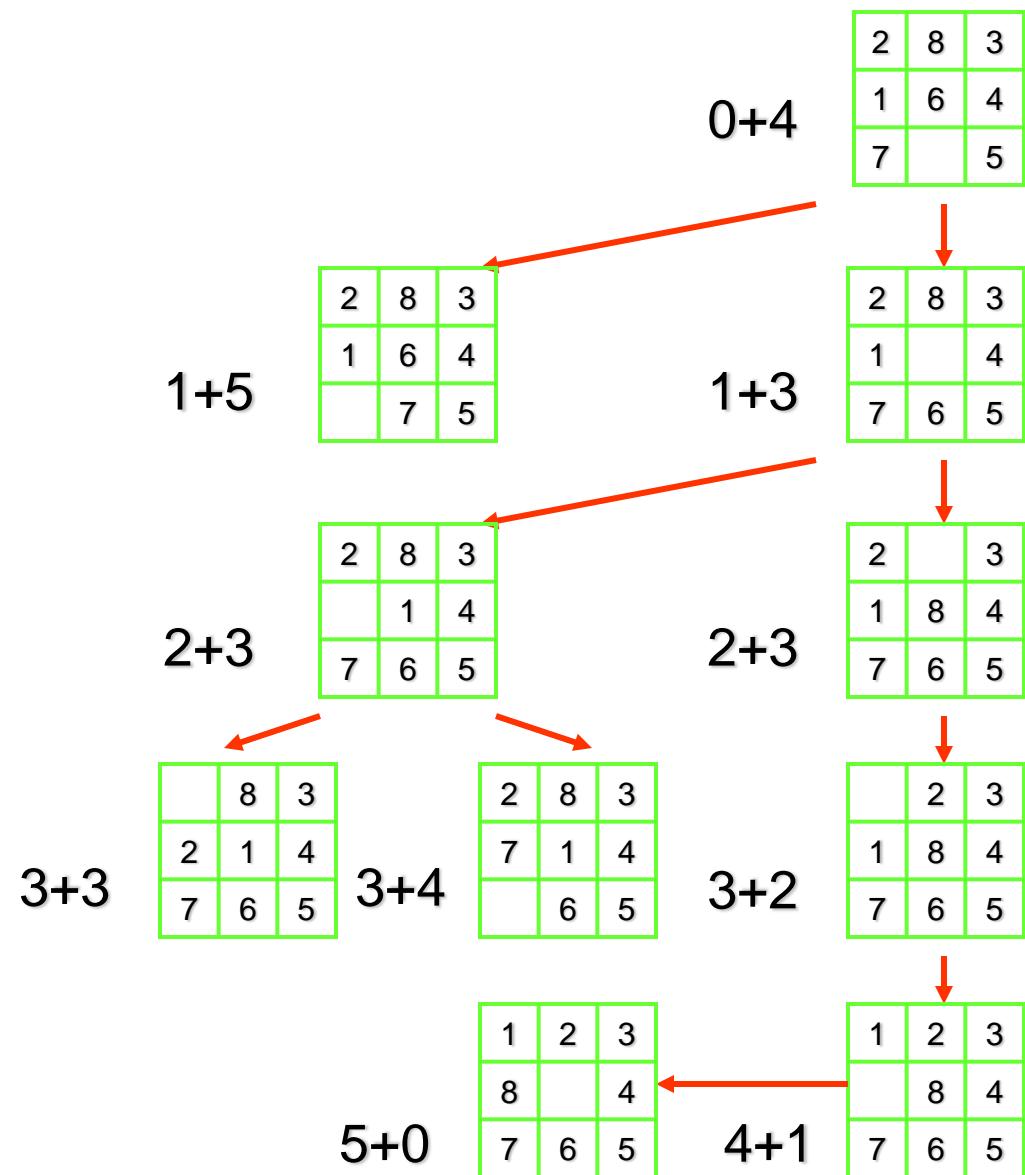
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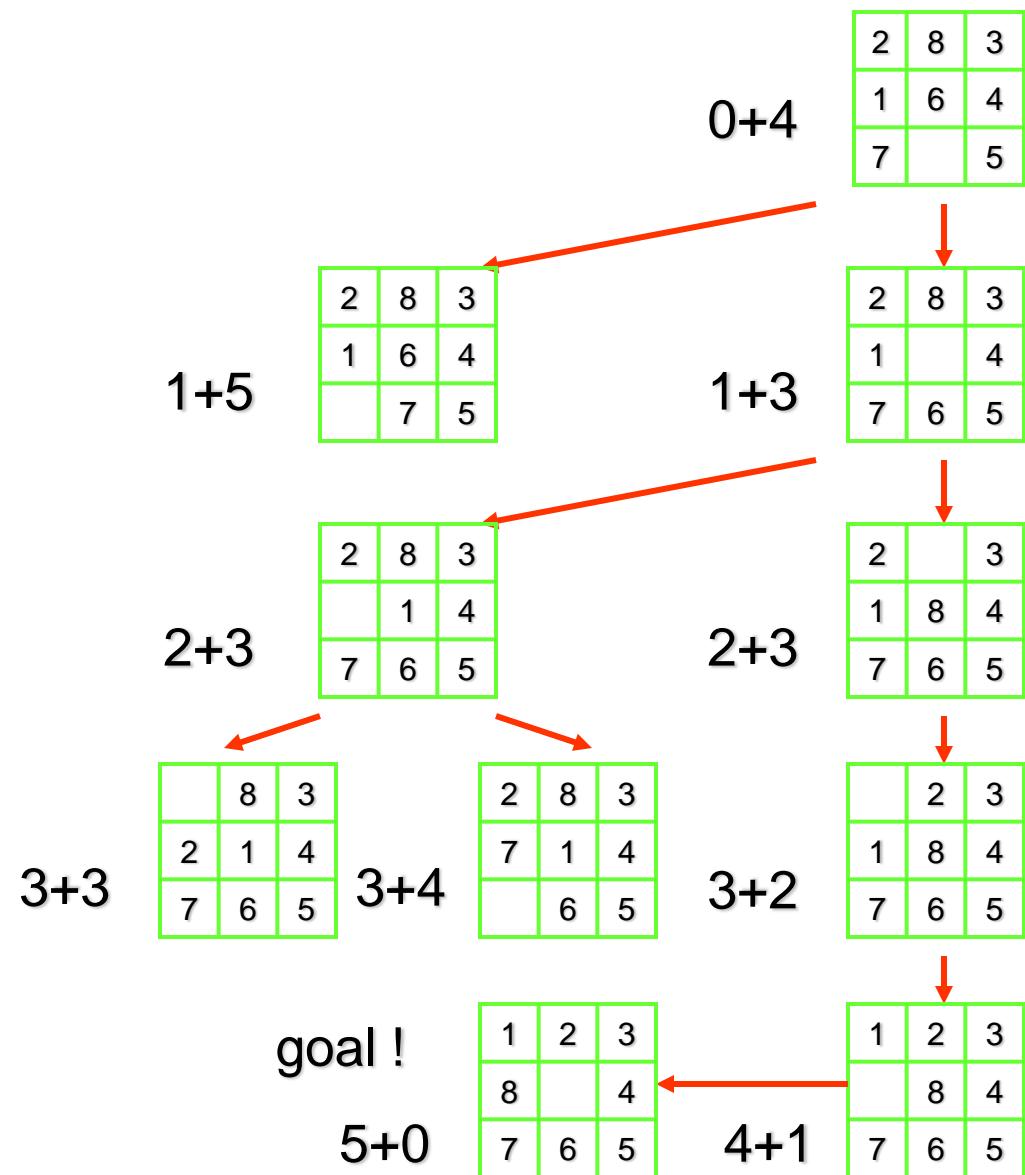
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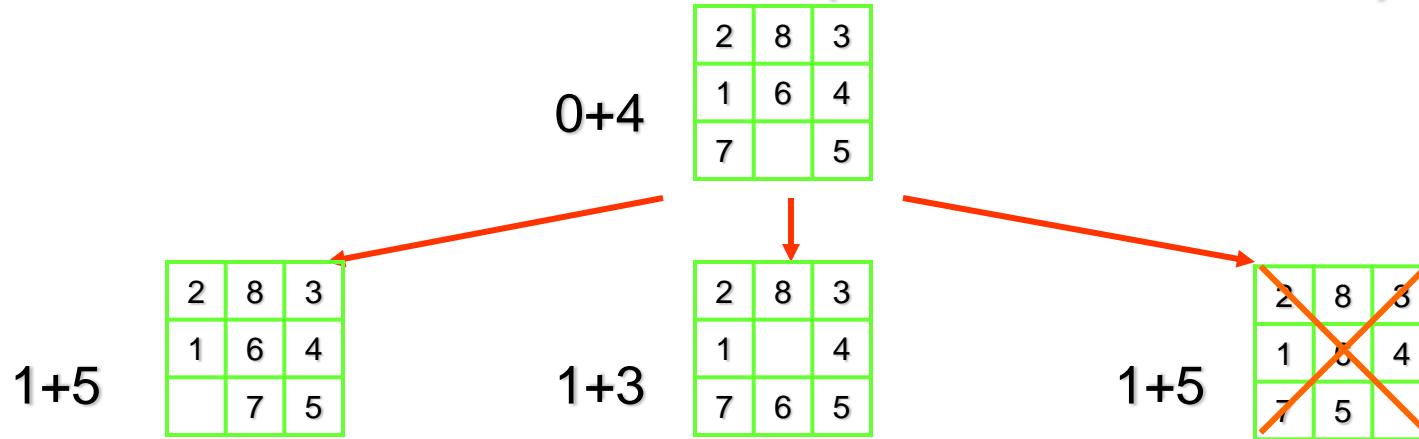


# Beam Search (with limit = 2)

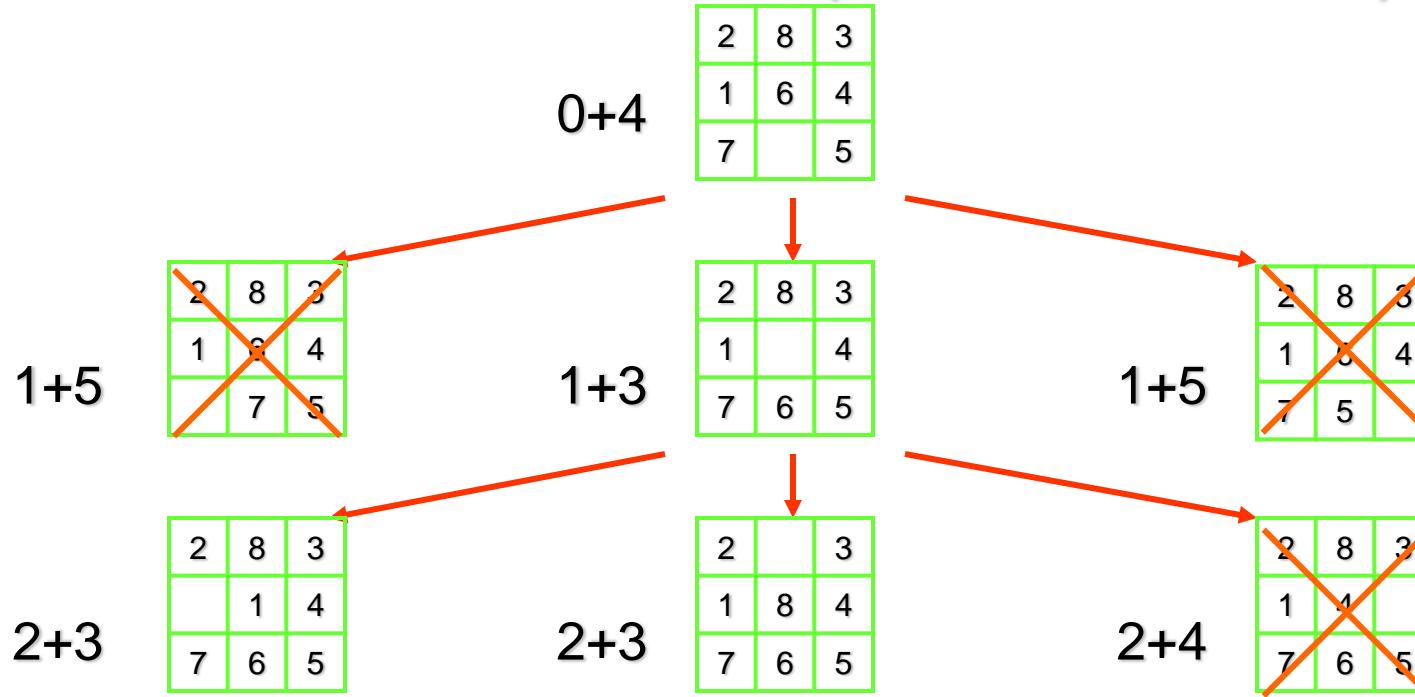
0+4

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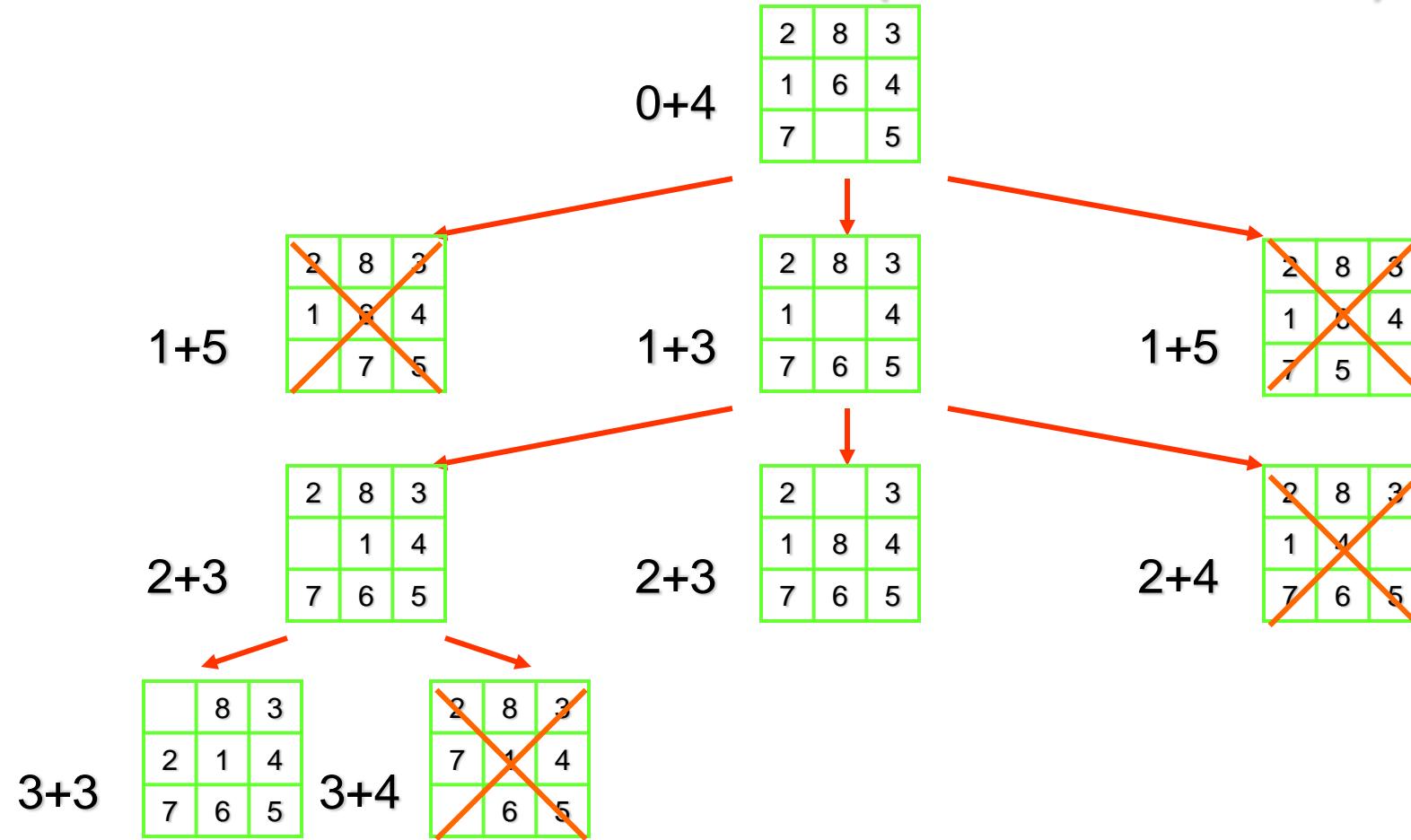
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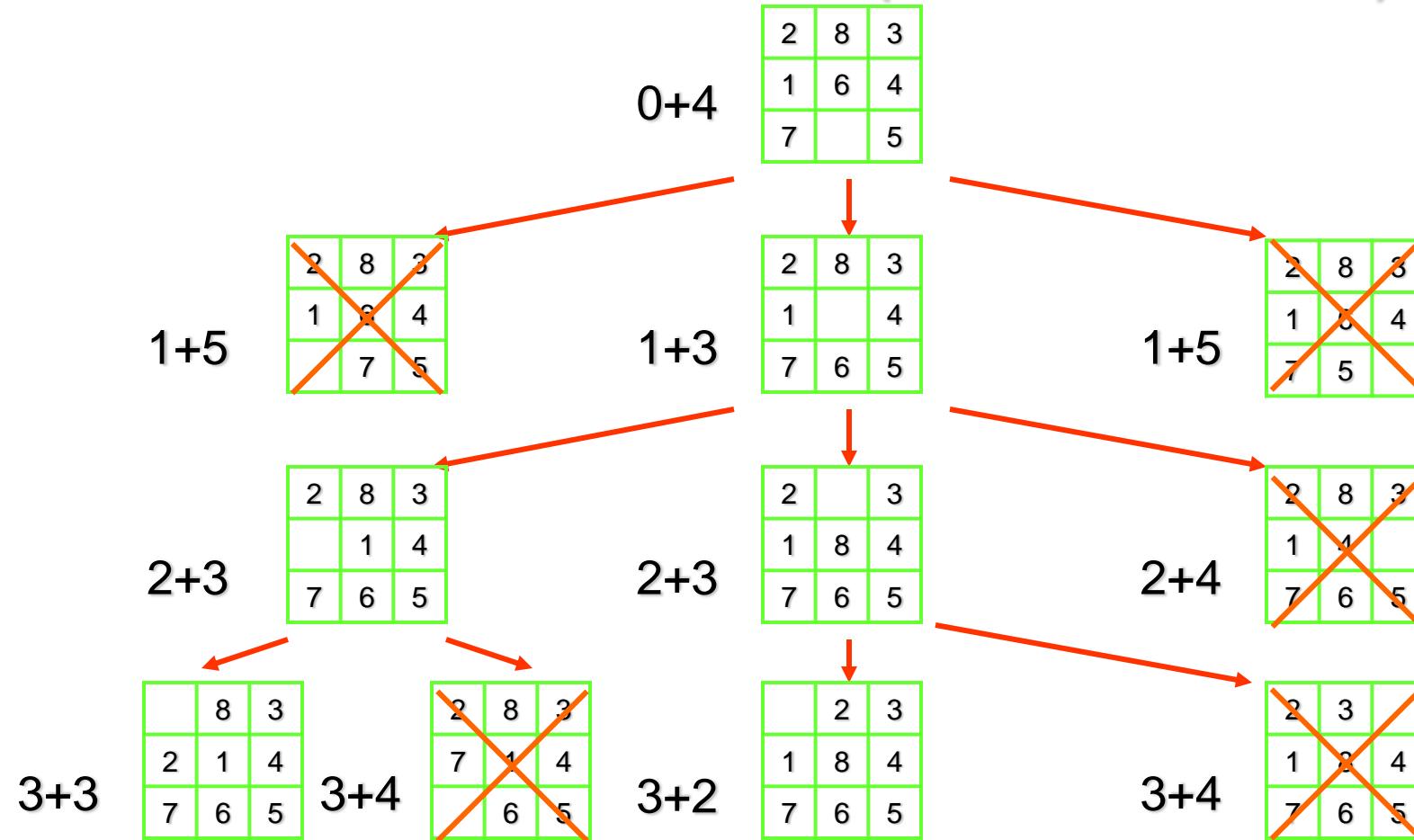
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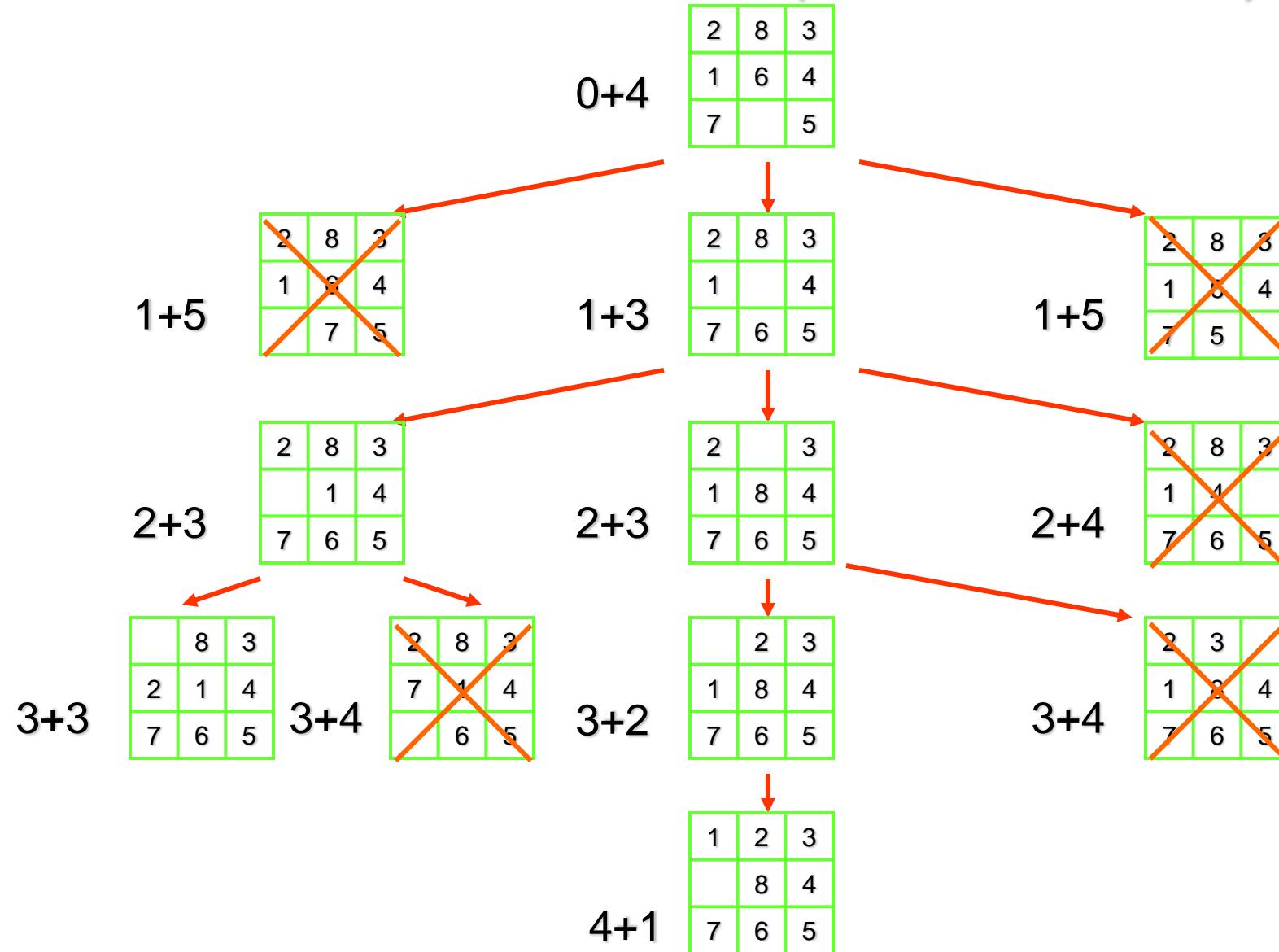
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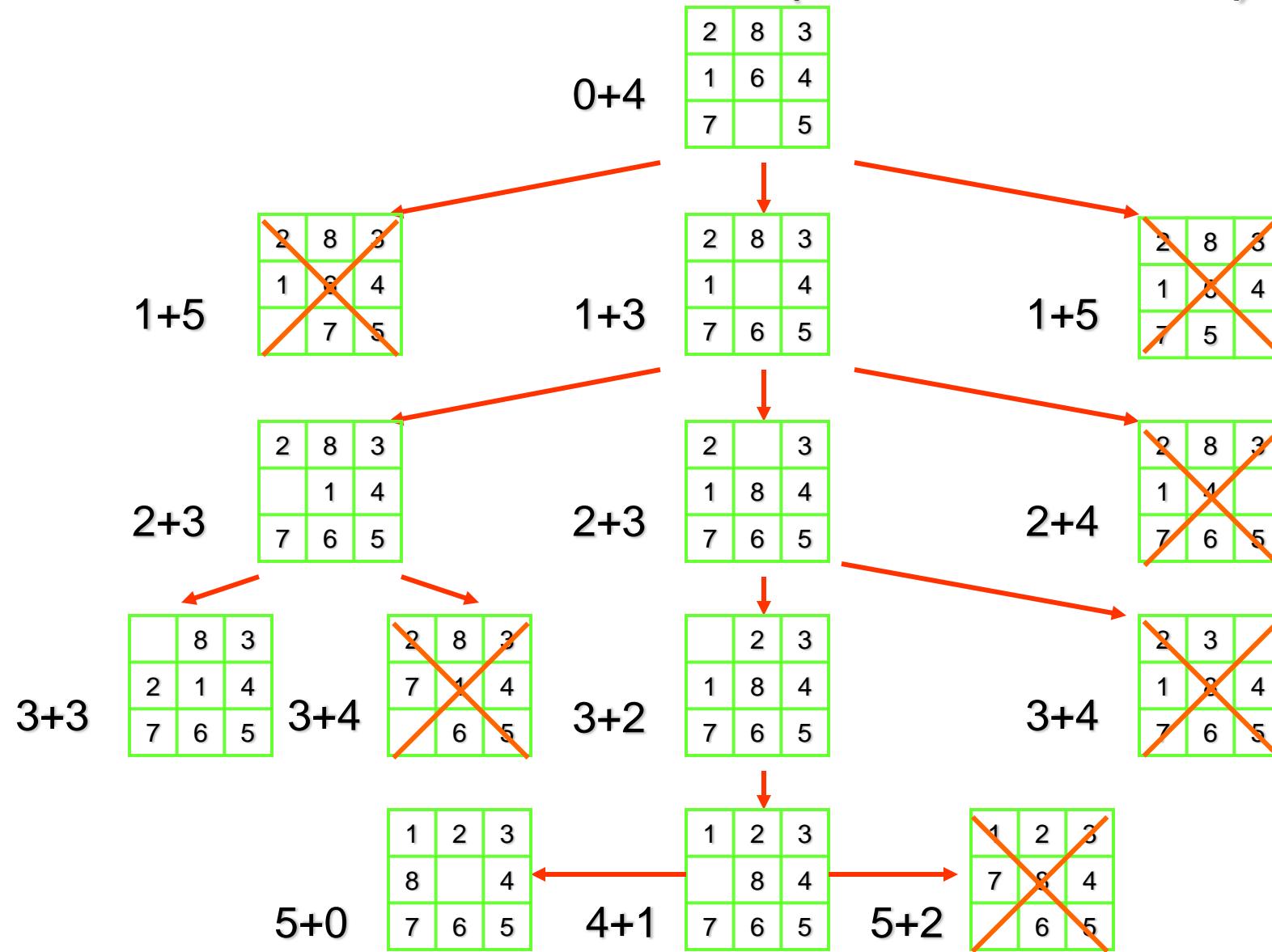
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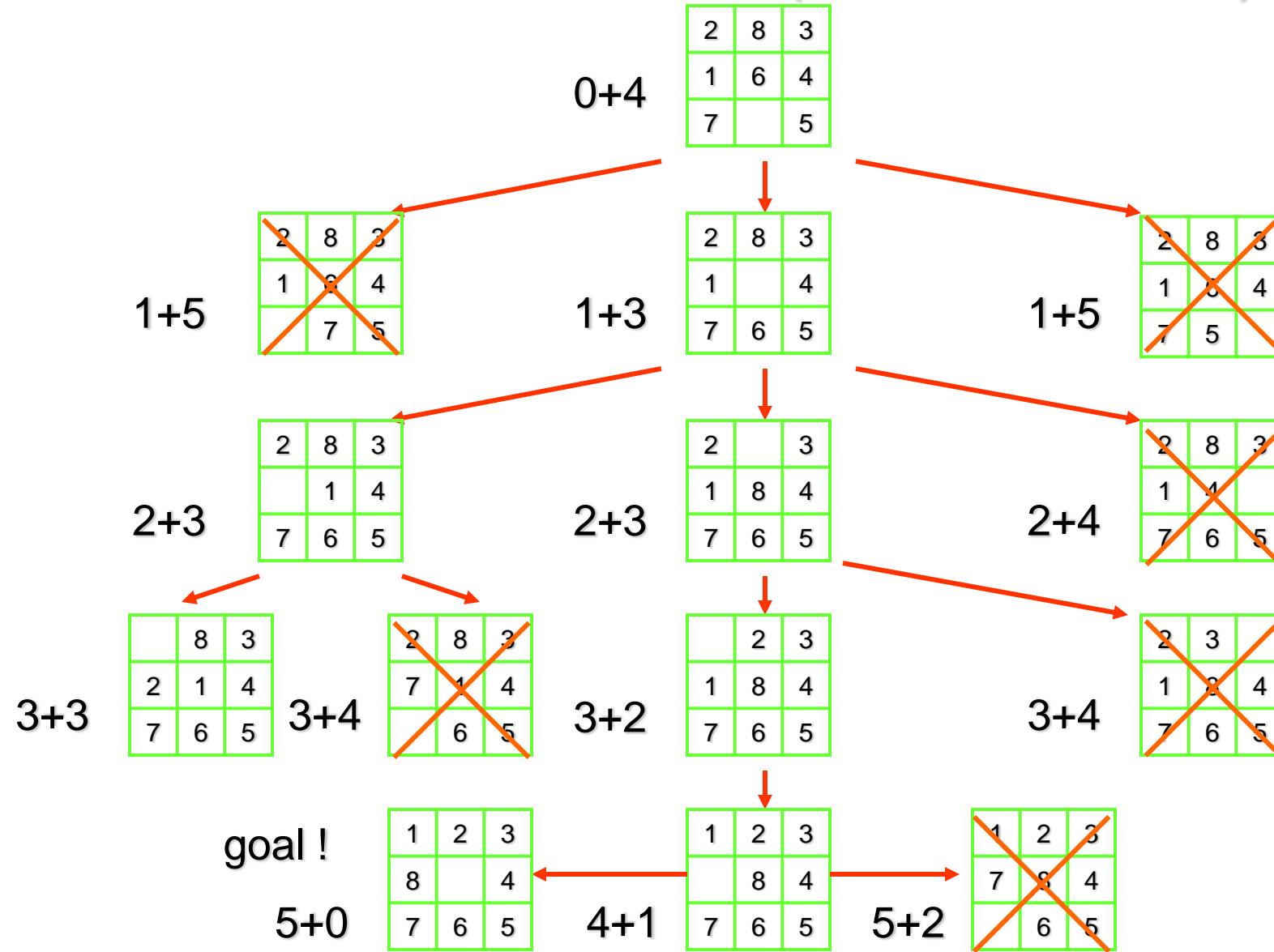
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# Generated and Test

- **Algorithm**
  1. Generate a (potential goal) state:
    - Particular point in the problem space, or
    - A path from a start state
  2. Test if it is a goal state
    - Stop if positive
    - go to step 1 otherwise
- **Systematic or Heuristic?**

# Generated and Test

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  1. Generate a (potential goal) state:
    - Particular point in the problem space, or
    - A path from a start state
  2. Test if it is a goal state
    - Stop if positive
    - go to step 1 otherwise
- **Systematic or Heuristic?**
  - It depends on “Generate”

# Local search algorithms

- In many problems, the **path** to the goal is irrelevant; the goal state itself is the solution
- State space = set of "complete" configurations
- Find configuration satisfying constraints, e.g., n-queens
- In such cases, we can use **local search algorithms**
- keep a single "**current**" state, try to improve it

# Hill Climbing

- Simple Hill Climbing
  - expand the current node
  - evaluate its children one by one (using the heuristic evaluation function)
  - choose the FIRST node with a better value
  
- Steepest Ascend Hill Climbing
  - expand the current node
  - Evaluate all its children (by the heuristic evaluation function)
  - choose the BEST node with the best value

# Hill-climbing

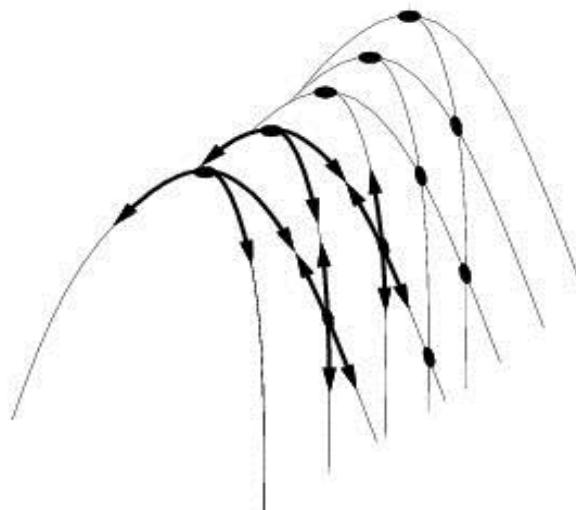
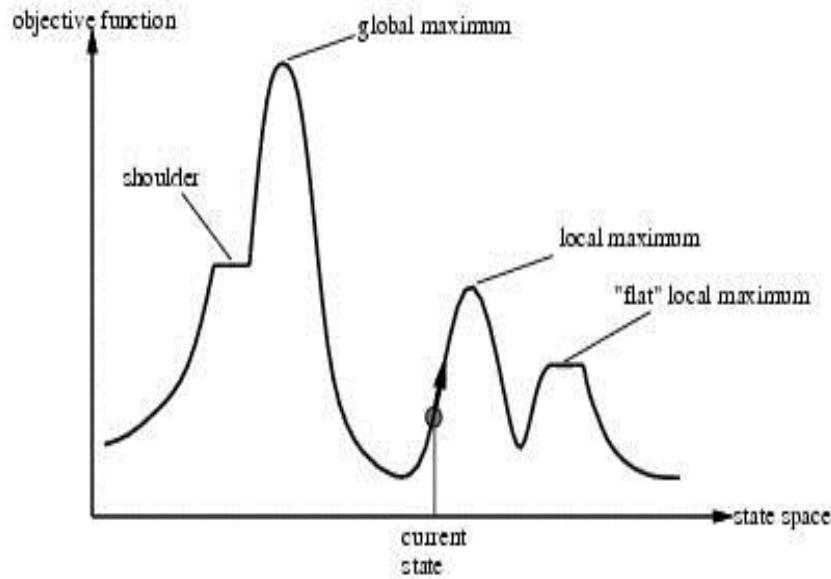
- Steepest Ascend

```
function HILL-CLIMBING(problem) returns a state that is a local maximum
  inputs: problem, a problem
  local variables: current, a node
                  neighbor, a node
  current  $\leftarrow$  MAKE-NODE(INITIAL-STATE[problem])
  loop do
    neighbor  $\leftarrow$  a highest-valued successor of current
    if VALUE[neighbor]  $\leq$  VALUE[current] then return STATE[current]
    current  $\leftarrow$  neighbor
```

# Hill Climbing

- Problems:
  - local maxima problem
  - plateau problem
  - Ridge

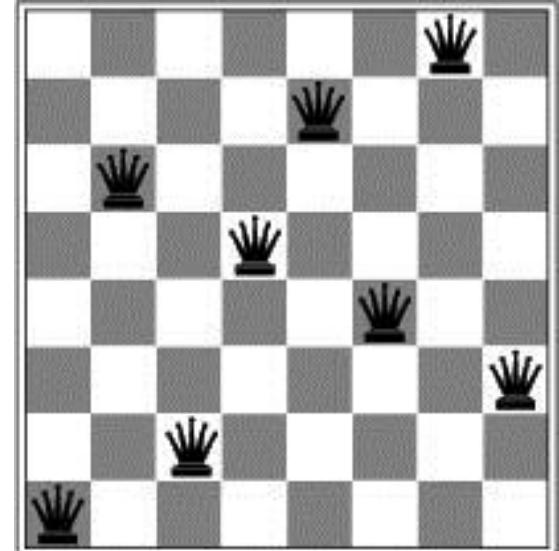
# Drawbacks



- Ridge = sequence of local maxima difficult for greedy algorithms to navigate
- Plateaux = an area of the state space where the evaluation function is flat.

# Hill-climbing example

18	12	14	13	13	12	14	14
14	16	13	15	12	14	12	16
14	12	18	13	15	12	14	14
15	14	14	14	18	13	16	13
18	14	17	15	15	14	16	16
17	18	16	18	15	18	15	18
18	14	18	15	15	14	18	16
14	14	13	17	12	14	12	18



- a) shows a state of  $h=17$  and the  $h$ -value for each possible successor.
- b) A local minimum in the 8-queens state space ( $h=1$ ).

# Hill-climbing variations

- Stochastic hill-climbing
  - Random selection among the uphill moves.
  - The selection probability can vary with the steepness of the uphill move.
- First-choice hill-climbing
  - Stochastic hill climbing by generating successors randomly until a better one is found.
- Random-restart hill-climbing
  - A series of Hill Climbing searches from randomly generated initial states
- Simulated Annealing
  - Escape local maxima by allowing some "bad" moves but **gradually decrease** their frequency

# Simulated Annealing

- $T = \text{initial temperature}$
- $x = \text{initial guess}$
- $v = \text{Energy}(x)$
- Repeat while  $T > \text{final temperature}$ 
  - Repeat  $n$  times
    - $X' \leftarrow \text{Move}(x)$
    - $V' = \text{Energy}(x')$
    - If  $V' < v$  then accept new  $x$  [  $x \leftarrow X'$  ]
    - Else accept new  $x$  with probability  $\exp(-(V' - v)/T)$
  - $T = 0.95T$  /\* Annealing schedule\*/
- At high temperature, most moves accepted
- At low temperature, only moves that improve energy are accepted

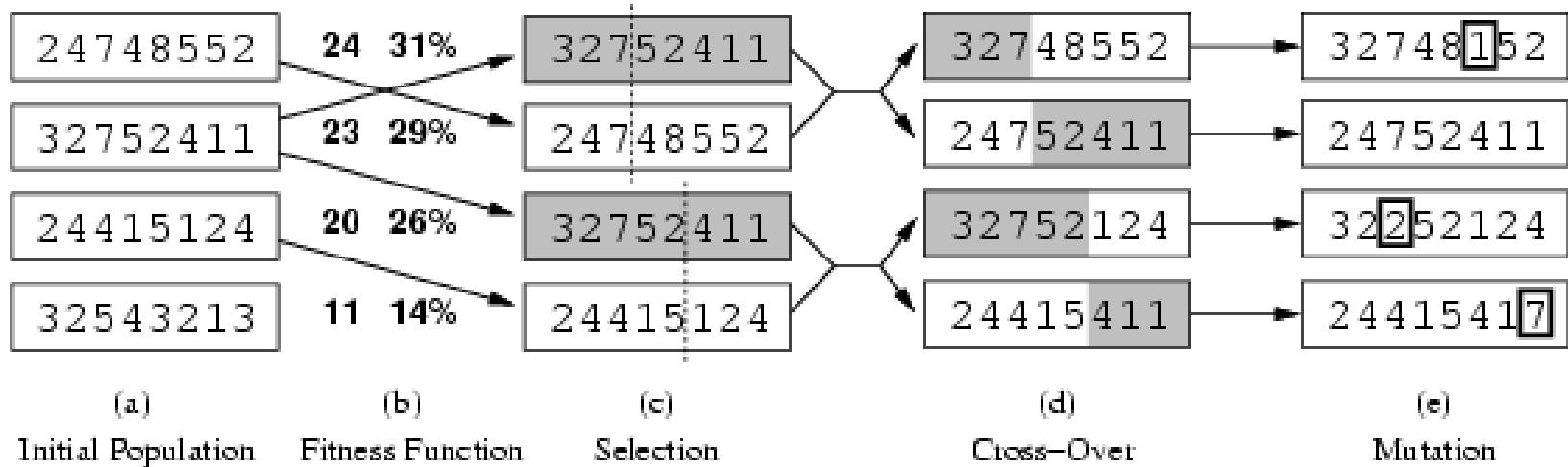
# Local beam search

- Keep track of  $k$  “current” states instead of one
  - Initially:  $k$  random initial states
  - Next: determine all successors of the  $k$  current states
  - If any of successors is goal → finished
  - Else select  $k$  best from the successors and repeat.
- Major difference with  $k$  random-restart search
  - Information is shared among  $k$  search threads.
- Can suffer from lack of diversity.
  - Stochastic variant: choose  $k$  successors at proportionally to the state success.

# Genetic algorithms

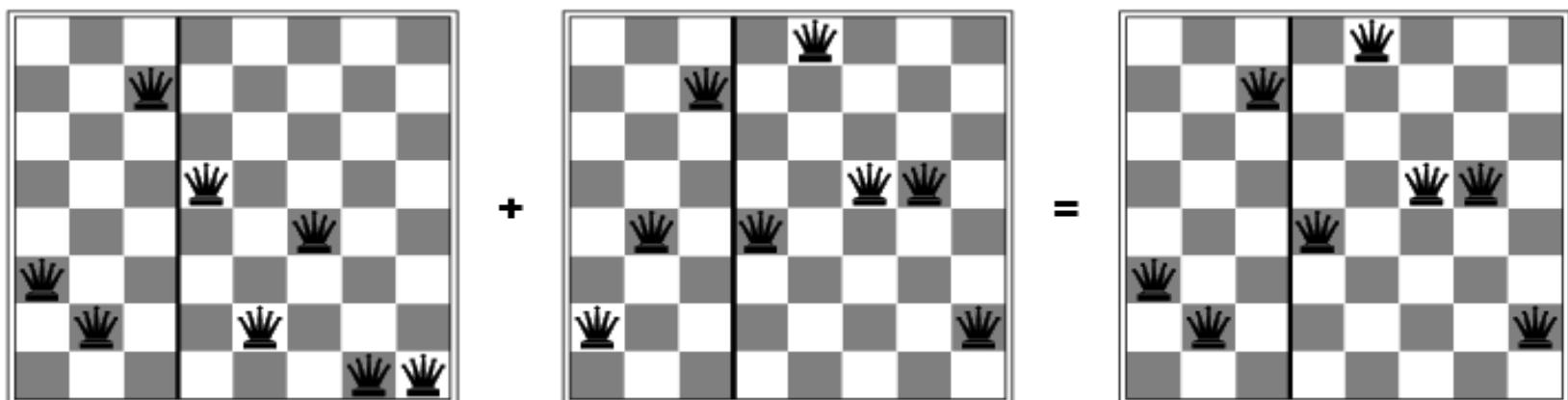
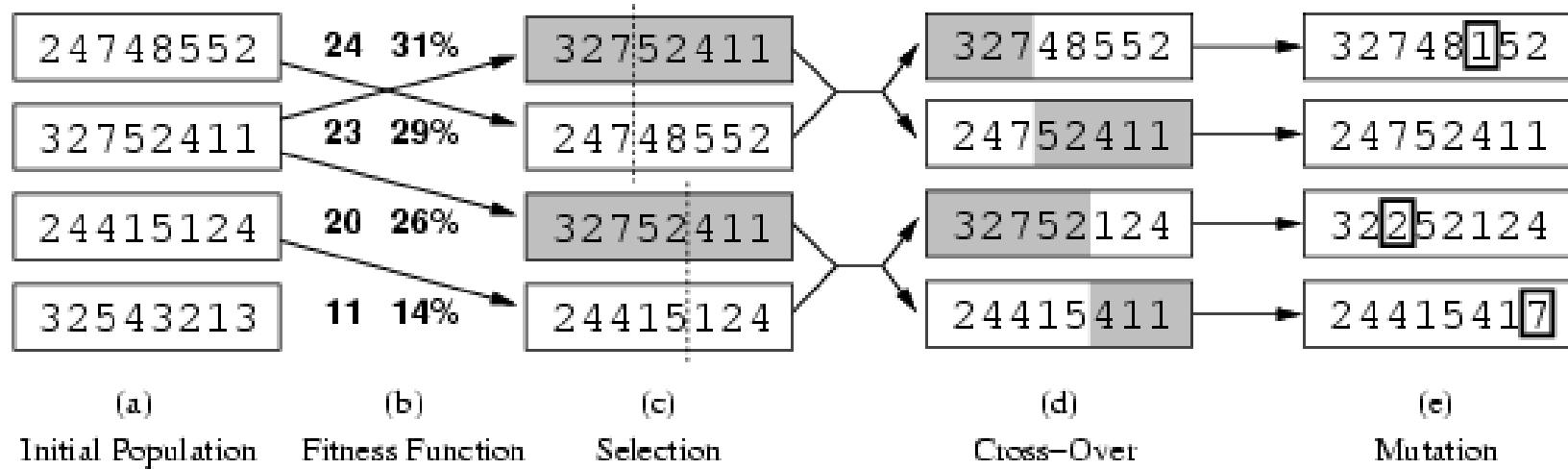
- Variant of local beam search with *sexual recombination*.
- A successor state is generated by combining two parent states
- Start with  $k$  randomly generated states (population)
- A state is represented as a string over a finite alphabet (often a string of 0s and 1s)
- Evaluation function (fitness function). Higher values for better states.
- Produce the next generation of states by selection, crossover, and mutation

# Genetic algorithms



- Fitness function: number of non-attacking pairs of queens  
(min = 0, max =  $8 \times 7/2 = 28$ )
- $24/(24+23+20+11) = 31\%$
- $23/(24+23+20+11) = 29\%$  etc

# Genetic algorithms



# A Genetic algorithm

**function** GENETIC\_ALGORITHM( *population*, FITNESS-FN) **return** an individual  
**input:** *population*, a set of individuals

FITNESS-FN, a function which determines the quality of the individual

**repeat**

*new\_population*  $\leftarrow$  empty set

**loop for** i **from** 1 **to** SIZE(*population*) **do**

*x*  $\leftarrow$  RANDOM\_SELECTION(*population*, FITNESS\_FN)

*y*  $\leftarrow$  RANDOM\_SELECTION(*population*, FITNESS\_FN)

*child*  $\leftarrow$  REPRODUCE(*x,y*)

**if** (small random probability) **then** *child*  $\leftarrow$  MUTATE(*child*)

add *child* to *new\_population*

*population*  $\leftarrow$  *new\_population*

**until** some individual is fit enough or enough time has elapsed

**return** the best individual

# A Genetic Algorithm (Cont.)

**function** REPRODUCE( *x*, *y*) **return** an individual

**input:** *x*, *y*, parent individuals

*n*  $\leftarrow$  LENGTH(*x*); *c*  $\leftarrow$  random number from 1 to *n*

**return** APPEND(SUBSTRING(*x*, 1, *c*), SUBSTRING(*y*, *c* + 1, *n*))

In this more popular version of GA, from each two parents, only one offspring is produced, not two.