# طراحی و تحلیل الگوریتم ها

دکتر امیر لکی زاده استادیار گروه مهندسی کامپیوتر دانشگاه قم

### ضرب ماتریس ها(۱)

```
SQUARE-MATRIX-MULTIPLY (A, B)

1  n = A.rows

2  let C be a new n \times n matrix

3  for i = 1 to n

4  for j = 1 to n

5  c_{ij} = 0

6  for k = 1 to n

7  c_{ij} = c_{ij} + a_{ik} \cdot b_{kj}

8  return C
```

#### ضرب ماتریس ها (۲)

$$A = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix}, \quad B = \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}, \quad C = \begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix},$$

so that we rewrite the equation  $C = A \cdot B$  as

$$\begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \cdot \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}.$$

Equation (4.10) corresponds to the four equations

$$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21} ,$$

$$C_{12} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22} ,$$

$$C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21} ,$$

$$C_{22} = A_{21} \cdot B_{12} + A_{22} \cdot B_{22} .$$

#### ضرب ماتریس ها (۲)

#### SQUARE-MATRIX-MULTIPLY-RECURSIVE (A, B)

```
1 n = A.rows
                                                T(n) = \begin{cases} \Theta(1) & \text{if } n = 1, \\ 8T(n/2) + \Theta(n^2) & \text{if } n > 1. \end{cases}
2 let C be a new n \times n matrix
3 if n == 1
   c_{11} = a_{11} \cdot b_{11}
   else partition A, B, and C as in equations (4.9)
        C_{11} = \text{SQUARE-MATRIX-MULTIPLY-RECURSIVE}(A_{11}, B_{11})
              + SQUARE-MATRIX-MULTIPLY-RECURSIVE (A_{12}, B_{21})
        C_{12} = \text{SQUARE-MATRIX-MULTIPLY-RECURSIVE}(A_{11}, B_{12})
              + SQUARE-MATRIX-MULTIPLY-RECURSIVE (A_{12}, B_{22})
        C_{21} = \text{SQUARE-MATRIX-MULTIPLY-RECURSIVE}(A_{21}, B_{11})
8
              + SQUARE-MATRIX-MULTIPLY-RECURSIVE (A_{22}, B_{21})
        C_{22} = \text{SQUARE-MATRIX-MULTIPLY-RECURSIVE}(A_{21}, B_{12})
9
              + SQUARE-MATRIX-MULTIPLY-RECURSIVE (A_{22}, B_{22})
    return C
```

### ضرب ماتریس ها (۳) Strassen Alg

$$S_1 = B_{12} - B_{22}$$
,  
 $S_2 = A_{11} + A_{12}$ ,  
 $S_3 = A_{21} + A_{22}$ ,  
 $S_4 = B_{21} - B_{11}$ ,  
 $S_5 = A_{11} + A_{22}$ ,  
 $S_6 = B_{11} + B_{22}$ ,  
 $S_7 = A_{12} - A_{22}$ ,  
 $S_8 = B_{21} + B_{22}$ ,  
 $S_9 = A_{11} - A_{21}$ ,  
 $S_{10} = B_{11} + B_{12}$ .

### ضرب ماتریس ها (۳) Strassen Alg

$$C_{11} = P_5 + P_4 - P_2 + P_6.$$

$$C_{12} = P_1 + P_2 \,,$$

$$C_{21} = P_3 + P_4$$

$$C_{22} = P_5 + P_1 - P_3 - P_7 ,$$

# ضرب ماتریس ها (۲) Strassen Alg

$$T(n) = \begin{cases} \Theta(1) & \text{if } n = 1, \\ 7T(n/2) + \Theta(n^2) & \text{if } n > 1. \end{cases}$$

$$T(n) = \Theta(n^{\lg 7}).$$

#### ضرب دو عدد n بیتی(۱)

```
Multiply(u.v){
n = \max(number \ of \ digits \ in \ u \ and \ v);
else if (n \leq threshold) return (u \times v); \rightarrow بودند کوچک بودند
else \{M = \lfloor n/2 \rfloor;
رخ می دهد \theta (n) \leftarrow \begin{cases} x = u \text{ divide } 10^M; y = u \text{ rem } 10^M; \\ w = V \text{ divide } 10^M; z = V \text{ rem } 10^M; \end{cases}
return (multiply (x, w) \times 10^{2M} + Multiply(x, z) \times 10^{M} + Multiply(y, w)
                    \times 10^M + Multiply(y,z);
T(n) = 4T(n/2) + 4T(n/2) + \theta(n) \rightarrow T(n) = \theta(n^2)
```

# ضرب دو عدد n بیتی (۲)

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