Sharif University of Technology - Big Dash - Notebook		k	7 Combinatorics         7.1 LP simplex
C	Contents		7.3 FFT
$\cup$	Officerius		7.4 NTT
			7.5 FWHT
1	Tester	1	7.6 Stirling 1
	1.1 Tester for Bash	1	7.7 Chinese remainder
			7.8 Stirling 2
2	Number theory	1	7.9 Duality of LP
2	Number theory		7.10 Extended catalan
	2.1 Primality Test	1	7.11 Find polynomial from it's points
	2.2 Linear Sieve	2	roce polyments of process and
	2.3 Integration	2	8 Constants
	2.4 Extended Euclidean Algorithm	2	
	2.5 Linear Diophantine Equation	2	8.1 Number of primes
			8.2 Most divisor
3	Geometry	3	
•	3.1 Geometry Primitives with Complex	_	1
			1 Tester
	3.2 Geometry Primitives	3	
	3.3 Line intersection	5	1.1 Tester for Bash
	3.4 Line and circle intersection	5	1.1 Tester for Bash
	3.5 Intersection of two circles	5	
	3.6 Rotating Calipers	6	# a.cpp -> code, naive.cpp -> naive solution, gen.cpp -> generator
	3.7 Delaunay Triangulation $O(n^2)$	6	g++ a.cpp -std=c++17 -02 -o a.out
	3.8 Delaunay $O(n \log^2 n)$	8	g++ naive.cpp -std=c++17 -02 -o naive.out
	3.9 Convex Hull 3D	9	g++ gen.cpp -std=c++17 -02 -o gen.out
			t=1000
	3.10 Half Plane Intersection	10	for ((i = 1; i <= t; i++)); do
	3.11 Minimum Enclosing Circle	11	<pre>printf "%s %d\r" "running on test" \$i</pre>
	3.12 Number of integer points inside polygon	12	./gen.out > test.txt
	3.13 Useful Geo Facts	12	./a.out < test.txt > out.txt
	3.14 Duality and properties	12	<pre>./naive.out &lt; test.txt &gt; ans.txt</pre>
			diff -bBq out.txt ans.txt
4	String	12	<pre>if [[ \$? != 0 ]]; then</pre>
	4.1 Suffix Array	12	echo wrong on test \$i
	4.2 Suffix Automata	13	echo input:
	4.3 Suffix Tree	13	cat test.txt
		13 13	echo output:
	4.4 Palindromic Tree	13	cat out.txt
_		4.4	echo answer:
5	Data structure	14	cat ans.txt
	5.1 Ordered Set	14	exit
	5.2 Treap	14	fi
	5.3 Dynimic convex hull	14	<pre>done printf "all tests passed"</pre>
			princi all tests passed
6	Graph	15	
	6.1 Maximum matching - Edmond's blossom	15	
	6.2 Biconnected components	16	2 Number theory
	6.3 Flow - Dinic	16	
			0.1 D. 11 / D. /
			2.1 Primality Test
	6.5 Maximum weighted matching - Hungarian	17	
	6.6 Ear decomposition	18	<pre>using u64 = uint64 t;</pre>
	6.7 Stoer-Wagner min cut $O(n^3)$	18	using u128 =uint128_t;
	6.8 Directed minimum spanning tree $O(m \log n)$	18	1164 binpower(1164 base, 1164 e. 1164 mod) {

u64 binpower(u64 base, u64 e, u64 mod) {

u64 result = 1;

base %= mod;

6.9

```
while (e) {
        if (e & 1)
            result = (u128) result * base % mod;
        base = (u128)base * base % mod;
        e >>= 1;
    return result;
bool check_composite(u64 n, u64 a, u64 d, int s) {
    u64 x = binpower(a, d, n);
    if (x == 1 | | x == n - 1)
        return false:
    for (int r = 1; r < s; r++) {
        x = (u128)x * x % n;
        if (x == n - 1)
            return false;
    return true;
bool MillerRabin (u64 n) { // true if n is prime, else false.
    if (n < 2)
        return false;
    int r = 0;
    u64 d = n - 1;
    while ((d & 1) == 0) {
        d >>= 1;
        r++;
    // for a 32 bit integer it is only necessary to check [2, 3, 5, 7]
    for (int a: {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}) {
        if (n == a)
            return true;
        if (check_composite(n, a, d, r))
            return false;
    return true;
```

### 2.2 Linear Sieve

```
const int N = 10000000;
vector<int> lp(N+1);
vector<int> pr;
for (int i=2; i <= N; ++i) {
    if (lp[i] == 0) {
        lp[i] = i;
        pr.push_back(i);
    }
    for (int j = 0; i * pr[j] <= N; ++j) {
        lp[i * pr[j]] = pr[j];
        if (pr[j] == lp[i]) {
            break;
        }
    }
}</pre>
```

### 2.3 Integration

```
const int N = 1000 * 1000; // must be even
double simpson_integration(double a, double b) {
    double h = (b - a) / N;
    double s = f(a) + f(b);
    for (int i = 1; i <= N - 1; ++i) {
        double x = a + h * i;
        s += f(x) * ((i & 1) ? 4 : 2);
    }
    s *= h / 3;
    return s;
}</pre>
```

### 2.4 Extended Euclidean Algorithm

```
int gcd(int a, int b, int& x, int& y) {
   if (b == 0) {
        x = 1;
        y = 0;
        return a;
   }
   int x1, y1;
   int d = gcd(b, a % b, x1, y1);
        x = y1;
        y = x1 - y1 * (a / b);
   return d;
}
```

### 2.5 Linear Diophantine Equation

```
int gcd(int a, int b, int& x, int& y) {
    if (b == 0) {
        x = 1;
        \mathbf{v} = 0;
        return a;
    int x1, y1;
    int d = gcd(b, a % b, x1, y1);
    y = x1 - y1 * (a / b);
    return d;
// ax + by = c
bool find_any_solution(int a, int b, int c, int &x0, int &y0, int &g)
    g = gcd(abs(a), abs(b), x0, y0);
    if (c % g) {
        return false;
    x0 \star = c / g;
    y0 \star = c / q;
    if (a < 0) x0 = -x0;
    if (b < 0) y0 = -y0;
    return true;
```

### 3 Geometry

### 3.1 Geometry Primitives with Complex

```
typedef complex<double> point;
#define X real()
#define Y imag()
a + b // vector addition
r * a // scalar multiplication
double dot(point a, point b) { return (conj(a) * b).X;}
double cross(point a, point b) { return (conj(a) * b).Y;}
double squared_dist(point a, point b) { return norm(a - b);}
double dist(point a, point b) {return abs(a - b);} // abs = rad(norm)
double angle(point a, point b) {return arg(b - a);}
double slope(point a, point b) {return tan(arg(b - a));}
polar(r, theta); // polar to cartesian
point(abs(p), arg(p)); // cartesian to polar
point rotate(point a, double theta) {return a * polar(1.0, theta);}
point rotate (point a, double theta, point p) {return p + rotate(a-p,
    theta);} // rotate about p
point angle (point a, point b, point c) {abs(remainder(arg(a-b) - arg(c
    -b), 2.0) * M_PI);
point project(point p, point v) {return v * dot(p, v) / norm(v);} //
   project p in vector v
point project(point p, point a, point b) {return a + project(p - a, b
    - a);} // project p in line (a, b)
point reflect(point p, point a, point b) {return a + conj((p - a) / (b
     -a)) * (b - a);}
point intersection (point a, point b, point p, point q) {
   double c1 = cross(p - a, b - a), c2 = cross(q - a, b - a);
   return (c1 * q - c2 * p) / (c1 - c2); // undefined if parallel
```

### 3.2 Geometry Primitives

```
const int ON = 0, LEFT = 1, RIGHT = -1, BACK = -2, FRONT = 2, IN = 3,
   OUT = -3;
bool byX(const pt &a, const pt &b) {
        if (Equ(a.x, b.x)) return Lss(a.v, b.v);
        return Lss(a.x, b.x);
bool byY(const pt &a, const pt &b) {
        if (Equ(a.y, b.y)) return Lss(a.x, b.x);
        return Lss(a.y, b.y);
struct cmpXY{ bool operator ()(const pt &a, const pt &b){ return byX(a
    , b); };
struct cmpYX{ bool operator () (const pt &a, const pt &b) { return byY(a
istream& operator >> (istream &in, pt &p) { in >> p.x >> p.y; return in
ostream& operator << (ostream &out, pt p) { out << p.x << ' ' << p.v;
    return out; }
pt rot(pt a) { return pt(-a.y, a.x); }
pt proj(pt a, pt b, pt c){
        b = b-a, c = c-a;
        return a + (b*c)/(b*b)*b;
pt reflect(pt a, pt b, pt c) {
```

```
pt d = c;
        b = b-a, c = c-a;
        return d + (c^b) /abs(b) *rot(unit(b)) *2;
bool intersect(pt a, pt b, pt c, pt d) {
        int as = dir(c, d, a), bs = dir(c, d, b),
                cs = dir(a, b, c), ds = dir(a, b, d);
        if (as && as == bs || cs && cs == ds)
                return false;
        else if (as || bs || cs || ds)
                return true;
        for (int j = 0; j < 2; j++) {
                cord mX = min(a.x, b.x), MX = max(a.x, b.x),
                         mY = min(a.y, b.y), MY = max(a.y, b.y);
                for (int k = 0; k < 2; k++) {
                        if (c.x + EPS > mX && c.x < MX + EPS && c.y +
                            EPS > mY && c.y < MY + EPS)
                                return true;
                        swap(c, d);
                swap(a, c); swap(b, d);
        return false;
pt intersection(pt a, pt b, pt c, pt d) {
        cord c1 = (b-a)^(c-a), c2 = (b-a)^(d-a);
        return (c1*d - c2*c)/(c1 - c2);
ld signedArea(vector<pt> &p) {
        int n = p.size();
        cord res = 0;
        for (int i = 0; i < n; i++)
              res += (p[i]^p[(i+1)%n]);
        return (ld) res/2;
//relative position of c toward ab
int relpos(pt a, pt b, pt c){
       b = b-a; c = c-a;
        if (Grt(b^c, 0)) return LEFT;
       if (Lss(b^c, 0)) return RIGHT;
        if (Lss(b*c, 0)) return BACK;
        if (Grt(b*c, abs(b))) return FRONT;
        return ON:
//distance of a point from a line segment
ld distLSP(pt a, pt b, pt c){
        int rpos = relpos(a, b, proj(a, b, c));
        if (rpos == BACK) return abs(c-a);
        if (rpos == FRONT) return abs(c-b);
        b = b-a, c = c-a;
        return abs(b^c/abs(b));
//distance between two line segments
ld distLS(pt a, pt b, pt c, pt d) {
        if (intersect(a, b, c, d)) return 0;
        return min(min(distLSP(a, b, c), distLSP(a, b, d)), min(
            distLSP(c, d, a), distLSP(c, d, b)));
//angles less than or equal to 180
bool isConvex(vector<pt> &p) {
        int n = p.size();
```

```
bool neg = false, pos = false;
                                                                                                      if (Grt(it->y, po[i].y+rad))
        for (int i = 0; i < n; i++) {
                                                                                                              break:
                int rpos = relpos(p[i], p[(i+1)%n], p[(i+2)%n]);
                                                                                                     ld cur = abs(po[i] - (*it));
                if (rpos == LEFT) pos = true;
                                                                                                     if (Lss(cur, rad)){
                if (rpos == RIGHT) neg = true;
                                                                                                            rad = cur;
                                                                                                             res = \{*it, po[i]\};
        return (neg&pos) == false;
int crossingN(vector<pt> &p, pt a) {
        int n = p.size();
                                                                                     return res;
        pt b = a;
        for (pt q : p)
                                                                             //Cuts polygon with line ab and returns the left cut polygon
                b.x = max(b.x, q.x);
                                                                             vector<pt> convexCut(vector<pt> &po, pt a, pt b) {
        int cn = 0;
                                                                                     int n = po.size();
        for (int i = 0; i < n; i++) {
                                                                                     vector<pt> res;
                                                                                     for (int i = 0; i < n; i++) {
                pt q1 = p[i], q2 = p[(i+1)%n];
                if (intersect(a, b, q1, q2) && (dir(a, b, q1) == 1 ||
                                                                                             if (dir(a, b, po[i]) >= 0)
                    dir(a, b, q2) == 1))
                                                                                                     res.push_back(po[i]);
                        cn++;
                                                                                             if (abs(dir(a, b, po[i]) - dir(a, b, po[(i+1)%n])) ==
        return cn;
                                                                                                     res.push_back(intersection(a, b, po[i], po[(i
                                                                                                          +1)%n]));
int windingN(vector<pt> &p, pt a){
        int n = p.size();
                                                                                     return res;
        pt b = a;
        for (pt q : p)
                                                                             //slightly line
                b.x = max(b.x, q.x);
                                                                             pair <pt, pt> get_segment(pt a, pt b){
                                                                                 const int deltax = b.x - a.x;
        int wn = 0;
        for (int i = 0; i < n; i++) {
                                                                                 const int deltay = b.v - a.v;
                pt q1 = p[i], q2 = p[(i+1)%n];
                                                                                 const int k = 100001;
                if (intersect(a, b, q1, q2)){
                                                                                 pt aa(a.x - k * deltax, a.y - k * deltay);
                        int ps = dir(a, b, q1), qs = dir(a, b, q2);
                                                                                 pt bb(b.x + deltax, b.v + deltay);
                        if (qs >= 0) wn++;
                                                                                 static const int dx[4] = \{ -1, 0, 1, 0 \};
                        if (ps >= 0) wn--:
                                                                                 static const int dv[4] = \{ 0, 1, 0, -1 \};
                                                                                 for (int d = 0; d < 4; ++d) {
                                                                                     pt aaa(aa.x + dx[d], aa.y + dy[d]);
        return wn;
                                                                                     pt bbb(bb.x, bb.y);
                                                                                     if (dir(aaa, bbb, a) >= 0) continue;
//returns IN, ON or OUT
int pointInPoly(vector<pt> &p, pt a) {
                                                                                     if (dir(aaa, bbb, b) == 0) continue;
        int n = p.size();
                                                                                     return {aaa, bbb};
        for (int i = 0; i < n; i++)
                if (relpos(p[i], p[(i+1)%n], a) == ON)
                        return ON;
                                                                             pair<int, int> tangent(vector <pt> &A0, vector <pt> &B0){
        return (crossingN(p, a)%2 ? IN : OUT);
                                                                                 vector <pair<pt, int>> A, B;
        //return (windingN(po, a) ? IN : OUT);
                                                                                 for (int i = 0; i < sz(A0); i++)
                                                                                     A.pb({A0[i], i});
pair <pt, pt> nearestPair(vector<pt> &po) {
                                                                                 for (int i = 0; i < sz(B0); i++)
        int n = po.size();
                                                                                     B.pb({B0[i], i});
        sort(po.begin(), po.end(), cmpXY());
                                                                                 sort(A.begin(), A.end());
        multiset <pt, cmpYX> s;
                                                                                 sort(B.begin(), B.end());
        ld rad = abs(po[1]-po[0]);
                                                                                 A = convex hull(A);
        pair <pt, pt> res = {po[0], po[1]};
                                                                                 B = convex_hull(B);
                                                                                 int ia = 0, ib = 0;
        int 1 = 0, r = 0;
        for (int i = 0; i < n; i++) {
                                                                                 //direction must be considered
                while (1 < r \&\& Geq(po[i].x - po[l].x, rad))
                                                                                 while (1) {
                        s.erase(po[1++]);
                                                                                     bool fin = true;
                while (r < i && Leq(po[r].x, po[i].x))</pre>
                                                                                     for (; dir(A[ia].F, B[ib].F, B[(ib+sz(B)-1)%sz(B)].F) < 0 | |</pre>
                                                                                         dir(A[ia].F, B[ib].F, B[(ib+1) %sz(B)].F) < 0; ib = (ib+1) %
                        s.insert(po[r++]);
                for (auto it = s.lower_bound(pt(po[i].x, po[i].y-rad))
                                                                                         sz(B))
                    ; it != s.end(); it++){
                                                                                         fin = false;
```

```
for (; dir(B[ib].F, A[ia].F, A[(ia+sz(A)-1)%sz(A)].F) > 0 | |
            dir(B[ib].F, A[ia].F, A[(ia+1) %sz(A)].F) > 0; ia = (ia+1) %
            sz(A))
            fin = false;
        if (fin) break;
    return {A[ia].S, B[ib].S};
//Sweep Line Example
int main(){
        vector <pt> adds, rems;
        vector <pair<pt, pt>> query;
        int it;
        cin >> it;
        for (int i = it; i; i--) {
                pt p, q;
                cin >> p >> q;
                if (Equ(p.y, q.y)) {
                         if (Lss(q.x, p.x))
                                 swap(p, q);
                         adds.push back(p);
                         rems.push_back(q);
                else{
                         if (Lss(q.y, p.y))
                                 swap(p, q);
                         query.push_back({p, q});
        sort(adds.begin(), adds.end());
        sort(rems.begin(), rems.end());
        sort(query.begin(), query.end());
        multiset <ld> vs;
        int iadd = 0, irem = 0;
        int ans = 0;
        for (auto p : query) {
                while (iadd < adds.size() && Leg(adds[iadd].x, p.F.x))</pre>
                        ys.insert(adds[iadd++].y);
                while (irem < rems.size() && Lss(rems[irem].x, p.F.x)</pre>
                     && ys.find(rems[irem].y) != ys.end())
                        ys.erase(ys.find(rems[irem++].y));
                int cur = distance(ys.lower_bound(p.F.y), ys.
                    upper_bound(p.S.y));
                ans += cur;
        cout << ans << endl;
        return 0;
//rotate b with center=a theta radian
PT rotate(PT a, PT b, ld theta) {
        return (b-a) *polar<ld>(1, theta) + a;
```

### 3.3 Line intersection

```
point intersection(point a, point b, point c, point d) {
    point ab = b - a;
    point cd = d - c;
```

```
point ac = c - a;
double alpha = cross(ac, cd) / cross(ab, cd);
return a + alpha * ab;
```

### 3.4 Line and circle intersection

```
// return pair<point, point> which is intersections point
// for each point if it's not exist, return (INF,INF)
typedef pair<point,point> ppp;
const ld INF = 1e18;
const ld eps = 1e-15;
ppp line_circle_intersection(point p1,point p2,point o,ld r) {
   point q = dot(o-p1,p2-p1)/dist(p1,p2)*(p2-p1) + p1;
   ld d = r*r - dist(o,q);
   if(d<eps && d>-eps) return ppp(q, point(INF,INF));
   if(d<0) return ppp(point(INF,INF), point(INF,INF));
   point dif = sqrt(d/dist(p1,p2))*(p1-p2);
   return ppp(q-dif,q+dif);
}</pre>
```

### 3.5 Intersection of two circles

```
#define USE MATH DEFINES
const int MAX_N = 2e5+10;
const int INF = 1e9;
const ld eps = 1e-8;
struct circle {
public:
  ld r;
  point o:
  circle(ld rr, ld x, ld y) {
   r = rr;
   o.x = x;
  ld S() {
    return M PI*r*r;
  ld distance(point p1, point p2) { return hypot(p2.x-p1.x,p2.y-p1.y);
   0 = other is inside this, zero point
   1 = other is tangent inisde of this, one point
   2 = other is intersect with this, two point
    3 = other is tangent outside of this, one point
    4 = other is outside of this, zero point
  pair<int, vector<point> > intersect(circle other) {
    vector<point> v;
    ld sumr = other.r + r:
    ld rr = r - other.r;
    ld dis = distance(o, other.o);
    ld = (r*r - other.r*other.r + dis*dis)/(2*dis);
    1d h = sqrt(r*r-a*a);
    point p2(o.x, o.v);
    p2.x = a*(other.o.x - o.x)/dis;
    p2.y = a*(other.o.y - o.y)/dis;
```

```
if(is_zero(sumr-dis)) {
      v.push back(p2);
      return make_pair(3, v);
    if(is_zero(rr - dis)) {
      v.push back(p2);
      return make_pair(1, v);
    if(dis <= rr)</pre>
      return make_pair(0, v);
    if(dis >= sumr)
      return make_pair(4, v);
    point p3(p2.x + h*(other.o.y - o.y)/dis, p2.y - h*(other.o.x - o.x
    point p4(p2.x - h*(other.o.y - o.y)/dis, p2.y + h*(other.o.x - o.x)
        )/dis);
    v.push_back(p3);
    v.push_back(p4);
    return make pair (2, v):
  ld f(ld l, ld r, ld R) {
    1d \cos a = (1*1 + r*r - R*R)/(2.0*r*1);
   ld a = acos(cosa);
    return r*r*(a - \sin(2*a)/2);
  ld intersection area(circle c2) {
    ld l = distance(o, c2.o);
    if(1 >= r + c2.r) return 0;
    else if (c2.r >= 1 + r) return S();
    else if(r \ge 1 + c2.r) return c2.S();
    return f(1, r, c2.r) + f(1, c2.r, r);
};
```

### 3.6 Rotating Calipers

```
vector <pair<pt, pt>> get_antipodals(vector <pt> &p) {
        int n = sz(p);
        sort(p.begin(), p.end());
        vector <pt> U, L;
        for (int i = 0; i < n; i++) {
                while (sz(U) > 1 \&\& side(U[sz(U)-2], U[sz(U)-1], p[i])
                      >= ())
                         U.pop_back();
                 while (sz(L) > 1 \&\& side(L[sz(L)-2], L[sz(L)-1], p[i])
                         L.pop_back();
                U.pb(p[i]);
                L.pb(p[i]);
        vector <pair<pt, pt>> res;
        int i = 0, j = sz(L)-1;
        while (i+1 < sz(U) | | j > 0) {
                res.pb({U[i], L[j]});
                if (i+1 == sz(U)) j--;
                else if ( \dot{j} == 0)
                                    <u>i</u>++;
                else if (cross(L[j]-L[j-1], U[i+1]-U[i]) >= 0) i++;
        return res;
```

### 3.7 Delaunay Triangulation $O(n^2)$

```
struct Delaunay{
    vector <pt> p;
    vector <int> to, nxt, perm;
    int add_edge(int q, int bef=-1) {
        int cnt = sz(to);
        to.pb(a);
        nxt.pb(-1);
        if (bef !=-1) {
            nxt[bef] = cnt;
            to.pb(to[bef]);
            nxt.pb(-1);
        return cnt;
   bool onconvex(int e) {
        if (nxt[nxt[e]]] != e) return true;
        if (dir(p[to[e^1]], p[to[e]], p[to[nxt[e]]]) < 0) return true;</pre>
        return false;
    int before(int e) {
        int cur = e, last = -1;
            last = cur;
            cur = nxt[cur^1];
        }while (cur != e);
        return last^1;
    void easy_triangulate(){
        to.clear():
        nxt.clear();
        perm = vector<int>(sz(p));
        for (int i = 0; i < sz(p); i++)
            perm[i] = i;
        sort(perm.begin(), perm.end(), [&] (int i, int j) {
            return p[i] < p[j]; });</pre>
        sort(p.begin(), p.end());
        if (dir(p[0], p[1], p[2]) > 0){
            swap(p[1], p[2]);
            swap(perm[1], perm[2]);
        int to0 = add edge(0), to0c = add edge(2),
        to1 = add edge(1), to1c = add edge(0),
        to2 = add_edge(2), to2c = add_edge(1);
        nxt[to1] = to2; nxt[to2] = to0;
        nxt[to0] = to1; nxt[to0c] = to2c;
        nxt[to2c] = to1c; nxt[to1c] = to0c;
        int e = to0:
       bool D2 = true;
        for (int i = 3; i < sz(p); i++) {</pre>
            pt q = p[i];
            if (D2) {
                int edge = e;
                    if (dir(q, p[to[edge^1]], p[to[edge]])){
                        D2 = false;
                        break;
```

```
edge = nxt[edge];
            } while (edge != e);
        vector <int> vis;
        if (D2) {
            while (p[to[e^1]] < p[to[e]])
                e = nxt[e];
            vis.pb(e);
            e = nxt[e];
        else{
            while (dir(q, p[to[e^1]), p[to[e]]) \le 0 \mid | dir(q, p[
                to[e^1], p[to[before(e)^1]] < 0)
                e = nxt[e];
            while (dir(q, p[to[e^1]), p[to[e])) > 0){
                vis.pb(e);
                e = nxt[e];
        int b = before(vis[0]);
        int ex = add edge(i, b);
        int last = ex^1;
        for (int edge : vis) {
            nxt[last] = edge;
            int eq = add_edge(i, edge);
            nxt[edge] = eg;
            nxt[eq] = last;
            last = eq^1;
        nxt[ex] = last;
        nxt[last] = e;
bool incircle(pt a, pt b, pt c, pt d) {
    if (dir(a, b, c) < 0)
        swap(b, c);
    return a.z() * (b.x * (c.y - d.y) - c.x * (b.y - d.y)
                    + d.x * (b.y - c.y)
    -b.z() * (a.x * (c.y - d.y) - c.x * (a.y - d.y) + d.x * (a.y
    + c.z() * (a.x * (b.y - d.y) - b.x * (a.y - d.y) + d.x * (a.y
        -b.y)
    -d.z() * (a.x * (b.y - c.y) - b.x * (a.y - c.y) + c.x * (a.y
        - b.y)) > 0;
bool locally(int e) {
    pt a = p[to[e^1]], b = p[to[e]], c = p[to[nxt[e]]], d = p[to[e^1]]
    if (onconvex(e)) return true;
    if (onconvex(e^1)) return true;
    if (incircle(a, b, c, d)) return false;
    if (incircle(b, a, d, c)) return false;
    return true;
void flip(int e) {
    int a = nxt[e], b = nxt[a],
    c = nxt[e^1], d = nxt[c];
    nxt[d] = a;
    nxt[b] = c;
    to[e] = to[c];
```

```
nxt[a] = e;
    nxt[e] = d;
    to[e^1] = to[a];
    nxt[c] = e^1;
    nxt[e^1] = b;
void delaunay_triangulate(){
    if (sz(to) == 0)
        easy_triangulate();
    bool *mark = new bool[sz(to)];
    fill (mark, mark + sz(to), false);
    vector <int> bad;
    for (int e = 0; e < sz(to); e++) {
        if (!mark[e/2] && !locally(e)){
            bad.pb(e);
            mark[e/2] = true;
    while (sz(bad)){
        int e = bad.back();
        bad.pop_back();
        mark[e/2] = false;
        if (!locally(e)){
            int to_check[4] = {nxt[e], nxt[nxt[e]], nxt[e^1], nxt[
                 nxt[e^1]]};
            flip(e);
            for (int i = 0; i < 4; i++)
                if (!mark[to_check[i]/2] && !locally(to_check[i]))
                    bad.pb(to_check[i]);
                    mark[to_check[i]/2] = true;
    for (int e = 0; e < sz(to); e^{++})
        assert(locally(e));
vector <tri> get_triangles() {
    vector <tri> res;
    bool *mark = new bool[sz(to)];
    fill(mark, mark + sz(to), false);
    for (int e = 0; e < sz(to); e++) {
        if (mark[e]) continue;
        if (onconvex(e)) continue;
        pt a = p[to[e^1]], b = p[to[e]], c = p[to[nxt[e]]];
        mark[e] = mark[nxt[e]] = mark[nxt[nxt[e]]] = true;
        res.pb(tri(perm[to[e^1]], perm[to[e]], perm[to[nxt[e]]]));
    }
    return res;
vector <pair<ls, pt>> get_voronoi_edges() {
    vector <pair<ls, pt>> res;
    for (int e = 0; e < sz(to); e++) {
        pt a = p[to[e^1]], b = p[to[e]], c = p[to[nxt[e]]], d = p[to[nxt[e]]]
            to[nxt[e^1]]];
        if (onconvex(e^1)){
            pt ol = center(a, b, c),
            o2 = (a+b)/2;
            pt ab = (b-a);
            pt per(ab.y, -ab.x);
            o2 = o2 + per*100000; //infinity
```

### 3.8 Delaunay $O(n \log^2 n)$

```
const int MAXN = 100 * 1000 + 10;
const int MAXLG = 20;
const int INF = 100 * 1000 * 1000 + 10;
const int MAXPOINTS = MAXN * MAXLG;
typedef pair<int ,int> point;
struct tria{
        int a, b, c;
        tria(int _a, int _b, int _c) {
                a = _a; b = _b; c = _c;
        tria() \{a = b = c = 0; \}
};
struct Delaunay {
        typedef pair<point, int> ppi;
        typedef pair<int, int> pii;
        typedef pair<pii, int> pip;
        tria t[MAXPOINTS]:
        bool mrk[MAXPOINTS];
        int last[MAXPOINTS]:
        int childs[MAXPOINTS][3];
        int cnt;
        vector<ppi> points;
        set<pip> edges;
        vector<tria> res;
        int n;
        inline void add_edge(int a, int b, int c) {
                edges.insert(pip(pii(min(a, b), max(a, b)), c));
        inline void remove edge(int a, int b, int c){
                edges.erase(pip(pii(min(a, b), max(a, b)), c));
        int add_triangle(int a, int b, int c){
                if (cross(points[b].first - points[a].first, points[c
                    ].first - points[a].first) == 0)
                        return -1:
                if (cross(points[b].first - points[a].first, points[c
                    ].first - points[a].first) < 0)</pre>
                        swap(b, c);
                add edge(a, b, cnt);
                add_edge(b, c, cnt);
                add_edge(c, a, cnt);
                t[cnt] = tria(a, b, c);
                childs[cnt][0] = childs[cnt][1] = childs[cnt][2] = -1;
                mrk[cnt] = false;
                last[cnt] = -1;
                cnt++;
```

```
return cnt - 1;
inline void remove triangle(int v) {
        childs[v][0] = childs[v][1] = childs[v][2] = -1;
        remove_edge(t[v].a, t[v].b, v);
        remove edge(t[v].b, t[v].c, v);
        remove_edge(t[v].c, t[v].a, v);
inline void relax edge (const int &a, const int &b) {
        pii key(min(a, b), max(a, b));
        set<pip>::iterator it = edges.lower_bound(pip(key, -1)
            );
        if (it == edges.end() || it->first != key)
                return:
        set<pip>::iterator it2 = it;
        it2++;
        if (it2 == edges.end() || it2->first != key)
                return:
        int c1 = t[it->second].a + t[it->second].b + t[it->
            secondl.c - a - b:
        int c2 = t[it2->second].a + t[it2->second].b + t[it2->
            second].c - a - b;
        if (c1 > n | | c2 > n)
                return;
        if (inCircle(points[a].first, points[b].first, points[
            c1].first, points[c2].first) < 0 ||</pre>
                        inCircle(points[a].first, points[b].
                            first, points[c2].first, points[c1
                            l.first) < 0)
                int v1 = it->second;
                int v2 = it2 -> second:
                remove_triangle(v1);
                remove_triangle(v2);
                mrk[v1] = mrk[v2] = true;
                childs[v1][0] = childs[v2][0] = add_triangle(a
                    , c1, c2);
                childs[v1][1] = childs[v2][1] = add_triangle(b
                    , c1, c2);
                relax(childs[v1][0]);
                relax(childs[v1][1]);
inline void relax(int v) {
        relax_edge(t[v].a, t[v].b);
        relax_edge(t[v].b, t[v].c);
        relax_edge(t[v].c, t[v].a);
inline bool inLine(int a, int b, int c) {
        return cross (points[b].first - points[a].first, points
            [c].first - points[a].first) >= 0;
inline bool inTriangle(int a, int b, int c, int d) {
        return inLine(a, b, d) && inLine(b, c, d) && inLine(c,
             a, d);
void find(int v, int p, int cl) {
       if (last[v] == cl)
                return;
        bool reached = false;
        last[v] = cl;
```

```
for (int i = 0; i < 3; i++)
                    int u = childs[v][i];
                    if (u == -1)
                            continue;
                    reached = true;
                    if (mrk[u] || inTriangle(t[u].a, t[u].b, t[u].
                        c, p))
                            find(u, p, cl);
            if (reached)
                    return :
            remove_triangle(v);
            childs[v][0] = add_triangle(p, t[v].a, t[v].b);
            childs[v][1] = add_triangle(p, t[v].b, t[v].c);
            childs[v][2] = add_triangle(p, t[v].c, t[v].a);
            relax(childs[v][0]);
            relax(childs[v][1]);
            relax(childs[v][2]):
    void getRes(int v, int cl){
            if (last[v] == cl)
                    return;
            last[v] = cl;
            bool reached = false;
            for (int i = 0; i < 3; i++)
                    int u = childs[v][i];
                    if (u == -1)
                            continue;
                    reached = true;
                    getRes(u, cl);
            if (!reached && t[v].a < n && t[v].b < n && t[v].c < n
                    res.push_back(t[v]);
vector<tria> delaunay(vector<point> v) {
            cnt = 0;
            int cl = 0;
            points.clear();
            for (int i = 0; i < v.size(); i++)</pre>
                    points.push_back(ppi(v[i], i));
            random_shuffle(points.begin(), points.end());
            n = points.size();
            points.push_back(ppi(point(INF, INF), n));
            points.push_back(ppi(point(-INF * 3, INF), n + 1));
            points.push_back(ppi(point(INF, -INF * 3), n + 2));
            int root = add triangle(n, n + 1, n + 2);
            for (int i = 0; i < n; i++) {
                    // cout << "" << inTriangle(n, n+1, n+2, i
                        ) << endl;
                    find(root, i, cl++);
            res.clear();
            getRes(root, cl++);
            for (int i = 0; i < res.size(); i++){</pre>
                    res[i].a = points[res[i].a].second;
                    res[i].b = points[res[i].b].second;
                    res[i].c = points[res[i].c].second;
```

```
return res;
};
long double getRadius(pointD a, pointD b, pointD c) {
        pointD v1 = norm(b - a) + ((a + b) / 2);
        pointD v2 = norm(c - b) + ((b + c) / 2);
        pointD center = intersect((a + b) / 2, v1, (b + c) / 2, v2);
        pointD ret = a - center;
        return sqrt(dot(ret, ret));
Delaunay d;
int main(){
        srand(2019);
        ios::sync_with_stdio(false);
        cin.tie(0);
        int n:
        cin >> n:
        vector<point> v;
        for (int i = 0; i < n; i++)
                int x, y;
                cin >> x >> v;
                v.push_back(point(x, y));
        vector<tria> ans = d.delaunay(v);
        long double res = 0;
        for (int i = 0; i < ans.size(); i++)</pre>
                res = max(res, getRadius(v[ans[i].a], v[ans[i].b], v[
                     ans[i].c]));
        cout.precision(6);
        cout << fixed << res << endl;</pre>
```

### 3.9 Convex Hull 3D

```
struct point{
  int X, Y, Z;
  point(int x=0,int y=0,int z=0) {
    X=x; Y=y; Z=z;
  bool operator==(const point& rhs) const {
    return (rhs.X==this->X && rhs.Y==this->Y && rhs.Z==this->Z);
  bool operator<(const point& rhs) const {</pre>
    return rhs.X > this->X || (rhs.X == this->X && rhs.Y > this->Y) ||
          (rhs.X==this->X && rhs.Y==this->Y && rhs.Z>this->Z);
};
const int maxn=1000:
int n;
point P[maxn];
vector<point>ans;
queue<pii>Q;
set<pii>mark:
int cross2d(point p,point q) { return p.X*q.Y-p.Y*q.X; }
point operator - (point p, point q) { return point (p.X-q.X,p.Y-q.Y,p.Z-q.
point _cross(point u, point v) { return point(u.Y*v.Z-u.Z*v.Y,u.Z*v.X-u.
    X \star v \cdot Z \cdot u \cdot X \star v \cdot Y - u \cdot Y \star v \cdot X); }
```

```
point cross(point o, point p, point q) { return _cross(p-o, q-o); }
point shift(point p) { return point(p.Y,p.Z,p.X);}
point norm(point p)
  if(p.Y<p.X || p.Z<p.X) p=shift(p);</pre>
  if(p.Y<p.X) p=shift(p);
  return p;
int main()
  cin>>n;
  int mn=0;
  for (int i=0;i<n;i++) {</pre>
   cin>>P[i].X>>P[i].Y>>P[i].Z;
   if(P[i]<P[mn]) mn=i;
  int nx=(mn==0);
  for(int i=0:i<n:i++)</pre>
    if(i!=mn && i!=nx && cross2d(P[nx]-P[mn],P[i]-P[mn])>0)
      nx=i:
  O.push(pii(mn,nx));
  while(!Q.empty())
      int v=Q.front().first,u=Q.front().second;
      if (mark.find(pii(v,u))!=mark.end()) continue;
      mark.insert(pii(v,u));
      int p=-1;
      for (int q=0; q< n; q++)
        if(q!=v && q!=u)
          if (p==-1 || dot(cross(P[v],P[u],P[p]),P[q]-P[v])<0)</pre>
      ans.push_back(norm(point(v,u,p)));
      Q.push(pii(p,u));
      Q.push (pii (v,p));
  sort(ans.begin(),ans.end());
  ans.resize(unique(ans.begin(),ans.end())-ans.begin());
  for (int i=0; i < ans.size(); i++)</pre>
    cout << ans[i]. X << " " << ans[i]. Y << " " << ans[i]. Z << endl;</pre>
```

### 3.10 Half Plane Intersection

```
typedef int T;
typedef long long T2;
typedef long long T4; // maybe int128_t

const int MAXLINES = 100 * 1000 + 10;
const int INF = 20 * 1000 * 1000;

typedef pair<T, T> point;
typedef pair<point, point> line;

#define X first
#define Y second
#define A first
#define B second
```

```
// REPLACE ZERO WITH EPS FOR DOUBLE
point operator - (const point &a, const point &b) {
        return point(a.X - b.X, a.Y - b.Y);
T2 cross(point a, point b) {
        return ((T2)a.X * b.Y - (T2)a.Y * b.X);
bool cmp(line a, line b) {
        bool aa = a.A < a.B;</pre>
        bool bb = b.A < b.B;</pre>
        if (aa == bb) {
                point v1 = a.B - a.A;
                point v2 = b.B - b.A;
                if (cross(v1, v2) == 0)
                        return cross(b.B - b.A, a.A - b.A) > 0;
                else
                        return cross(v1, v2) > 0;
        else
                return aa;
bool parallel(line a, line b) {
        return cross(a.B - a.A, b.B - b.A) == 0;
pair<T2, T2> alpha(line a, line b) {
        return pair<T2, T2>(cross(b.A - a.A, b.B - b.A),
                                                 cross(a.B - a.A, b.B -
                                                      b.A));
bool fcmp(T4 f1t, T4 f1b, T4 f2t, T4 f2b) {
        if (f1b < 0) {
                f1t *= -1;
                f1b *= -1:
        if (f2b < 0) {
                f2t *= -1:
                f2b *= -1;
        return f1t * f2b < f2t * f1b; // check with eps
bool check(line a, line b, line c) {
        bool crs = cross(c.B - c.A, a.B - a.A) > 0;
        pair<T2, T2> a1 = alpha(a, b);
        pair<T2, T2> a2 = alpha(a, c);
        bool alp = fcmp(a1.A, a1.B, a2.A, a2.B);
        return (crs ^ alp);
bool notin(line a, line b, line c) { // is intersection of a and b in
    ccw direction of c?
        if (parallel(a, b))
                return false;
        if (parallel(a, c))
```

```
return cross(c.B - c.A, a.A - c.A) < 0;
        if (parallel(b, c))
                return cross (c.B - c.A, b.A - c.A) < 0;
        return ! (check(a, b, c) && check(b, a, c));
void print(vector<line> lines) {
        cerr << " " << endl; for (int i = 0; i < lines.size();</pre>
            i++)cerr << lines[i].A.X << " " << lines[i].A.Y << " -> "
            << lines[i].B.X << " " << lines[i].B.Y << endl;cerr << "
 << endl<< endl;
line dq[MAXLINES];
vector<line> half plane(vector<line> lines) {
        lines.push_back(line(point(INF, -INF), point(INF, INF)));
        lines.push_back(line(point(-INF, INF), point(-INF, -INF)));
        lines.push_back(line(point(-INF, -INF), point(INF, -INF)));
        lines.push_back(line(point(INF, INF), point(-INF, INF)));
        sort(lines.begin(), lines.end(), cmp);
        int ptr = 0;
        for (int i = 0; i < lines.size(); i++)</pre>
                if (i > 0 \& \&
                         (lines[i-1].A < lines[i-1].B) == (lines[i
                             ].A < lines[i].B) &&
                         parallel(lines[i - 1], lines[i]))
                         continue;
                else
                         lines[ptr++] = lines[i];
        lines.resize(ptr);
        if (lines.size() < 2)</pre>
                return lines;
        //print(lines):
        int f = 0, e = 0;
        dq[e++] = lines[0];
        dq[e++] = lines[1];
        for (int i = 2; i < lines.size(); i++) {</pre>
                while (f < e - 1 && notin(dq[e - 2], dq[e - 1], lines[
                    i1))
                //print (vector<line>(dq + f, dq + e));
                if (e == f + 1) {
                        T2 crs = cross(dq[f].B - dq[f].A, lines[i].B -
                              lines[i].A);
                         if (crs < 0)
                                 return vector<line>();
                         else if (crs == 0 && cross(lines[i].B - lines[
                             i].A, dq[f].B - lines[i].A) < 0)
                                return vector<line>();
                while (f < e - 1 \&\& notin(dq[f], dq[f + 1], lines[i]))
                dq[e++] = lines[i];
        while (f < e - 1 \& \& notin(dg[e - 2], dg[e - 1], dg[f]))
        while (f < e - 1 \&\& notin(dq[f], dq[f + 1], dq[e - 1]))
```

```
f++;
        vector<line> res;
        res.resize(e - f);
        for (int i = f; i < e; i++)
                res[i - f] = dq[i];
        return res;
int main() {
        int n;
        cin >> n;
        vector<line> lines;
        for (int i = 0; i < n; i++) {
                 int x1, y1, x2, y2;
                 cin >> x1 >> y1 >> x2 >> y2;
                 lines.push_back(line(point(x1, y1), point(x2, y2)));
        lines = half plane(lines):
        cout << lines.size() << endl;</pre>
        for (int i = 0; i < lines.size(); i++)</pre>
                 cout << lines[i].A.X << " " << lines[i].A.Y << " " <<</pre>
                     lines[i].B.X << " " << lines[i].B.Y << endl;</pre>
```

### 3.11 Minimum Enclosing Circle

```
const int N = 1000 * 100 + 10;
struct point {
    11 x, y, z;
typedef vector<point> circle;
bool ccw(point a, point b, point c) {
    return (b.x - a.x) * (c.y - a.y) - (c.x - a.x) * (b.y - a.y) >= 0;
bool incircle( circle a , point p ) {
    if( sz(a) == 0 ) return false;
    if(sz(a) == 1)
        return a[0].x == p.x && a[0].y == p.y;
    if(sz(a) == 2) {
        point mid = \{a[0].x+a[1].x, a[0].y+a[1].y\};
        return sq(2*p.x-mid.x) + sq(2*p.y-mid.y) \le sq(2*a[0].x-mid.x
            ) + sq(2*a[0].v-mid.v);
    if( !ccw(a[0], a[1], a[2]) )
        swap(a[0], a[2]);
    return incircle(a[0],a[1],a[2], p) >= 0;
point a[N];
circle solve(int i, circle curr) {
    assert(curr.size() <= 3);</pre>
    if( i == 0 )
        return curr:
    circle ret = solve(i-1, curr);
    if( incircle(ret, a[i-1]) )
        return ret;
    curr.pb(a[i-1]);
    return solve(i-1, curr);
int n;
```

```
void gg(circle c) {
    if( sz(c) == 1 ) {
        cout << ld(a[0].x) << " " << ld(a[0].y) << endl;</pre>
        cout << 0.1 << endl;
        return;
    if( sz(c) == 2 ) {
        point mid = \{c[0].x+c[1].x, c[0].y+c[1].y\};
        1d ret = sqrt(sq(2*c[0].x-mid.x) + sq(2*c[0].y-mid.y))/2;
        cout << ld(mid.x) / 2 << " " << ld(mid.y) /2 << endl;</pre>
        cout << ret << endl;</pre>
    } else {
        lpt a[3];
        for (int i = 0; i < 3; i++)
            a[i] = lpt(c[i].x, c[i].y);
        lpt A = ld(0.5) * (a[0] + a[1]), C = ld(0.5) * (a[1] + a[2]);
        lpt B = A + (a[1] - a[0]) * lpt(0, 1), D = C + (a[2] - a[1]) *
             lpt(0, 1);
        lpt center = intersection( A , B , C , D );
        ld ret = abs(a[0] - center);
        cout << center.real() << " " << center.imag() << endl;</pre>
        cout << ret << endl;
int main(){
    cin >> n;
    for(int i = 0; i < n; i++) {
        cin >> a[i].x >> a[i].y;
        a[i].z = sq(a[i].x) + sq(a[i].y);
    srand(time(NULL));
    for(int i = 1; i < n; i++)</pre>
        swap(a[i], a[rand()%(i+1)]);
    circle ans = solve(n, circle());
    cout << fixed << setprecision(3);</pre>
    gg(ans);
    return 0;
```

### 3.12 Number of integer points inside polygon

```
S = I + B / 2 - 1
```

### 3.13 Useful Geo Facts

```
Area of triangle with sides a, b, c: \operatorname{sqrt}(S * (S-a) * (S-b) * (S-c)) where S = (a+b+c)/2
Area of equilateral triangle: s^2 * \operatorname{sqrt}(3) / 4 where is side lenght Pyramid and cones volume: 1/3 area (base) * height

if p1=(x1, x2), p2=(x2, y2), p3=(x3, y3) are points on circle, the center is
x = -((x2^2 - x1^2 + y2^2 - y1^2) * (y3 - y2) - (x2^2 - x3^2 + y2^2 - y3^2) * (y1 - y2)) / (2*(x1 - x2) * (y3 - y2) - 2*(x3 - x2) * (y1 - y2))
y = -((y2^2 - y1^2 + x2^2 - x1^2) * (x3 - x2) - (y2^2 - y3^3 + x2^2 - x3^2) * (x1 - x2)) / (2*(y1 - y2) * (x3 - x2) - 2*(y3 - y2) * (x1 - x2))
```

### 3.14 Duality and properties

duality of point (a, b) is y = ax - b and duality of line y = ax + b is (a, -b) Properties:

- 1. p is on l iff l\* is in p\*
- 2. p is in intersection of l1 and l2 iff l1\* and l2\* lie on p\*
- 3. Duality preserve vertical distance
- 4. Translating a line in primal to moving vertically in dual
- 5. Rotating a line in primal to moving a point along a non-vertical line
- 6.  $li \cap lj$  is a vertex of lower envelope  $\iff$  (li\*, lj\*) is an edge of upper hull in dual

### 4 String

### 4.1 Suffix Array

```
const int maxn = 2e5 + 10;
int suf[maxn], rnk[maxn << 1], cnt[maxn], tmp[maxn], lcp[maxn];</pre>
bool scmp(int x, int y) {
        return (rnk[x] == rnk[y]? rnk[x+w] < rnk[y+w]: rnk[x] < rnk[y</pre>
void build_suffix_array(string &s){
        memset(rnk, 0, sizeof rnk);
        int n = s.size();
        for (int i = 0; i < n; i++) {
                rnk[i] = s[i] - 'a' + 1;
                 suf[i] = i;
        w = 1;
        sort(suf, suf + n, scmp);
        for (;; w <<= 1) {
                 memset(cnt, 0, sizeof cnt);
                 for (int i = 0; i < n; i++) cnt[rnk[i]]++;</pre>
                 for (int i = 1; i \le max(26, n); i++) {
                         cnt[i] += cnt[i-1];
                 for (int i = n-1; ~i; i--) {
                         if (suf[i] < w) continue;</pre>
                         tmp[--cnt[rnk[suf[i]-w]]] = suf[i]-w;
                 for (int i = n-w; i < n; i++) {
                         tmp[--cnt[rnk[i]]] = i;
                 for (int i = 0; i < n; i++) suf[i] = tmp[i];</pre>
                 tmp[suf[0]] = 1;
                 for (int i = 1; i < n; i++) {
                         tmp[suf[i]] = tmp[suf[i-1]] + scmp(suf[i-1],
                              suf[i]);
```

### 4.2 Suffix Automata

```
const int maxn = 2 e5 + 42; // Maximum amount of states
map < char , int > to [ maxn ]; // Transitions
int link [ maxn ]; // Suffix links
int len [ maxn ]; // Lengthes of largest strings in states
int last = 0; // State corresponding to the whole string
int sz = 1; // Current amount of states
void add_letter ( char c ) { // Adding character to the end
   int p = last ; // State of string s
   last = sz ++; // Create state for string sc
   len [ last ] = len [ p ] + 1;
    for (; to [ p ][ c ] == 0; p = link [ p ]) // (1)
        to [p][c] = last; // Jumps which add new suffixes
   if ( to [ p ][ c ] == last ) { // This is the first occurrence of
        c if we are here
        link [ last ] = 0;
        return ;
   int q = to [ p ] [ c ];
    if ( len [ q ] == len [ p ] + 1) {
        link [ last ] = q;
        return ;
    // We split off cl from q here
   int cl = sz ++;
   to [cl] = to [q]; // (2)
   link [ cl ] = link [ q ];
   len [ cl ] = len [ p ] + 1;
   link [ last ] = link [ q ] = cl ;
    for (; to [p][c] == q; p = link [p]) // (3)
        to [ p ][ c ] = cl ; // Redirect transitions where needed
```

### 4.3 Suffix Tree

```
#define fpos adla
const int inf = 1e9;
const int maxn = 1e4;
char s[maxn];
map<int, int> to[maxn];
```

```
int len[maxn], fpos[maxn], link[maxn];
int node, pos;
int sz = 1, n = 0;
int make_node(int _pos, int _len) {
    fpos[sz] = \_pos;
    len [sz] = len;
    return sz++;
void go edge() {
    while(pos > len[to[node][s[n - pos]]]) {
        node = to[node][s[n - pos]];
        pos -= len[node];
void add_letter(int c) {
    s[n++] = c;
    pos++;
    int last = 0;
    while(pos > 0) {
        go_edge();
        int edge = s[n - pos];
        int &v = to[node][edge];
        int t = s[fpos[v] + pos - 1];
        if(v == 0) {
            v = make_node(n - pos, inf);
            link[last] = node;
            last = 0;
        } else if(t == c) {
            link[last] = node;
            return;
        } else {
            int u = make_node(fpos[v], pos - 1);
            to[u][c] = make_node(n - 1, inf);
            to[u][t] = v;
            fpos[v] += pos - 1;
            len [v] -= pos - 1;
            v = u:
            link[last] = u;
            last = u;
        if(node == 0)
            pos--;
        else
            node = link[node];
int main() {
    len[0] = inf;
    string s;
    cin >> s;
    int ans = 0;
    for(int i = 0; i < s.size(); i++)</pre>
        add_letter(s[i]);
    for(int i = 1; i < sz; i++)
        ans += min((int)s.size() - fpos[i], len[i]);
    cout << ans << "\n";
```

### 4.4 Palindromic Tree

```
int n, last, sz;
void init() {
    s[n++] = -1;
   link[0] = 1;
    len[1] = -1;
    sz = 2;
int get_link(int v) {
    while (s[n - len[v] - 2] != s[n - 1]) v = link[v];
    return v;
void add_letter(int c) {
    s[n++] = c;
    last = get_link(last);
    if(!to[last][c]) {
        len [sz] = len[last] + 2;
        link[sz] = to[get_link(link[last])][c];
        to[last][c] = sz++;
    last = to[last][c];
```

### 5 Data structure

### 5.1 Ordered Set

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
typedef tree<int, null_type, less<int>, rb_tree_tag,
    tree_order_statistics_node_update> ordered_set; //less_equal<int>
    for multiset
ordered_set s;
s.insert(1);
s.find_by_order(idx) //iterator to the idx
s.order_of_key(key) // index of the key
//implementation with fenwick
const int maxn = 1e5 + 10;
int f[maxn];
void change (int idx, int x) { // x = 1 insert, 0 erase
        for (; idx < maxn; idx += idx & -idx) f[idx] += x;</pre>
int lg = 20;
int find by order(int idx) {
        int res = 0;
        for (int i = lq-1; ~i; i--) {
                int tmp = res + (1 << i);</pre>
                if (tmp < maxn && f[tmp] < idx) {</pre>
                        idx -= f[tmp]:
                         res = tmp;
        return res + 1; // maxn if not exist
int order_of_key(int key) {
        int res = 0;
        for (; key; key -= key & -key) res += f[key];
        return res;
```

### 5.2 Treap

```
struct item {
   int key, prior;
    item * 1, * r;
    item() { }
    item (int key, int prior) : key(key), prior(prior), l(NULL), r(
        NULL) { }
typedef item * pitem;
void split (pitem t, int key, pitem & l, pitem & r) {
   if (!t)
       1 = r = NULL;
    else if (key < t->key)
        split (t->1, key, 1, t->1), r = t;
        split (t->r, key, t->r, r), l = t;
void insert (pitem & t, pitem it) {
    if (!t)
       t = it;
    else if (it->prior > t->prior)
        split (t, it->key, it->l, it->r), t = it;
        insert (it->key < t->key ? t->l : t->r, it);
void merge (pitem & t, pitem l, pitem r) {
   if (!l || !r)
       t = 1 ? 1 : r;
    else if (l->prior > r->prior)
       merge (1->r, 1->r, r), t = 1;
    else
       merge (r->1, 1, r->1), t = r;
void erase (pitem & t, int key) {
   if (t->kev == kev)
       merge (t, t->1, t->r);
    else
        erase (key < t->key ? t->1 : t->r, key);
pitem unite (pitem 1, pitem r) {
   if (!1 || !r) return 1 ? 1 : r;
    if (l->prior < r->prior) swap (l, r);
   pitem lt, rt;
   split (r, 1->key, lt, rt);
   1->1 = unite (1->1, 1t);
   1->r = unite (1->r, rt);
   return 1;
```

### 5.3 Dynimic convex hull

```
const ld is_query = -(1LL << 62);
struct Line {
   ld m, b;
   mutable std::function<const Line *()> succ;
   bool operator<(const Line &rhs) const {
      if (rhs.b != is_query) return m < rhs.m;</pre>
```

```
const Line *s = succ();
        if (!s) return 0;
        ld x = rhs.m;
        return b - s -> b < (s -> m - m) * x;
struct HullDynamic : public multiset<Line> { // dynamic upper hull +
    max value query
   bool bad(iterator y) {
        auto z = next(y);
        if (y == begin()) {
            if (z == end()) return 0;
            return y->m == z->m && y->b <= z->b;
        auto x = prev(y);
       if (z == end()) return y->m == x->m && y->b <= x->b;
        return (x-b-y-b) * (z-m-y-m) >= (y-b-z-b) * (y-m)
   void insert_line(ld m, ld b) {
        auto y = insert({m, b});
        v -> succ = [=] { return next(v) == end() ? 0 : &*next(v); };
        if (bad(y)) {
            erase(v);
            return;
        while (next(v) != end() && bad(next(v))) erase(next(v));
        while (y != begin() && bad(prev(y))) erase(prev(y));
   ld best(ld x) {
        auto l = *lower_bound((Line) {x, is_query});
        return 1.m * x + 1.b;
};
```

### 6 Graph

### 6.1 Maximum matching - Edmond's blossom

```
/*
GETS:
n->number of vertices
you should use add edge (u, v) and
add pair of vertices as edges (vertices are 0..n-1)
(note: please don't add multiple edge)
output of edmonds() is the maximum matching in general graph
match[i] is matched pair of i (-1 if there isn't a matched pair)
O(mn^2)
*/
struct struct edge{int v:struct edge* nxt;};
typedef struct_edge* edge;
const int MAXN=500;
struct Edmonds{
        struct_edge pool[MAXN*MAXN*2];
        edge top=pool,adj[MAXN];
        int n, match[MAXN], qh, qt, q[MAXN], father[MAXN], base[MAXN];
        bool inq[MAXN], inb[MAXN];
```

```
void add_edge(int u,int v) {
        top->v=v,top->nxt=adj[u],adj[u]=top++;
        top->v=u,top->nxt=adj[v],adj[v]=top++;
int LCA(int root,int u,int v) {
        static bool inp[MAXN];
        memset(inp, 0, sizeof(inp));
        while(1){
                inp[u=base[u]]=true;
                if (u==root) break;
                u=father[match[u]];
        while(1){
                if (inp[v=base[v]]) return v;
                else v=father[match[v]];
void mark_blossom(int lca,int u) {
        while (base[u]!=lca){
                int v=match[u];
                inb[base[u]]=inb[base[v]]=true;
                u=father[v];
                if (base[u]!=lca) father[u]=v;
void blossom_contraction(int s,int u,int v) {
        int lca=LCA(s,u,v);
        memset(inb, 0, sizeof(inb));
        mark blossom(lca,u);
        mark blossom(lca,v);
        if (base[u]!=lca)
                 father[u]=v;
        if (base[v]!=lca)
                father[v]=u;
        for (int u=0; u< n; u++)
                if (inb[base[u]]){
                         base[u]=lca:
                         if (!ing[u])
                                 ing[q[++qt]=u]=true;
int find_augmenting_path(int s){
        memset(ing, 0, sizeof(ing));
        memset(father,-1,sizeof(father));
        for (int i=0;i<n;i++) base[i]=i;</pre>
        inq[q[qh=qt=0]=s]=true;
        while (gh<=gt) {</pre>
                int u=q[qh++];
                for (edge e=adj[u];e;e=e->nxt) {
                         int v=e->v;
                         if (base[u]!=base[v] && match[u]!=v){
                                 if (v==s || (match[v]!=-1 &&
                                      father[match[v]]!=-1))
                                          blossom contraction(s,
                                              u, v);
                                 else if (father[v]==-1) {
                                          father[v]=u;
                                          if (match[v] == -1)
                                                  return v:
                                          else if (!ing[match[v
                                              ]])
```

```
inq[q[++qt] =
                                                       match[v]]=
                                                       true;
        return -1;
int augment_path(int s,int t){
        int u=t, v, w;
        while (u!=-1) {
                v=father[u];
                w=match[v];
                match[v]=u;
                match[u]=v;
                u=w;
        return t!=-1;
int edmonds(){
        int matchc=0;
        memset (match, -1, sizeof (match));
        for (int u=0; u<n; u++)
                if (match[u] == -1)
                         matchc+=augment_path(u,
                              find_augmenting_path(u));
        return matchc;
```

### 6.2 Biconnected components

};

```
vector<int> adj[maxn];
bool vis[maxn];
int dep[maxn], par[maxn], lowlink[maxn];
vector<vector<int> > comp;
stack<int> st;
void dfs(int u, int depth = 0, int parent = -1) {
        vis[u] = true;
        dep[u] = depth;
        par[u] = parent;
        lowlink[u] = depth;
        st.push(u);
        for (int i = 0; i < adj[u].size(); i++) {</pre>
                int v = adj[u][i];
                if (!vis[v])
                         dfs(v, depth + 1, u);
                         lowlink[u] = min(lowlink[u], lowlink[v]);
                else
                         lowlink[u] = min(lowlink[u], dep[v]);
        if (lowlink[u] == dep[u] - 1){
                comp.push_back(vector<int>());
                while (st.top() != u)
                         comp.back().push_back(st.top());
                         st.pop();
```

### 6.3 Flow - Dinic

```
struct FlowEdge {
    int v, u;
    long long cap, flow = 0;
    FlowEdge(int v, int u, long long cap) : v(v), u(u), cap(cap) {}
};
struct Dinic {
    const long long flow_inf = 1e18;
    vector<FlowEdge> edges;
    vector<vector<int>> adj;
    int n, m = 0;
    int s, t;
    vector<int> level, ptr;
    queue<int> q;
    Dinic(int n, int s, int t) : n(n), s(s), t(t) {
        adj.resize(n);
        level.resize(n);
        ptr.resize(n);
    void add_edge(int v, int u, long long cap) {
        edges.emplace_back(v, u, cap);
        edges.emplace_back(u, v, 0);
        adj[v].push_back(m);
        adi[u].push back(m + 1);
        m += 2;
    bool bfs() {
        while (!q.empty()) {
            int v = q.front();
            q.pop();
            for (int id : adj[v]) {
                if (edges[id].cap == edges[id].flow)
                    continue;
                if (level[edges[id].u] != -1)
                    continue;
                level[edges[id].u] = level[v] + 1;
                q.push(edges[id].u);
        return level[t] != -1;
    long long dfs(int v, long long pushed) {
```

```
if (pushed == 0)
            return 0:
        if (v == t)
            return pushed;
        for (int& cid = ptr[v]; cid < (int)adj[v].size(); cid++) {</pre>
            int id = adi[v][cid];
            int u = edges[id].u;
            if (level[v] + 1 != level[u])
                continue;
            long long tr = dfs(u, min(pushed, edges[id].cap - edges[id
                ].flow));
            if (tr == 0)
                continue;
            edges[id].flow += tr;
            edges[id ^ 1].flow -= tr;
            return tr;
        return 0;
   long long flow() {
        long long f = 0;
        while (true) {
            fill(level.begin(), level.end(), -1);
            level[s] = 0;
            q.push(s);
            if (!bfs())
                break;
            fill(ptr.begin(), ptr.end(), 0);
            while (long long pushed = dfs(s, flow_inf)) {
                f += pushed;
        return f;
};
```

### 6.4 Min Cost Max Flow

```
template <typename F, typename C, int MAXN, int MAXM> struct
   MinCostMaxFlow {
    struct Edge {
        int from, to;
        F cap;
        C cost;
    };
   Edge E[2 * MAXM];
   int m, par[MAXN], s, t;
   C dist[MAXN], cost = 0, pot[MAXN];
   F mn[MAXN], flow = 0;
   vector<int> adj[MAXN];
   MinCostMaxFlow() { memset(pot, 0, sizeof pot); }
   inline void add_edge(int u, int v, F cap, C cost) {
        adj[u].push_back(m);
        E[m++] = \{u, v, cap, cost\};
        adj[v].push_back(m);
        E[m++] = \{v, u, 0, -cost\};
   inline void dijkstra() {
        fill(dist, dist + MAXN, numeric_limits<C>::max());
```

```
fill(par, par + MAXN, -1);
    priority_queue<pair<C, int>, vector<pair<C, int>>, greater
        pair<C, int>>> pq;
    dist[s] = 0;
   mn[s] = numeric_limits<F>::max();
    pg.push({dist[s], s});
    while (!pq.empty()) {
        int v = pq.top().Y;
        C d = pq.top().X;
        pq.pop();
        if (d != dist[v]) continue;
        for (int id : adj[v]) {
            int u = E[id].to;
            if (!E[id].cap) continue;
            C d_u = d + pot[v] - pot[u] + E[id].cost;
            if (d_u < dist[u]) {</pre>
                dist[u] = d_u;
                pq.push({d_u, u});
                par[u] = id;
                mn[u] = min(mn[v], E[id].cap);
inline F solve() {
   dijkstra();
    if (par[t] == -1) return 0;
    F c = min(mn[t], flow), v = t;
   cost += c * (dist[t] + pot[t]);
    while (v != s) {
        int id = par[v];
        E[id].cap -= c;
        E[id ^1].cap += c;
        v = E[id].from;
    for (int v = 0; v < MAXN; v++)
        if (par[v] != -1)
           pot[v] += dist[v];
    return c:
inline C min_cost(int _s, int _t, F _flow) {
    s = \_s, t = \_t, flow = \_flow;
    while (true) {
        F c = solve();
        if (c == 0)
            break;
    return cost;
```

### 6.5 Maximum weighted matching - Hungarian

};

```
const int N = 2002;
const int INF = 1e9;
int hn, weight[N][N];
int x[N], y[N];
int hungarian() // maximum weighted perfect matching{
    int n = hn;
```

```
int p, q;
vector<int> fx(n, -INF), fy(n, 0);
fill(x, x + n, -1);
fill(y, y + n, -1);
for (int i = 0; i < n; ++i)
        for (int j = 0; j < n; ++j)
                fx[i] = max(fx[i], weight[i][j]);
for (int i = 0; i < n; ) {</pre>
        vector<int> t(n, -1), s(n+1, i);
        for (p = 0, q = 1; p < q && x[i] < 0; ++p) {
                int k = s[p];
                for (int j = 0; j < n \&\& x[i] < 0; ++j)
                        if (fx[k] + fy[j] == weight[k][j] && t
                             [\dot{1}] < 0) {
                                 s[q++] = y[j], t[j] = k;
                                 if (y[j] < 0) // match found!
                                         for (int p = j; p >=
                                              0; j = p
                                                  y[j] = k = t[j]
                                                      ], p = x[k]
                                                      ], x[k] =
                                                      j;
        if (x[i] < 0) {
                int d = INF;
                for (int k = 0; k < q; ++k)
                        for (int j = 0; j < n; ++j)
                                 if (t[j] < 0) d = min(d, fx [s
                                     [k]] + fy[j] - weight[s]k
                                     ]][j]);
                for (int j = 0; j < n; ++j) fy[j] += (t[j] <
                    0? 0: d);
                for (int k = 0; k < q; ++k) fx[s[k]] -= d;
        } else ++i;
for (int i = 0; i < n; ++i) ret += weight[i][x[i]];</pre>
return ret:
```

### 6.6 Ear decomposition

- 1- Find a spanning tree of the given graph and choose a root for the tree.
- 2- Determine, for each edge uv that is not part of the tree, the distance between the root and the lowest common ancestor of u and v.
- 3- For each edge uv that is part of the tree, find the corresponding "master edge", a non-tree edge wx such that the cycle formed by adding wx to the tree passes through uv and such that, among such edges, w and x have a lowest common ancestor that is as close to the root as possible (with ties broken by edge identifiers).
- 4- Form an ear for each non-tree edge, consisting of it and the tree edges for which it is the master, and order the ears by their master edges' distance from the root (with the same tie-breaking rule).

### 6.7 Stoer-Wagner min cut $O(n^3)$

```
const int N = -1, MAXW = -1;
int g[N][N], v[N], w[N], na[N];
bool a[N];
int minCut( int n ) // initialize g[][] before calling!{
    for ( int i = 0; i < n; i++ ) v[i] = i;
    int best = MAXW * n * n;
    while (n > 1)
        // initialize the set A and vertex weights
        a[v[0]] = true:
        for( int i = 1; i < n; i++ ) {</pre>
            a[v[i]] = false;
            na[i - 1] = i;
            w[i] = q[v[0]][v[i]];
        // add the other vertices
        int prev = v[0];
        for( int i = 1; i < n; i++ ) {</pre>
            // find the most tightly connected non-A vertex
            int zi = -1;
            for ( int j = 1; j < n; j++ )
                if( !a[v[j]] \&\& (zj < 0 || w[j] > w[zj] ) )
                    zj = j;
            // add it to A
            a[v[zj]] = true;
            // last vertex?
            if(i == n - 1){
                // remember the cut weight
                best = min(best, w[zj]);
                // merge prev and v[zj]
                for ( int j = 0; j < n; j++ )
                    g[v[j]][prev] = g[prev][v[j]] += g[v[zj]][v[j]];
                v[zj] = v[--n];
                break:
            prev = v[zi];
            // update the weights of its neighbors
            for ( int j = 1; j < n; j++ ) if ( !a[v[j]] )
                w[j] += g[v[zj]][v[j]];
    return best;
```

### **6.8** Directed minimum spanning tree $O(m \log n)$

```
const int INF = 2e7;
struct MinimumAborescense{
        struct edge {
                int src, dst, weight;
        };
        struct union find {
                vector<int> p;
                union_find(int n) : p(n, -1) { };
                bool unite(int u, int v) {
                         if ((u = root(u)) == (v = root(v))) return
                         if (p[u] > p[v]) swap(u, v);
                         p[u] += p[v]; p[v] = u;
                         return true;
                bool find(int u, int v) { return root(u) == root(v); }
                int root(int u) { return p[u] < 0 ? u : p[u] = root(p[</pre>
                int size(int u) { return -p[root(u)]; }
        };
        struct skew_heap {
                 struct node {
                         node *ch[2];
                         edge kev;
                         int delta;
                 } *root;
                skew heap() : root(0) { }
                void propagate(node *a) {
                         a->kev.weight += a->delta;
                         if (a->ch[0]) a->ch[0]->delta += a->delta;
                         if (a->ch[1]) a->ch[1]->delta += a->delta;
                         a \rightarrow delta = 0:
                node *merge(node *a, node *b) {
                         if (!a || !b) return a ? a : b;
                         propagate(a); propagate(b);
                         if (a->key.weight > b->key.weight) swap(a, b);
                         a - ch[1] = merge(b, a - ch[1]);
                         swap(a->ch[0], a->ch[1]);
                         return a:
                void push(edge kev) {
                         node *n = new node():
                         n \rightarrow ch[0] = n \rightarrow ch[1] = 0;
                         n\rightarrow kev = kev; n\rightarrow delta = 0;
                         root = merge(root, n);
                void pop() {
                         propagate(root);
                         node *temp = root;
                         root = merge(root->ch[0], root->ch[1]);
                edge top() {
                         propagate(root);
                         return root->key;
                bool empty() {
                         return !root:
                void add(int delta) {
                         root->delta += delta:
```

```
void merge(skew heap x) {
                root = merge(root, x.root);
};
vector<edge> edges;
void add_edge(int src, int dst, int weight) {
        edges.push_back({src, dst, weight});
int n;
void make_graph(int _n) {
        n = _n;
        edges.clear();
int dmst(int r) {
        union find uf(n):
        vector<skew_heap> heap(n);
        for (auto e: edges)
                heap[e.dst].push(e);
        double score = 0;
        vector<int> seen(n, -1);
        seen[r] = r;
        for (int s = 0; s < n; ++s) {</pre>
                vector<int> path;
                for (int u = s; seen[u] < 0;) {</pre>
                        path.push_back(u);
                        seen[u] = s;
                        if (heap[u].empty()) return INF;
                        edge min e = heap[u].top();
                        score += min_e.weight;
                        heap[u].add(-min_e.weight);
                        heap[u].pop();
                        int v = uf.root(min_e.src);
                        if (seen[v] == s) {
                                 skew_heap new_heap;
                                 while (1) {
                                         int w = path.back();
                                         path.pop_back();
                                         new_heap.merge(heap[w
                                         if (!uf.unite(v, w))
                                             break:
                                 heap[uf.root(v)] = new_heap;
                                 seen[uf.root(v)] = -1;
                        u = uf.root(v);
        return score;
```

### **6.9** Directed minimum spanning tree O(nm)

};

```
/*
    GETS:
        call make_graph(n) at first
        you should use add_edge(u,v,w) and
        add pair of vertices as edges (vertices are 0..n-1)
```

```
GIVES:
                output of dmst(v) is the minimum arborescence with
                    root v in directed graph
                 (-1 if it hasn't a spanning arborescence with root v)
        (mn)
const int INF = 2e7;
struct MinimumAborescense{
        int n;
        struct edge {
                int src, dst;
                int weight;
        vector<edge> edges;
        void make_graph(int _n) {
                n=\underline{n};
                edges.clear();
        void add edge(int u, int v, int w) {
                edges.push_back({u, v, w});
        int dmst(int r) {
                int N = n;
                for (int res = 0; ;) {
                         vector<edge> in(N, \{-1, -1, (int) INF\});
                         vector<int> C(N, -1);
                         for (auto e: edges)
                                 if (in[e.dst].weight > e.weight)
                                         in[e.dst] = e;
                         in[r] = \{r, r, 0\};
                         for (int u = 0; u < N; ++u) { // no comming
                             edge ==> no aborescense
                                 if (in[u].src < 0) return -1;</pre>
                                 res += in[u].weight;
                         vector<int> mark(N, -1); // contract cycles
                         int index = 0;
                         for (int i = 0; i < N; ++i) {
                                 if (mark[i] != -1) continue;
                                 int u = i;
                                 while (mark[u] == -1) {
                                         mark[u] = i;
                                         u = in[u].src;
                                 if (mark[u] != i || u == r) continue;
                                 for (int v = in[u].src; u != v; v = in
                                     [v].src) C[v] = index;
                                 C[u] = index++;
                         if (index == 0) return res; // found
                             arborescence
                         for (int i = 0; i < N; ++i) // contract</pre>
                                 if (C[i] == -1) C[i] = index++;
                         vector<edge> next;
                         for (auto &e: edges)
                                 if (C[e.src] != C[e.dst] && C[e.dst]
                                     != C[r])
                                         next.push_back({C[e.src], C[e.
                                              dst], e.weight - in[e.dst
```

```
| .weight});
edges.swap(next);
N = index; r = C[r];
}
};
```

### 6.10 Dominator tree

```
struct DominatorTree{
        vector<int> adj[MAXN], radj[MAXN], tree[MAXN], bucket[MAXN];
            // SET MAXIMUM NUMBER OF NODES
        int sdom[MAXN], par[MAXN], idom[MAXN], dsu[MAXN], label[MAXN];
        int arr[MAXN], rev[MAXN], cnt;
        void clear(){
                for (int i = 0; i < MAXN; i++) {</pre>
                        adj[i].clear();
                         radj[i].clear();
                         tree[i].clear();
                         sdom[i] = idom[i] = dsu[i] = label[i] = i;
                        arr[i] = -1;
                cnt = 0;
        void add_edge(int u, int v) {
                adj[u].push_back(v);
        void dfs(int v){
                arr[v] = cnt;
                rev[cnt] = v;
                cnt++;
                for (int i = 0; i < adj[v].size(); i++) {</pre>
                         int u = adj[v][i];
                         if (arr[u] == -1) {
                                 dfs(u);
                                 par[arr[u]] = arr[v];
                         radj[arr[u]].push_back(arr[v]);
        int find(int v, int x = 0) {
                if (dsu[v] == v)
                         return (x ? -1 : v);
                int u = find(dsu[v], x + 1);
                if (u < 0)
                         return v;
                if (sdom[label[dsu[v]]] < sdom[label[v]])</pre>
                        label[v] = label[dsu[v]];
                dsu[v] = u;
                return (x ? u : label[v]);
        void merge(int u, int v) {
                dsu[v] = u;
        void build(int root) {
                dfs(root):
                int n = cnt;
                for (int v = n - 1; v >= 0; v--) {
                         for (int i = 0; i < radj[v].size(); i++) {</pre>
```

int u = radj[v][i];

```
sdom[v] = min(sdom[v], sdom[find(u)]);
                         if (v > 0)
                                 bucket[sdom[v]].push_back(v);
                         for (int i = 0; i < bucket[v].size(); i++){</pre>
                                 int u = bucket[v][i];
                                 int w = find(u);
                                 if (sdom[u] == sdom[w])
                                         idom[u] = sdom[u];
                                 else
                                         idom[u] = w;
                         if (v > 0)
                                 merge(par[v], v);
                for (int v = 1; v < n; v++) {
                         if (idom[v] != sdom[v])
                                 idom[v] = idom[idom[v]];
                        tree[rev[v]].push_back(rev[idom[v]]);
                         tree[rev[idom[v]]].push_back(rev[v]);
        DominatorTree() {
                clear();
};
```

### 7 Combinatorics

### 7.1 LP simplex

```
// Two-phase simplex algorithm for solving linear programs of the form
       maximize c^T x
11
       subject to Ax <= b
                   x >= 0
// INPUT: A -- an m x n matrix
        b -- an m-dimensional vector
         c -- an n-dimensional vector
         x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution (infinity if unbounded
          above, nan if infeasible)
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
  int m, n;
 VI B, N;
 VVD D;
  LPSolver (const VVD &A, const VD &b, const VD &c) :
   m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2, VD(n + 2)) {
    for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D[i][j] =
   for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n] = -1; D[i][n + i]
         1] = b[i];
```

```
for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
    N[n] = -1; D[m + 1][n] = 1;
  void Pivot(int r, int s) {
    double inv = 1.0 / D[r][s];
    for (int i = 0; i < m + 2; i++) if (i != r)
      for (int j = 0; j < n + 2; j++) if (j != s)
        D[i][j] -= D[r][j] * D[i][s] * inv;
    for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] *= inv;
    for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] *= -inv;
    D[r][s] = inv;
    swap(B[r], N[s]);
  bool Simplex(int phase) {
    int x = phase == 1 ? m + 1 : m;
    while (true) {
      int s = -1;
      for (int j = 0; j \le n; j++) {
        if (phase == 2 && N[i] == -1) continue;
        if (s == -1 \mid | D[x][j] < D[x][s] \mid | D[x][j] == D[x][s] && N[j]
              < N[s]) s = j;
      if (D[x][s] > -EPS) return true;
      for (int i = 0; i < m; i++) {</pre>
        if (D[i][s] < EPS) continue;</pre>
        if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s]
          (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r][s]) && B[i] <
               B[r]) r = i;
      if (r == -1) return false;
      Pivot(r, s);
  DOUBLE Solve (VD &x) {
    for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n + 1]) r = i;
    if (D[r][n + 1] < -EPS) {
      Pivot(r, n):
      if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return -</pre>
          numeric_limits<DOUBLE>::infinity();
      for (int i = 0; i < m; i++) if (B[i] == -1) {
        int s = -1;
        for (int j = 0; j <= n; j++)</pre>
          if (s == -1 \mid \mid D[i][j] < D[i][s] \mid \mid D[i][j] == D[i][s] && N[
              j] < N[s]) s = j;
        Pivot(i, s);
    if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
    x = VD(n);
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];
    return D[m][n + 1];
};
```

### 7.2 LP for game theory

```
ld solve(int n, int m) {
```

```
VVD A;
VD B;
VD C;
for (int j = 0; j < m; j++) {
    VD v;
    for (int i = 0; i < n; i++)
        v.pb(-mat[i][j]);
    v.pb(1);
    A.pb(v);
VD v;
for (int i = 0; i < n; i++)
    v.pb(1);
v.pb(0);
A.pb(v);
v.clear();
for (int i = 0; i < n; i++)
    v.pb(-1);
(0) dq.v
A.pb(v);
for (int i = 0; i < m; i++)</pre>
    B.pb(0);
B.pb(1);
B.pb (-1);
for (int i = 0; i < n; i++)
    C.pb(0);
C.pb(1);
LPSolver solver (A, B, C);
VD x;
ld res = solver.Solve(x);
return res;
```

### 7.3 FFT

```
typedef complex<double> point;
const double pi=acos(-1);
const int mod=998244353;
const int N=(1<<20);
int rev[N];
void FFT(point *A, int n, bool inv) {
    int lg=__builtin_ctz(n);
    for (int i=1; i<n; i++) {</pre>
         rev[i] = (rev[i>>1]>>1) | ((i&1) << (lq-1));
         if (rev[i]<i) swap(A[i], A[rev[i]]);</pre>
    for (int len=1; len<n; len<<=1) {</pre>
         double theta=pi/len;
         if (inv) theta*=-1;
         point wn=point(cos(theta), sin(theta));
         for (int i=0: i<n: i+=2*len) {</pre>
             point w=1;
             for (int j=i; j<i+len; j++) {</pre>
                 point x=A[j], y=w*A[j+len];
                 A[j]=x+y;
                 A[j+len]=x-v;
                 w \star = wn;
```

```
}
if (inv) {
    for (int i=0; i<n; i++) A[i]/=n;
}
</pre>
```

### 7.4 NTT

```
11 powmod(ll a, ll b) {
    11 res=1;
    for (; b; b>>=1, a=a*a%mod) if (b&1) res=res*a%mod;
    return res;
inline void fix(ll &x){
    if (x<0) x+=mod:
    if (x \ge mod) x = mod:
void NTT(ll *A, int n, bool inv) {
    int lg=__builtin_ctz(n);
    for (int i=1; i<n; i++) {</pre>
         rev[i] = (rev[i>>1]>>1) | ((i&1) << (lq-1));
        if (rev[i]<i) swap(A[i], A[rev[i]]);</pre>
    for (int len=1; len<n; len<<=1) {</pre>
        11 wn=powmod(5, mod/2/len);
        if (inv) wn=powmod(wn, mod-2);
        for (int i=0; i<n; i+=2*len) {</pre>
             11 w=1;
             for (int j=i; j<i+len; j++) {</pre>
                 ll x=A[j], y=w*A[j+len]%mod;
                 fix(A[j]=x+y);
                 fix(A[j+len]=x-y);
                 w=w*wn%mod;
    if (inv) {
        11 nn=powmod(n, mod-2);
         for (int i=0; i<n; i++) A[i]=A[i]*nn%mod;</pre>
```

### **7.5 FWHT**

```
A[i] >>= m;
}
```

### 7.6 Stirling 1

```
const int mod = 998244353;
const int root = 15311432;
const int root_1 = 469870224;
const int root pw = 1 << 23:
const int N = 400004;
vector<int> v[N];
11 modInv(11 a, 11 mod = mod) {
        11 \times 0 = 0, \times 1 = 1, \times 0 = 0, \times 1 = a;
        while(r1){
                11 q = r0 / r1;
                 x0 -= q * x1; swap(x0, x1);
                 r0 = q * r1; swap(r0, r1);
        return x0 < 0 ? x0 + mod : x0;
void fft (vector<int> &a, bool inv) {
        int n = (int) a.size();
        for (int i=1, j=0; i<n; ++i) {</pre>
                int bit = n >> 1;
                 for (; j>=bit; bit>>=1)
                         j -= bit;
                 i += bit:
                 if (i < j)
                         swap (a[i], a[j]);
        for (int len=2; len<=n; len<<=1) {</pre>
                 int wlen = inv ? root_1 : root;
                 for (int i=len; i<root_pw; i<<=1)</pre>
                         wlen = int (wlen * 111 * wlen % mod);
                 for (int i=0; i<n; i+=len) {</pre>
                         int w = 1;
                         for (int j=0; j<len/2; ++j) {</pre>
                                  int u = a[i+j], v = int (a[i+j+len/2]
                                       * 111 * w % mod);
                                  a[i+j] = u+v < mod ? u+v : u+v-mod;
                                  a[i+j+len/2] = u-v >= 0 ? u-v : u-v+
                                      mod;
                                  w = int (w * 111 * wlen % mod);
        if(inv) {
                 int nrev = modInv(n, mod);
                 for (int i=0; i<n; ++i)</pre>
                         a[i] = int (a[i] * 111 * nrev % mod);
void pro(const vector<int> &a, const vector<int> &b, vector<int> &res)
        vector<int> fa(a.begin(), a.end()), fb(b.begin(), b.end());
        int n = 1:
        while (n < (int) max(a.size(), b.size())) n <<= 1;</pre>
        n <<= 1:
        fa.resize (n), fb.resize (n);
```

```
fft(fa, false), fft (fb, false);
        for (int i = 0; i < n; ++i)
                fa[i] = 1LL * fa[i] * fb[i] % mod;
        fft (fa, true);
        res = fa;
int S(int n, int r) {
        int nn = 1;
        while(nn < n) nn <<= 1;</pre>
        for(int i = 0; i < n; ++i) {
                v[i].push_back(i);
                v[i].push_back(1);
        for (int i = n; i < nn; ++i) {
                v[i].push_back(1);
        for (int j = nn; j > 1; j >>= 1) {
                int hn = i >> 1:
                for (int i = 0; i < hn; ++i) {
                        pro(v[i], v[i + hn], v[i]);
        return v[0][r];
int fac[N], ifac[N], inv[N];
void prencr() {
        fac[0] = ifac[0] = inv[1] = 1;
        for (int i = 2; i < N; ++i)
                inv[i] = mod - 1LL * (mod / i) * inv[mod % i] % mod;
        for (int i = 1; i < N; ++i) {
                fac[i] = 1LL * i * fac[i - 1] % mod;
                ifac[i] = 1LL * inv[i] * ifac[i - 1] % mod;
int C(int n, int r) {
        return (r \ge 0 \&\& n \ge r) ? (1LL * fac[n] * ifac[n - r] % mod
            * ifac[r] % mod) : 0;
int main(){
        prencr();
        int n, p, q;
        cin >> n >> p >> q;
        11 ans = C(p + q - 2, p - 1);
        ans *= S(n - 1, p + q - 2);
        ans %= mod;
        cout << ans;
```

### 7.7 Chinese remainder

```
long long GCD(long long a, long long b) { return (b == 0) ? a : GCD(b,
   a % b); }
inline long long LCM(long long a, long long b) { return a / GCD(a, b)
   * b; }
```

```
inline long long normalize(long long x, long long mod) { x %= mod; if
    (x < 0) x += mod; return x; }
struct GCD_type { long long x, y, d; };
GCD_type ex_GCD(long long a, long long b) {
    if (b == 0) return {1, 0, a};
    GCD type pom = ex GCD(b, a % b);
    return {pom.y, pom.x - a / b * pom.y, pom.d};
int testCases;
int t;
long long r[N], n[N], ans, lcm;
int main(){
    cin >> t;
    for(int i = 1; i \le t; i++) cin >> r[i] >> n[i], normalize(r[i], n
        [i]);
    ans = r[1];
    lcm = n[1];
    for(int i = 2; i <= t; i++) {
        auto pom = ex_GCD(lcm, n[i]);
        int x1 = pom.x;
        int d = pom.d;
        if((r[i] - ans) % d != 0) return cerr << "No solutions" <</pre>
        ans = normalize(ans + x1 * (r[i] - ans) / d % (n[i] / d) * lcm
          , lcm * n[i] / d);
        lcm = LCM(lcm, n[i]); // you can save time by replacing above
            lcm * n[i] /d by lcm = lcm * n[i] / d
    cout << ans << " " << lcm << endl;
    return 0;
```

### 7.8 Stirling 2

$$\left\{\begin{array}{c} \mathbf{n} \\ \mathbf{k} \end{array}\right\} = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} \binom{k}{j} j^{n}$$

### 7.9 Duality of LP

primal: Maximize  $c^Tx$  subject to  $Ax \le b, x \ge 0$  dual: Minimize  $b^Ty$  subject to  $A^Ty \ge c, y \ge 0$ 

### 7.10 Extended catalan

number of ways for going from 0 to A with k moves without going to -B:

$$\binom{k}{\frac{A+k}{2}} - \binom{k}{\frac{2B+A+k}{2}}$$

### 7.11 Find polynomial from it's points

$$P(x) = \sum_{i=1}^{n} y_i \prod_{j=1, j \neq i}^{n} \frac{x - x_j}{x_i - x_j}$$

### 8 Constants

### 8.1 Number of primes

30: 10 60: 17 100: 25 1000: 168 10000: 1229 1000000: 78498 10000000: 664579

### 8.2 Most divisor

```
<= 1e2: 60 with 12 divisors
<= 1e3: 840 with 32 divisors
<= 1e4: 7560 with 64 divisors
<= 1e5: 83160 with 128 divisors
<= 1e6: 720720 with 240 divisors
<= 1e7: 8648640 with 448 divisors
<= 1e8: 73513440 with 768 divisors
<= 1e9: 735134400 with 1344 divisors
<= 1e10: 6983776800 with 2304 divisors
<= 1e11: 97772875200 with 4032 divisors
<= 1e12: 963761198400 with 6720 divisors
<= 1e13: 9316358251200 with 10752 divisors
<= 1e14: 97821761637600 with 17280 divisors
<= 1e15: 866421317361600 with 26880 divisors
<= 1e16: 8086598962041600 with 41472 divisors
<= 1e17: 74801040398884800 with 64512 divisors
<= 1e18: 897612484786617600 with 103680 divisors
```

### Useful formulas

 $\binom{n}{k} = \frac{n!}{k!(n-k)!}$  objects out of n- number of ways to choose k

 $\binom{n+k-1}{k-1}$  — number of ways to choose k objects out of n with repetitions  $\binom{n}{n}$  — Stirling numbers of the first kind: number of

permutations of n elements with k cycles  ${n+1\brack m}=n{n\brack m}+{n\brack m-1}$  $\lfloor \frac{n}{m} \rfloor$  — Stirling numbers of the first kind; number of

$${n+1 \brack m} = n{n \brack m} + {n \brack m-1}$$

$$(x)_n = x(x-1)\dots x - n + 1 = \sum_{k=0}^n (-1)^{n-k} {n \brack k} x^k$$

of partitions of set  $1, \ldots, n$  into k disjoint subsets.  ${n+1 \brace m} = k \begin{Bmatrix} n \end{Bmatrix} + \begin{Bmatrix} n \cr k-1 \end{Bmatrix}$  ${n \brace m} - ext{Stirling numbers of the second kind; number}$ 

$${\binom{n+1}{m}} = k {\binom{n}{k}} + {\binom{n}{k-1}}$$

$$\sum_{k=0}^{n} {n \brace k}(x)_k = x^n$$

$$C_n = \frac{1}{n+1} {2n \choose n} - \text{Catalan numbers}$$

$$C(x) = \frac{1-\sqrt{1-4x}}{2x}$$

## Binomial transform

If 
$$a_n = \sum_{k=0}^{n} {n \choose k} b_k$$
, then  $b_n = \sum_{k=0}^{n} (-1)^{n-k} {n \choose k} a_k$ 

• 
$$a = (1, x, x^2, ...), b = (1, (x+1), (x+1)^2, ...)$$

• 
$$a_i = i^k, b_i = {n \brace i} i!$$

### Burnside's lemma

shifts of array, rotations and symmetries of  $n \times n$ matrix, ...) Let G be a group of action on set X (Ex.: cyclic

action f that transforms x to y: f(x) = y. Call two objects x and y equivalent if there is an

The number of equivalence classes then can be calculated as follows:  $C = \frac{1}{|G|} \sum_{f \in G} |X^f|$ , where  $X^f$ 

is the set of fixed points of  $f: X^f = \{x | f(x) = x\}$ 

## Generating functions

sequence  $a_0, a_1, \dots, a_n, \dots$  is  $A(x) = \sum_{n=0}^{\infty} a_i x^i$  $a_0, a_1, \dots, a_n, \dots$  is  $A(x) = \sum_{i=1}^{\infty} a_i x^i$ Ordinary generating function (o.g.f.) for sequence  $B(x) = A'(x), b_{n-1} = n \cdot a_n$ Exponential generating function (e.g.f.)

$$c_n = \sum_{k=0}^{n} a_k b_{n-k} \text{ (o.g.f. convolution)}$$

$$c_n = \sum_{k=0}^{n} \binom{n}{k} a_k b_{n-k} \text{ (e.g.f. convolution, compute}$$
with FFT using  $\widetilde{a_n} = \frac{a_n}{n!}$ )

# General linear recurrences

If 
$$a_n = \sum_{k=1}^{\infty} b_k a_{n-k}$$
, then  $A(x) = \frac{a_0}{1-B(x)}$ . We also can compute all  $a_n$  with Divide-and-Conquer algorithm in  $O(n\log^2 n)$ .

# Inverse polynomial modulo x'

Given 
$$A(x)$$
, find  $B(x)$  such that  $A(x)B(x)=1+x^l\cdot Q(x)$  for some  $Q(x)$ 

1. Start with 
$$B_0(x) = \frac{1}{a_0}$$

2. Double the length of 
$$B(x)$$
:
$$B_{k+1}(x) = (-B_k(x)^2 A(x) + 2B_k(x)) \mod x^{2^{k+1}}$$

# Fast subset convolution

Given array  $a_i$  of size  $2^k$ , calculate  $b_i =$ 

for 
$$b = 0..k-1$$
  
for  $i = 0..2^k-1$   
if  $(i & (1 << b)) != 0:$   
 $a[i + (1 << b)] += a[i]$ 

## Hadamard transform

size  $2 \times 2 \times \ldots \times 2$ , calculate FFT of that array: Treat array a of size  $2^k$  as k-dimentional array