



Sharif University of Technology
Department of Computer Engineering

Low Power Digital System Design

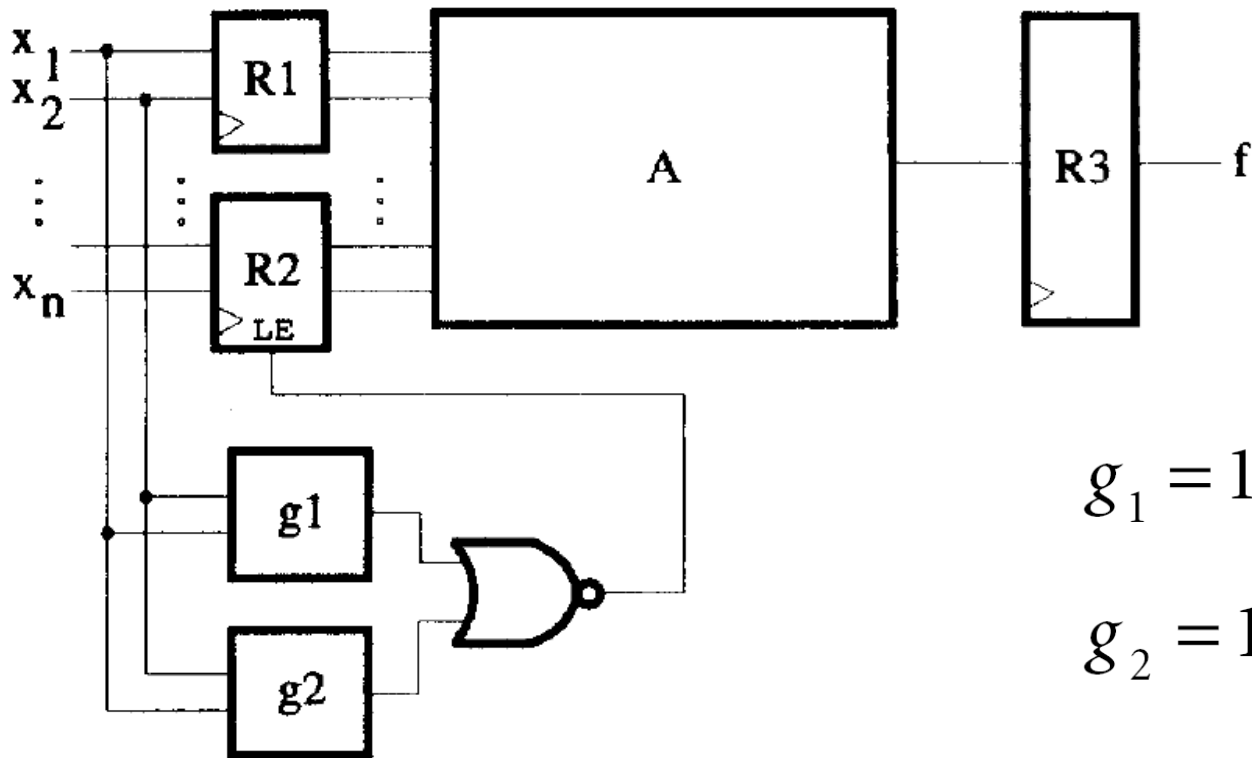
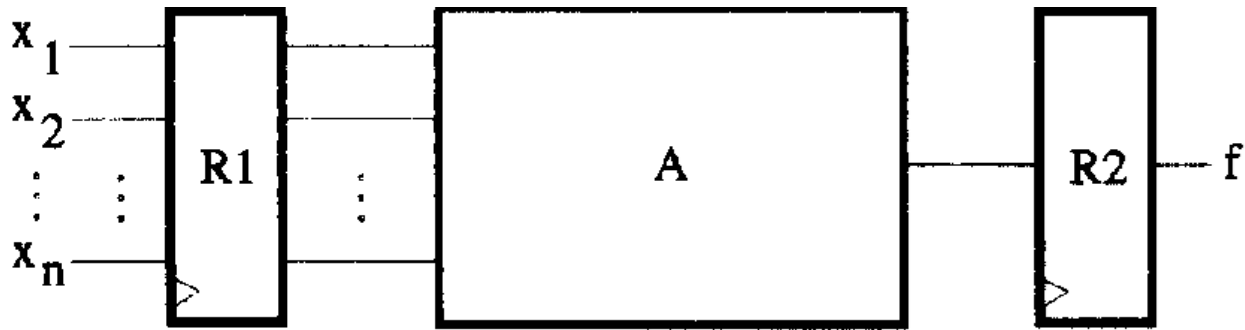
Precomputation

A. Ejlali

Precomputation

- Basic Ideal:
 - **Selectively** precomputing the output logic values of a circuit one clock cycle before they are required.
 - Use the precomputed values to reduce the **internal switching activity** of the combinational logic in the successive clock cycle.

Precomputation Architecture



$$g_1 = 1 \Rightarrow f = 1$$

$$g_2 = 1 \Rightarrow f = 0$$

Predictor Functions: g_1 and g_2

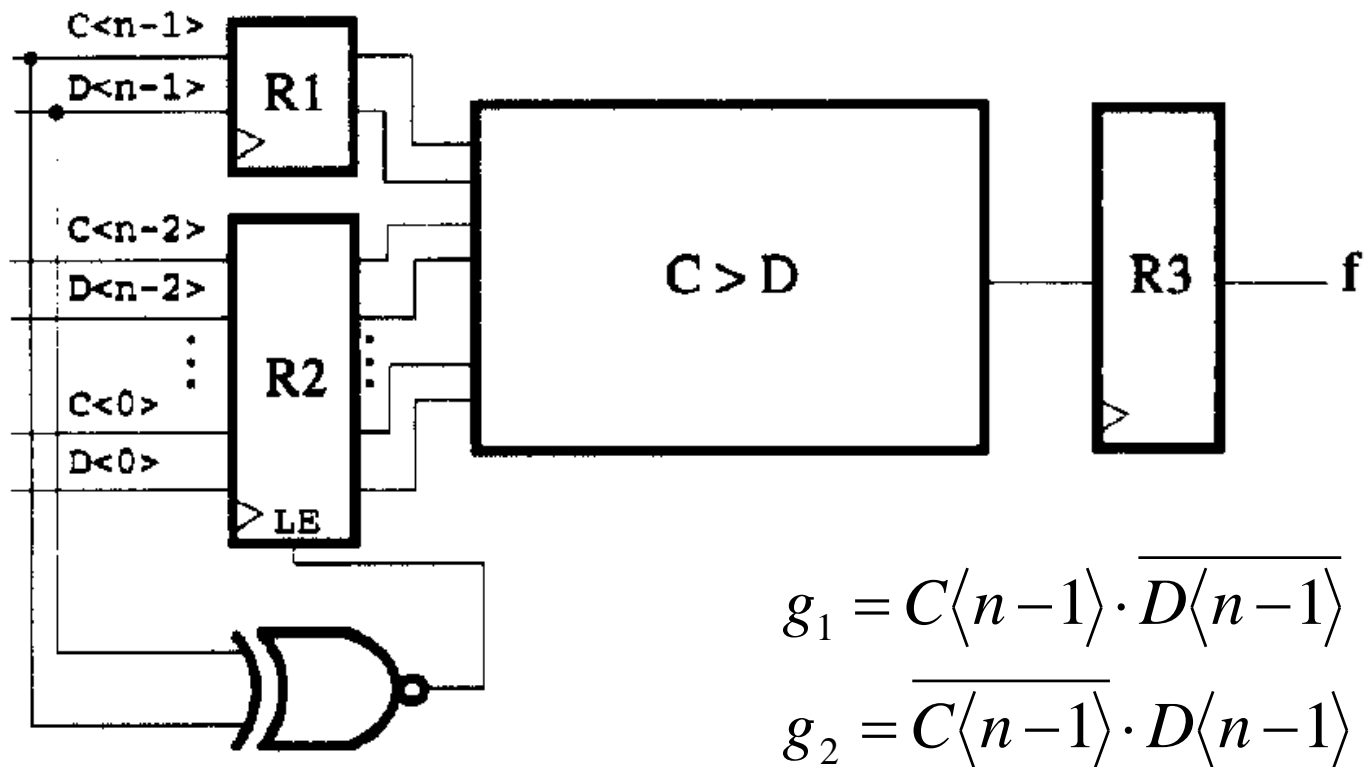
- A power reduction is obtained because for a subset of input conditions corresponding to g_1+g_2 , R2 does not change implying reduced switching activity.
- The probability of g_1+g_2 being a 1 should be high.
- g_1 and g_2 have to be significantly less complex than f .
- g_1 and g_2 cannot both be 1 during the same clock cycle due to the conditions imposed by:

$$g_1 = 1 \Rightarrow f = 1$$

$$g_2 = 1 \Rightarrow f = 0$$

An Example of Precomputation

- Example: a n-bit comparator that compares two n-bit numbers C and D and computes the function $C > D$.



An Example of Precomputation (Cont.)

- In this example, the probability of $g_1 + g_2$ being a 1 is 0.5.
- If we add the inputs $C(n-2)$ and $D(n-2)$ to g_1 and g_2 it is possible to achieve a power reduction close to 75%.
 - Suitable for large n

Precomputation Logic

- Observability don't care set for input x :

$$ODC_x = \overline{f_x} \oplus \overline{f_x^-} = f_x \cdot f_x^- + \overline{f_x} \cdot \overline{f_x^-}$$

- f_x and f_x^- are the cofactors of f .
- What does ODC mean?
 - For all the input combinations in ODC_x , we do not need the value of x in order to know what the value of f is.
- If we wish to disable the inputs $x_{m+1} \sim x_n$ we will have to implement the predictor function:

$$g = \prod_{i=m+1}^n ODC_{x_i}$$

Precomputation Logic (Cont.)

- Universal Quantification of a function:

$$U_x f = f_x \cdot f_{\bar{x}}$$

- What does $U_x f$ mean?
 - all the input combinations that result in $f=1$ such that the value of x , does not matter.
- Given a subset of inputs $S=\{x_1, x_2, \dots, x_n\}$, it has been proven that the best choices for g_1 and g_2 are:

$$g_1 = \prod_{i=m+1}^n U_{x_i} f$$

$$g_2 = \prod_{i=m+1}^n U_{x_i} \bar{f}$$

Precomputation Logic (Cont.)

- Problem: Selecting a subset of inputs for the precomputation logic.
 - Exhaustive check
 - Time-consuming
 - The simplest prime implicants in the simplified SOP and POS representations are suitable to be used as g_1 and g_2 functions.
 - For large combinational blocks it is very time-consuming to find simplified SOP and POS representations.

Reference

M. Alidina, "Precomputation-Based Sequential Logic Optimization for Low Power", IEEE Trans. VLSI 1994.