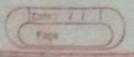
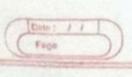


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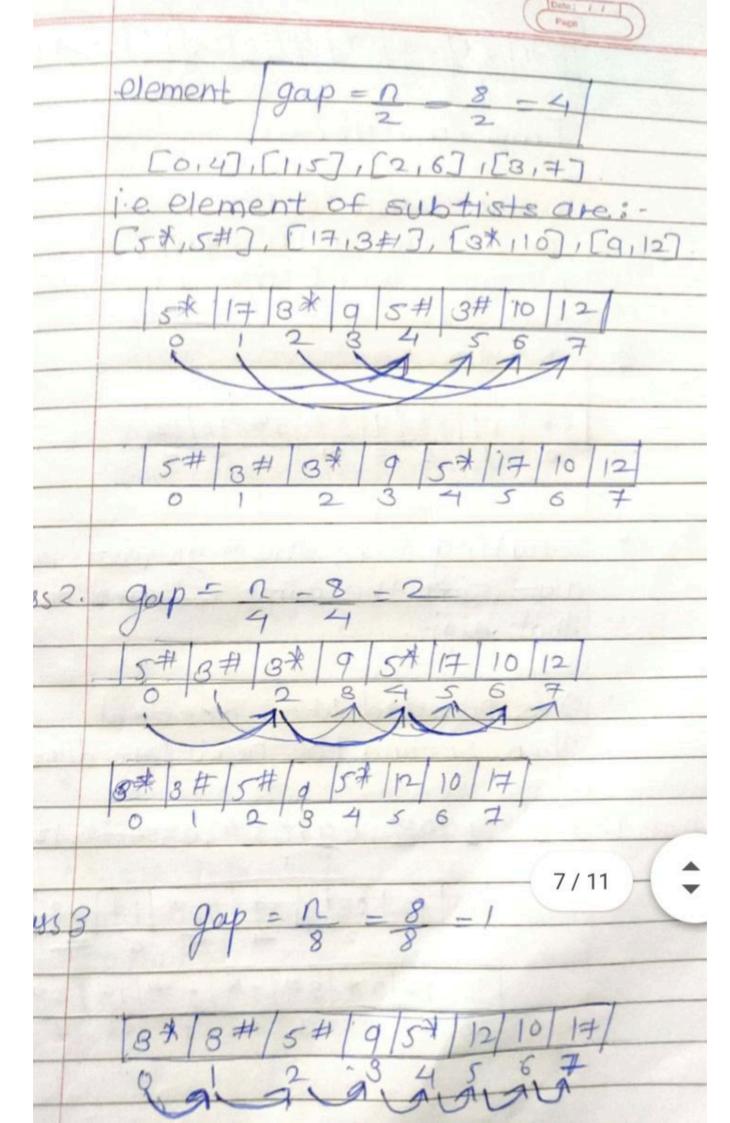


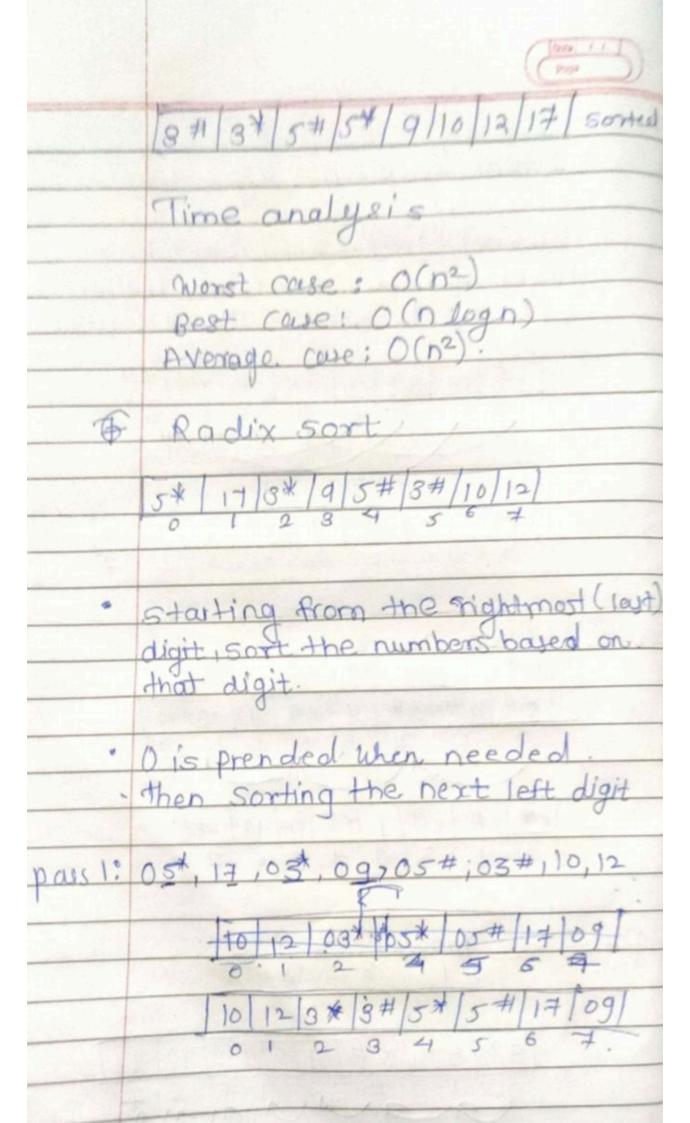
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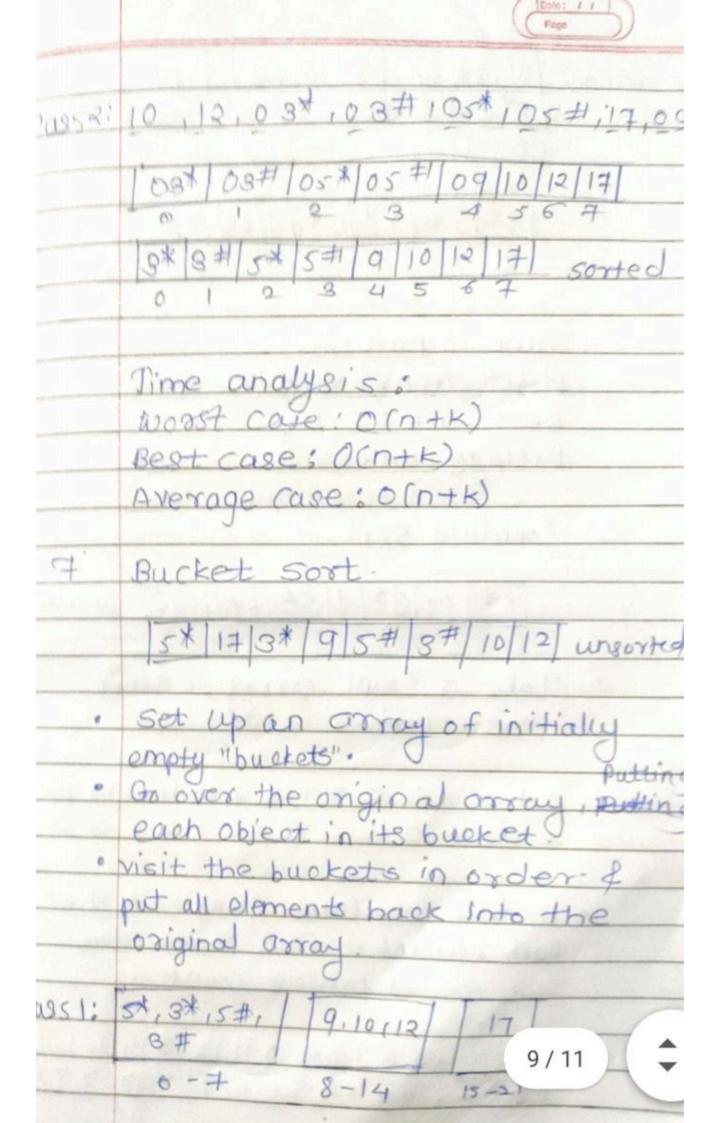


Time Analysis Worst case : O(n2) Best case ; o(n) Average Case: O(n2) 4. Quick sort. 5 17 3 4 9 5 # 3 # 10 12 unsortal · These are two indices if if at the very tegigning of the portition algorithm i points to the first element in the array & i points to the last one * Then i moves forward until ar element with value greaters or Bround to me pint is Round. · Index i is moved back word urtil an element with vave hosses or equal to the pirot is found I fisj then they swapped & i step to the next postion (i+1) is steps to the previous ore (j-1). · It otops when i be comes greaters them on After position, all values before ith

element are less or equal than the pivet







Arrange: (34,3#1) 19,10,12 M,5# 8-14 0 - 7 3 + 3 + 5 + 9 10 12 17 Time analysis: Worst cases - O(p2) Best case; O(n) Average case: O(n) 8. Counting sort : 5x 17 13 x 9 5# 3# 10 13 Unsorte Take a count array to store the count. Index-01234567891011121314 court 000202000110100 · modify the count array such that lack element it Puch indere store the sum of previous counts. Index 01234567891011121314 1516 0002244445667