

Unit

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Date ___/___/___

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Operating system Assignment - 2

Q₁

Ans Address translation in modern system.

- Each process generates logical (virtual) address.
- MMU (Memory management unit) translates these into physical address.

Translation Steps:

- CPU generates logical address.
- MMU checks page table for corresponding frame number.
- Concatenation frame + offset → physical address.

Q₂

Ans Process A (100kb of 120 kb block) / free 30 kb /
Process B (200 kb)

Internal → 20 kb wasted space inside allocated block.

External → 30 kb free, but too small for 40 kb.

Mitigation: Paging (removes external), segmentation with buddy system, slab allocators, or dynamic allocation with coalescing free space.

Q₃

Ans Memory is split into fixed size frames, process into pages, page table store mapping.

- Pros: No external fragmentation simple allocation.
- Cons: Page tables use extra memory, TLB missed add delay, last page may waste some space.

Q4.

Ans OS hardware interaction in virtual memory.

- Page-Table in memory.
- MMU translates virtual.
- TLB caches recent translations
- Protection bits

Q5

Ans 16-bit virtual address, 11kb page size.

- Address space = $2^{16} = 65536$ bytes
- Page size = $2^{10} = 1024$ bytes
- Number of Pages = $\frac{65536}{1024} = 64$
- Page table size = $64 \times 2 \text{ bytes} = 128 \text{ bytes}$

Part B

Q6

Ans

Process	Size (kb)
P ₁	212
P ₂	417
P ₃	112
P ₄	426

- First-fit

P ₁	P ₂	P ₃	P ₄
0	212	629	741
			1167

Unused memory = 259 kb

- Best-fit

P ₁	P ₂	P ₃	P ₄
212	417	112	

P₄ still can't fit

Unused = 259 kb

- Worst-fit

P_4 (426) can't fit, unused = 259 kb

All three methods leave 259 kb unused in this process.

Q7.

Ans Page replacement.

- FIFO = 10 faults
- Optimal = 7 faults
- LRU = 9 faults

Optimal is best, LRU is practical, FIFO may show Belady's anomaly.

Q8.

Ans

Disk write = 10ms

Memory write = 100ms

Dirty pages = 30% of 1000 = 300

$$\begin{aligned} \text{① Overhead} &= 300 \times 10 \\ &= 3000 \text{ ms} = 3 \text{ seconds.} \end{aligned}$$

- ② Optimization: write-back caching with dirty bit tracking or pre-clearing (background flush) reduce blocking time.

Q9.

Ans

Autonomous vehicle memory.

- Use working set model to allocate enough frames for real time (object detection) to prevent thrashing.

- Replacement policy : LRU or WBlock with priority.
- Strategy : Reserve fixed memory for real-time processes, share remaining dynamically. This ensures responsiveness and efficient utilization.

finds
2/11/25