

CS & IT ENGINEERING

Logical Equivalence



Lecture No. 02



By- SATISH YADAV SIR

TOPICS TO BE COVERED

01 Type 1

02 Logical equivalence

03 Conditional Equivalence

04 Biconditional equivalence

05 Type 2 Questions in Logic

* $A \equiv B$

1. A, B are having same behaviour

	A	B	$A \leftrightarrow B$
	-	-	T
	-	-	T
	-	-	T
	-	-	T
	-	-	T

2. A, B are having same columns

3.

$A \leftrightarrow B$
tautology.
valid.

we will try \rightarrow false

Type-I. :

$$\frac{\frac{\frac{P \wedge (P \rightarrow q) \wedge (\neg q \vee n) \rightarrow n}{\frac{(T \rightarrow F)}{\text{false}}}}{(T \rightarrow F)}}{(T \vee F)} \quad \begin{array}{l} n = F \\ P = T \\ \neg q = T \\ q = F \end{array}$$

Time

false

$$\rightarrow \gamma q = f$$

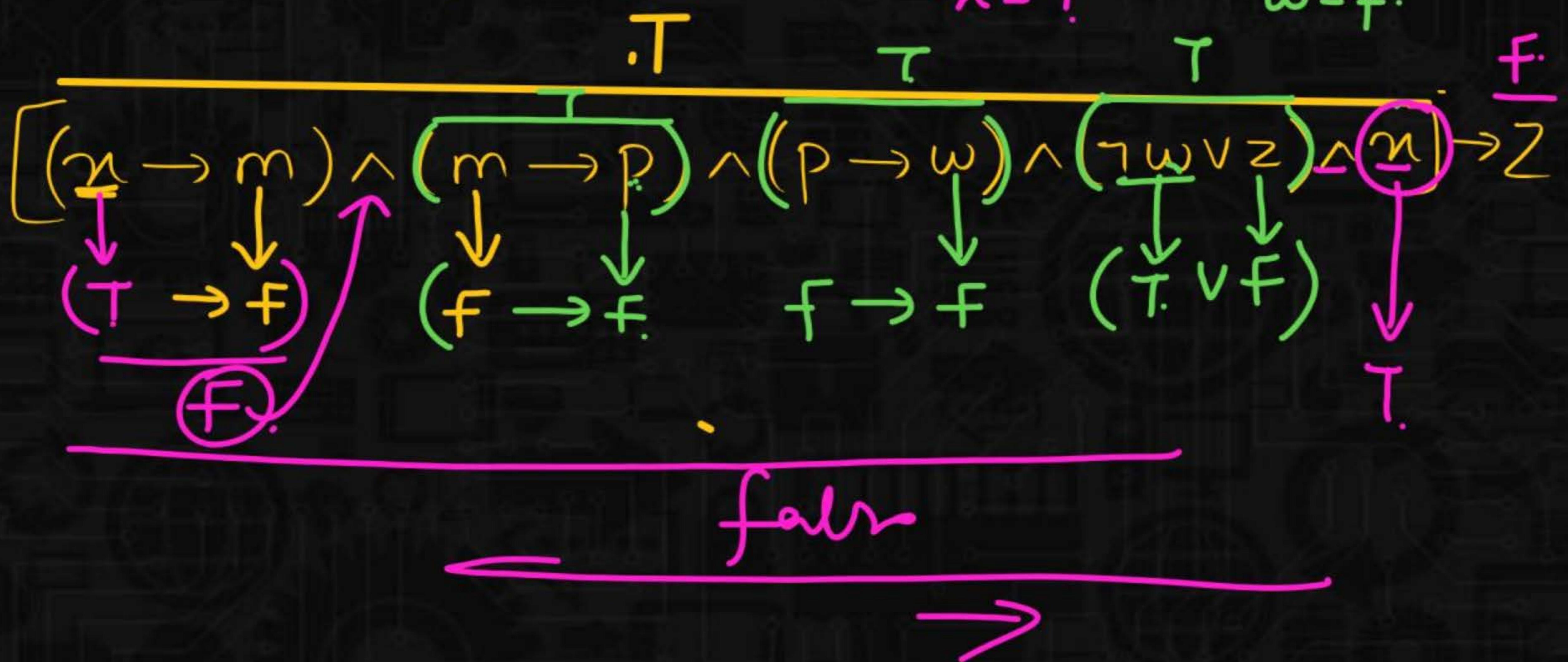
m = f.

$$P = F.$$

$$\underline{z = f}.$$

$$\begin{aligned} \neg w &= T \\ w &= F. \end{aligned}$$

七



$$\frac{\frac{\frac{\frac{\frac{\top}{(n \rightarrow m) \wedge (m \rightarrow p) \wedge (p \rightarrow \omega) \wedge ((\neg \omega \vee z) \wedge m)} \rightarrow z}{(n \rightarrow m) \wedge (m \rightarrow p) \wedge (p \rightarrow \omega) \wedge ((\neg \omega \vee z) \wedge m)} \rightarrow z}{(n \rightarrow m) \wedge (m \rightarrow p) \wedge (p \rightarrow \omega) \wedge ((\neg \omega \vee z) \wedge m)} \rightarrow z}{(n \rightarrow m) \wedge (m \rightarrow p) \wedge (p \rightarrow \omega) \wedge ((\neg \omega \vee z) \wedge m)} \rightarrow z}{(n \rightarrow m) \wedge (m \rightarrow p) \wedge (p \rightarrow \omega) \wedge ((\neg \omega \vee z) \wedge m)} \rightarrow z}$$

$z = f$

$n = \top$

$\omega = f$

$p = f$

$m = f$

Falsze

$$\frac{\overline{T} \wedge (\overline{u \rightarrow p}) \wedge \overline{m \rightarrow p}}{\overline{m}} \rightarrow \overline{f}$$

$$\frac{\downarrow}{\underline{(T \rightarrow T)}}$$

$$\frac{\downarrow}{\underline{(F \rightarrow T)}}$$

$$\frac{\overline{T} \wedge \overline{T} \wedge \overline{T}}{\overline{T}}$$

$$m = F.$$

$$x = T.$$

$$p = T.$$

$$\overline{T} \longrightarrow f.$$

false

Type-1.: Substⁿ method:

(\wedge \wedge \wedge) \rightarrow

Check valid? \checkmark_0 .

Type-2 logical equivalence:

$$\begin{aligned} 2: \quad P \wedge T &\equiv P \\ P \vee F &\equiv P \end{aligned}$$

1.

$$P \wedge P \equiv P$$

$$P \vee P \equiv P$$

2.

$$P \wedge \underline{T} \equiv P$$

$$P = T$$

$$\textcircled{T} \wedge T \equiv T$$

$$P = F$$

$$\textcircled{F} \wedge T \equiv F$$

$$P \vee \underline{F} \equiv P$$

$$P = T$$

$$\textcircled{T} \vee F \equiv T$$

$$P = F$$

$$\textcircled{F} \vee F \equiv F$$

Associative. \rightarrow same operator.

3:

$$P \vee \top \equiv \top.$$

$$P \wedge \text{f} \equiv \text{f}$$

5. $P \Delta (q \Delta R) \equiv (P \Delta q) \Delta R$

$$P \vee (q \vee R) \equiv (P \vee q) \vee R.$$

4. $P \vee q = q \vee P$

$$P \wedge q = q \wedge P.$$

6: Distributive laws. : diff operator.

$$P \wedge (q \vee R) \equiv [P \wedge q] \vee [P \wedge R]$$



$$P \vee (q \wedge R) \equiv (P \vee q) \wedge (P \vee R)$$

$$P \wedge (q \vee R) \equiv (P \wedge q) \vee (P \wedge R) \xrightarrow{F \vee (F \wedge -)} \equiv F.$$

Absorption law :

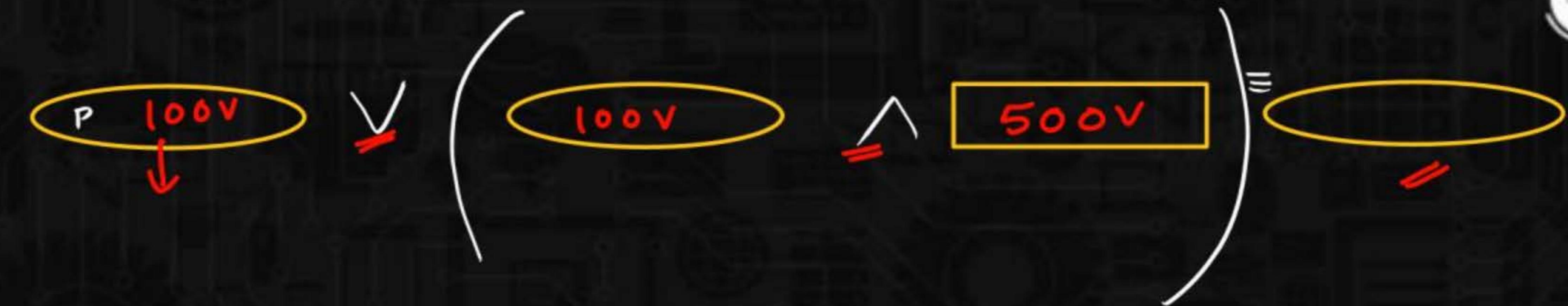
$$\underline{P} \vee (\underline{P} \wedge q) \equiv P.$$

$$P \wedge (P \vee q) \equiv P.$$

$$\underline{P} \vee (\underline{P} \wedge q) \equiv P.$$

$$P = T.$$

$$\underline{T} \vee (T \wedge -) \xrightarrow{T}$$



$$(P \rightarrow Q) \vee R) \vee ((P \rightarrow Q) \vee R) \wedge (t \rightarrow (m \vee n)) \Rightarrow (P \rightarrow Q) \vee R.$$

$$\neg(a \vee b) \equiv \neg a \wedge \neg b.$$

$$\neg(a \wedge b) \equiv \neg a \vee \neg b.$$

$$P \rightarrow q \equiv \neg q \rightarrow \neg P.$$

$$P \rightarrow q \equiv \neg P \vee q.$$

$$\neg(P \rightarrow q) \equiv \neg(\neg P \vee q)$$

$$\neg(P \rightarrow q) \equiv P \wedge \neg q$$

Good's Rule

$$\boxed{P \rightarrow q \equiv \neg q \rightarrow \neg P.}$$
$$P \rightarrow q \equiv \neg P \vee q$$

$$\neg P \rightarrow q = P \vee q.$$

$$\neg(P \rightarrow q)$$

$$\neg(\neg P \vee q)$$

$$\underline{P \wedge \neg q.}$$

$$\begin{aligned} p \leftrightarrow q &\equiv (p \rightarrow q) \wedge (q \rightarrow p) \equiv (p \rightarrow q) \wedge (q \rightarrow p) \\ &\equiv (\neg p \vee q) \wedge (q \rightarrow p) \equiv (p \rightarrow q) \wedge (\neg q \vee p) \\ &\equiv (\neg p \vee q) \wedge (\neg p \rightarrow \neg q) \equiv (\neg q \rightarrow \neg p) \wedge (\neg q \vee p) \end{aligned}$$

$$(P \rightarrow q) \wedge (P \rightarrow r) \equiv P \rightarrow (q \wedge r)$$

$$(\neg P \vee q) \wedge (\neg P \vee r)$$

$$\left\{ \begin{array}{l} (P \rightarrow q) \wedge (P \rightarrow r) \equiv P \rightarrow (q \wedge r) \\ \hline (P \rightarrow q) \vee (P \rightarrow r) \equiv P \rightarrow (q \vee r) \end{array} \right.$$

$$\not\models P \vee (q \wedge r)$$

$$P \rightarrow (q \wedge r)$$

P
W

$$(P \rightarrow R) \wedge (\varnothing \rightarrow R)$$

$$(\neg P \boxed{\vee R}) \wedge (\neg \varnothing \boxed{\vee R})$$

$$(\neg P \wedge \neg \varnothing) \vee R$$

$$\neg(P \vee \varnothing) \vee R$$

$$(P \vee \varnothing) \rightarrow R$$

$$\underline{(P \rightarrow R)} \wedge \underline{(P \rightarrow \varnothing)} \equiv P \rightarrow (\varnothing \wedge R)$$

$$\underline{(P \rightarrow R)} \vee \underline{(P \rightarrow \varnothing)} \equiv P \rightarrow (\varnothing \vee R)$$

$$(P \rightarrow \underline{R}) \wedge (\varnothing \rightarrow \underline{R}) \equiv$$

$$(P \vee \varnothing) \rightarrow R$$

$$(P \rightarrow R) \vee (\varnothing \rightarrow R) \equiv (P \wedge \varnothing) \rightarrow R$$

$$(P \rightarrow R) \wedge (Q \rightarrow R)$$

$$\neg P \vee R \quad \wedge \quad \neg Q \vee R$$

$$(\neg P \wedge \neg Q) \vee R$$

$$\neg(P \vee Q) \vee R \equiv (P \vee Q) \rightarrow R$$

$$(P \rightarrow Q) \wedge (P \rightarrow R) \equiv P \rightarrow (Q \wedge R)$$

$$(P \rightarrow Q) \vee (P \rightarrow R) \equiv P \rightarrow (Q \vee R)$$

God Rule
&

$$(P \rightarrow R) \wedge (Q \rightarrow R) \equiv (P \vee Q) \rightarrow R.$$

$$(P \rightarrow R) \vee (Q \rightarrow R) \equiv (P \wedge Q) \rightarrow R.$$

Demorgan's

$$(P \rightarrow R) \vee (Q \rightarrow R)$$

$$(\neg P \vee R) \vee (\neg Q \vee R)$$

$$\underline{(\neg P \vee \neg Q) \vee R}$$

$$\neg(P \wedge Q) \vee R = (P \otimes Q) \rightarrow R.$$

The Simplest form of $(p \wedge (\sim r \vee q \vee \sim q)) \vee ((r \vee t \vee \sim r) \wedge \sim q)$ is

- (a) $p \wedge \sim q$
- (b) $p \vee \sim q$
- (c) t
- (d) $(p \rightarrow \sim q)$

The Simplest form of

$(p \vee (p \wedge q) \vee (p \wedge q \wedge \sim r)) \wedge ((p \wedge r \wedge t) \vee t)$ is

(a) $p \wedge t$

(b) $q \wedge t$

(c) $p \wedge r$

(d) $p \wedge q$

Which one of the following is NOT equivalent to $p \leftrightarrow q$?

(GATE-15-Set1)

- (a) $(\sim p \vee q) \wedge (p \vee \sim q)$
- (b) $(\sim p \vee q) \wedge (q \rightarrow p)$
- (c) $(\sim p \wedge q) \vee (p \wedge \sim q)$
- (d) $(\sim p \wedge \sim q) \vee (p \wedge q)$

P and Q are two propositions. Which of the following logical expressions are equivalent?

- I. $P \vee \sim Q$
- II. $\sim(\sim P \wedge Q)$
- III. $(P \wedge Q) \vee (P \wedge \sim Q) \vee (\sim P \wedge \sim Q)$
- IV. $(P \wedge Q) \vee (P \wedge \sim Q) \vee (\sim P \wedge Q)$

(GATE-08)

- (a) Only I and II
- (b) Only I, II and III
- (c) Only I, II and IV
- (d) All of I, II, III and IV

\wedge	\cap	\cdot
\vee	\cup	$+$
T	\top	1
F	\emptyset	0

$$P \wedge (P \vee a) = P.$$

$$\underline{a} \cap (\underline{a} \cup b) = \underline{a}.$$

$$\underline{a} \vee (\underline{a} \wedge b) = \underline{a}.$$

