

Q1]

→ i) Activity selection of least duration will not work

Example:- if we consider these are and activities

1 to 4 which has different duration for each activity start time and finish time is depend as below

Activity	start time	finish time	Duration
1	3	5	2
2	4	8	4
3	1	4	3
4	1	5	4

Solution for least duration activity is $\{1, 3\}$ that is activity is duration is 2 but optimal solution is $\{2, 3\}$ that is activity 2 and 3

ii) Always selecting the compatible activity that overlaps the fewest other remaining activities this will not work, consider below example as there are 6 different activities each activity has start and finish time. these activities may or may not overlap with other activities start and finish time is defined in below take also overlap count will be added with other activities

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Activities	Start time	Finish time	Overlap
1	2	4	3
2	3	5	3
3	4	7	2
4	1	3	2
5	2	4	3
6	6	8	2
7	7	10	3
8	9	11	2
9	7	10	3

Solution for above approach is {3, 4, 8} or {4, 6, 8}

but optimal solution is {2, 4, 6, 8}

iii) Selecting the compatible remaining activity with earliest start time. This will not work. Consider below example.

Activity	Start time	Finish time	Duration
1	1	10	9
2	2	4	2
3	5	7	2
4	4	10	3

Solution for earliest start time is {1}

but optimal solution is {2, 3, 4}

P-③

Q2) if we consider problem of making change for n cents using fewer number of coins

- a)
- penny - 1
 - Nickle - 5
 - Dime - 10
 - quarter - 25

Consider 'n' be the no of cents. Now first divide the number of cents (n) with quarters (25) the floor value will be no of quarters if remainder is greater than 0 then remainder value to be divide by dime here floor value will be number of dime if remainder is greater than 0 then again divide remainder with Nickle if floor value is zero then that will be number of Nickle if remainder is non zero then we will divide remainder with consider remainder as no of cents penny

$$\text{Number of quarters} = (n/25)$$

$$n = n \bmod 25$$

if $n(n \bmod 25)$ is greater than 0 then

$$\text{no dime} = (n/10)$$

$$n = n \bmod 10$$

if $n(n \bmod 10)$ is greater than 0 then

$$\text{no of Nickles} = (n/5)$$

$$n = n \bmod 5$$

P-⑤

if $n \pmod{5}$ is gre

Number of pennies = $p = n$

here we have 5 denominations

- if we have $n = 52$, then there are different soln
consider 52 pennies or 5 dime 2 penny etc
but optimal solution will be 2 quarters / 2 penny
- if we have $n = 23$ then are more than
1 solution as we can consider 13 pennies
or 1 dime 3 pennies but optimal solution
will be 1 dime 3 pennies

hence there is always a greedy optimal solution
for every value of n

b)

if available coins are in the denomination
that are power of c ,

denominator are c^0, c^1, \dots, c^k

for integer $c > 1$ and $k > 1$

No of cents = n

$c > 1$ and $k > 1$

then $c^0, c^1, c^2, \dots, c^k$ are denominations
maximum value for denominations has to be
less than value of n , we have to make

1-⑤

change for number of cents. we can calculate value of k starting from 0 to c' ($i=0; i \leq c'$), comparing each c' value with n if $c' \leq n$ then we should break the loop the $i-1 = k'$

Now we have traverse from c^0, c^1, \dots, c^k from the

$$\text{Number of coins for } c^k = \left\lfloor \frac{c^k}{n} \right\rfloor$$

$$\text{Number of coins for } c^{k-1} = \left\lfloor \frac{c^k \bmod n}{c^{k-1}} \right\rfloor$$

repeat until $k=0$

after $k=0$ we will get number of coins required to make change of n cents

Eg. if $n=25$ cents & $c=3$

then max denomination is 3^2

3^3 is 27 which greater than quarter so $k=2$

$$3^0 = 1$$

$$3^1 = 3$$

$$3^2 = 9$$

The optimal solution is $3^2=2, 3^1=2, 3^0=1$
algorithm also give $\{3^0=25\}$ or $\{3^1=6, 3^2=7\}$ etc
which are not optimal solution

p-⑥

c]

for example denominations $arr = \{1, 5, 8\}$ and
cents = 20

So $n = 20$

greedy algorithm give $\{4, 0, 2\}$ coins respectively
above solⁿ give $4+2=6$ coins but optimal solⁿ
is $\{0, 4, 0\}$ only 4 coins can give end
result instead of 6 coins

d]

```
PowerChange(int cents, int den-arr[], int size)
```

```
int num-unit;
```

```
for (i=size, i>=0 && cents>0; i--)
```

```
{
```

```
    if (den-arr[i] < cents)
```

```
    {
```

```
        num-unit = cents / den-arr[i];
```

```
        cents = cents - num-unit * den-arr[i];
```

```
        printf("Denominations %d : den-arr[i];",
```

```
    }
```

```
}
```

den-array is containing denominations when one coin
is penny & size is length of denominations

above algorithm runs in $O(k)$ time
duration takes $O(1)$ time