

Assembly Programing with NASM (and GAS)

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NASM Data Directives

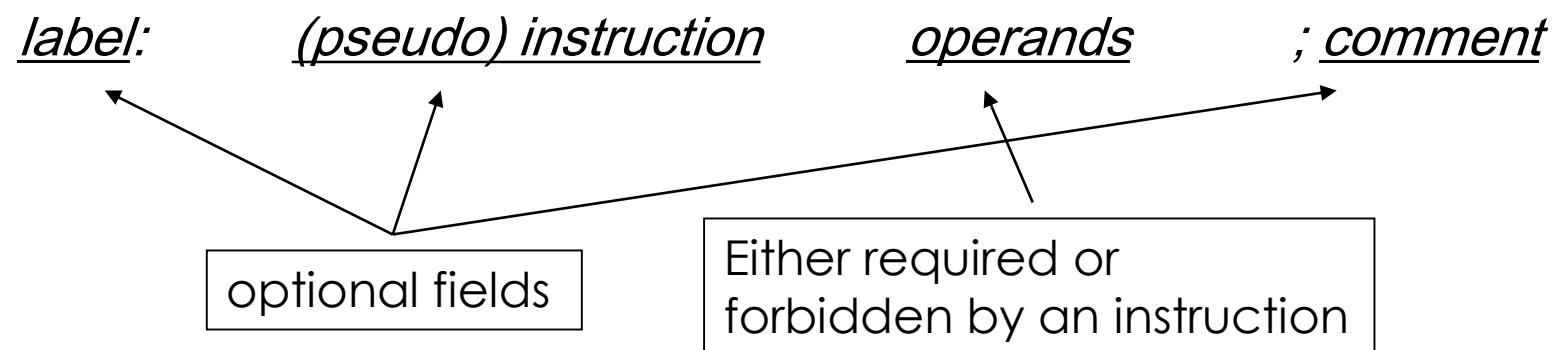
| | | | |
|-------|--------------|-------------------|----------------------|
| ▶ L1 | db | 0 | ; byte |
| ▶ L2 | dw | 1000 | ; word |
| ▶ L3 | db | 110101b | ; byte |
| ▶ L4 | db | 12h | ; byte |
| ▶ L5 | db | 17o | ; byte |
| ▶ L6 | dd | 1A92h | ; double word |
| ▶ L7 | resb | 1 | ; uninitialized byte |
| ▶ L8 | db | 'A' | ; ascii code = 'A' |
| ▶ L9 | db | 0,1,2,3 | ; 4 bytes |
| ▶ L10 | db | 'w','o','r','d',0 | ; string |
| ▶ L11 | db | 'word', 0 | |
| ▶ L12 | times 100 db | 0 | ; 100 bytes of zero |
| ▶ L13 | resw | 100 | ; 100*2(word bytes) |

Data Types in IA32 (AT&T format)

| C declaration | Intel data type | Assembly code suffix | Size (bytes) |
|---------------|--------------------|----------------------|--------------|
| char | Byte | b | 1 |
| short | Word | w | 2 |
| int | Double word | l | 4 |
| long int | Double word | l | 4 |
| long long int | — | — | 4 |
| char * | Double word | l | 4 |
| float | Single precision | s | 4 |
| double | Double precision | l | 8 |
| long double | Extended precision | t | 10/12 |

Instruction Basics

Each NASM standard source line contains a combination of the 4 fields:



1. **Backslash (\)** uses as the line continuation character: if a line ends with backslash, the next line is considered to be a part of the backslash-ended line.
2. **No restrictions on white space** within a line.
3. A colon after a label is optional.

Example Program

| | | | | |
|---|----------------------------|---|---------------------|--------------------------|
| 1 | int simple(int *xp, int y) | 1 | simple: | |
| 2 | { | 2 | pushl %ebp | Save frame pointer |
| 3 | int t = *xp + y; | 3 | movl %esp, %ebp | Create new frame pointer |
| 4 | *xp = t; | 4 | movl 8(%ebp), %edx | Retrieve xp |
| 5 | return t; | 5 | movl 12(%ebp), %eax | Retrieve y |
| 6 | } | 6 | addl (%edx), %eax | Add *xp to get t |
| | | 7 | movl %eax, (%edx) | Store t at xp |
| | | 8 | popl %ebp | Restore frame pointer |
| | | 9 | ret | Return |

Operand Modes

| | Source | Destination | C Analog |
|------|--------|-------------|----------------|
| movl | Imm | Reg | temp = 0x4; |
| | | Mem | *p = -147; |
| | Reg | Reg | temp2 = temp1; |
| | Mem | Mem | *p = temp; |
| | Mem | Reg | temp = *p; |

NASM Data Directives

| | | | |
|-------|--------------|-------------------|----------------------|
| ▶ L1 | db | 0 | ; byte |
| ▶ L2 | dw | 1000 | ; word |
| ▶ L3 | db | 110101b | ; byte |
| ▶ L4 | db | 12h | ; byte |
| ▶ L5 | db | 17o | ; byte |
| ▶ L6 | dd | 1A92h | ; double word |
| ▶ L7 | resb | 1 | ; uninitialized byte |
| ▶ L8 | db | 'A' | ; ascii code = 'A' |
| ▶ L9 | db | 0,1,2,3 | ; 4 bytes |
| ▶ L10 | db | 'w','o','r','d',0 | ; string |
| ▶ L11 | db | 'word', 0 | |
| ▶ L12 | times 100 db | 0 | ; 100 bytes of zero |
| ▶ L13 | resw | 100 | ; 100*2(word bytes) |

Examples

Data directives (different to MASM)

- ▶ **Mov al, [L1]** ; copy byte at L1
- ▶ **Mov eax, L1** ; eax = address of byte at L1
- ▶ **Mov [L1], ah** ; copy ah into byte at L1
- ▶ **Mov eax, [L6]** ; copy double word
- ▶ **Add eax, [L6]** ; eax = eax + double word at L6
- ▶ **Add [L6], eax** ; double word at L6 += eax
- ▶ **Mov al, [L6]** ; copy first byte of double word at L6 into al
- ▶ **Mov [L6], 1** ; operation size is not specified
- ▶ **Mov dword [L6], 1** ; store a 1 at L6

Instructions: Setup

- ▶ STACK
 - ▶ can be used as a convenient place to store data temporarily
 - ▶ Also used for making subprogram calls, passing parameters and local variables.
- ▶ Data can only be added in double word units
- ▶ PUSH
 - ▶ inserts a double word on the stack by **subtracting 4 from ESP**
 - ▶ And then stores the double word at [ESP]

Instructions: Setup

- ▶ POP
 - ▶ reads the double word at [ESP]
 - ▶ And then adds 4 to ESP
- ▶ CALL
 - ▶ Call subprogram
 - ▶ Make an unconditional jump to a subprogram
 - ▶ And pushes the address of the next instruction on the stack

Instructions: Setup

- ▶ RET
 - ▶ Pops off an address
 - ▶ And jumps to that address.
- ▶ When using this inst.: It is very important that one manage the stack correctly so that **the right number is popped off by the RET**.

Example

```
void swap(int *xp, int *yp)
{
    int t0 = *xp;
    int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

swap:

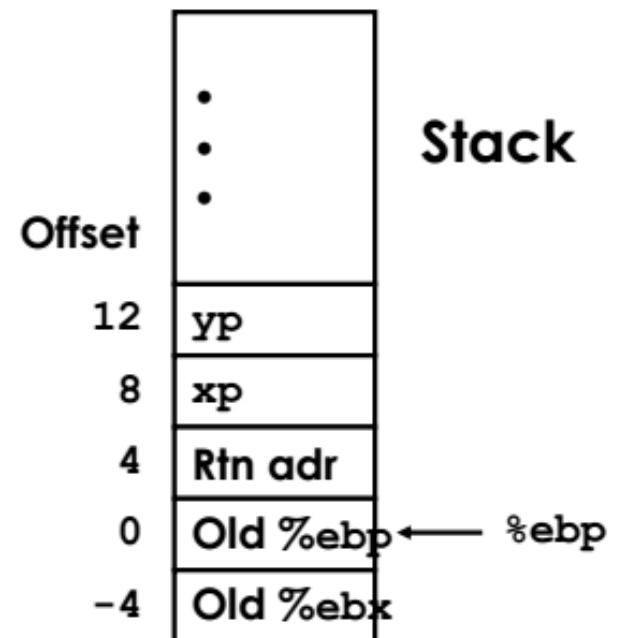
```
pushl %ebp  
movl %esp,%ebp  
pushl %ebx } Set Up  
  
movl 12(%ebp),%ecx  
movl 8(%ebp),%edx  
movl (%ecx),%eax  
movl (%edx),%ebx  
movl %eax,(%edx)  
movl %ebx,(%ecx) } Body  
  
movl -4(%ebp),%ebx  
movl %ebp,%esp  
popl %ebp  
ret } Finish
```

Example (contd.)

| Register | Variable |
|----------|----------|
| %ecx | yp |
| %edx | xp |
| %eax | t1 |
| %ebx | t0 |

```
void swap(int *xp, int *yp)
{
    int t0 = *xp;
    int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

```
movl 12(%ebp),%ecx # ecx = yp
movl 8(%ebp),%edx # edx = xp
movl (%ecx),%eax # eax = *yp (t1)
movl (%edx),%ebx # ebx = *xp (t0)
movl %eax,(%edx) # *xp = eax
movl %ebx,(%ecx) # *yp = ebx
```



Example (contd.)

| | |
|------|-------|
| %eax | |
| %edx | |
| %ecx | |
| %ebx | |
| %esi | |
| %edi | |
| %esp | |
| %ebp | 0x104 |

```

movl 12(%ebp), %ecx # ecx = yp
movl 8(%ebp), %edx # edx = xp
movl (%ecx), %eax # eax = *yp (t1)
movl (%edx), %ebx # ebx = *xp (t0)
movl %eax, (%edx) # *xp = eax
movl %ebx, (%ecx) # *yp = ebx

```

| Offset | |
|--------|---------|
| 12 | 0x120 |
| 8 | 0x124 |
| 4 | Rtn adr |
| 0 | 0x108 |
| -4 | 0x104 |
| | 0x100 |

%ebp → 0

Example (contd.)

| | |
|------|-------|
| %eax | |
| %edx | |
| %ecx | 0x120 |
| %ebx | |
| %esi | |
| %edi | |
| %esp | |
| %ebp | 0x104 |

```

movl 12(%ebp), %ecx # ecx = yp
movl 8(%ebp), %edx # edx = xp
movl (%ecx), %eax # eax = *yp (t1)
movl (%edx), %ebx # ebx = *xp (t0)
movl %eax, (%edx) # *xp = eax
movl %ebx, (%ecx) # *yp = ebx

```

Offset

| | | |
|----|---------|-------|
| 12 | 0x120 | 0x124 |
| 8 | 0x124 | 0x120 |
| 4 | Rtn adr | 0x11c |
| 0 | | 0x118 |
| -4 | | 0x114 |
| | | 0x110 |
| | | 0x10c |
| | | 0x108 |
| | | 0x104 |
| | | 0x100 |

%ebp → 0

Example (contd.)

| | |
|------|-------|
| %eax | |
| %edx | 0x124 |
| %ecx | 0x120 |
| %ebx | |
| %esi | |
| %edi | |
| %esp | |
| %ebp | 0x104 |

```

movl 12(%ebp),%ecx # ecx = yp
movl 8(%ebp),%edx # edx = xp
movl (%ecx),%eax # eax = *yp (t1)
movl (%edx),%ebx # ebx = *xp (t0)
movl %eax,(%edx) # *xp = eax
movl %ebx,(%ecx) # *yp = ebx

```

Offset

| | |
|-----|-----------|
| 123 | 0x124 |
| 456 | 0x120 |
| | 0x11c |
| | 0x118 |
| | 0x114 |
| | 12 0x120 |
| | 8 0x124 |
| | 0x10c |
| | 4 Rtn adr |
| | 0x108 |
| | 0x104 |
| | 0x100 |

%ebp → 0 -4

Example (contd.)

| | |
|------|-------|
| %eax | 456 |
| %edx | 0x124 |
| %ecx | 0x120 |
| %ebx | |
| %esi | |
| %edi | |
| %esp | |
| %ebp | 0x104 |

```

movl 12(%ebp),%ecx # ecx = yp
movl 8(%ebp),%edx # edx = xp
movl (%ecx),%eax # eax = *yp (t1)
movl (%edx),%ebx # ebx = *xp (t0)
movl %eax,(%edx) # *xp = eax
movl %ebx,(%ecx) # *yp = ebx

```

Offset

| | |
|-----|-------|
| 123 | 0x124 |
| 456 | 0x120 |
| | 0x11c |
| | 0x118 |
| | 0x114 |
| | 0x110 |
| | 0x10c |
| | 0x108 |
| | 0x104 |
| | 0x100 |

%ebp → 0

Example (contd.)

| | |
|------|-------|
| %eax | 456 |
| %edx | 0x124 |
| %ecx | 0x120 |
| %ebx | 123 |
| %esi | |
| %edi | |
| %esp | |
| %ebp | 0x104 |

```

movl 12(%ebp),%ecx # ecx = yp
movl 8(%ebp),%edx # edx = xp
movl (%ecx),%eax # eax = *yp (t1)
movl (%edx),%ebx # ebx = *xp (t0)
movl %eax,(%edx) # *xp = eax
movl %ebx,(%ecx) # *yp = ebx

```

Offset

| | |
|---------|-------|
| 123 | 0x124 |
| 456 | 0x120 |
| | 0x11c |
| | 0x118 |
| | 0x114 |
| | 0x110 |
| | 0x10c |
| 123 | 0x108 |
| Rtn adr | 0x108 |
| | 0x104 |
| | 0x100 |

%ebp → 0

-4

Example (contd.)

| | |
|------|-------|
| %eax | 456 |
| %edx | 0x124 |
| %ecx | 0x120 |
| %ebx | 123 |
| %esi | |
| %edi | |
| %esp | |
| %ebp | 0x104 |

```

movl 12(%ebp),%ecx # ecx = yp
movl 8(%ebp),%edx # edx = xp
movl (%ecx),%eax # eax = *yp (t1)
movl (%edx),%ebx # ebx = *xp (t0)
movl %eax,(%edx) # *xp = eax
movl %ebx,(%ecx) # *yp = ebx

```

Offset

| | |
|-----|-------|
| 456 | 0x124 |
| 456 | 0x120 |
| | 0x11c |
| | 0x118 |
| | 0x114 |
| | 0x110 |
| | 0x10c |
| | 0x108 |
| | 0x104 |
| | 0x100 |

yp 12 0x120

xp 8 0x124

%ebp → 0 Rtn adr 0x108

Example (contd.)

| | |
|------|-------|
| %eax | 456 |
| %edx | 0x124 |
| %ecx | 0x120 |
| %ebx | 123 |
| %esi | |
| %edi | |
| %esp | |
| %ebp | 0x104 |

```

movl 12(%ebp), %ecx # ecx = yp
movl 8(%ebp), %edx # edx = xp
movl (%ecx), %eax # eax = *yp (t1)
movl (%edx), %ebx # ebx = *xp (t0)
movl %eax, (%edx) # *xp = eax
movl %ebx, (%ecx) # *yp = ebx

```

Offset

| | |
|---------|-------|
| 456 | 0x124 |
| 123 | 0x120 |
| | 0x11c |
| | 0x118 |
| | 0x114 |
| | 0x110 |
| | 0x10c |
| Rtn adr | 0x108 |
| | 0x104 |
| | 0x100 |

%ebp → 0

Some Integer Instructions

Format

Computation

Two-Operand Instructions

| | | |
|-----------------------------|---------------------|-------------------------------|
| <code>addl Src,Dest</code> | $Dest = Dest + Src$ | |
| <code>subl Src,Dest</code> | $Dest = Dest - Src$ | |
| <code>imull Src,Dest</code> | $Dest = Dest * Src$ | |
| <code>sall k,Dest</code> | $Dest = Dest \ll k$ | Also called <code>shll</code> |
| <code>sarl k,Dest</code> | $Dest = Dest \gg k$ | Arithmetic |
| <code>shrl k,Dest</code> | $Dest = Dest \gg k$ | Logical |

k is an immediate value or contents of %cl

Some Integer Instructions

Format Computation

Two-Operand Instructions

xorl Src,Dest *Dest = Dest ^ Src*

andl Src,Dest *Dest = Dest & Src*

orl Src,Dest *Dest = Dest | Src*

Some Integer Instructions

Format Computation

One-Operand Instructions

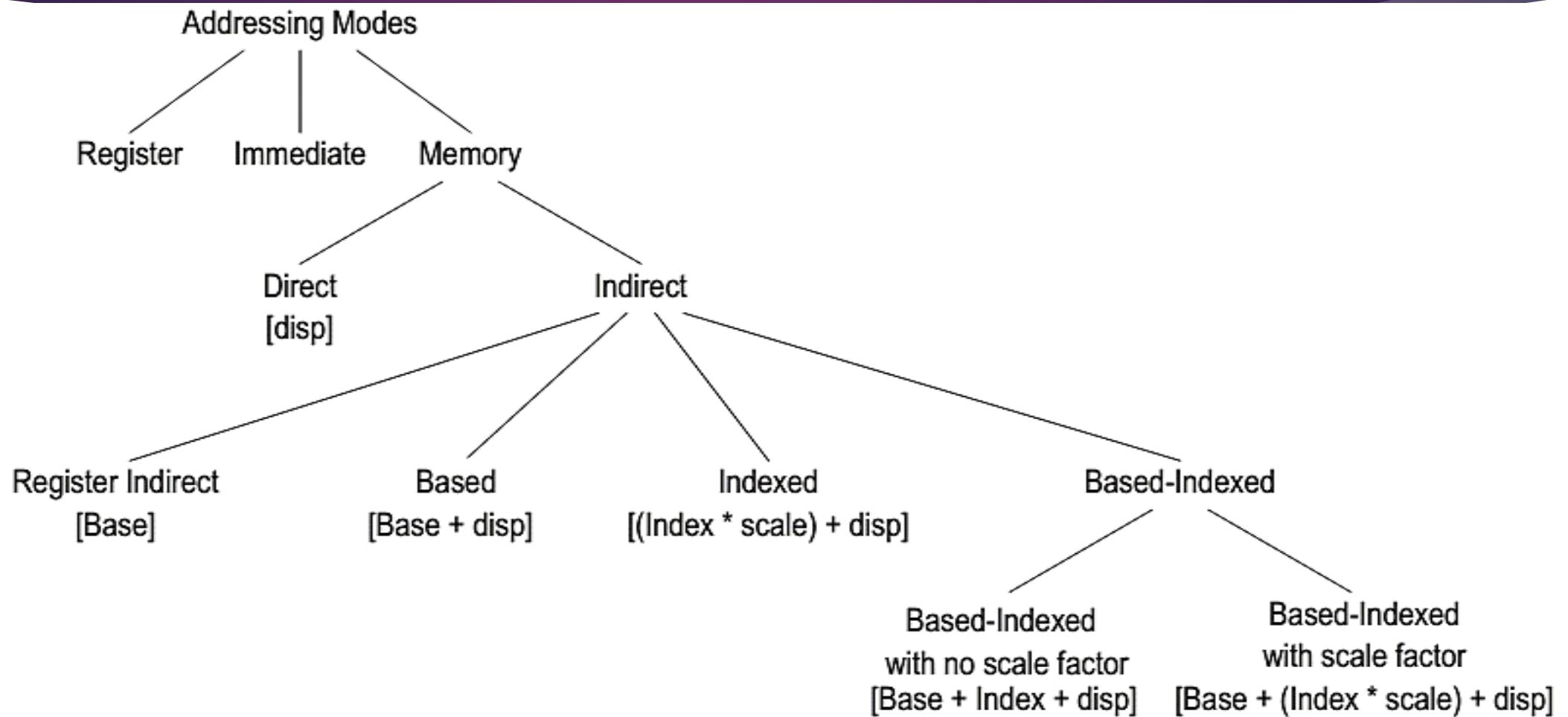
incl Dest $Dest = Dest + 1$

decl Dest $Dest = Dest - 1$

negl Dest $Dest = -Dest$

notl Dest $Dest = \sim Dest$

Addressing modes (32-bits)



Operands in X86 (NASM)

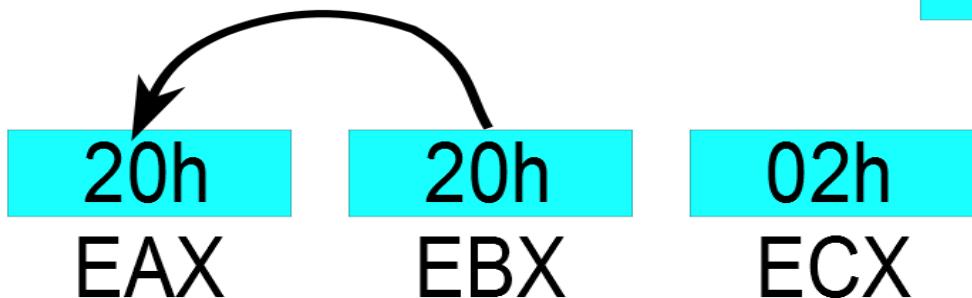
- ▶ Register: **MOV EAX, EBX**
 - ▶ Copy content from one register to another
- ▶ Immediate: **MOV EAX, 10h**
 - ▶ Copy constant to register
- ▶ Memory: different addressing modes
 - ▶ Typically at most one memory operand
 - ▶ Complex address computation supported

Addressing modes

- ▶ Direct: **MOV EAX, [10h]**
 - ▶ Copy value located at address 10h
- ▶ Indirect: **MOV EAX, [EBX]**
 - ▶ Copy value pointed to by register BX
- ▶ Indexed: **MOV AL, [EBX + ECX * 4 + 10h]**
 - ▶ Copy value from array (BX[4 * CX + 0x10])
- ▶ Pointers can be associated to type
 - ▶ **MOV AL, byte ptr [BX]**

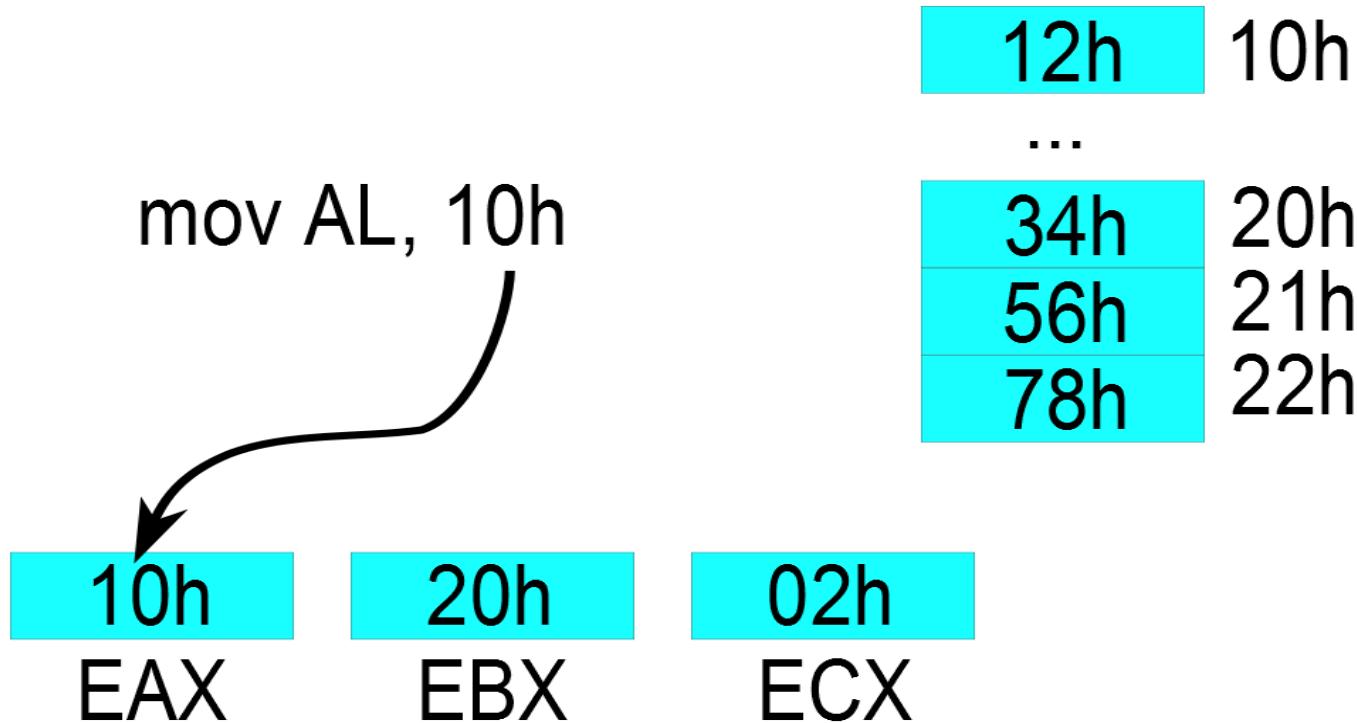
Operands and addressing modes: Register

mov AL, BL

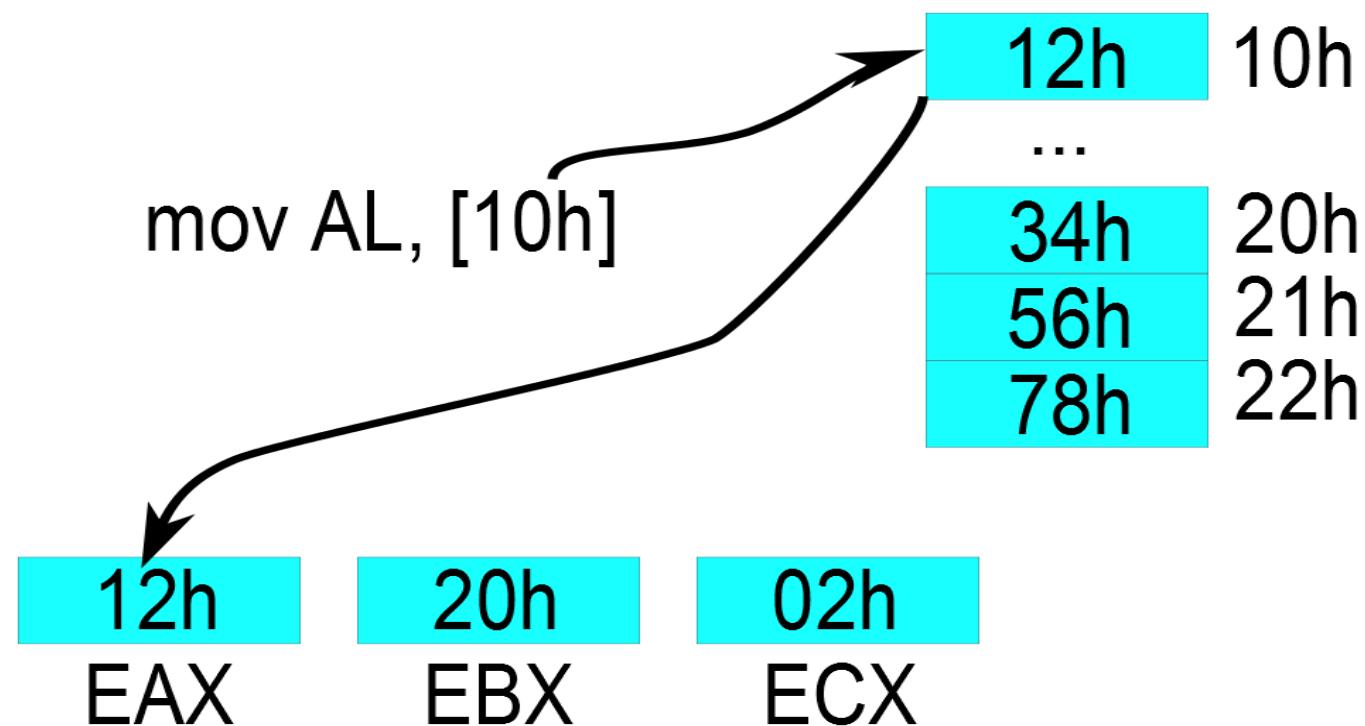


| | |
|------------------|------------------|
| <code>12h</code> | <code>10h</code> |
| ... | |
| <code>34h</code> | <code>20h</code> |
| <code>56h</code> | <code>21h</code> |
| <code>78h</code> | <code>22h</code> |

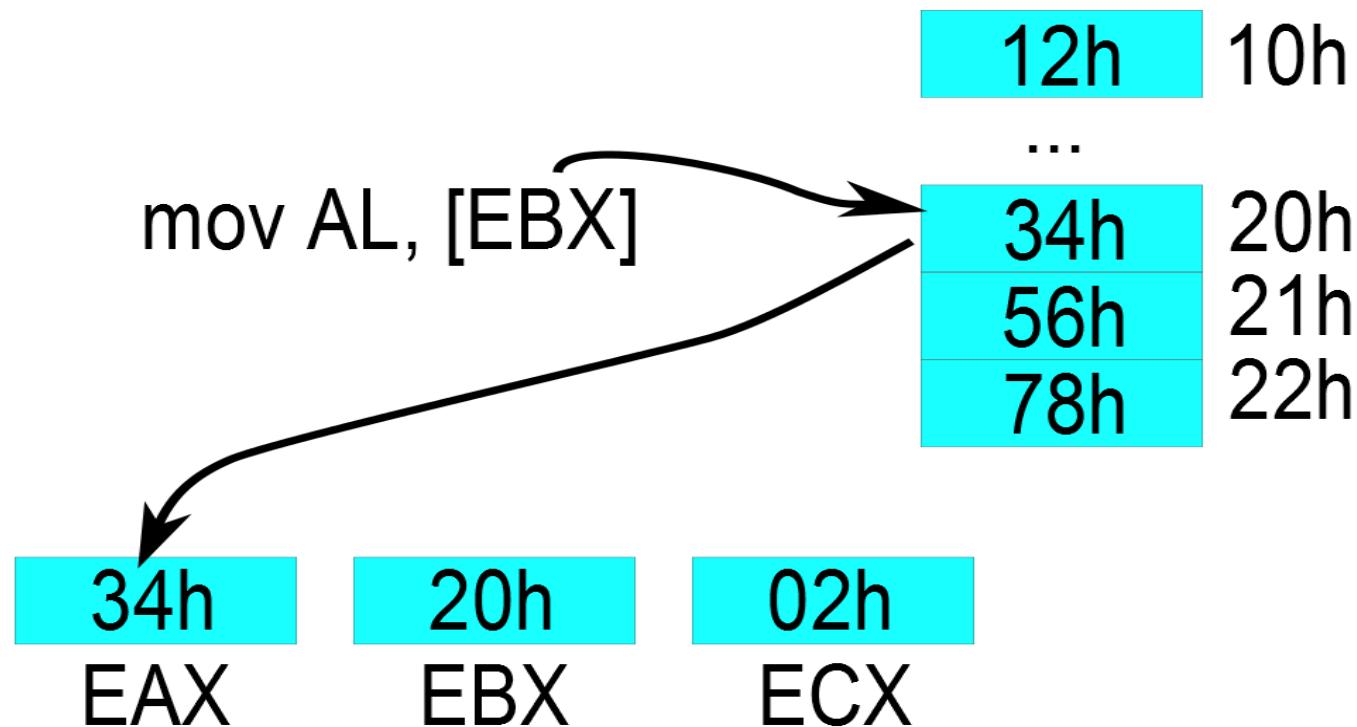
Operands and addressing modes: Immediate



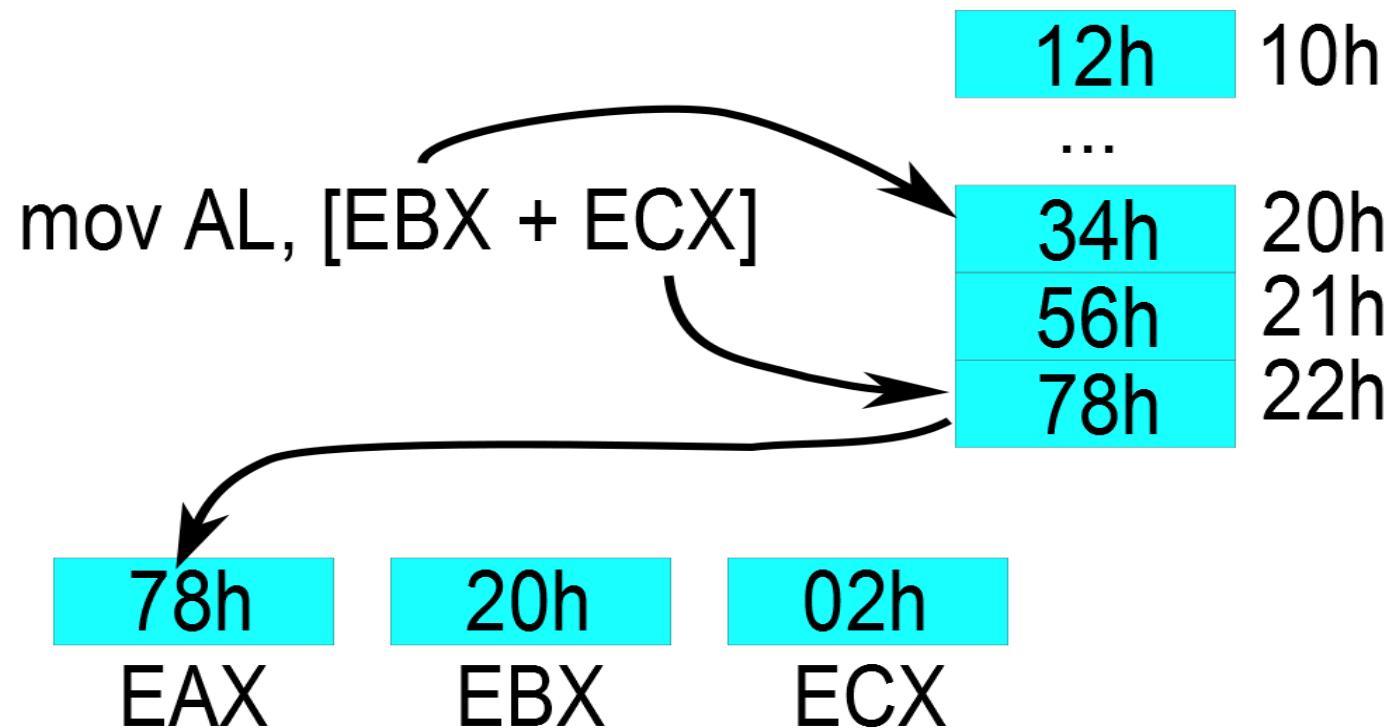
Operands and addressing modes: Direct



Operands and addressing modes: Indirect



Operands and addressing modes: Indexed



Generic 32bit Addressing Mode

- ▶ **Segment + Base + (Index * Scale) + displacement**

| | | | | |
|----|-----|-----|---|---------------------|
| CS | EAX | EAX | 1 | no displacement |
| SS | EBX | EBX | 2 | 8-bit displacement |
| DS | ECX | ECX | 4 | 32-bit displacement |
| ES | EDX | EDX | 8 | |
| FS | ESI | ESI | | |
| GS | EDI | EDI | | |
| | EBP | EBP | | |
| | ESP | | | |

Addressing Mode GNU-AS

► **D(Rb, Ri, S)**

Mem [Reg[Rb] + S*Reg[Ri] + D]

- ▶ D: Constant "displacement" 1, 2, or 4 bytes
- ▶ Rb: Base register: Any of 8 integer registers
- ▶ Ri: Index register: Any, **except for %esp**
 - ▶ Unlikely you'd use %ebp, either
- ▶ S: Scale: 1, 2, 4, or 8

GAS - Indexed Addressing Modes

▶ Special Cases

- ▶ (Rb, Ri) **Mem [Reg[Rb] + Reg[Ri]]**
- ▶ D(Rb, Ri) **Mem [Reg[Rb] + Reg[Ri] + D]**
- ▶ (Rb, Ri, S) **Mem [Reg[Rb] + S*Reg[Ri]]**

Solve the Followings:

| Expression |
|----------------------------|
| <code>0x8(%edx)</code> |
| <code>(%edx,%ecx)</code> |
| <code>(%edx,%ecx,4)</code> |
| <code>0x80(,%edx,2)</code> |

► Given:

| | |
|-------------------|---------------------|
| <code>%edx</code> | <code>0xf000</code> |
| <code>%ecx</code> | <code>0x100</code> |

Another Programming Example

Objective:

- ▶ Take an integer through Command-line
- ▶ Print the value
- ▶ Add 1 to that
- ▶ Print the result

Another Program (1)

```
SECTION .text
    global main
main:
; set-up phase
    push ebp
    mov ebp, esp

; finish phase
    mov esp, ebp
    pop ebp
    ret
```

Another Program (2)

```
SECTION .text
    global main
main:
; set-up phase
    push ebp
    mov ebp, esp

; get the command-line data
    mov ebx, DWORD [esp + 12] ; get argv starting address
    mov ebx, [ebx + 4]          ; get the second argument data

; finish phase
    mov esp, ebp
    pop ebp
    ret
```

Another Program (3)

```
SECTION .data
msg:    db      "You Entered - %s", 10, 0 ; print argv[1] data

SECTION .text
        extern printf
        global main

main:
; set-up phase
        push ebp
        mov ebp, esp
```

Another Program (3)

```
; get the command-line data
    mov ebx, DWORD [esp + 12] ; get argv starting address
    mov ebx, [ebx + 4]          ; get the second argument data

; print the value
    push ebx                   ; put data on stack for call
    push msg                   ; print the value
    call printf

; finish phase
    add esp, 8                 ; esp back to start
    mov esp, ebp
    pop ebp
    ret
```

Another Program (4)

```
SECTION .data
msg:    db      "You Entered - %s", 10, 0 ; print argv[1] data
msg2:   db      "This is int - %d", 10, 0 ; print INT equivalent

SECTION .text
extern printf
extern atoi
global main

main:
; set-up phase
push ebp
mov ebp, esp
```

Another Program (4)

```
; get the command-line data
    mov ebx, DWORD [esp + 12] ; get argv starting address
    mov ebx, [ebx + 4]          ; get the second argument data

; print the value
    push ebx                   ; put data on stack for call
    push msg                   ; print the value
    call printf

; convert to integer
    add esp, 4                 ; stack points to entry edx
    call atoi                  ; call atoi - return in EAX?
```

Another Program (4)

```
; get return value and add 1
add esp, 4                      ; esp points to start
inc eax                          ; increase eax for testing

; print the result
push eax                         ; push arg for print
push msg2                         ; push print message
call printf

; finish phase
add esp, 8                        ; esp back to start
mov esp, ebp
pop ebp
ret
```

Another Program (final)

```
SECTION .data
msg:    db      "You Entered - %s", 10, 0 ; print argv[1] data
msg2:   db      "This is int - %d", 10, 0 ; print INT equivalent

SECTION .text
extern printf
extern atoi
global main

main:
; set-up phase
push ebp
mov ebp, esp
```

Another Program (final)

```
; get the command-line data
    mov ebx, DWORD [esp + 12] ; get argv starting address
    mov ebx, [ebx + 4]          ; get the second argument data

; print the value
    push ebx                  ; put data on stack for call
    push msg                  ; print the value
    call printf

; convert to integer
    add esp, 4                ; stack points to entry edx
    call atoi                 ; call atoi - return in EAX?
```

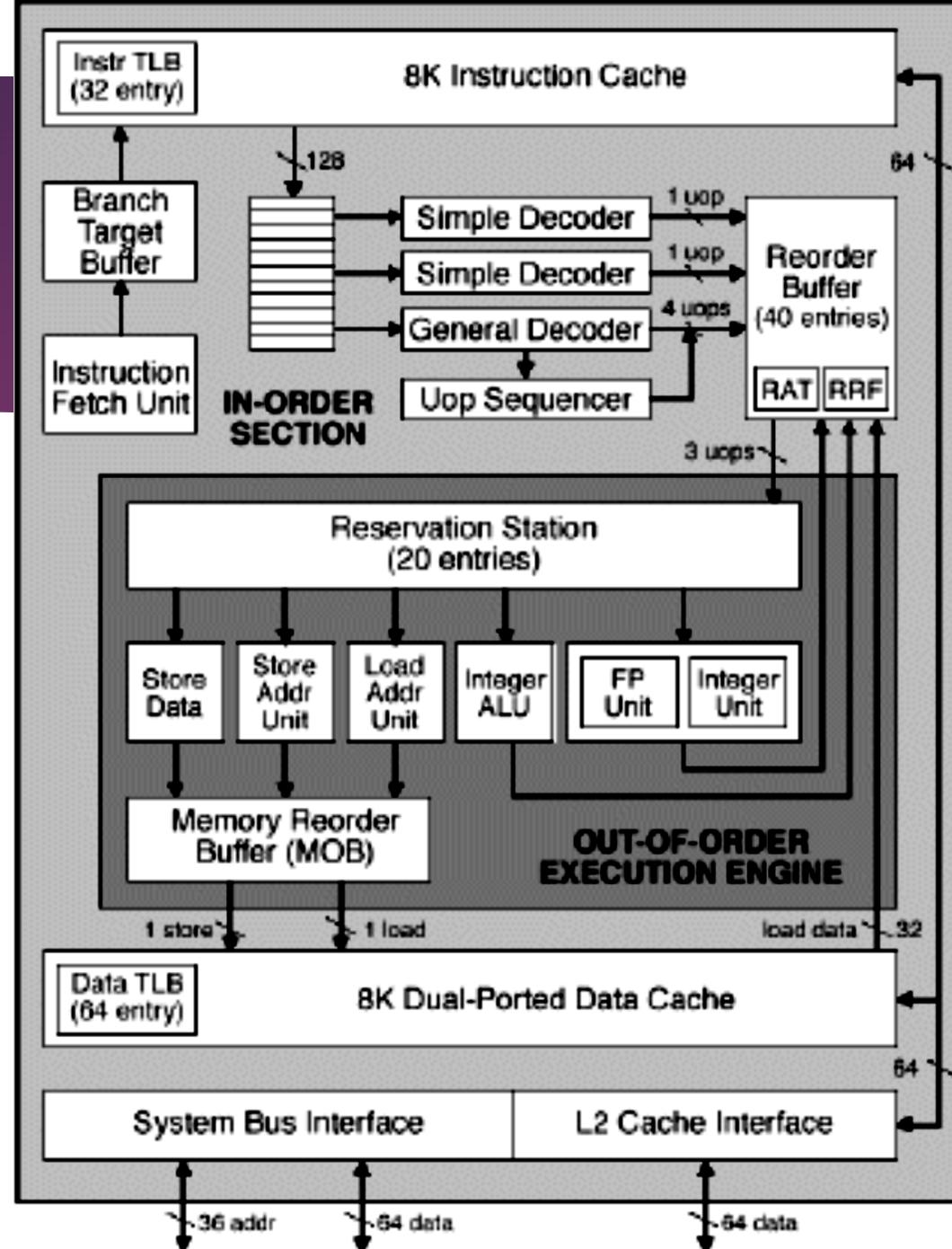
Another Program (final)

```
; get return value and add 1
add esp, 4                      ; esp points to start
inc eax                          ; increase eax for testing

; print the result
push eax                         ; push arg for print
push msg2                         ; push print message
call printf

; finish phase
add esp, 8                        ; esp back to start
mov esp, ebp
pop ebp
ret
```

PentiumPro Block Diagram



Operations

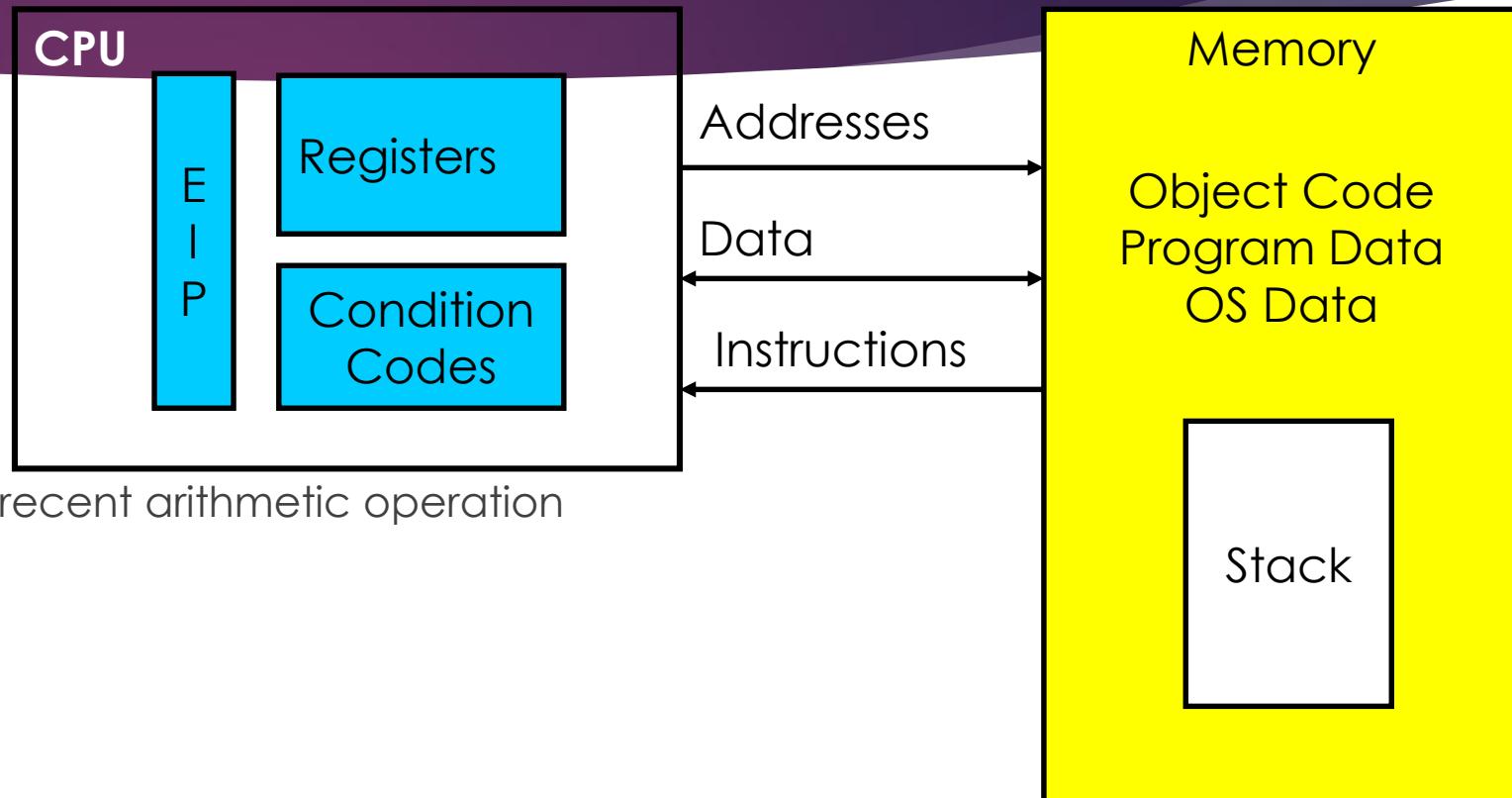
- ▶ Translates instructions dynamically into “Uops” / “μops”
 - ▶ 128 bits wide
- ▶ Holds
 - ▶ Operation 32
 - ▶ Two sources, and + 32 + 32
 - ▶ Destination + 32 = 128bits **WHY??**
- ▶ Executes Uops with “**Out of Order**” engine
 - ▶ Uop executed when
 1. Operands are available
 2. Functional unit available

Operations

- ▶ Execution controlled by “Reservation Stations”
 - ▶ Keeps track of data **dependencies** between uops
 - ▶ Allocates resources
- ▶ Consequences
 - ▶ Indirect relationship between
 - ▶ IA32 code &
 - ▶ What actually gets executed
 - ▶ Tricky to predict / optimize performance at assembly level

Machine View

- ▶ EIP (Program Counter)
 - ▶ Address of next instruction
- ▶ Register File
 - ▶ Heavily used program data
- ▶ Condition Codes
 - ▶ Store status information about most recent arithmetic operation
 - ▶ Used for conditional branching
- ▶ Memory
 - ▶ Byte addressable array
 - ▶ Code, user data, (some) OS data
 - ▶ Includes stack used to support procedures

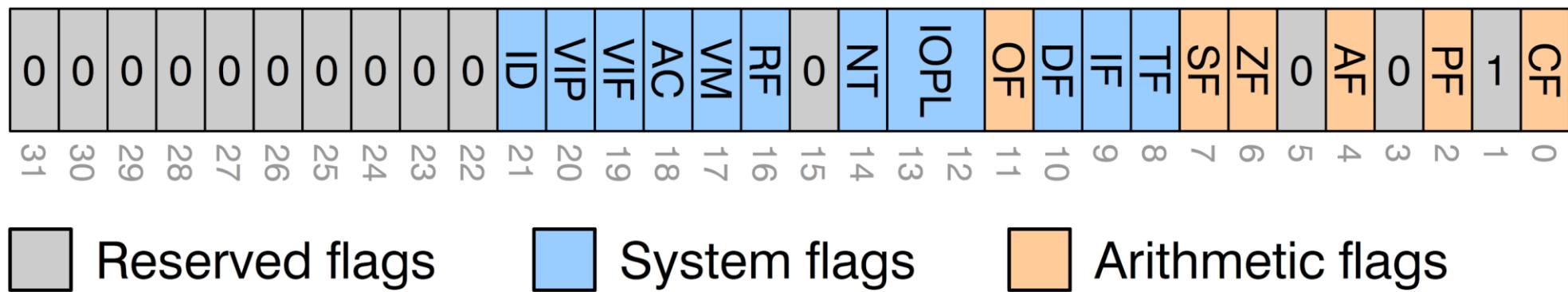


Flow Control Instructions

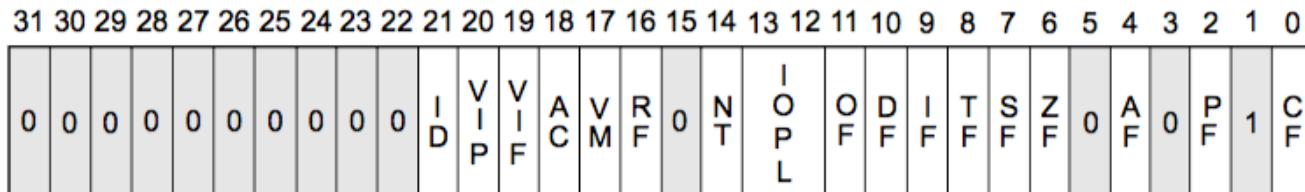
- ▶ JMP
- ▶ Jcc
- ▶ CALL
- ▶ RET

Revisit EFLAGS

eflags register



Revisit EFLAGS



- X ID Flag (ID) _____
- X Virtual Interrupt Pending (VIP) _____
- X Virtual Interrupt Flag (VIF) _____
- X Alignment Check / Access Control (AC) _____
- X Virtual-8086 Mode (VM) _____
- X Resume Flag (RF) _____
- X Nested Task (NT) _____
- X I/O Privilege Level (IOPL) _____
- S Overflow Flag (OF) _____
- C Direction Flag (DF) _____
- X Interrupt Enable Flag (IF) _____
- X Trap Flag (TF) _____
- S Sign Flag (SF) _____
- S Zero Flag (ZF) _____
- S Auxiliary Carry Flag (AF) _____
- S Parity Flag (PF) _____
- S Carry Flag (CF) _____

S Indicates a Status Flag

C Indicates a Control Flag

X Indicates a System Flag

Reserved bit positions. DO NOT USE.
Always set to values previously read.

JMP: jump

► Syntax:

JMP *dest*

► Operation (absolute jump):

EIP \leftarrow *dest*

► Operation (relative jump):

EIP \leftarrow EIP + *dest*

| | | | | | | | |
|----|----|----|----|----|----|----|---|
| - | - | - | - | - | - | - | - |
| of | df | sf | zf | af | pf | cf | |

Unconditional Jumps

- ▶ Jmp statement label
- ▶ We have two types of jumps,
 - ▶ Intersegment
 - ▶ Intrasegment
- ▶ Address can be in a register, variable or label.

Unconditional Jumps

- ▶ Example:

```
Start:    MOV AX, 0  
          INC AX,  
          JMP Start
```

Jcc: short jump conditional

► Syntax:

`Jcc dest`

► Operation:

```
if(cc)
    EIP ← EIP + dest
endif
```

► **Notes:** *cc* is any of the condition codes. *dest* must be within a signed 8-bit range (-128 to 127).

| | | | | | | | |
|----|----|----|----|----|----|----|---|
| - | - | - | - | - | - | - | - |
| of | df | sf | zf | af | pf | cf | |

Condition Codes

► Implicitly Set By Arithmetic Operations

addl Src, Dest

C analog: $t = a + b$

► CF set if carry out from most significant bit

► Used to detect unsigned overflow

► ZF set if $t == 0$

► SF set if $t < 0$

► OF set if two's complement overflow

$(a>0 \ \&\& \ b>0 \ \&\& \ t<0) \ || \ (a<0 \ \&\& \ b<0 \ \&\& \ t>=0)$

► **Not** Set by leal instruction

Condition Codes

| Sufix | Meaning | Flags |
|-------|----------------------------|-------|
| O | <i>Overflow</i> | OF=1 |
| NO | <i>No Overflow</i> | OF=0 |
| C | <i>Carry</i> | |
| B | <i>Below</i> | CF=1 |
| NAE | <i>Not Above nor Equal</i> | |
| NC | <i>No Carry</i> | |
| NB | <i>Not Below</i> | CF=0 |
| AE | <i>Above or Equal</i> | |

Condition Codes (continued)

| Sufix | Meaning | Flags |
|-------|----------------------------|---------------|
| Z | <i>Zero</i> | ZF=1 |
| E | <i>Equal</i> | |
| NZ | <i>Not Zero</i> | ZF=0 |
| NE | <i>Not Equal</i> | |
| BE | <i>Below or Equal</i> | CF=1 OR ZF=1 |
| NA | <i>Not Above</i> | |
| A | <i>Above</i> | CF=0 AND ZF=0 |
| NBE | <i>Not Below nor Equal</i> | |

Condition Codes (continued)

| Sufix | Meaning | Flags |
|-------|--------------------|-------|
| S | <i>Sign</i> | SF=1 |
| NS | <i>Not Sign</i> | SF=0 |
| P | <i>Parity</i> | PF=1 |
| PE | <i>Parity Even</i> | |
| NP | <i>Not Parity</i> | PF=0 |
| PO | <i>Parity Odd</i> | |

Condition Codes (continued)

| Sufix | Meaning | Flags |
|-------|------------------------------|----------------|
| L | <i>Less</i> | SF<>OF |
| NGE | <i>Not Greater nor Equal</i> | |
| GE | <i>Greater or Equal</i> | SF=OF |
| NL | <i>Not Less</i> | |
| LE | <i>Less or Equal</i> | ZF=1 OR SF<>OF |
| NG | <i>Not Greater</i> | |
| G | <i>Greater</i> | ZF=0 AND SF=OF |
| NLE | <i>Not Less nor Equal</i> | |

Condition Codes (continued)

- ▶ **Above** and **Below** are used for unsigned integer comparisons.
- ▶ **Greater** and **Less** are used for signed integer comparisons.

Conditional Jumps

- ▶ Dependent on condition codes.
- ▶ Example:

JZ → jump if zero flag is set.

Conditional Codes

- ▶ Code the following C routine using assembly language instructions.

Add a value X to eax;

If x < 0

Then

... (body for negative condition)

Else if x = 0

... (body for zero condition)

Else

... (body for positive condition)

End if

Conditional Codes

► Solution

```
Add eax, X      ;add a value to eax  
JNS elseifZero ;jump if eax is not negative  
...              ; code for negative condition  
JMP endCheck
```

elseifZero:

```
JNZ elsePos     ; jump if x is not zero  
...              ; code for zero condition  
JMP endCheck
```

elsePos: ...

; code for positive balance

endCheck:

Conditional Codes

```
Add eax, X          ;add a value to eax  
JNS elseifZero    ;jump if eax is not negative  
...                 ; code for negative condition  
JMP endCheck  
elseifZero:  
  JNZ elsePos       ; jump if x is not zero  
  ...                 ; code for zero condition  
  JMP endCheck  
elsePos:  
  ...                 ; code for positive balance  
endCheck:
```

Add a value X to eax;
If $x < 0$
Then
... (**body for negative condition**)
Else if $x = 0$
... (**body for zero condition**)
Else
... (**body for positive condition**)
End if

Setting Codes

Explicit Setting by Compare Instruction

`cmp Src1, Src2`

- ▶ like computing **Src1-Src2** without setting destination

- ▶ **CF** set if carry out from most significant bit
 - ▶ Used for unsigned comparisons
- ▶ **SF** set if $(a-b) < 0$
- ▶ **ZF** set if $a == b$
- ▶ **OF** set if two's complement overflow
$$(a>0 \ \&\& \ b<0 \ \&\& \ (a-b)<0) \ \|\ \ (a<0 \ \&\& \ b>0 \ \&\& \ (a-b)>0)$$

Compare Examples

- ▶ OP1= 3B
- ▶ OP2= 3B

- ▶ CF=OF=SF=0
- ▶ ZF=1

- ▶ OP1==OP2 signed and unsigned

Compare Examples

- ▶ OP1= 3B
- ▶ OP2= 15

- ▶ OP1-OP2= 26
- ▶ CF=OF=SF=ZF=0

- ▶ OP1>OP2 signed and unsigned

Compare Examples

- ▶ OP1= 15
- ▶ OP2= F6
- ▶ OP1-OP2= 1F

$$15 - F6 \Rightarrow 15 + 0A$$

$$\begin{array}{r} 0001\ 0101 \\ +0000\ 1010 \\ \hline 0001\ 1111 = 1F \end{array}$$

- ▶ SF= OF= ZF= 0s
- ▶ Signed operation = op1>op2
- ▶ CF= (1 – borrow)
- ▶ Unsigned operation =op1 < op2

Compare Examples

Legal Examples

Cmp eax, 356

cmp value, 03dh

Cmp bh, '\$'

Illegal examples

Cmp 1000, total

Compare Programming Ex.

- ▶ Code the following routine in assembly language.

If val < 10

Then

add 1 to xcount;

Else

add 1 to ycount;

End if;

Compare Programming Ex

- ▶ Solution: assume 'val' is in ebx

```
cmp ebx, 10      ;value < 10
jnl Elsey
inc xcount      ;add 1 to xcount
jmp endVal
```

Elsey:

```
inc ycount      ;add 1 to ycount
```

endVal:

Compare Examples for JNL

- ▶ OP1=7F $7F - FF \Rightarrow 7F + 01$
 - ▶ OP2= FF
 - ▶ OP1-OP2=80
 - ▶ SF=1
 - ▶ OF=1
 - ▶ ZF=0
 - ▶ Signed operation: **op1>=op2**
 - ▶ Unsigned operation =op1 < op2
- | | |
|-------|------|
| 0111 | 1111 |
| +0000 | 0001 |
| <hr/> | |
| 1000 | 0000 |
| = 80 | |

Instruction: Test

Explicit Setting by Test instruction

`testl Src2, Src1`

`test Src1, Src2`

- ▶ Sets condition codes based on value of: **Src1 & Src2**
 - ▶ Useful to have one of the operands be a mask
- ▶ **testl b,a**
 - ▶ like computing **a&b** without setting destination

Instruction: Test

- ▶ *Src1*

- ▶ AL/AX/EAX (only if **Src2** is immediate)
- ▶ Register
- ▶ Memory

- ▶ *Src2*

- ▶ Register
- ▶ Immediate

Instruction: Test

- ▶ Modified flags (Assume **Temp = a & b**)
 - ▶ **SF** = **MostSignificantBit(Temp)**
 - ▶ **If (Temp == 0) ZF = 1 else ZF = 0**
 - ▶ **PF** = **BitWiseXorNor(Temp [Max-1:0])**
 - ▶ **CF** = 0
 - ▶ **OF** = 0
 - ▶ **AF** is undefined

Set on Condition (setcc reg8)

| <u>Instruction</u> | <u>Description</u> | <u>Condition</u> | <u>Comments</u> |
|--------------------|--------------------|------------------|-------------------------|
| SETC | Set if carry | Carry = 1 | Same as SETB, SETNAE |
| SETNC | Set if no carry | Carry = 0 | Same as SETNB, SETAE |
| SETZ | Set if zero | Zero = 1 | Same as SETE |
| SETNZ | Set if not zero | Zero = 0 | Same as SETNE |
| SETS | Set if sign | Sign = 1 | - |
| SETNS | Set if no sign | Sign = 0 | - |
| SETO | Set if overflow | Overflow=1 | - |
| SETNO | Set if no overflow | Overflow=0 | - |
| SETP | Set if parity | Parity = 1 | Same as SETPE |
| SETPE | Set if parity even | Parity = 1 | Same as SETP |
| SETNP | Set if no parity | Parity = 0 | Same as SETPO |
| SETPO | Set if parity odd | Parity = 0 | Same as SETNP |

Condition Code Example

```
int max(int x, int y)
{
    if (x > y)
        return x;
    else
        return y;
}
```

```
_max:
    pushl %ebp
    movl %esp,%ebp
    movl 8(%ebp),%edx
    movl 12(%ebp),%eax
    cmpl %eax,%edx
    jle L9
    movl %edx,%eax
L9:
    movl %ebp,%esp
    popl %ebp
    ret
```

Set Up

Body

Finish

Condition Code Example

```
int max(int x, int y)
{
    if (x > y)
        goto Flag;

    return y;

Flag:
    return x;
}
```

```
_max:
    pushl %ebp
    movl %esp,%ebp
    movl 8(%ebp),%edx
    movl 12(%ebp),%eax
    cmpl %eax,%edx
    jle L9
    movl %edx,%eax
L9:
    movl %ebp,%esp
    popl %ebp
    ret
```

Set Up

Body

Finish

“Do-While” Example

C Code

```
int fact_do
    (int x)
{
    int result = 1;
    do {
        result *= x;
        x = x-1;
    } while (x > 1);
    return result;
}
```

Goto Version

```
int fact_goto(int x)
{
    int result = 1;
loop:
    result *= x;
    x = x-1;
    if (x > 1)
        goto loop;
    return result;
}
```

“Do-While” Example

Goto Version

```
int fact_goto
    (int x)
{
    int result = 1;
loop:
    result *= x;
    x = x-1;
    if (x > 1)
        goto loop;
    return result;
}
```

- Registers

| | |
|------|--------|
| %edx | x |
| %eax | result |

Assembly

```
_fact_goto:
    pushl %ebp          # Setup
    movl %esp,%ebp      # Setup
    movl $1,%eax        # eax = 1
    movl 8(%ebp),%edx  # edx = x

L11:
    imull %edx,%eax    # result *= x
    decl %edx           # x--
    cmpl $1,%edx        # Compare x : 1
    jg L11              # if > goto loop

    movl %ebp,%esp      # Finish
    popl %ebp            # Finish
    ret                 # Finish
```

Multiplication

- ▶ **mul arg** Unsigned multiplication
 - ▶ This multiplies "arg" **by** the value of corresponding byte-length in the **EAX register**.

| | | | |
|---|--------|---------|---------|
| operand size | 1 byte | 2 bytes | 4 bytes |
| other operand | AL | AX | EAX |
| higher part of result stored in: | AH | DX | EDX |
| lower part of result stored in: | AL | AX | EAX |

- ▶ Signed Multi[pli]cation: **imul arg**

Division

- ▶ Unsigned: **div arg**
 - ▶ This divides the value in the dividend register(s) by "arg", see table below

| divisor size: | 1 byte | 2 bytes | 4 bytes |
|-----------------------------|--------|---------|---------|
| dividend | AX | DX:AX | EDX:EAX |
| remainder stored in: | AH | DX | EDX |
| quotient stored in: | AL | AX | EAX |

- ▶ Signed: **idiv arg**

Sign extended Register

- ▶ CDQ: Convert Double to Quad-Word (Intel)
- ▶ CLTD: AT&T

- ▶ Sign-extends EAX into EDX
- ▶ Forms the quad-word **EDX:EAX**.
- ▶ Since (I)DIV uses EDX:EAX as its input, CDQ must be called after setting EAX if EDX is not manually initialized (as in 64/32 division) before (I)DIV.

Load Effective Address

- ▶ AT&T: **leal** src, dest
- ▶ Intel: **lea** dest, src

- ▶ Calculates the **address** of the **src** operand and **loads** it into the **dest** operand
- ▶ **Src:** Immediate / Register / Memory
- ▶ **Dest:** Register

- ▶ No FLAGS are modified by this instruction

Load Effective Address

- ▶ Powerful operation
- ▶ Can also be used for calculating general-purpose unsigned integer arithmetic
- ▶ SRC can use **General Addressing Mode**
- ▶ **[eax + edx*4 -4] (Intel syntax)**
 - ▶ or
- ▶ **-4(%eax, %edx, 4) (GAS syntax)**

Load Effective Address Example

- ▶ Suppose you want to multiply the content of **EBX** with **18**
- ▶ LEA can be used as follows

```
lea ebx, [ebx*2] ; Multiply ebx by 2
```

```
lea ebx, [ebx*8+ebx] ; Multiply ebx by 9, which totals ebx*18
```

Jump Table

- ▶ To implement a “**switch-case**”.
- ▶ To select a **function** to be evoked.

- ▶ We will construct a **array** of the **jump addresses** .
- ▶ For each number will jump to the **corresponding entry** in the jump table

Example

- ▶ Main.c

```
extern void jumper(int);

int main (int argc , char* argv)

{

    jumper (0);

    jumper (1);

    jumper (2);

    return 0;

}
```

jump.s

section .data

jt:

| | |
|----|---------|
| dd | label_1 |
| dd | label_2 |

str0: db "Got the number 0",10,0

str1: db "Got the number 1",10,0

str2: db "Out of bound",10,0

str3: db "num = %d",10,0

section .text

align 16

global jumper

extern printf

jump.s

jumper:

| | | |
|-------|-------------------|--------------------------------------|
| push | ebp | |
| mov | ebp,esp | |
| pusha | | |
| mov | ebx,dword [ebp+8] | |
| push | ebx | |
| push | str3 | |
| call | printf | ; Print num |
| add | esp, 8 | |
| cmp | ebx,0 | ; Check if num is in bounds |
| jb | out_of | |
| cmp | ebx ,1 | |
| ja | out_of | |
| shl | ebx,2 | ; num = num * 4 |
| jmp | dword [ebx+jt] | ; Jump according to address in table |

jump.s

label_1:

```
push    str0
call    printf
add     esp, 4
jmp     end
```

label_2:

```
push    str1
call    printf
add     esp, 4
jmp     end
```

out_of:

```
push    str2
call    printf
add     esp, 4
jmp     end
```

end:

```
popa
pop    ebp
ret
```

Output

num = 0

Got the number 0

num = 1

Got the number 1

num = 2

Out of bound

jump.s

section .data

jt:

| | |
|----|---------|
| dd | label_1 |
| dd | label_2 |

str0: db "Got the number 0",10,0
str1: db "Got the number 1",10,0
str2: db "Out of bound",10,0
str3: db "num = %d",10,0

section .text

| | |
|--------|--------|
| align | 16 |
| global | jumper |
| extern | printf |

jump.s

jumper:

```
push    ebp
mov     ebp,esp
pusha
mov     ebx,dword [ebp+8]
push    ebx
push    str3
call    printf          ; Print num
add    esp, 8
cmp    ebx,0            ; Check if num is in bounds
jb     out_of
cmp    ebx ,1
ja     out_of
shl    ebx,2
jmp    dword [ebx+jt]
```

; num = num * 4
; Jump according to address in table

jump.s

label_1:

```
push    str0
call    printf
add     esp, 4
jmp    end
```

label_2:

```
push    str1
call    printf
add     esp, 4
jmp    end
```

out_of:

```
push    str2
call    printf
add     esp, 4
jmp    end
```

end:

```
popa
pop    ebp
ret
```

Switch Statement Example (IA32)

Setup:

```
switch_eg:  
    pushl %ebp          # Setup  
    movl %esp, %ebp     # Setup  
    pushl %ebx          # Setup  
    movl $1, %ebx        # w = 1  
    movl 8(%ebp), %edx  # edx = x  
    movl 16(%ebp), %ecx # ecx = z  
    cmpl $6, %edx      # x:6  
    ja .L61            # if > goto default  
    jmp * .L62(,%edx,4) # goto JTab[x]
```

```
long switch_eg  
(long x, long y, long z)  
{  
    long w = 1;  
    switch(x) {  
        . . .  
    }  
    return w;  
}
```

Jump Table

100

Table Contents

```
.section .rodata
.align 4
.L62:
.long .L61 # x = 0
.long .L56 # x = 1
.long .L57 # x = 2
.long .L58 # x = 3
.long .L61 # x = 4
.long .L60 # x = 5
.long .L60 # x = 6
```

```
switch(x) {
    case 1:          // .L56
        w = y*z;
        break;
    case 2:          // .L57
        w = y/z;
        /* Fall Through */
    case 3:          // .L58
        w += z;
        break;
    case 5:
    case 6:          // .L60
        w -= z;
        break;
    default:         // .L61
        w = 2;
}
```

Code Blocks (Partial)

```
switch(x) {  
    . . .  
    case 2:      // .L57  
        w = y/z;  
        /* Fall Through */  
    case 3:      // .L58  
        w += z;  
        break;  
    . . .  
    default:     // .L61  
        w = 2;  
}
```

```
.L61: // Default case  
    movl $2, %ebx    # w = 2  
    movl %ebx, %eax # Return w  
    popl %ebx  
    leave  
    ret  
.L57: // Case 2:  
    movl 12(%ebp), %eax # y  
    cltd                # Div prep  
    idivl %ecx          # y/z  
    movl %eax, %ebx # w = y/z  
# Fall through  
.L58: // Case 3:  
    addl %ecx, %ebx # w+= z  
    movl %ebx, %eax # Return w  
    popl %ebx  
    leave  
    ret
```

Code Blocks (Rest)

```
switch(x) {  
    case 1:          // .L56  
        w = y*z;  
        break;  
        . . .  
    case 5:  
    case 6:          // .L60  
        w -= z;  
        break;  
        . . .  
}
```

```
.L60: // Cases 5&6:  
    subl %ecx, %ebx # w -= z  
    movl %ebx, %eax # Return w  
    popl %ebx  
    leave  
    ret  
.L56: // Case 1:  
    movl 12(%ebp), %ebx # w = y  
    imull %ecx, %ebx      # w*= z  
    movl %ebx, %eax # Return w  
    popl %ebx  
    leave  
    ret
```