

Binary Search.

- Real life example
- Coding Problem Example
- Iterative Code
- Recursive Code
- Time Complexity
- Overflow.

Binary Search:-

Binary Search:-
→ If you find from the first till last page → is known as LINEAR SEARCH.
i.e. - have a sorted arr. "ja".

→ Binary Search is applicable → if we have a sorted array. "May".
the Search Space is Sorted.

- Binary Search is applicable →
- Binary Search applied → where the Search Space is Sorted.

Ex: \rightarrow $\begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ \underline{3}, & \underline{4}, & \underline{6}, & 7, & 9, & 12, & 16, & 17 \end{matrix}$ $n=8$.
array is Sorted.

target = 6.

→ then it is a 3 step.

- But suppose we want 17. → then it is a 28 step process.
- where TIME COMPLEXITY = $O(N)$. because you have iterated the entire array.

→ Now if we take Binary Search.

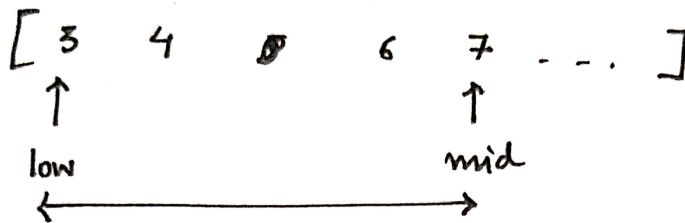
[3, 4, 6, 7, 9, 12, 16, 17]

↑ ↑
low high

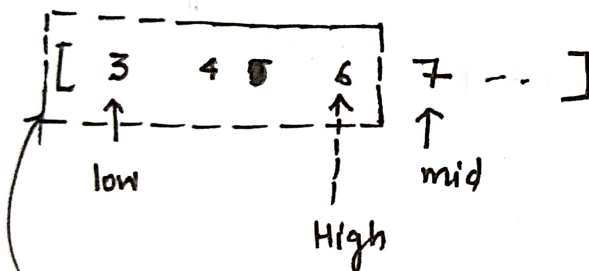
target = 6.

$$\text{mid} = \frac{0+7}{2} = 3.5 = 3$$

at $\text{mid} = 3 \Rightarrow a[\text{mid}] = a[3] = 7$.



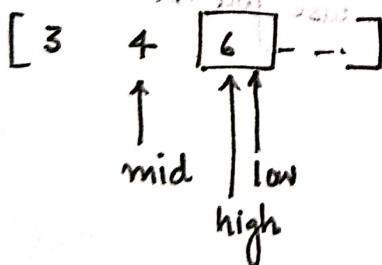
$a[\text{mid}] > a[\text{target}] \rightarrow$ so now Binary Search will be on LEFT HAND SIDE



now this much is only the Search Space.

$$\text{mid} = \frac{\text{low} + \text{high}}{2} = \frac{0+2}{2} = 1 =$$

$a[\text{mid}] = a[1] = 4$.



now move low at the one place right of mid. and low and high pointing to the same index.

$$\text{mid} = \frac{\text{low} + \text{high}}{2} = 2$$

$a[\text{mid}] = a[2] = 6$.

Iterative Soln. C++.

```
int search (vector<int> & nums, int target) {  
    int n = nums.size();  
    int low = 0, high = n-1;  
    while (low <= high) {  
        int mid = (low+high)/2;  
        if (nums[mid] == target) return mid;  
        else if (target > nums[mid])  
            low = mid+1;  
        else high = mid-1;  
    }  
    return -1;  
}
```

Recursive C++.

Recursion used when we are doing repetitive steps.

$f(arr, \overset{0}{low}, \overset{7}{high})$

↓
 $f(arr, 4, 7)$

↓
 $f(arr, 6, 7)$

RECURSIVE

[3, 4, 6, 7, 9, 12, 16, 17]

Target = 13.

f(arr, low, high, target)

{

if (low > high)

return -1;

mid = $\frac{(low + high)}{2}$

if (a[mid] == target)

return mid

else if (target > a[mid])

return f(arr, mid+1, high, target)

else

return f(arr, low, mid-1, target)

}

ultimately this will return -1.

[3, 4, 6, 7, 9, 12, 16, 17]

0 7 0 1 2 3 4 5 6 7

f(arr, low, high, target)

```

{
  if (low > high)
    return -1;
  mid = (low + high) / 2; (7)
  if (a[mid] == target)
    return mid;
  else if (target > a[mid])
    return f(arr, mid + 1, high, target);
  else
    return f(arr, low, mid - 1, target);
}

```

f(arr, low, high, target)

```

{
  if (low > high)
    return -1;
  mid = (low + high) / 2; (12)
  if (a[mid] == target)
    return mid;
  else if (target > a[mid])
    return f(arr, mid + 1, high, target);
  else
    return f(arr, low, mid - 1, target);
}

```

Index 6 7

f(arr, low, high, target)

```

{
  if (low > high)
    return -1;
  mid = (low + high) / 2; (16)
  if (a[mid] == target)
    return mid;
  else if (target > a[mid])
    return f(arr, mid + 1, high, target);
  else
    return f(arr, low, mid - 1, target);
}

```

f(arr, low, high, target)

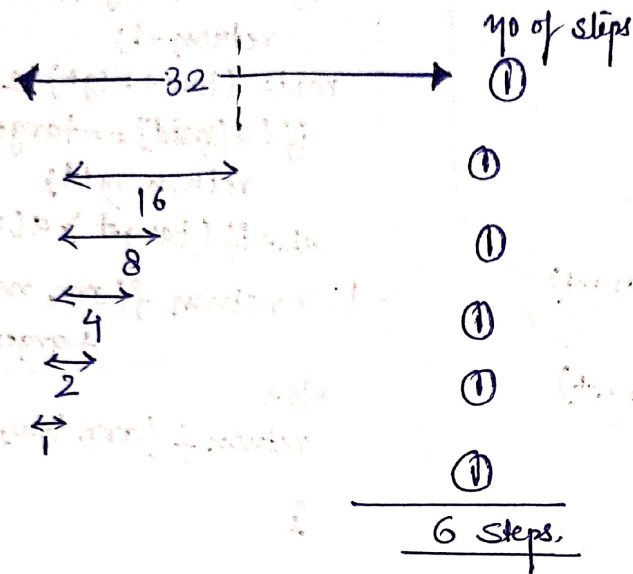
```

{
  if (low > high)
    return -1;
  mid = (low + high) / 2;
  if (a[mid] == target)
    return mid;
  else if (target > a[mid])
    return f(arr, mid + 1, high, target);
  else
    return f(arr, low, mid - 1, target);
}

```

Base Case executed.

TIME COMPLEXITY OF BINARY SEARCH.



$$\text{for } 32 = 2^5$$

$$64 = 2^6$$

\therefore So Time Complexity is somewhat near about $O(\log_2 n)$.

* Overflow case of Binary Search.



$$\boxed{\text{mid} = \frac{\text{low} + \text{high}}{2}}$$

if we keep on reducing, so low and high both are low and high = $\frac{\text{INT-MAX} + \text{INT-MAX}}{2}$

So it becomes $2 * \text{INT-MAX}$, thus it can't be stored in the Integers.

Solution of this problem is \rightarrow either Take long long (data type)

\therefore Other Option will be, if you don't want to use long long or long.

$$\rightarrow \boxed{\text{mid} = \text{low} + \frac{\text{high} - \text{low}}{2}} \quad \begin{array}{l} \text{it is same as} \rightarrow \\ \boxed{\text{mid} = \frac{\text{low} + \text{high}}{2}} \end{array}$$

So here if low and high both become equal to INT-MAX then $\frac{(\text{INT-MAX} - \text{INT-MAX})}{2} = 0$.
So it will prevent overflow.