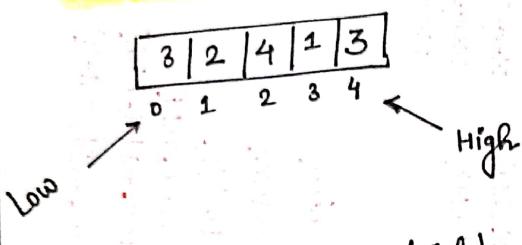


## Data Structure and Algorithms

By

Malay Tripathi

### MERGE SORT



Merge Sort (arr, low, high)

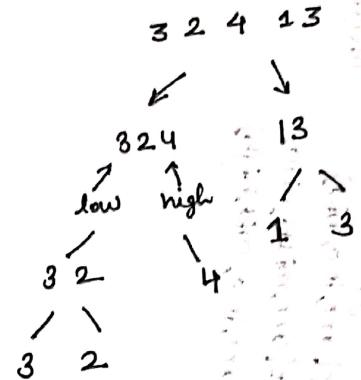
{ if (low >= high) return ; } Base condition.

$$\text{mid} = (\text{low} + \text{high}) / 2$$

merge sort (arr, low, mid)

merge sort (arr, mid+1, high)

merge (arr, low, mid, high)



1	2	3	3	4
3	2	4	1	3

0 1 2 3 4  
2 3 4  
2 3  
mergesort (arr, low, high)

```
{
    if (low >= high) return;
    mid = (low + high) / 2;
    MergeSort (arr, low, mid);
    MergeSort (arr, mid+1, high);
    merge (arr, low, mid, high);
}
```

temp = 1, 1, 2, 2, 3, 4, 4, 5, 6.

```
* merge (arr, low, mid, high)
temp [] temp []
left = low;
right = mid + 1;
while (left <= mid) && right <= high)
{
    if (arr [left] <= arr [right])
        temp.add (arr [left]);
    left++;
    else
        temp.add (arr [right]);
    right++;
}
```

## RECURSION OF MERGESORT

mergesort (arr, low, high)

```
{
    if (low >= high) return;
```

```
    mid = (low + high) / 2;
```

MergeSort (arr, low, mid)

MergeSort (arr, mid+1, high)

Merge (arr, low, mid, high)



mergesort (arr, low, high)

```
{
    if (low >= high) return;
```

```
    mid = (low + high) / 2;
```

mergesort (arr, low, mid)

mergesort (arr, mid+1, high)

Merge (arr, low, mid, high)



mergesort (arr, low, high)

```
{
    if (low >= high) return;
```

```
    mid = (low + high) / 2;
```

mergesort (arr, low, mid)

mergesort (arr, mid+1, high)

Merge (arr, low, mid, high)

mergesort (arr, low, high)

```
{
    if (low >= high) return;
```

```
    mid = (low + high) / 2;
```

MergeSort (arr, low, mid)

MergeSort (arr, mid+1, high)

Merge (arr, low, mid, high)



mergesort (arr, low, high)

```
{
    if (low >= high) return;
```

```
    mid = (low + high) / 2;
```

MergeSort (arr, low, mid)

MergeSort (arr, mid+1, high)

Merge (arr, low, mid, high)

mergesort (arr, low, high)

```
{
    if (low >= high) return;
```

```
    mid = (low + high) / 2;
```

MergeSort (arr, low, mid)

MergeSort (arr, mid+1, high)

Merge (arr, low, mid, high)

mergesort (arr, low, high)

```
{
    if (low >= high) return;
```

```
    mid = (low + high) / 2;
```

MergeSort (arr, low, mid)

MergeSort (arr, mid+1, high)

Merge (arr, low, mid, high)

mergesort (arr, low, high)

```
{
    if (low >= high) return;
```

```
    mid = (low + high) / 2;
```

MergeSort (arr, low, mid)

MergeSort (arr, mid+1, high)

Merge (arr, low, mid, high)

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(7)

MULTIPLE RECURSION CALLS

$f()$   
{  
   $f()$   
  {  
     $f()$  } 2 Times.  
 $f()$   
}

different multiple recursion calls

FIBONACCI SERIES.

0 1 1 2 3 5 8 13 21 34  
1st 2nd 3rd 4th ←----- numbers of fibonacci.

Write Recursive function to tell you the  $N^{th}$  Fibonacci Number.

$$f(3) = 2$$

$$f(4) = 3$$

$$f(5) = f(4) + f(3)$$

SIMPLE.

$$f[0] = 0$$

$$f[1] = 1$$

for ( $i=2 \rightarrow n$ )

$$f[i] = f[i-1] + f[i-2]$$

6.1.1

6.1.2

6.1.3

6.1.4

6.1.5

6.1.6

6.1.7

6.1.8

6.1.9

6.1.10

6.1.11

6.1.12

6.1.13

6.1.14

6.1.15

6.1.16

6.1.17

6.1.18

6.1.19

6.1.20

6.1.21

6.1.22

6.1.23

6.1.24

OKEN

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Now WITH THE HELP OF RECURSION.

```
f(n)
{
    if (n<=1)
        return n;
    return f(n-1)+f(n-2);
}
```

```
f(n)
{
    if (n<=1)
        return n;
    waiting for this line
    last = f(n-1);
    secondlast = f(n-2);
    return last+secondlast;
}

main
{
    n=4
    print (f(4));
}
```

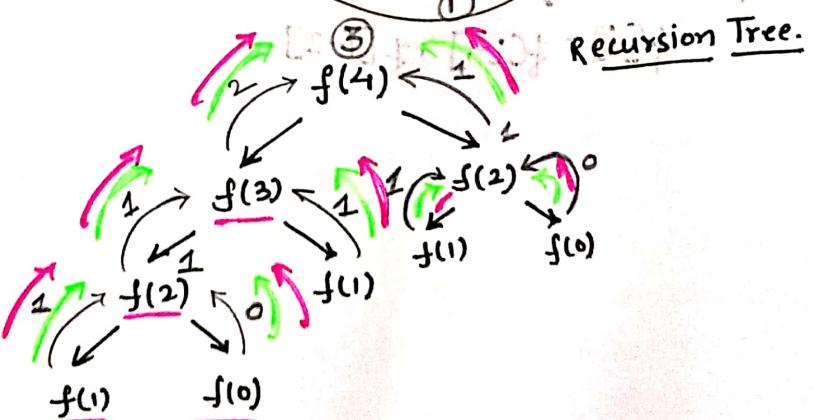
Diagram illustrating the execution flow:

```
f(n)
{
    if (n<=1)
        return n;
    last = f(n-1);
    slast = f(n-2);
    return last+slast;
}
main
{
    n=4
    print (f(4));
}
```

The diagram shows the state of variables at each step:

- Step 1:  $n=4$ ,  $last=f(3)$ ,  $slast=f(2)$ .
- Step 2:  $n=3$ ,  $last=f(2)$ ,  $slast=f(1)$ .
- Step 3:  $n=2$ ,  $last=f(1)$ ,  $slast=f(0)$ .
- Step 4:  $n=1$ ,  $last=f(0)$ ,  $slast=f(0)$ .

final output = 3.



$Tc = (2^n)$  approx (exponential)  $\rightarrow$  as you are calling 2 func for each function.



# Data Structure and Algorithms

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Malay Tripathi

⑨

```

o 1 2
arr [3, 1, 2]
o [ ]
f(ind, [ ])
{
    if (ind >= n)
        print([ ])
        return;
    [3]. add(arr[i]);
    f(ind+1, [ ]);
    [ ]. remove(arr[i]);
    f(ind+1, [ ]);
}
main
{
    arr → [312];
    f(0, [ ]);
}

```

```

→ f(1, [3])
{
    if (ind >= n)
        print([ ])
        return;
    [3]. add(arr[1]);
    f(2, [31])
}
[31]. remove(arr[1])
→ f(2, [31])
{
    if (ind >= n)
        print([ ])
        return;
    [31]. add(arr[2]);
    f(3, [312])
}
[312]. remove(arr[2])
→ f(3, [312])
{
    if (ind >= n)
        print([ ])
        return;
}
---:
3
f(2, [3])
{
    if (ind >= n)
        print([ ])
        return;
    [3]. add(arr[2]);
    f(3, [3, 2])
}
[32]. remove(arr[2])
f(3, [3, 2]).
```

o/p      312

31

32

3

1 2

1

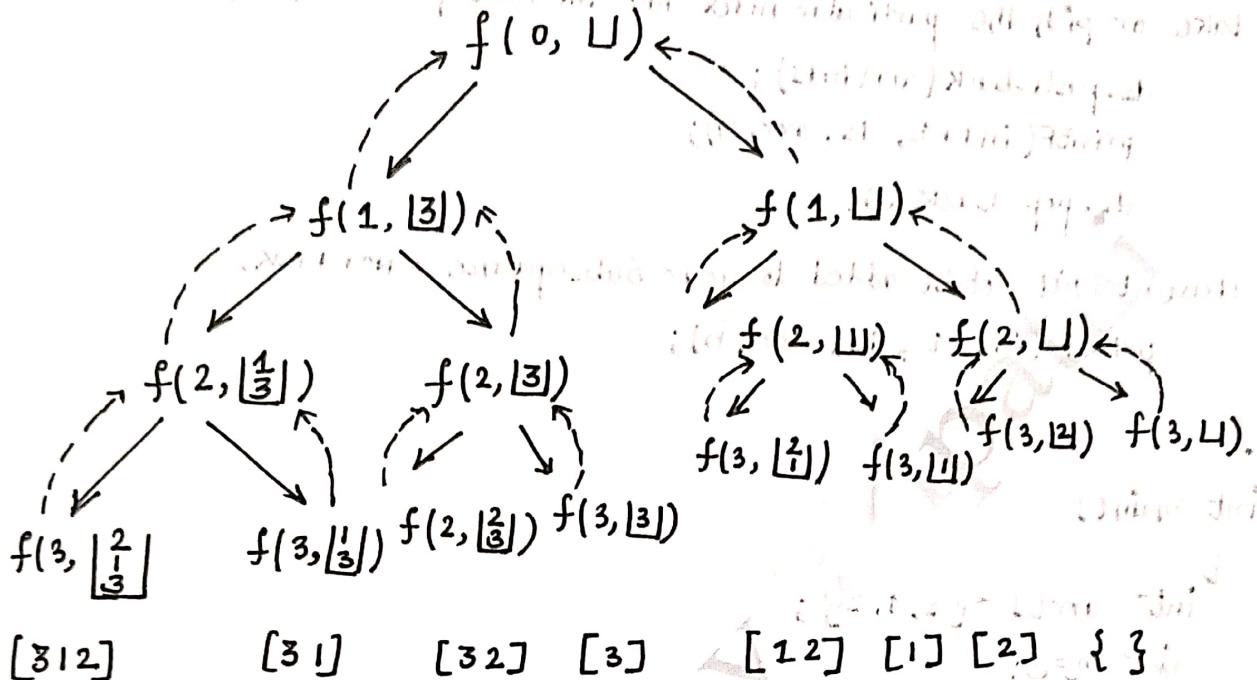
2

{ }

RECURSION TREE

arr = [3 1 2]

## RECURSION ON SUBSEQUENCE.



```
#include <bits/stdc++.h>
using namespace std;
void printf(int ind, vector<int>&ds, int arr[], int n)
{
    if(ind == n)
    {
        for(auto it : ds)
        {
            cout << it << " ";
        }
        if(ds.size() == 0)
        {
            cout << "{}" << endl;
        }
        cout << endl;
        return;
    }
}
```

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// take or pick the particular index into the Subsequence = PICK

```
ds.push-back(arr[ind]);  
printf(ind+1, ds, arr, n);  
ds.pop-back();
```

// element will not be added to your Subsequence = NOT PICK.

```
printf(ind+1, ds, arr, n);
```

int main()

```
{  
    int arr[] = {3, 1, 2};  
    int n = 3;  
    vector<int> ds;  
    printf(0, ds, arr, n);
```

\* TIME COMPLEXITY - for every index there is couple of answer/options

$$O(2^n) * n$$

\* AUXILIARY STACK SPACE = Not more than 3 Recursion  
SPACE COMPLEXITY call waiting at the stack  
space. because the Depth of the  
Recursion is THREE.

$$O(N).$$



# Data Structure and Algorithms

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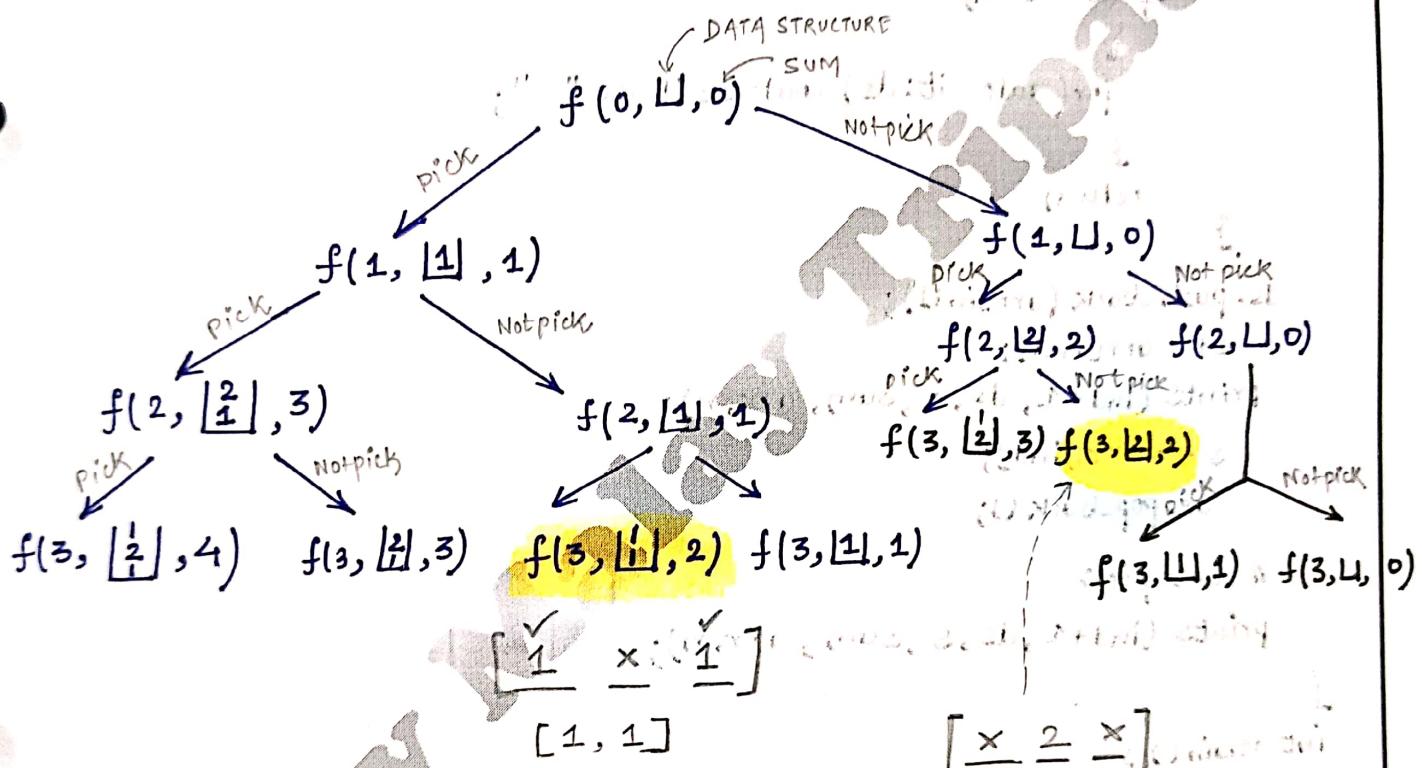
## PRINTING SUBSEQUENCE WHOSE SUM IS K

→ Technique of Take/Not Take.

arr → [1, 2, 1] sum = 2

[1, 1]

[2]



Pseudo code:-

```
f(i, s)
{
    if(i == n)
    {
        if(s == sum)
            print(ds)
    }
    return;
}
```

```
ds.add(arr[i]);
s += arr[i];
f(i+1, ds, s)
```

```
ds.remove(arr[i])
s -= arr[i]
f(i+1, ds, s)
```

PICK  
NOT  
PICK



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O/P = [2 2]  
[2]

Data Structure and Algorithms  
By  
Malay Tripathi

Code

```
#include <bits/stdc++.h>
using namespace std;

void prints(int ind, vector<int> &ds, int s, int sum, int arr[], int n)
{
    if (ind == n)
    {
        if (s == sum)
        {
            for (auto it : ds) cout << it << " ";
        }
        return;
    }
    ds.push_back(arr[ind]);
    prints(ind + 1, ds, s + arr[ind], sum, arr, n);
    s -= arr[ind];
    ds.pop_back();
}

// not pick
prints(ind + 1, ds, s, sum, arr, n);
}

int main()
{
    int arr[] = {1, 2, 1};
    int n = 3;
    int sum = 2;
    vector<int> ds;
    prints(0, ds, 0, sum, arr, n);
    return 0;
}
```