

working with numbers in the system has helped much more self obviously with no problem presented, or

INFIX NOTATION.

→ To add A, B, we write ~~addition plus multiplication~~ $A + B$.

→ To multiply A, B we write ~~addition plus multiplication~~ $A * B$.

→ The operators ('+' or '*') go in between the Operands ('A' and 'B'). This is "Infix" Notation.

PREFIX NOTATION

→ Instead of saying "A plus B", we could say "add A, B" and

write $+AB$

→ "Multiply A, B" would be written $*AB$

Postfix Notation.

→ Another alternative is to put the operators after the operands as in

$AB+$

$(C D E F G H I J K L M N O P Q R S T U V W X Y Z) A B C D E F G H I J K L M N O P Q R S T U V W X Y Z$

and

$AB*$

→ This is Postfix Notation.

The terms infix, prefix and postfix tell us whether the operators go, between, before or after the operands.

Precedence and Associativity of Operators

Operator	Priority	Associativity
\wedge	High	$R \rightarrow L$
$*, /$	Medium	$L \rightarrow R$
$+, -$	Low	$L \rightarrow R$

INFIX TO POSTFIX CONVERSION

• Parathesis Method

- Step 1: Fully parenthesis the expression

- Step 2: Replace right parenthesis (closing parenthesis) by moving operator it.

- Delete left parenthesis (opening parenthesis).

- e.g Given Expression: $A + B * C / D ^ E - F$

$$\wedge ((A + ((B * C) / (D ^ E))) - F)$$

- Moving operator to their right parenthesis

$$((A ((B C * (D E ^) / F) - ;$$

- Delete all left parenthesis

- Note it requires multiple traversals of the expression.

"Here are the steps of the algorithm to convert INFIX to POSTFIX using stack in C".

- Scan all the symbols one by one from left to right in the given Infix Expression.
- If the reading symbol is an operand, then immediately append it to the Postfix Expression.
- If the reading symbol is left Parenthesis '(', thus push it onto stack.
- If the reading symbol is right Parenthesis ')', thus pop all the contents of the stack until the respective left Parenthesis is popped and append each popped symbol to Postfix Expression.
- If the reading symbol is an operator (+, -, *, /), then Push it onto the stack. However, first, pop the operator which are already on the stack that have higher precedence than the current operator and append them to the postfix. If an open parenthesis is there on top of the stack then push the operator into the stack.
- If the input is over, pop all the remaining symbols from the stack and append them to the postfix.

return of mail drivers at multiple sites to make all sort of
arrangements for delivery of packages and parcels.

Still no update yet from our govt. regarding rate the post
and delivery of packages and parcels. We are awaiting their response.

Truckers (truckers) will take care of delivery services and we
are awaiting their response.

Now the stamp duty, TDS, documents and delivery charges will be
increased after 1st Oct.

Stamp duty will increase by 10% and delivery charges will be
increased by 10%.

Delivery charges will increase by 10% and stamp duty will be
increased by 10%.

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Dry Run to Infix to Postfix Conversion Using Stack inc.

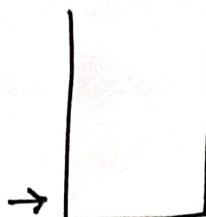
Infix Expression $\rightarrow a * (b + c + d)$

(1) Infix

$a * (b + c + d)$

'a' is an operand
append it to the
postfix

Stack



Postfix

a

$(b + c + d) * \dots$ (n)

(2) $a * (b + c + d)$

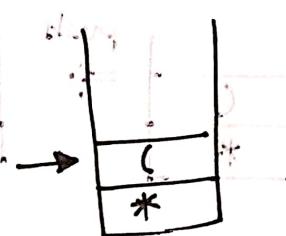
"*" operator push it onto
stack after removing all
higher or equal precedence
operator



a

(3) $a * (b + c + d)$

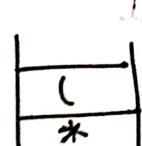
"(" → push left
Parenthesis onto the
stack



Postfix

a

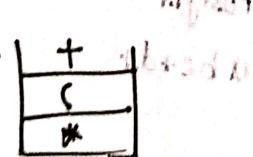
(4) $a * (b + c + d)$



Postfix

ab

(5) $a * (b + c + d)$



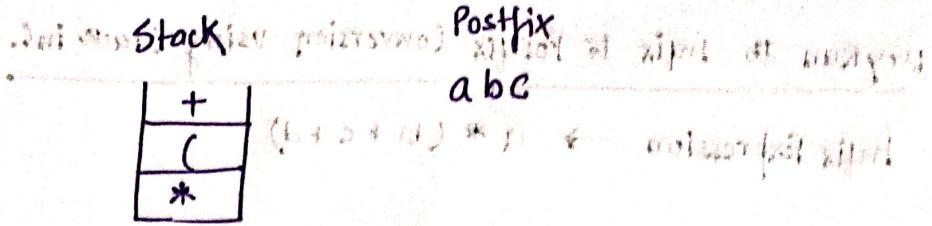
Postfix

ab

$(b + c + d) * \dots$ (n)



Infix
(6) $a * (b + c + d)$



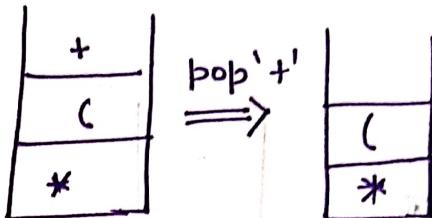
(7) $a * (b + c + d)$



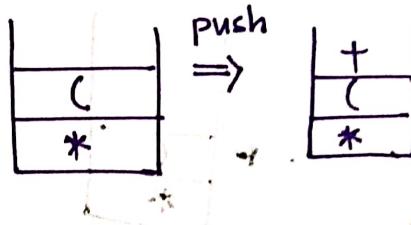
$\Rightarrow '+'$ is an operator,
so push it into stack
after removing higher or
equal priority operators
from top to stack.

\Rightarrow current top of stack is '+', it
has equal precedence with
the input symbol '+'.
So pop it & append to Postfix

\Rightarrow Now the top of stack is '(', so
push the input operator.

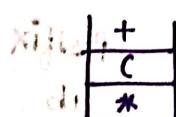


Postfix $\Rightarrow abc +$ (pop ' + ' from stack)



(8) $a * (b + c + d)$

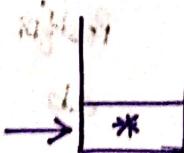
'd' is an operand.
append it to postfix



Postfix
 $abc + d$

(9) $a * (b + c + d)$

') ' → right parenthesis,
follow steps of an algo
& pop all content of stack
until respective.



Postfix
 $abc + d +$

(10) $a * (b + c + d)$

i/p → is over, so pop all remaining



Postfix
 $abc + d + *$
FINAL EXPRESSION.

CONVERT POSTFIX TO INFIX USING A STACK.

1. Initialize an empty stack.
2. Traverse the Postfix expression from left to Right
 - If the symbol is an Operand (eg A,B,I,2) push it onto the stack.
 - If the symbol is an Operator (eg +,-,*,/)
 - pop the top elements from the stack.
let's call them Op1 and Op2.
 - form a new infix expression by enclosing Op1 and Op2 in parenthesis, placing the operator between
 $((Op_1 \text{ Operator } Op_2))$.
 - Push this new expression back onto the stack.
3. At the end of the traversal, the stack will contain a single element, which is the final infix expression.

CONVERT INFIX TO PREFIX EXPRESSION

↳ Convert Infix to Prefix Expression using stacks.

↳ steps to converting infix to Prefix using stacks

1. Reverse the Infix Expression:

→ traverse the expression from right to left.

→ reverse the order of Operators and Operands.

→ Swap opening Parenthesis with closing
parenthesis. and vice-versa.

2. Convert the reversed expression to postfix notation:

→ use the Infix-to-Postfix algorithm (precedence
and associativity rules using a stack).

3. Reverse the postfix expression : to obtain the final
prefix expression.

↳ Example, : $A + B * C$

→ Step-by-Step Solution.

1. Reverse the Infix expression

$C * B + A$

2. Apply Infix to Postfix expression

Postfix

$C * B + A$



C

$C * B + A$



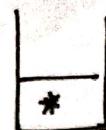
C

$C * B + A$



CB

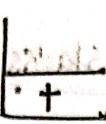
$C * B + A$



$CB *$

Topological Order of Stack Operations

$C * B + A$



$CB *$

$C * B + A$



$CB * A +$

$C * B + A$



$CB * A +$

Postfix of Reversed Expression and $CB * A +$ will demand of
stack and multiply addition will occur.

Reverse the Postfix Expression $\Rightarrow + A * B C$

Infix to Prefix $\Rightarrow + A * B C$

FINAL ANSWER.

Data Structures
Topological Order of Stack Operations
Stack and Stack Operations
Final Answer

Topological Order of Stack Operations
Infix



PREFIX TO INFIX.

1. Initialize an empty stack.
2. Traverse the prefix expression from RIGHT to LEFT.
 - If the symbol is an operand (e.g A, B, 1, 2), push it onto the stack.
 - If the symbol is an operator (e.g +, -, *, /)
 - ↳ Pop the two elements from the stack.
let's call them op1 and op2.
 - ↳ Form a new infix expression by placing op1 and op2 below the operator, enclosed in parenthesis
 $(op1 \text{ operator } op2)$.
 - ↳ Push this new expression back onto the stack.
3. At the end of the Traversal, the stack will contain a single element, which is the final Infix expression.

POSTFIX TO PREFIX USING A STACK.

1. Initialize an empty stack
2. Traverse the postfix expression from left to right
 - If the symbol is an operand (e.g. A, B, 1, 2), push it onto the stack.
 - If the symbol is an operator (e.g. +, -, *, /)
 - pop the top two elements from the stack.
 - pop the top two elements from the stack.
 - create a new expression by concatenating the operator with op1 and op2 in prefix form (operator op1 op2).
 - push this new expression back onto the stack.
3. At the end of the Traversal → the stack will contain a single element, which is the prefix expression.

EXAMPLE

$A \ B \ C \ * \ + \ D \ -$

Step-by-step Solution:-

1. Initialize an empty stack.

2. Traverse each character in the postfix expression:-

• Operand A → Push to Stack: $["A"]$

• Operand B → Push to Stack: $["A", "B"]$

• Operand C → Push to Stack: $["A", "B", "C"]$

• Operator * : $A \ B \ C \ * \ + \ D \ -$
 ↳ pop C and B from the stack

• Operator + : $A \ B \ C \ * \ + \ D \ -$
 ↳ pop *BC and A from the stack

• Operator + : $A \ B \ C \ * \ + \ D \ -$
 ↳ pop *BC and A from the stack
 ↳ from the expression $+A * BC$ and push back on stack

on the stack: stack $["+A * BC"]$

• Operand D :
 ↳ Push to stack: stack $["+A * BC", "D"]$

$A \ B \ C \ * \ + \ D \ -$

• Operator -
 ↳ Pop D and $+A * BC$ from the stack

 ↳ from the expression $- +A * BCD$.

Final Prefix expression.

PREFIX TO POSTFIX.

1. Initialize an empty stack.
2. Traverse the prefix expression from Right to Left.
 - If the symbol is an Operand, push on to stack.
 - If the symbol is an Operator
 - ↳ Pop the top two elements from the stack
 - ↳ form new postfix expression by concatenating op₁ and op₂ with the operator at the end
(op₁ and op₂ operator)
 - ↳ Push this new expression back onto stack
3. At the end of the Traversal, the stack will contain a single element, which is the final postfix expression.