

Data Structure and Algorithm

By

Malay Tripathi

SPARSE MATRIX

A Sparse Matrix \rightarrow is a two dimensional array where the majority of the elements have the value null. for ex \rightarrow

$$\begin{bmatrix} - & x & - & - & - & - \\ x & - & - & - & - & - \\ - & - & x & - & - & - \\ - & - & - & x & - & - \end{bmatrix}$$

Where 'x' = represents NON-NULL VALUES.

Memory Representation of lower triangular

$$\begin{bmatrix} a_{11} & & & \\ a_{21} & a_{22} & & \\ a_{31} & a_{32} & a_{33} & \\ \vdots & \vdots & \vdots & \\ a_{n1} & a_{n2} & a_{n3} & \dots a_{nn} \end{bmatrix}_{n \times n}$$

Row-major order.

According to row major order, the address of any element a_{ij} , $1 \leq i, j \leq n$, can be obtained as \rightarrow

Address (a_{ij}) = No. of elements upto the a_{ij} elements

= Total no. of elements in the first $i-1$ rows

+ no. of elements upto j th column in i th row

$$= 1 + 2 + 3 + \dots + (i-1) + j$$

$$= \frac{i(i-1)}{2} + j$$

$$\text{Address } (a_{ij}) = M + \frac{i(i-1)}{2} + j - 1.$$

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Column Major Order.

According to column-major order, the address of any element a_{ij} , $1 \leq i, j \leq n$ can be obtained as \rightarrow

Address $(a_{ij}) = \text{No. of elements upto the } a_{ij} \text{ elements}$
 $= \text{Total number of elements in the first } j-1 \text{ columns}$
 $+ \text{number of elements upto the } i^{\text{th}} \text{ row in the } j^{\text{th}}$

$$= [n + (n-1) + (n-2) + \dots + (n-j+2)] + (i-j+1)$$

$$= \{ n \times [j-1] - [1+2+3+\dots+(j-2)+(j-1)] + i \}$$

$$\Rightarrow n \times (j-1) - \frac{j(j-1)}{2} + i$$

$$\Rightarrow (j-1) \left[n - \frac{j}{2} \right] + i$$

If the starting location of the first element (that is, of a_{11}) is M then the address of a_{ij} , $1 \leq i, j \leq n$ will be

$$\text{Address}(a_{ij}) = M + (j-1) \times \left(n - \frac{j}{2} \right) + i - 1$$

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Memory Representation of an Upper-Triangular

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & \dots & a_{1n} \\ & a_{22} & a_{23} & \dots & a_{2n} \\ & & a_{33} & \dots & a_{3n} \\ & & & \dots & a_{nn} \end{bmatrix}_{n \times n}$$

(a). Row-Major Order

Acc. to Row-Major Order, the address of any element a_{ij} , $1 \leq i, j \leq n$ can be obtained as \Rightarrow

Address (a_{ij}) = Number of elements upto the a_{ij} element
= Total no. of elements in the first $(i-1)$ rows +
Number of elements upto the j^{th} column in the i^{th} row

$$= n + (n-1) + (n-2) + (n-3) + \dots + (j-i+1)$$

$$= n + (n-1) + (n-2) + \dots + (n-i+2) + (j-i+1)$$

$$= n \times (i-1) - [1+2+3+\dots+(i-2)+(i-1)] + j$$

$$= n \times (i-1) - \frac{i(i-1)}{2} + i$$

$$= (i-1) \times \left(n - \frac{i}{2}\right) + j$$

$$\boxed{\text{Address } (a_{ij}) = M + (i-1) \times \left(n - \frac{i}{2}\right) + j - 1}$$

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Column Major Order.

Acc. to column major order, the address of any element a_{ij} , $1 \leq i, j \leq n$ can be obtained as \rightarrow

Address (a_{ij}) = No. of elements up to the a_{ij} elements
= Total no. of elements in the first $(j-1)$ columns
+ no. of elements up to the i th row in the j th column

$$= [1 + 2 + 3 + \dots + (j-1)] + i$$

$$= \frac{j(j-1)}{2} + i$$

If the starting location of the first element, i.e. of a_{11} is M , then the address of a_{ij} , $1 \leq i, j \leq n$ will be

$$\text{Address } (a_{ij}) = M + \frac{j(j-1)}{2} + i - 1$$

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MEMORY REPRESENTATION OF A TRIDIAGONAL MATRIX

In sparse matrix having the elements only on the diagonal, the following points are evident:

Number of elements in an $n \times n$ square matrix = n

Any elements a_{ij} can be retrieved in memory using the formula

$$\text{Address}(a_{ij}) = i[\text{or } j]$$

$$\begin{bmatrix} a_{11} & a_{12} & & & \\ a_{21} & a_{22} & a_{23} & & \\ & a_{32} & a_{33} & a_{34} & \\ & & a_{43} & a_{44} & a_{45} \\ & & & \vdots & \\ & & & a_{(n-1)(n-2)} & a_{(n-1)(n-1)} & a_{(n-1)n} \\ & & & & a_{n(n-1)} & a_{nn} \end{bmatrix}$$

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Row Major Order.

The address of the element, a_{ij} , $1 \leq i, j \leq n$ can be obtained as:-

Address (a_{ij}) = Number of elements up to the a_{ij} element

= Total number of elements in the first $(i-1)$ rows
+ number of elements up to the j^{th} column in the
 i^{th} row

$$= \{ 2 + [3 + 3 + 3 + \dots + \text{upto } (i-2) \text{ terms}] \} + (j-i+2)$$

$$= 2 + 3(i-2) + j - (i-2)$$

$$= 2 + 2(i-2) + j \Rightarrow 2 + 2(i-2) + j$$

If the starting location of the first element, i.e. of a_{11} is M ,
then the address of a_{ij} , $1 \leq i, j \leq n$ will be

$$\text{Address } (a_{ij}) = M + 2(i-2) + j + 1.$$

Column Major Order

According to column major order, the address of any element a_{ij} , $1 \leq i, j \leq n$

Address (a_{ij}) = Number of elements upto the a_{ij} element

= Total no. of elements in the first $(j-1)$ columns
+ number of elements upto the i^{th} row in the j^{th} column.

$$= \{ 2 + [3 + 3 + \dots + \text{upto } (j-2) \text{ terms}] \} + (i-j+2)$$

$$= 2 + (j-2) \times 3 + i - (j-2)$$

$$= 2 + 2*(j-2) + i$$

If the starting location of the first element, i.e. of a_{11} is M , then the
address of a_{ij} , $1 \leq i, j \leq n$ will be

$$\text{Address } (a_{ij}) = M + 2*(j-2) + i + 1.$$