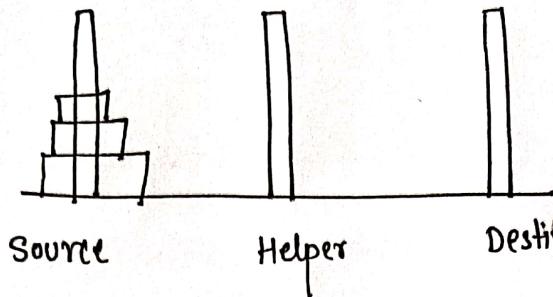


TOWER OF HANOI PROBLEM.



Good Job

(Tower of Hanoi)

of hanoi

(Solve Tower of Hanoi)

Steps 1. The tower of hanoi problem is to find a sequence of disk moves so that all the disks are moved from stand 1 (source) to stand 3 (Destination), adhering to the following rules.

1. Move only one disk at a time.
2. A larger disk cannot be placed on top of a smaller disk.
3. All disks except the one being moved should be on a stand.

Recurrence Relation for solving the Tower of Hanoi problem can be written as :-

$$\text{Tower of Hanoi (disks)} = \begin{cases} \text{move the disk} & \text{if } \text{disks} = 1 \\ \text{Tower of Hanoi (disks-1)} & \text{if } \text{disks} > 1 \end{cases}$$

```
#include < >  
void move (int, int, int, int);
```

```
void main()
```

```
{  
    int disks = 3;
```

```
    printf ("follow these moves:\n");
```

```
    move (disks, 1, 3, 2)
```

```
    move (count - 1, start, temp, finish);
```

```
    move (count - 1, temp, finish, start);
```

```
    if (count > 0)
```

```
        move (count - 1, start, temp, finish);
```

```
        move (count - 1, temp, finish, start);
```

```
        move (count - 1, start, temp, finish);
```

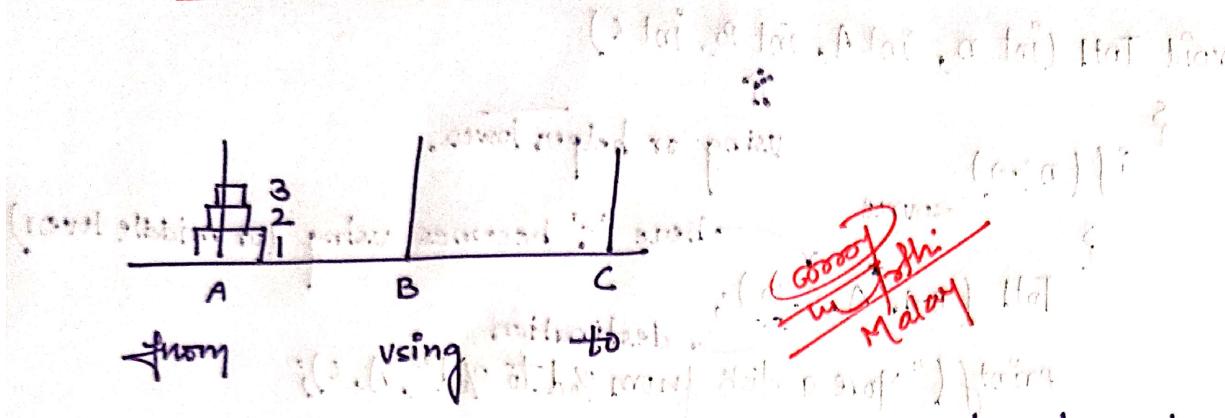
```
        move (count - 1, temp, finish, start);
```

```
}
```

```
} // main ()
```

```
move (count - 1, start, temp, finish);
```

Solution for 3 Disks.



- Move 2 Disks from A to B using C.
- Move a Disk from A to C \rightarrow
- Move 2 Disks from B to C using A.

Now if there are n Disks

- Move $n-1$ disks from A to B using C
- Move a disk from A to C.
- Move $n-1$ disks from B to C using A.

Recall example of and our last comment on base case
at the bottom of previous page note

Algorithm.

```
void ToH (int n, int A, int B, int C)
```

↑

using or helper towers.

{ if ($n > 0$)

{ source

ToH ($n-1, A, C, B$);

here 'C' becomes destination

printf ("Move a disk from %d to %d", A, C);

ToH ($n-1, B, A, C$); here 'A' becomes destination

here 'B' becomes destination

~~Copy this
not today.~~

Source

no. of disk

no. of towers

main & void

ToH ($3, 1, 2, 3$)

call → $A \leftarrow B$
ToH ($2, 1, 3, 2$)

$B \leftarrow C$
 $A \leftarrow C$
 $A \leftarrow B$ ← call
ToH ($2, 2, 1, 3$)

ToH ($1, 1, 2, 3$)

$A \leftarrow C$
 $B \leftarrow C$

1 → 3

ToH ($1, 2, 3, 1$)

ToH ($1, 1, 2, 3$)

2 → 3

$A \leftarrow C$

$B \leftarrow C$

1 → 3

$A \rightarrow D$

2 → 3

3 → 2

$S \rightarrow D$

2 → 1

$S \rightarrow D$

1 → 3

$S \rightarrow D$

1 → 3

1 → 2

3 → 2

1 → 3

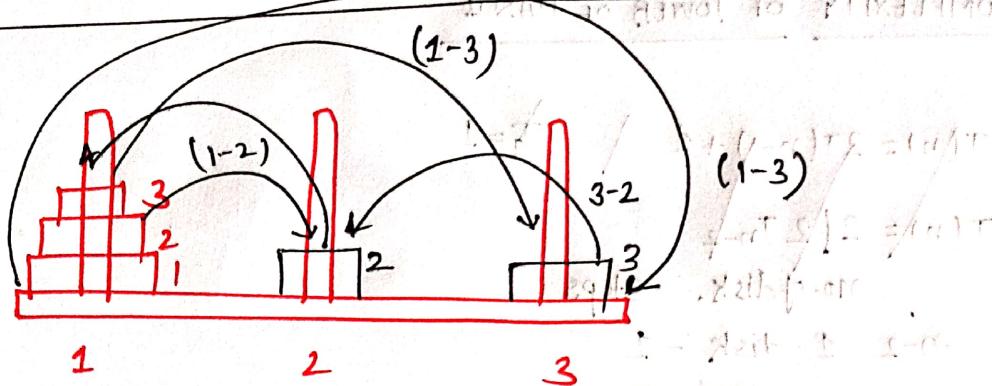
2 → 1

2 → 3

1 → 3

These are moves that we have to take to move disks from Source to destination.

$$(1-3) \rightarrow (1-2) \rightarrow (3-2) \rightarrow (1-3) \rightarrow (2-1) \rightarrow (2-3) \rightarrow (1-3)$$



check for this order towers of hanoi.

$$[1 + 1 \cdot n^T S = (n)T]$$

$$1 + (1 \cdot n) T S = (n)T$$

$$1 + [(1 \cdot n) + S] T = (n)T$$

$$1 + (C \cdot n) T^2 S =$$

$$1 + [1 + (C \cdot n) T^2 S] T =$$

$$1 + (C \cdot n) T^3 S =$$

$$[1 + N_S + (N \cdot n) T^N S = (n)T]$$

$$1 + N_S + (N \cdot n) T^N S = (n)T$$

$$1 + N_S + (N \cdot n) T^N S =$$

$$[1^N S + (n)T^N S = (n)T]$$

TIME COMPLEXITY OF TOWER OF HANOI

$$T(n) = 2T(n-1) + 1$$

$\uparrow n=1$

$$T(n) = 2[2T(n-2) + 1] + 1$$

\downarrow
no. of disk. steps

$$n-2 \quad 1 \text{ disk} - 1$$

$$n-1 = 2 \text{ disk} - 3$$

$$n = 3 \text{ disk} - 7$$

So the Recurrence Relation on the basis of Recurrence Relation

$$T(n) = 2T(n-1) + 1$$

Recurrence Relation

~~cost of
moving
disk~~

$$T(n) = 2T(n-1) + 1$$

$$T(n) = 2[2T(n-2) + 1] + 1$$

$$= 2^2 T(n-2) + 3$$

$$= 2^2 [2T(n-3) + 1] + 3$$

$$= 2^3 T(n-3) + 7$$

or K number of disk

$$T(n) = 2^K T(n-K) + 2^K - 1$$

Calculate : \rightarrow total number of steps for N disks.

$$T(n) = 2^K T(n-K) + 2^K - 1$$

if $K \rightarrow n$

$$= 2^n T(n-n) + 2^n - 1$$

$$\Rightarrow T(n) = 2^n T(0) + 2^n - 1$$

Another Method to calculate

Recurrence relation

$$T = CT + T$$

$$T = CT + T$$

$$T = CT + T$$

ANALYSIS OF RECURSION.

$$T(n) = 2T(n-1) + 1 \rightarrow ①$$

Good
method
Malay.

Solving it using Back-Substitution:

$$T(n-1) = 2T(n-2) + 1 \rightarrow ②$$

$$T(n-2) = 2T(n-3) + 1 \rightarrow ③$$

($T(n-1)$ = $2T(n-2) + 1$) \Rightarrow $T(n-2) = \frac{T(n-1) - 1}{2}$ \therefore $T(n-2)$ = $T(n-1) - 1$

Put the value of $T(n-2)$ in the equation 2 with the help
of equation 3.

$$T(n-1) = 2[2T(n-3) + 1] + 1 \rightarrow ④$$

Now put the value $T(n-1)$ in equation 1 with the help of
equation 4.

$$\begin{aligned} T(n) &= 2[2[2T(n-3) + 1] + 1] + 1 \\ &= [4[2T(n-3) + 1] + 2 + 1] \\ &\Rightarrow [4 \cdot 2T(n-3) + 4 + 2 + 1] \\ &= 2^3 \cdot T(n-3) + 2^2 + 2^1 + 2^0 \end{aligned}$$

Good
method
Malay.

So if there is k disks

$$= 2^k \cdot T(n-k) + 2^{k-1} + 2^{k-2} + \dots + 2^2 + 2^1 + 1$$

Base Condition.

$$T(1) = 1$$

$$n-k = 1$$

$$k = n-1$$

$$T(n) = 2^{n-1} T(1) + \dots + 2^{(n-2)} + \dots + 2^2 + 2 + 1$$

$$(1) \leftarrow 1 + (n-1) T \Rightarrow (n) T$$

$$(2) \leftarrow 1 + (n-1) T \Rightarrow (1-n) T$$

it is a GP Series, and the sum is $2^n - 1$

$$(3) \leftarrow 1 + (n-1) T \Rightarrow (2^n - 1) T$$

$$T(n) = O(2^n - 1), \text{ or its Time Complexity} = O(2^n)$$

which is exponential.

~~unathi~~ ~~Malay~~ ~~Exponentially~~

$$(4) \leftarrow 1 + [1 + (n-1) T]^n \Rightarrow (1+n) T$$

THE SPACE COMPLEXITY OF THE TOWER OF HANOI

$O(N) \rightarrow$ where n is number of disks.

This is because for each recursion, the

disks take up $N-1$ recursive stack space.

$$\dots + [1 + (n-1) T] \cdot P \Rightarrow$$

$$[1 + 1 + [1 + (n-1) T] \cdot P] \cdot P \Rightarrow$$

$$[2 + 1 + [1 + (n-1) T] \cdot P] \cdot P \Rightarrow$$

~~stack~~

~~disk~~ ~~of stack~~ ~~of disk~~

$$[2 + 2 + [1 + (n-1) T] \cdot P] \cdot P \Rightarrow$$

