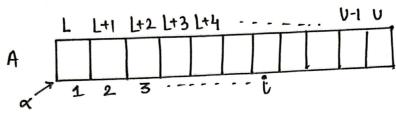
ADDRESS CALCULATION IN 1-D ARRAY

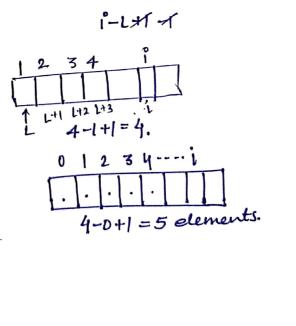


- → A[L:V]
 - → How many elements = V-L+1
- → Address of A[1] = d A[2] = 0+1 A(3) = d+2 A [u] = 0 + 3 A [i] = <+ (i-1)
- · first Assumption Index starts as 1.
- . Second Assumption → every element is of 1 bytes. or requires 1 byte.

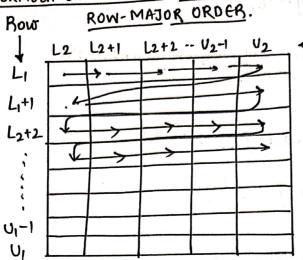
Now suppose every element requires n bytes.

A[1]=
$$q$$

A[2] = $x+1\cdot n$
A[3] = $x+2\cdot n$.
A[4] = $x+3\cdot n$.
:
A[i] = $x+(i-1)n$
A[i] = $x+(i-1)n$
A[i] = $x+(i-(1-1)\pi i)n$
A[i] = $x+(i-(1-1)\pi i)n$
Base Address.



INDEX FORMULA COMPUTATION - 20 ARRAY.



- Column.

A[3][4]

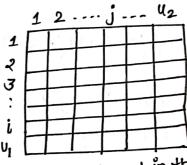
Column. [L2: U2]

Assumption 1). 1 Byte storage per element.

A[1: V1. 1: U2]

-> immediately after U2, (L1+1: L2) will get stored.

$$A[1, U_2] = d + (U_2^{-1})$$



 $A[2,1] = \alpha + (U_2-1) + 1 = \alpha + U_2 = \alpha + 1 \cdot U_2$ $A[2,2] = \alpha + (U_2-1) + 1 = \alpha + U_2 = \alpha + 1 \cdot U_2$

$$A[2,2] = x + U_2+1.$$

. U_2 elements of first crows. $A[3,1] = classed + U_2 + U_2 = classed + 2 \cdot U_2$.

$$A (1,2) = x + (1-1) 1 + 1$$

$$A[1,3] = x + (1-1) L_2 + 2$$

$$A(i,j) = \alpha + (i-1)\mu_2 + (j-1)$$

Ati,j] =
$$\alpha + [(i-1)U_2 + (j-1)] * \gamma$$
.

 $U_2 - Show no. of elements present in the first column. 50 N$

$$A[i,j] = d + [(i-L_1+1-1)(U_2-L_2+1) + (j-L_2+1-1)] \times n$$

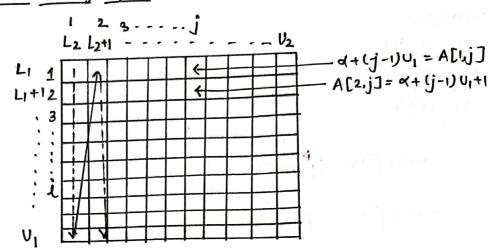
$$A[i,j] = d + [(i-L_1)(U_2-L_2+1) + (j-L_2)] n.$$

A L 37(6)=
$$1430 + [(3-(-4))(0-3+1)+(6-3)]4$$

= $[430 + [7*6+3] 4$
= $[430 + [45*4]]$

2-D ARRAY ADDRESS CALCULATION

COLUMN MAJOR ORDER



$$A[1,1] = 0$$

 $A[2,1] = 0 + 1$
 $A[3,1] = 0 + 2$
 $A[0_1,1] = 0 + (0_1-1)$

$$A[1,2] = d + (u_1-1)+1 = d + u_1$$

$$A[1,3] = d + U_1 + u_1 = d + 2u_1$$

$$A[1,4] = d + U_1 + u_1 + u_1 = d + 3u_1$$

$$A[1,j] = d + (j-1)u_1$$

$$A[2,j] = d + (j-1)u_1 + 1$$

$$A[3,j] = d + (j-1)u_1 + 2$$

A [i + j] = x + (j-1) u, + (i-1)

$$A[i,j] = \alpha + \left[(j-L_2+1-1)(U_1-L_1+1) + (i-L_1+1-1) \right] *$$

```
column major ordes.
         ACI, j) = d+ [f-12) (U1-1+1) + (1-1) mas assuda yazar una
Ques
        B[10][20]
         n=2 bytes
         B[2] [1] = 2140.
         B(5) (4)
             2/40+[14-1)(10)+(5-2)]2
            2140 + 5 30+3 ] 2
            2140+ 66
                                  Address = Base Address + Wx[(1-LBR)+(j-LBC)*M
 Ques - arr[15][20]
                    LBR LBL
     Bars address Arr [1][1] 4000
       address of ar [6][0]=44400 4440
4400 = 4000 + W[(6-1)+(8-1)+15]
      4400 = 4000 + W[5+4*15]
     44400 = 4000+ HOW 110W
      A[m][m]
 aues
                  UBR LBC
       W=4
      Base Address A [I][1] = 1500
        adof A[4][5] = 1608
m=?
           1608 = 1500 +4 ((4-1) + (5-1) m)
           1608 = 1500 + 4 (3+4m) = 1500 + 12 + 16m
              108-12 = 1670
                96 = 16m = 6 = m=6 Answer
```

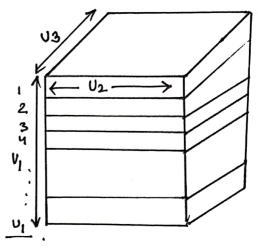
INDEX ARRAY CALCULATION.

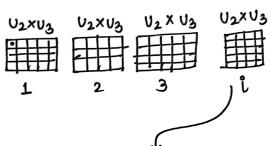
ROW MAJOR ORDER.

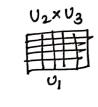
A[LI:UI, L2:U2, L3:U3]

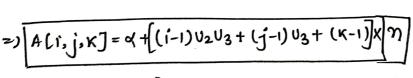
=)
$$A[3,1,1] = 0 + 2(02 \times 03)$$

=) $A[i,j,1] = \alpha + (i-1) U_2 U_3 + (j-1) U_3 - - - A[i,2,1]$







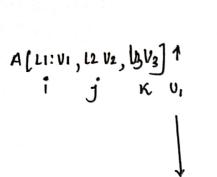


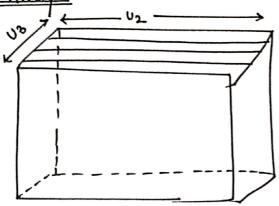
Index formula Computation - 3-0 Array.

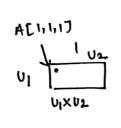
$$A[-1:5, -1:6, 0:10]$$
 $L_1 U_1 L_2 U_2 L_3 U_3$

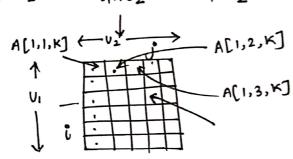
$$= 1000 + \left[2 - (-1)(6 - (-1) + 1)(10 - 0 + 1) + (4 - (-1))(10 - 0 + 1) + (6 - 0)\right] +$$

> Column mayor Order - 3D-Array









$$A[1,1,K] = &+(K-1)U_1U_2$$

$$A[3,j,\kappa] = 4+(\kappa-1)U_1U_2+(j-1)U_1+2$$

$$A[i,j,K] = x + (K-1)U_1U_2 + (j-1)U_1 + (i-1)$$

$$j = J - L_1 + 1$$

 $j = j - L_2 + 1$
 $K = K - L_3 + 1$
 $U_1 = U_1 - L_1 + 1$

$$A(i,j,K) = x + \left[(K-L_3 + 1/2) (U_1-L_1+1) (U_2-L_2+1) + (j-L_2+1-1) (U_1-L_1+1) + (j-L_1+1-1) (U_1-L_1+1-1) + (j-L_1+1-1) (U_1-L_1+1-1) + (j-L_1+1-1) (U_1-L_1+1-1) + (j-L_1+1-1) + (j-L_1+1-1$$