

# Principles of Information Security & Privacy

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ITIS-6200/8200

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# Syllabus ...

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# Chapter 1: An Overview of Computer Security

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- Components of computer security
- Threats
- Policies and mechanisms
- The role of trust
- Assurance
- Operational Issues
- Human Issues



# Basic Components: the CIA

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- Confidentiality
  - ❑ Definition: is the avoidance of the unauthorized disclosure of information
  - ❑ Example
  - ❑ How to achieve this: cryptography, access control
- Integrity
  - ❑ Trustworthiness of data or resources, or, improper or unauthorized modifications can be prevented or detected
  - ❑ Example
  - ❑ How to achieve this: prevention (via authentication and access control) and detection (analyze system events or the data itself)
- Availability
  - Definition: the property that information is accessible and modifiable in a timely fashion by those authorized to do so
  - Denial of Service attacks

# Classes of Threats

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- Threat: a potential violation of security
- Attacks: actions that actually cause a violation
- Disclosure: unauthorized access to information
- Deception: acceptance of false data
- Disruption: interruption or prevention of correct operation
- Usurpation: unauthorized control of some part of a system



# Case Studies of Threats

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- Snooping → passive wiretapping, packet sniffing
  - a form of disclosure
- Phishing: an impersonation of one entity (e.g., legitimate web site) by another (e.g., fake web site)
  - a form of deception
- Other examples on the Web

# Policies and Mechanisms

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- Policy says what is, and is not, allowed
  - This defines “security” for the site/system/*etc.*
- Mechanism: a method, tool, or procedure for enforcing a security policy
- Example: access to your project source files on the lab computers



# Goals of Security

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- Prevention
  - Prevent attackers from violating security policy
- Detection
  - Detect attackers' violation of security policy
- Recovery
  - Stop attack, assess and repair damage
  - Continue to function correctly even if attack succeeds



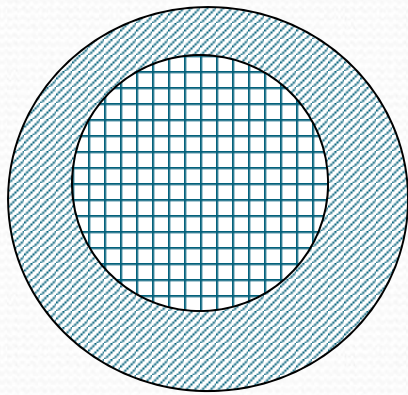
# Trust and Assumptions

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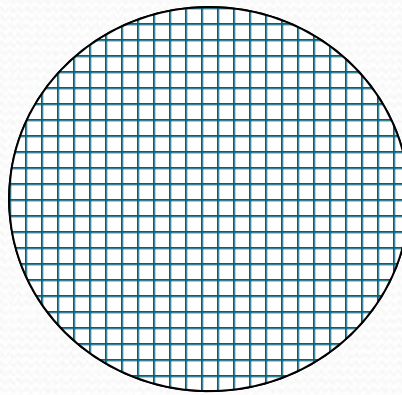
- Underlie *all* aspects of security
- Policies
  - Unambiguously partition system states
  - Correctly capture security requirements
- Mechanisms
  - Assumed to enforce policy, i.e., must be appropriate
    - ✓ For example, using cryptography to ensure that a web site is available won't work
  - Support mechanisms work correctly

# Types of Mechanisms

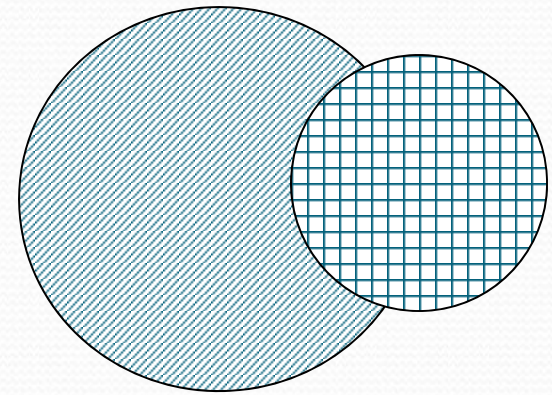
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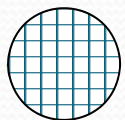
secure



precise



broad



set of reachable states



set of secure states



# Assurance

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- The degree to which the policies meet the requirements of the organizations using the system
- The degree to which the mechanisms correctly implement the policies
- More details in Chapter 17

# Operational Issues

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- Cost-Benefit Analysis
  - Is it cheaper to prevent or recover?



# Operational Issues

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- Cost-Benefit Analysis
  - Is it cheaper to prevent or recover?
- Risk Analysis:
  - Should we protect something?
  - How much should we protect this thing?

Case study: salary database

# Operational Issues

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# Operational Issues

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- Cost-Benefit Analysis
  - Is it cheaper to prevent or recover?
- Risk Analysis:
  - Should we protect something?
  - How much should we protect this thing?
  - A function of the environment and dynamically changing.
- Laws and Customs
  - Are desired security measures illegal? Crypto software
  - Will people do them? DNA samples for authentication, SSN as a password unacceptable

# Human Issues

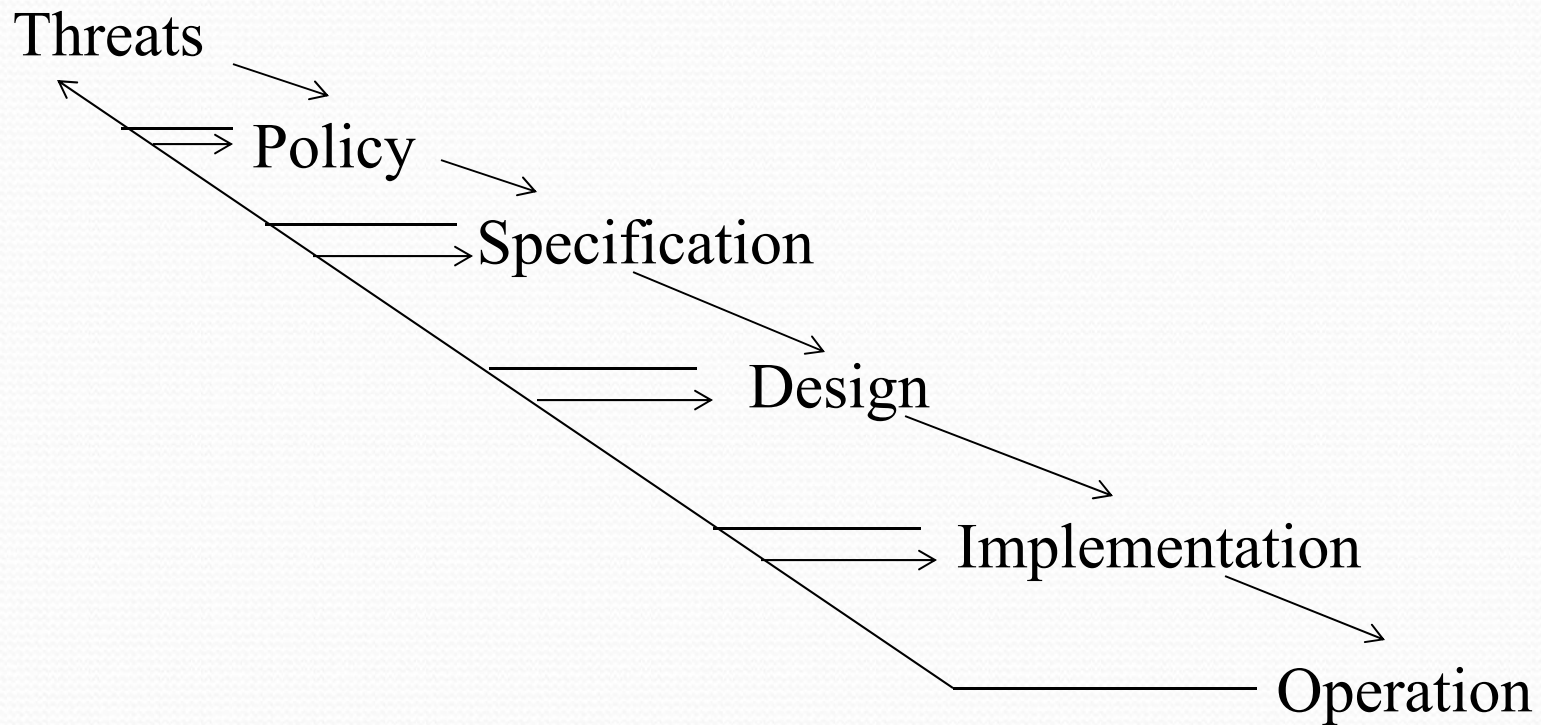
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- Organizational Problems
  - Financial benefits
- People problems
  - Outsiders and insiders (e.g., disgruntled employees)
  - Social engineering



# Tying Together

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# Key Points

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- Policy defines security, and mechanisms enforce security
  - Confidentiality
  - Integrity
  - Availability
- Trust and knowing assumptions
- Importance of assurance
- The human factor



# Chapter 2: Access Control Matrix

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- Overview
- Access Control Matrix Model
- Protection State Transitions
  - Commands
  - Conditional Commands

# Overview

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- Protection state of system
  - Describes current settings, values of system relevant to protection
- Access control matrix
  - Describes protection state precisely
  - Matrix describing rights of subjects
  - State transitions change elements of matrix



# Description

|          |       | objects (entities) |     |       |       |     |       |
|----------|-------|--------------------|-----|-------|-------|-----|-------|
|          |       | $o_1$              | ... | $o_m$ | $s_1$ | ... | $s_n$ |
| subjects | $s_1$ |                    |     |       |       |     |       |
|          | $s_2$ |                    |     |       |       |     |       |
|          | ...   |                    |     |       |       |     |       |
|          | $s_n$ |                    |     |       |       |     |       |

- Subjects  $S = \{ s_1, \dots, s_n \}$
- Objects  $O = \{ o_1, \dots, o_m \}$
- Rights  $R = \{ r_1, \dots, r_k \}$
- Entries  $A[s_i, o_j] \subseteq R$
- $A[s_i, o_j] = \{ r_x, \dots, r_y \}$   
means subject  $s_i$  has  
rights  $r_x, \dots, r_y$  over object  
 $o_j$

The triple  $(S, O, A)$  represents the protection state

# Example

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- Processes  $p, q$
- Files  $f, g$
- Rights  $r, w, x, a, o$

|     | $f$   | $g$  | $p$    | $q$    |
|-----|-------|------|--------|--------|
| $p$ | $rwo$ | $r$  | $rwxo$ | $w$    |
| $q$ | $a$   | $ro$ | $r$    | $rwxo$ |

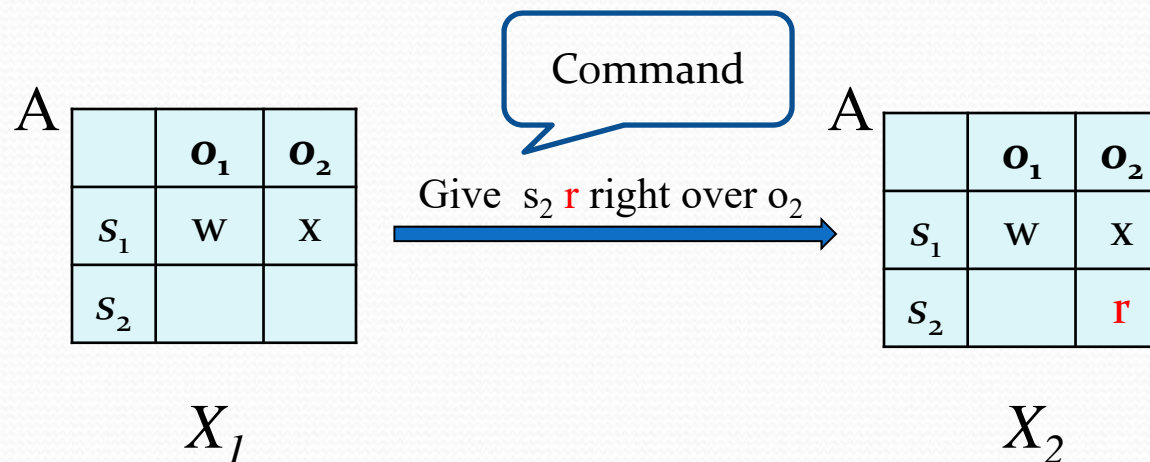


# State Transitions

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- Change the protection state  $X=(S,O,A)$  of system
- $|-$  represents transition
  - $X_i |-\tau X_{i+1}$ : command  $\tau$  moves system from state  $X_i$  to  $X_{i+1}$
  - $X_i |-\ast X_{i+1}$ : a sequence of commands moves system from state  $X_i$  to  $X_{i+1}$
- Commands often called *transformation procedures*
- $X_i=(S_i, O_i, A_i)$

# State Transition (An Example)





# Primitive Operations

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- **create subject  $s$ ; create object  $o$** 
  - Creates new row, column in ACM; creates new column in ACM
- **destroy subject  $s$ ; destroy object  $o$** 
  - Deletes row, column from ACM; deletes column from ACM
- **enter  $r$  into  $A[s, o]$** 
  - Adds  $r$  rights for subject  $s$  over object  $o$
- **delete  $r$  from  $A[s, o]$** 
  - Removes  $r$  rights from subject  $s$  over object  $o$
- ACM: Access Control Matrix

# Creating File

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- Process  $p$  creates file  $f$  with  $r$  and  $w$  permission

```
command create•file( $p, f$ )  
  create object  $f$ ;  
  enter  $own$  into  $A[p, f]$ ;  
  enter  $r$  into  $A[p, f]$ ;  
  enter  $w$  into  $A[p, f]$ ;  
end
```



# Mono-Operational Commands

- Definition: single primitive operation in a command

- Make process  $p$  the owner of file  $g$

```
command makeowner( $p$ ,  $g$ )  
    enter own into  $A[p, g]$ ;  
end
```

# Conditional Commands

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- Let  $p$  give  $q$   $r$  rights over  $f$ , if  $p$  owns  $f$   
**command**  $grant \bullet read \bullet file \bullet 1(p, f, q)$   
    **if**  $own$  **in**  $A[p, f]$   
    **then**  
        **enter**  $r$  **into**  $A[q, f];$   
    **end**
- Mono-conditional command
  - Single condition in this command



# Multiple Conditions

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- Let  $p$  give  $q$   $r$  and  $w$  rights over  $f$ , if  $p$  owns  $f$  and  $p$  has  $c$  rights over  $q$

```
command grant•read•file•2( $p, f, q$ )  
    if own in  $A[p, f]$  and  $c$  in  $A[p, q]$   
    then  
        enter  $r$  into  $A[q, f]$ ;  
        enter  $w$  into  $A[q, f]$ ;  
end
```

# Key Points

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- Access control matrix simplest abstraction mechanism for representing protection state
- Transitions alter protection state
- 6 primitive operations alter matrix
  - Transitions can be expressed as commands composed of these operations and, possibly, conditions





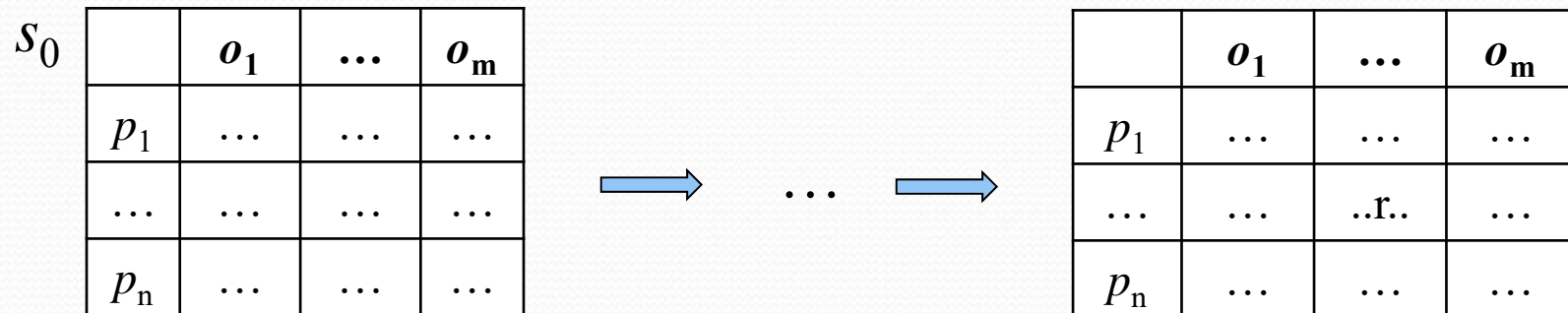
# Chapter 3: Foundational Results

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- Safety Question

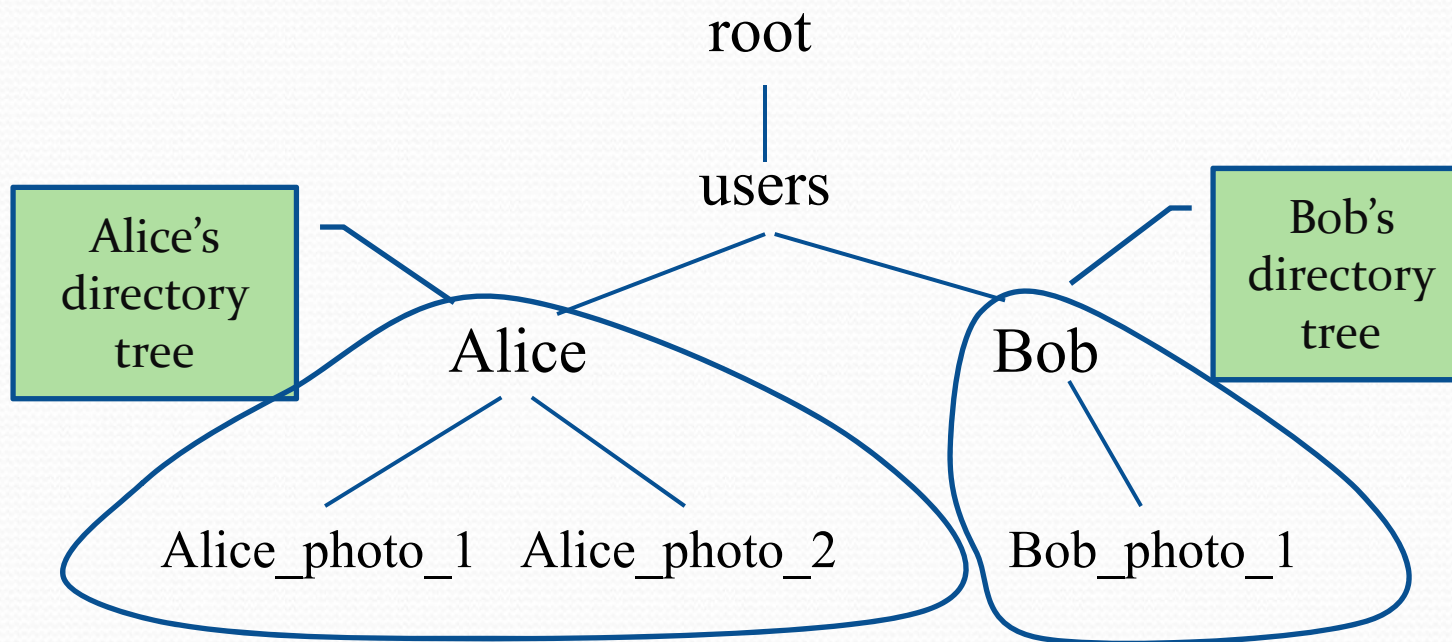
# What Is “Secure”?

- Adding a generic right  $r$  where there was not one is “leaking”
- If a system  $S$ , beginning in initial state  $s_0$ , cannot leak right  $r$ , it is *safe with respect to the right  $r$* .
  - Example: Bob should not **see** Alice’s photo in Facebook, but if he somehow is able to due to an implementation flaw (directory traversal) in the Facebook web server.





# Example Insecure System



- Alice access url: [http://facebook.com/Alice\\_photo\\_1](http://facebook.com/Alice_photo_1)
- Bob access url: [http://facebook.com/../../Alice/Alice\\_photo\\_1](http://facebook.com/../../Alice/Alice_photo_1), violating security policy!

# Safety Question

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- Does there exist an algorithm for determining whether a protection system  $S$  with initial state  $s_0$  is safe with respect to a generic right  $r$ ?
  - Here, “safe” = “secure” for an abstract model
  - No authorized transfer of rights



# Mono-Operational Commands

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- Answer: yes
- Sketch of proof:

Consider minimal sequence of commands  $c_1, \dots, c_k$  to leak the right.

  - Can omit **delete**, **destroy**
  - Can merge all **creates** into one

Worst case: insert every right into every entry; with  $s$  subjects and  $o$  objects initially, and  $n$  rights, upper bound is  $k \leq n(s+1)(o+1)$
- We can show that the length of this sequence is bounded. Therefore, we can enumerate all possible states and determine whether the system is safe.

# General Case

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- Answer: *no*
- Proof idea:
  - Reduce halting problem to safety problem
- If safety question decidable, then represent Turing Machine using protection states and determine if  $q_f$  leaks
  - Implies halting problem decidable
- Conclusion: safety question undecidable



# Key Points

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- Safety problem undecidable
- Limiting scope of systems can make problem decidable