

Chapter 5: Confidentiality Policies

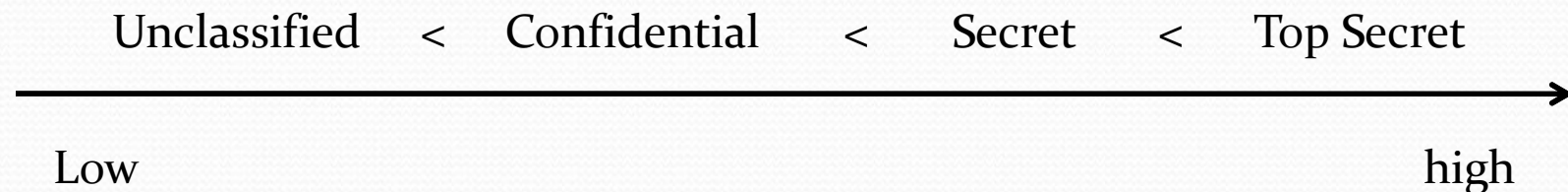
- Overview
 - What is a confidentiality model
- Bell-LaPadula Model
 - General idea
 - Informal description of rules
- A case study (DG/UX)

Confidentiality Policy

- Also called an *information flow policy*
- Goal: prevent the unauthorized *disclosure* of information
 - Deals with information flow
 - Integrity is secondary
 - E.g., military security policy
- Multi-level security models are best-known examples
 - Bell-LaPadula Model basis for many, or most, of these

Bell-LaPadula Model (Preliminary Version)

- Security *levels* arranged in linear ordering, e.g.



- Levels consist of *security clearance*
 - Subjects have *security clearances* $L(s)$
 - Objects have *security classification* $L(o)$

Example

<i>security level</i>	<i>subject</i>	<i>object</i>
Top Secret	Tamara	Personnel Files
Secret	Samuel	E-Mail Files
Confidential	Claire	Activity Logs
Unclassified	Ulaley	Telephone Lists

- Tamara can read all files
- Claire cannot read Personnel or E-Mail Files
- Ulaley can only read Telephone Lists

Reading Information

- Information flows *up*, not *down*
 - “Reads up” disallowed, “reads down” allowed
- Simple Security Condition (Preliminary Version)
 - Subject s can read object o iff $L(o) \leq L(s)$ and s has permission to read o
 - Note: combines mandatory control (relationship of security levels) and discretionary control (the required permission)
 - Sometimes called “no reads up” rule

Writing Information

- Information flows up, not down
 - “Writes up” allowed, “writes down” disallowed
- *-Property (Preliminary Version)
 - Subject s can write object o iff $L(s) \leq L(o)$ and s has permission to write o
 - Note: combines mandatory control (relationship of security levels) and discretionary control (the required permission)
 - Sometimes called “no writes down” rule

Basic Security Theorem, Preliminary Version

- If a system is initially in a secure state, and every transition of the system satisfies the simple security condition, preliminary version, and the $*$ -property, preliminary version, then every state of the system is secure

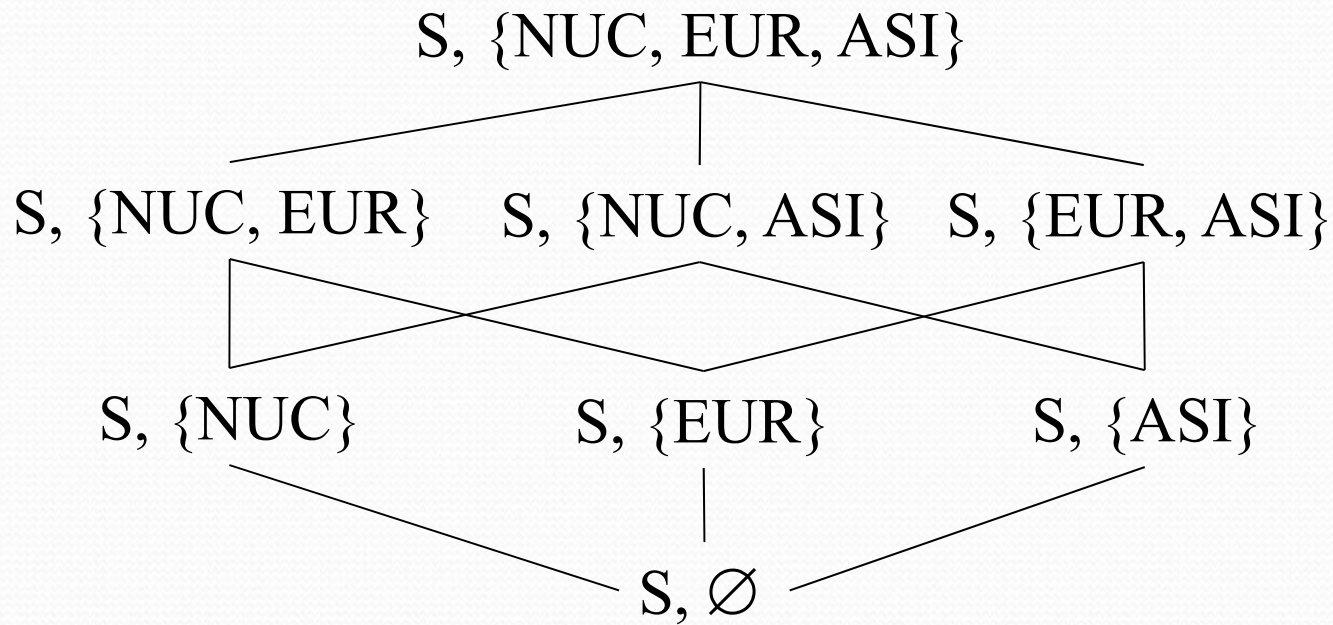
Bell-LaPadula Model, Full Version

- *Expand* notion of security level to include *categories*
→ the “need to know” principle
- Security level is now (*clearance, category set*)
- Examples
 - (Top Secret, { NUC, EUR, ASI })
 - (Confidential, { EUR, ASI })
 - (Secret, { NUC, ASI })

Levels and Lattices

- How do we compare the new security levels?
- The *dom* relationship
 - $(A, C) \text{ dom } (A', C') \text{ iff } A' \leq A \text{ and } C' \subseteq C$
- Examples
 - $(\text{Top Secret}, \{\text{NUC}, \text{ASI}\}) \text{ dom } (\text{Secret}, \{\text{NUC}\})$
 - $(\text{Secret}, \{\text{NUC}, \text{EUR}\}) \text{ dom } (\text{Confidential}, \{\text{NUC}, \text{EUR}\})$
 - $(\text{Top Secret}, \{\text{NUC}\}) \neg \text{dom } (\text{Confidential}, \{\text{EUR}\})$
- The *dom* relationship forms a lattice

Lattice Example



- S represents Secret
- The categories are NUC, EUR, and ASI
- The relationship is *dom*

Levels and Ordering

- Security levels partially ordered
 - Any pair of security levels may (or may not) be related by *dom*
- “dominates” serves the role of “greater than” in the preliminary version
 - “greater than” is a total ordering, though

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 - Note: combines mandatory control (relationship of security levels) and discretionary control (the required permission)
 - Sometimes called “no writes down” rule

Basic Security Theorem, Full Version

- If a system is initially in a secure state, and every transition of the system satisfies the simple security condition, full version, and the $*$ -property, full version, then every state of the system is secure

A Practical Problem

- Colonel has (Secret, {NUC, EUR}) security level
- Major has (Secret, {EUR}) security level
- Major can talk to colonel (“writes up” or “reads down”)
- Colonel cannot talk to major (“no reads up” or “no writes down”)
- Clearly does not make sense

Solution

- Define *maximum* and *current* levels for subjects
 - $\text{maxlevel}(s) \text{ dom } \text{curlevel}(s)$
- Allow subjects to adjust their current levels (thus power) as needed
- Example
 - Treat Major as an object (Colonel is writing to him/her)
 - Colonel has *maxlevel* (Secret, { NUC, EUR })
 - Colonel sets *curlevel* to (Secret, { EUR })
 - Now $L(\text{Major}) \text{ dom } \text{curlevel}(\text{Colonel})$
 - Colonel can write to Major without violating “no writes down”

The Data General/Unix (DG/UX) System

- Provides mandatory access controls
 - Was a Unix operating system
 - MAC label identifies security level
- Initially
 - Subjects assigned MAC label of parent
 - Initial label (at login time) is the label assigned to the user, kept in Authorization and Authentication database
 - Object assigned label at creation, and the label may be
 - Explicit: stored as part of attributes
 - Implicit: determined from parent directory

MAC Regions in the DG/UX Lattice

Hierarchy levels	↑	A&A database, audit	Administrative Region
		User data and applications	User Region
VP-1		Site executables	Virus Prevention Region
VP-2		Trusted data	
VP-3		Executables not part of the TCB	
VP-4		Executables part of the TCB	
VP-5		Reserved for future use	
		Categories	

IMPL_HI is “maximum” (least upper bound) of all levels

IMPL_LO is “minimum” (greatest lower bound) of all levels

Object Labels

- Requirement: every file system object must have a MAC label
 1. Roots of file systems have explicit MAC labels
 - ❑ If mounted file system has no label, it gets label of mount point
 2. Object with implicit MAC label inherits label of parent
 3. Creating hard link requires explicit label
 - ❑ If target object label implicit, it is made explicit
 - ❑ Moving a file makes label explicit

Object Labels (Cont.)

- 4. Change to directory label makes child labels explicit *before* the change
- Symbolic links are files, and treated as such, so ...
- 5. When resolving symbolic link, label of object is label of target of the link
 - System needs access to the symbolic link itself

Using MAC Labels

- Simple security condition implemented
- *-property not fully implemented
 - Process MAC must equal object MAC
 - Writing allowed only at same security level
- Overly restrictive in practice

MAC Tuples

- Up to 3 MAC ranges (one per region)
- MAC range is a set of labels with upper, lower bound
 - Upper bound must dominate lower bound of range

MAC Range Examples

1. [(Secret, {NUC}), (Top Secret, {NUC})]
 2. [(Secret, \emptyset), (Top Secret, {NUC, EUR, ASI})]
 3. [(Confidential, {ASI}), (Secret, {NUC, ASI})]
- (Top Secret, {NUC}) in ranges 1, 2
 - (Secret, {NUC, ASI}) in ranges 2, 3
 - [(Secret, {ASI}), (Top Secret, {EUR})] not valid range
 - as (Top Secret, {EUR}) $\neg dom$ (Secret, {ASI})

Objects and Tuples

- Objects must have MAC labels
 - May also have a MAC tuple
 - If both, tuple overrides label
- Example
 - Paper has MAC range:
[(Secret, {EUR}), (Top Secret, {NUC, EUR})]

Read Control Based on MAC Tuples

- Process can read object when:
 - Object MAC range (lr, hr) ; process MAC label pl
 - $pl \text{ dom } hr$
 - Process MAC label grants read access to upper bound of range
- Example
 - Paper has MAC range: $[(\text{Secret}, \{\text{EUR}\}), (\text{Top Secret}, \{\text{NUC}, \text{EUR}\})]$
 - Can Peter, with label $(\text{Secret}, \{\text{EUR}\})$, read paper?
 - No, because $(\text{Secret}, \{\text{EUR}\}) \neg \text{dom } (\text{Top Secret}, \{\text{NUC}, \text{EUR}\})$
 - Can Paul, with label $(\text{Top Secret}, \{\text{NUC}, \text{EUR}, \text{ASI}\})$, read paper?
 - Yes, $(\text{Top Secret}, \{\text{NUC}, \text{EUR}, \text{ASI}\}) \text{ dom } (\text{Top Secret}, \{\text{NUC}, \text{EUR}\})$

Write Control Based on MAC Tuples

- Process can write object when:
 - Object MAC range (lr, hr) ; process MAC label pl
 - $pl \in (lr, hr)$
 - Process MAC label grants write access to any label in range
- Example
 - Paper has MAC range: $[(\text{Secret}, \{\text{EUR}\}), (\text{Top Secret}, \{\text{NUC}, \text{EUR}\})]$
 - Can Peter, with label $(\text{Secret}, \{\text{EUR}\})$, write paper?
 - Yes, because $(\text{Top Secret}, \{\text{NUC}, \text{EUR}\}) \text{ dom } (\text{Secret}, \{\text{EUR}\})$ and $(\text{Secret}, \{\text{EUR}\}) \text{ dom } (\text{Secret}, \{\text{EUR}\})$

Key Points

- Confidentiality models restrict flow of information
- Bell-LaPadula models multilevel security
 - Cornerstone of much work in computer security