





Nitish Kumar Gupta

Course: GATE Computer Science Engineering(CS)

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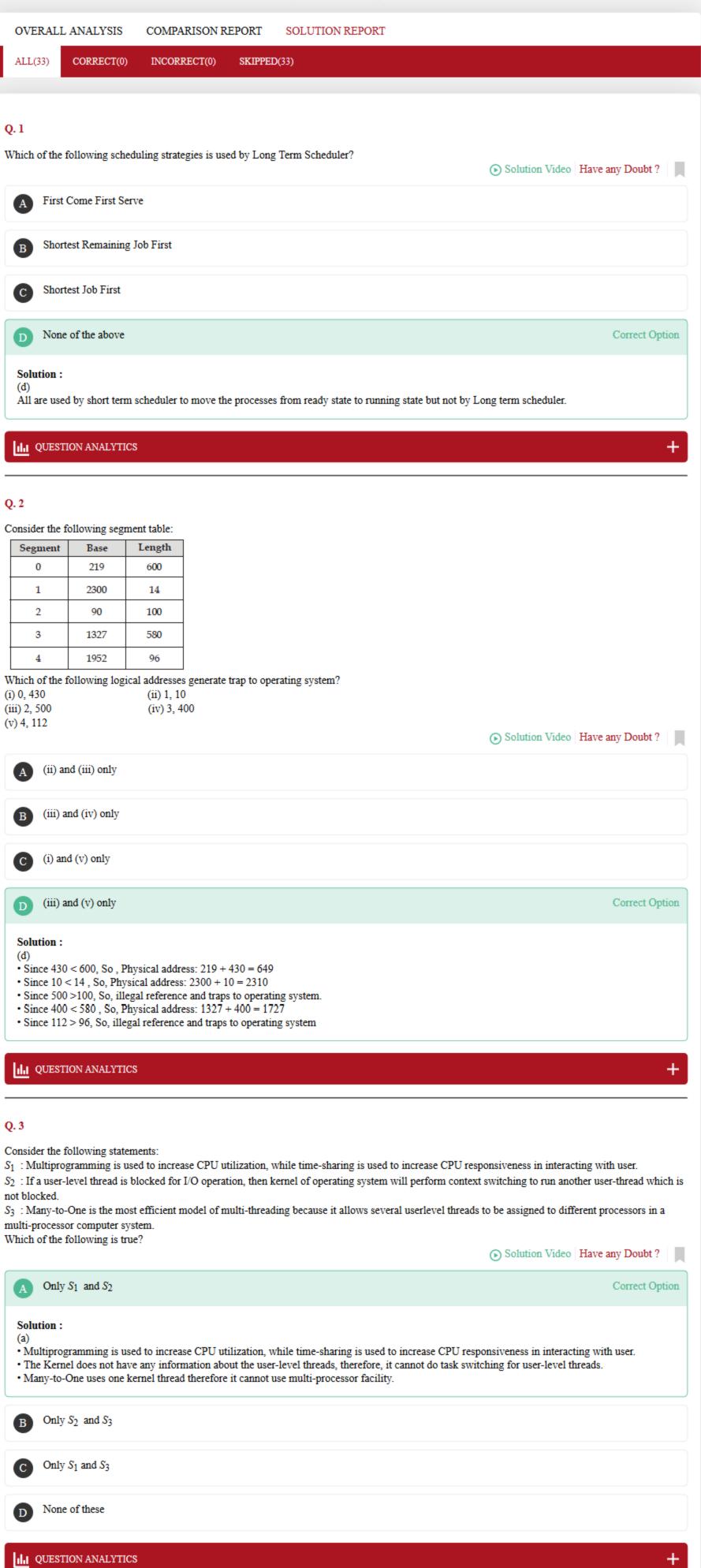
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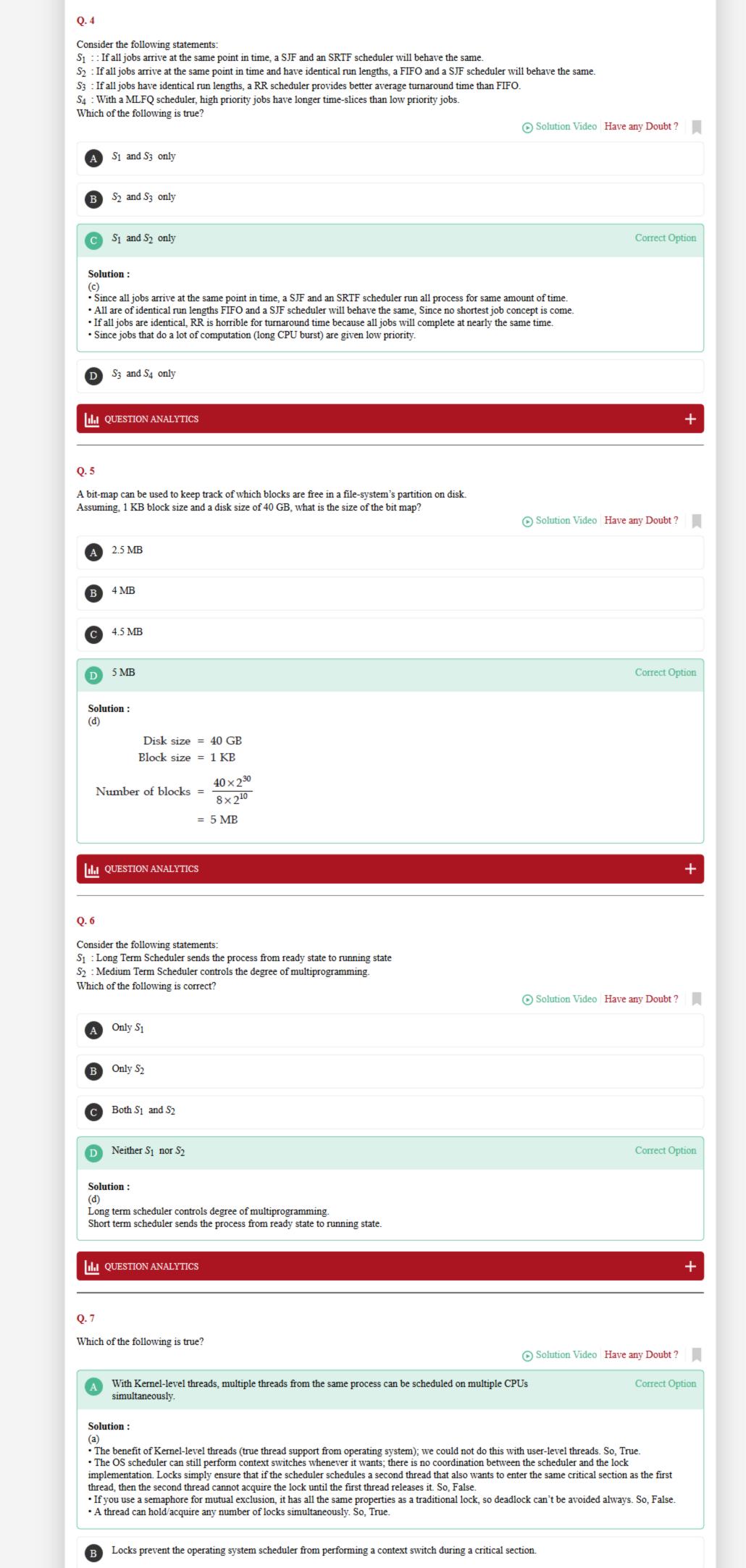
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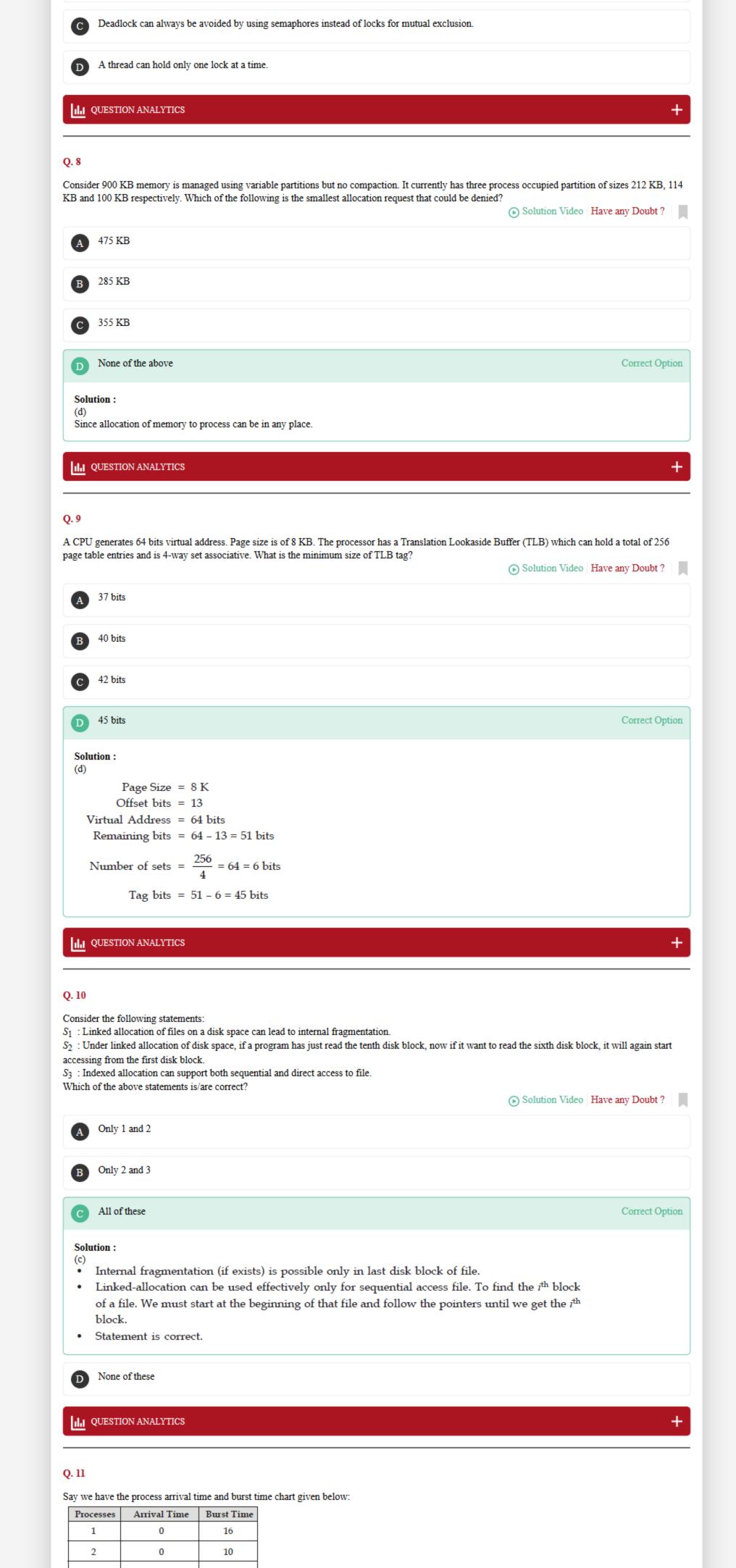
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SINGLE SUBJECT : OPERATING SYSTEM (GATE - 2019) - REPORTS







3	6	4
4	7	6
5	8	10

The average waiting time if round robin scheduling (Time quantum = 5 units) is used is ___ (Upto 1 decimal place)

Solution Video Have any Doubt?

Correct Option

23.4 (23.3 - 23.5)

Solution:

23.4 (23.3 - 23.5) Following is the Grantt chart:

	P_1	P ₂	P_1	P_3	P_4	P_5	P ₂	P_1	P_4	P_5	P_1
() !	5 1	0 1	5 1	9 2	4 2	9 3	34 3	19 4	0 4	15 4

Waiting queue: p_1' p_2' p_1' p_3' p_4' p_5' p_2' p_1' p_4' p_5' p_1'

Turnaround time = Completion Time - Arrival Time

Waiting time = Turn Around Time - Burst Time

Updation of burst times:

Processes	Burst time updation sequence
1	$16 \rightarrow 11 \rightarrow 6 \rightarrow 1$
2	$10 \rightarrow 5 \rightarrow 0$
3	$4 \rightarrow 0$
4	$6 \rightarrow 1 \rightarrow 0$
5	$10 \rightarrow 5 \rightarrow 0$

Processes	Arrival Time	Burst Time	Completion Time	Turn Around Time	Waiting Time
1	0	16	46	46	30
2	0	10	34	34	24
3	6	4	19	13	9
4	7	6	40	33	27
5	8	10	45	37	27
					Average = $\frac{117}{5}$ = 23.4



Q. 12

Let P and V be semaphore operations. P represents wait and V represents signal operation. Counting semaphore variable S is initialized to 1 and no blocked processes are present in the system. If the following operations are performed in the given order then the value of S is ___

Solution Video Have any Doubt?



5

Solution:

7 P operations, 11 V operations $S = 1 \implies S = 1 - 7 + 11 = 5$

 \therefore S = 5 after the given operations.

ILI QUESTION ANALYTICS

Correct Option

Correct Option

Q. 13

Consider a computer system with 34 bit logical address and 30 bit physical address. If the page size is 4 KB then the number of bits required to represent number of entries is an inverted page table is ______.

Solution Video Have any Doubt?



18

Solution:

18

Physical memory size = 1 GB

Physical address Number of entries in inverted page table = Page size

$$= \frac{2^{30}}{2^{12}} = 2^{18}$$

Number of bits = 18

III QUESTION ANALYTICS

Q. 14

A multiprocessor with 8 processors has 20 attached tape drives. There is a large number of jobs submitted to the system and each requires 3 tape drives to complete execution. Assume that each job starts running with only 3 tape drives. The maximum number of jobs that can execute at once are

Solution Video Have any Doubt?

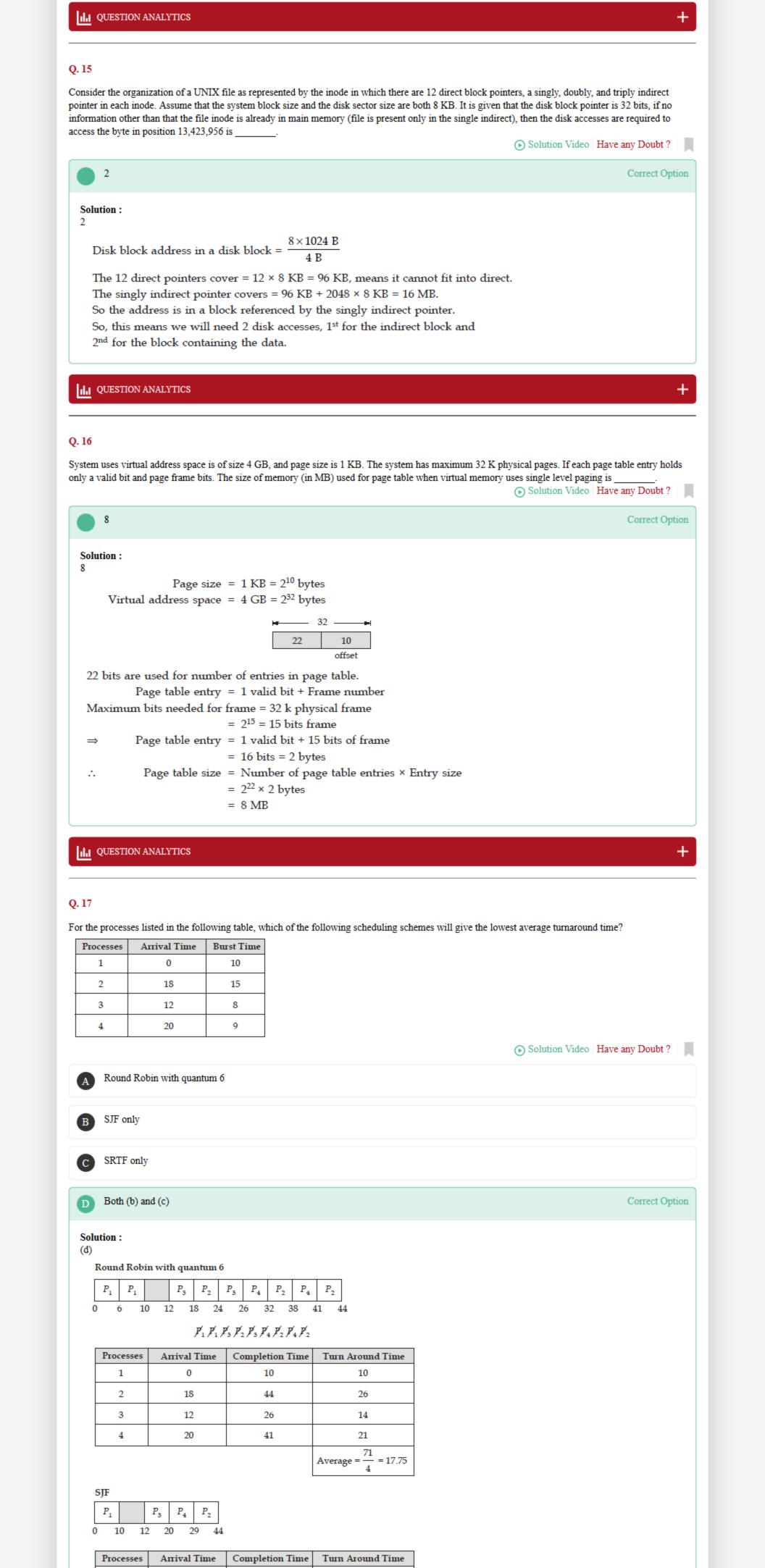


Solution:

Correct Option

If each process can be initially allocated with three tape drives. The fourth one will be allocated

on demand then maximum $\left\lfloor \frac{20}{3} \right\rfloor = 6$ processes can be active simultaneously.



1	0	10	10
2	18	44	26
3	12	20	8
4	20	29	9
			Average = $\frac{53}{4}$ = 13.25

SRTF

	P_1		P_3	P.	£	P_2
0	1	0 1	2	20	29	4

Processes	Arrival Time	Completion Time	Turn Around Time
1	0	10	10
2	18	44	26
3	12	20	8
4	20	29	9
			Average = $\frac{53}{}$ = 13.25

QUESTION ANALYTICS

+

Q. 18

A computer system uses the Banker's Algorithm to deal with deadlocks. Its current state is shown in the table below, where P_0 , P_1 , P_2 , P_3 , P_4 are processes and R_0 , R_1 , R_2 are resources types.

Process	Maxi	imum	need	Curre	nt allo	cation	Available			
	R_0	R_1	R ₂	R ₀	R_1	R ₂	R ₀	R_1	R ₂	
P_0	7	7	4	7	5	3	1	3	2	
P_1	3	5	6	3	2	2				
P_2	9	8	2	3	6	2				
P_3	2	2	2	2	1	1				
P_4	3	3	4	2	0	1				

From the perspective of deadlock avoidance, which one of the following is true?

Solution Video | Have any Doubt ?

A

The system is not in safe state, but would be safe if one more instance of R_0 were available.

B The system is in safe state.

Correct Option

Solution:

(b)

Process	Maxi	mum	need	Current allocation			Available			Remaining need		
	R_0	R_1	R_2	R_0	R_1	R ₂	R_0	R_1	R_2	R_0	R_1	R_2
P_0	7	7	4	7	5	3	1	3	2	0	2	1
P_1	3	5	6	3	2	2				0	3	4
P_2	9	8	2	3	6	2				6	2	0
P_3	2	2	2	2	1	1				0	1	1
P_4	3	3	4	2	0	1				1	3	3

The system is not in safe state, but would be safe if two more instance of R_2 were available.

The system would be in unsafe state, if 1 instance of R_1 were available rather than 3.

QUESTION ANALYTICS

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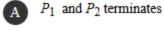
Q. 19

Consider the following 'C' code:

$$x = 0;$$
 $y = 1;$
 $z = \text{true};$
initialization
 P_2
while (z)

 $\{x = x + 1;\}$ $\{y = y + 1; z = (x! = y);\}$ Assume P_1 and P_2 are two concurrent processes and sharing the global variables x, y and z. Assignments and tests are atomic. Which is the following is false about P_1 and P_2 ?

Solution Video | Have any Doubt ?



 P_1

while (x < y)

P₂ terminates but P₁ does not

 P_1 terminates but P_2 does not

Correct Option

Solution:

(c) Initially: x = 0, y = 1, z = true;

 $P_1: 1.$ while (x < y)2. $\{x = x + 1\}$

 $P_2:1.$ while (z)

2. $\{y = y + 1;$

3. z = (x ! = y):

1. Both threads terminate : is possible.

 $P_2: 1, 2$

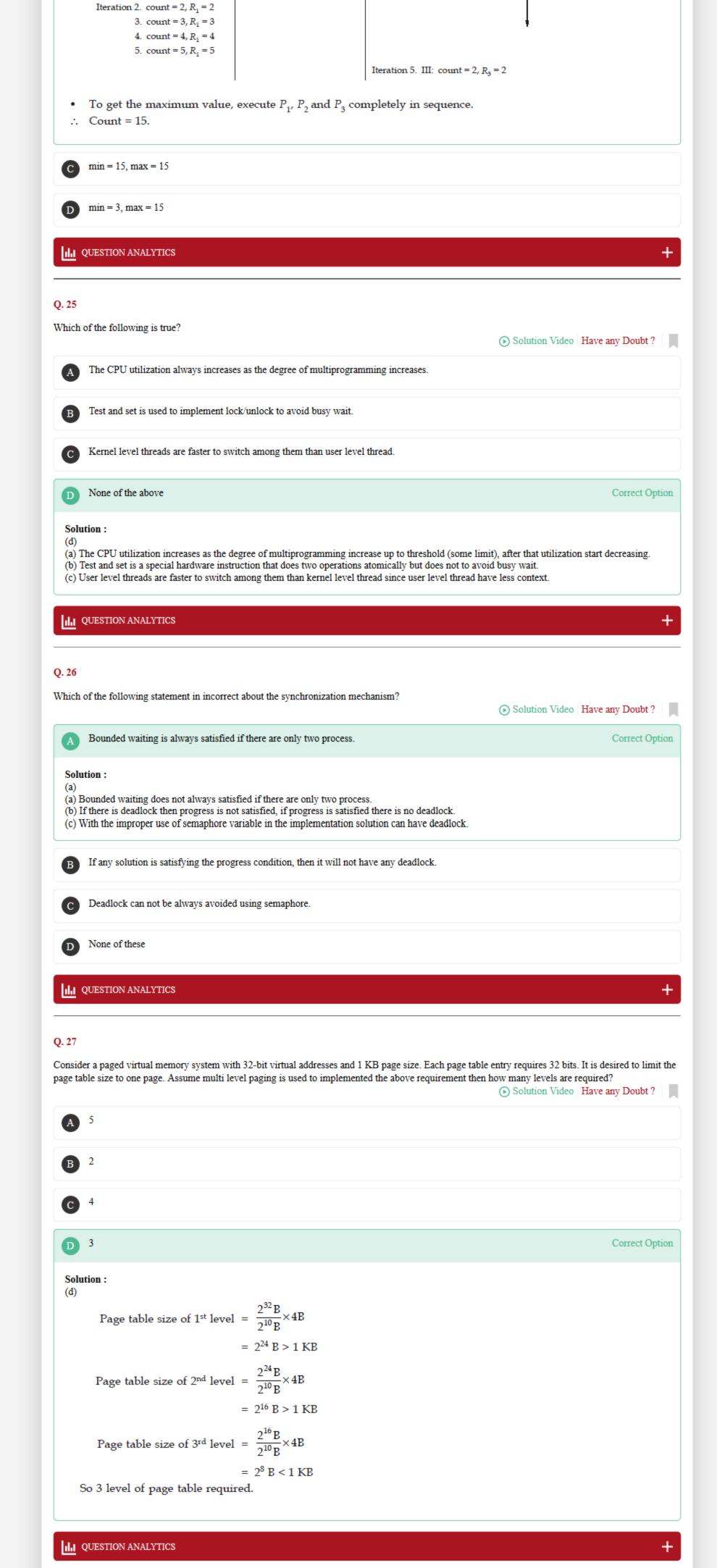
 $P_1: 1, 2, 1, 2, 1$ (terminated)

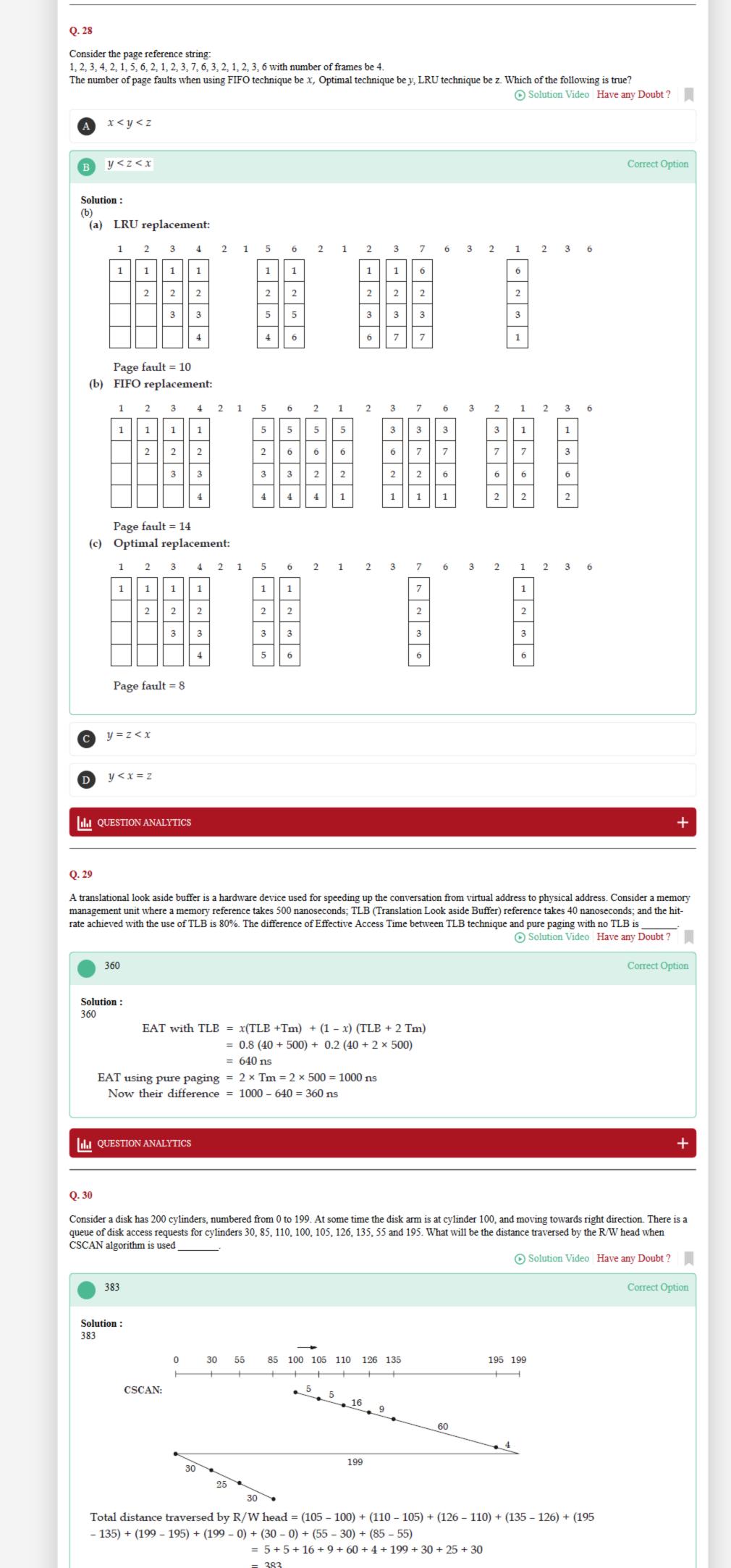
 P_2 : 3, 1 (terminated) 2. P_1 terminates but P_2 does not: is possible

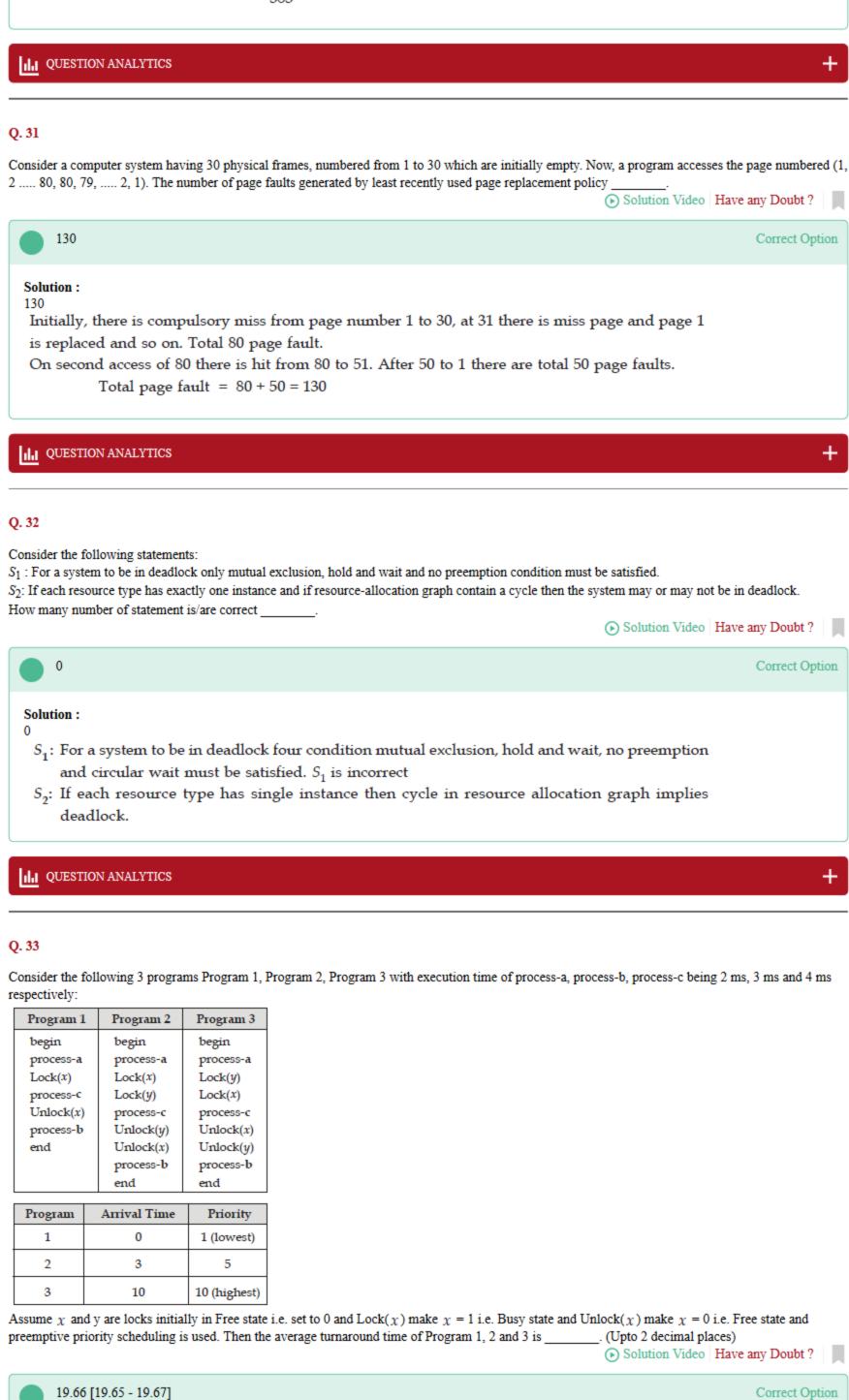
 P_1 terminates but P_2 does not: is poss $P_1 \cdot 1 \cdot 2 \cdot 1$ (terminated)

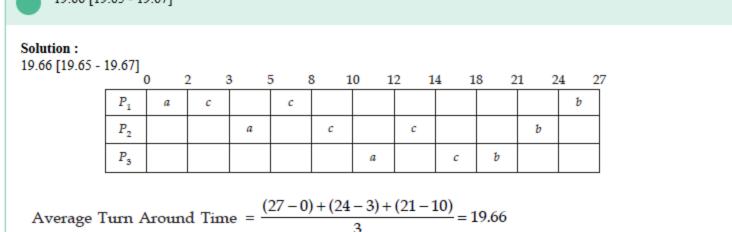
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P_2: 1, 2, 3, 1, 2, 3 (infinite loop)
     3. P_2 terminates but P_1 does not : is not possible
     4. Neither P_1 nor P_2 terminates : is possible
         P_2: 1, 2, 3
         P_1^-: 1, 2
          P_2: 1, 2, 3 \\ P_1: 1, 2  Infinite iterations (P_2 \text{ followed by } P_1)
       Neither P_1nor P_2 terminates
 III QUESTION ANALYTICS
Q. 20
Consider the following set of processes that need to be scheduled on a single CPU operating system uses preemptive shortest remaining time first
algorithm. What is the average waiting time of these processes?
     Process
                   Arrival Time
                                       Burst Time
        P_0
                                           3
                        2
        P_1
                                            2
                        0
       P_2
                        3
                                            4
       P_3
                        7
                                           1
       P_4
                        8
                                           3
                              (All time in milliseconds)
                                                                                                  Solution Video Have any Doubt?
 A 1
                                                                                                                           Correct Option
  Solution:
  (a)
                  Waiting Time = Turn Around Time - Burst Time
       Average Waiting Time = \sum_{i=0}^{n} \frac{\text{Waiting time of } P_i}{\text{Total number of process}}
                                          Process
                                                        Waiting Time
                                             P_0
                                             P_1
                                                              0
                                             P_2
                                                              3
                                             P_3
                                                              0
                                                              2
       Average Waiting Time = \frac{0+0+3+0+2}{5} = 1 ms
 B 1.2
 C 1.4
 D
       1.5
 III QUESTION ANALYTICS
Q. 21
Consider the two process P_0 and P_1 using binary semaphore R, S, T initial value of R = 1, S = 0, T = 0.
           P_0
       While (1)
                                                While (1)
           P(R);
                                                    P(S);
       printf("1");
                                               printf("1");
      printf("4");
                                                    P(T);
           V(S);
                                                printf("3");
       printf("2");
                                               printf("2");
           V(T)
                                                    V(S);
What will be the output of program when process P_0 and P_1 concurrently executing?
                                                                                                  Solution Video Have any Doubt?
       141322
       142132
       141232
  D Both (b) and (c)
                                                                                                                           Correct Option
  Solution:
    Value of semaphore R=1 and S=0 so first P_0 is executed and 14 will be printed. After V(S) both
    P_{\mathrm{0}} and P_{\mathrm{1}} will concurrently executing and 21 or 12 can be printed. Now 32 will be printed.
    Both 1 4 2 1 3 2 and 1 4 1 2 3 2 can be printed.
    So option (d) is correct.
  III QUESTION ANALYTICS
Q. 22
```

Consider the following statements: S_1 : Pages that are shared between two or more processes can never be swapped out to the disk. S₂: Demand paging requires the programmer to instruct the operating system to load a particular virtual memory page. S_3 : The translation look aside buffer is a software data structure that supports the virtual memory address translation operation. Which of the following is correct? Solution Video Have any Doubt? Only S_1 Only S_1 and S_2 Only S_1 , S_2 and S_3 Correct Option None of these Solution: · Pages that are shared between two or more processes can be swapped out to disk when demand paging is applied and we have to swap in new pages and main memory is full. The operating system automatically loads pages from disk when necessary when it is needed • The translation look aside buffer is a hardware data structure. III QUESTION ANALYTICS Q. 23 To obtain better memory utilization, dynamic loading is used. A routine is not loaded until it is called. For implementing dynamic loading, which of the following is correct? Solution Video Have any Doubt? All routine are kept on memory in a relocatable load format. All routines become part of the code and are loaded at the same time regardless of whether they are called or not. Both (a) and (b) Correct Option Dynamic loading optimizes memory space utilization. Solution: The purpose of dynamic loading is optimal utilization of memory. Option (a) all routine are kept on disk in a relocatable load format. Option (b) is property of static loading **III** QUESTION ANALYTICS Q. 24 Consider the following segment: int count = 0;void tally() for (int i = 1; i <=5; i++) count = count +1; main() parbegin tally(); tally(); tally(); parend Note: Assume the count = count + 1; will execute in '3' different instructions. I. Load R_i, m[count] II. INC R_i III. Store m[count], R; where m[count] refers to memory value of count variable. Preemption can occur while executing the above instructions. Each of the tally function has a separate register R allocated to it in which the value of count is stored. After completion of main function, what will be the minimum and maximum values of count in the end of the program? Solution Video Have any Doubt? min = 5, max = 15min = 2, max = 15Correct Option Solution: To get the minimum value of count. Tally2() Tally3() Tally1() Iteration 1. I: count = 0, $R_1 = 0$ II: count = 0, $R_1 = 0$ Iteration 1. count = 1, $R_2 = 1$ 2. count = 2, $R_2 = 2$ 3. count = 3, R_2 = 3 4. count = 4, $R_2 = 4$ 5. count = 5, $R_2 = 5$ Iteration 1. count = 6, $R_3 = 6$ 2. count = 7, $R_3 = 7$ 3. count = 8, $R_3 = 8$ Time 4. count = 9, $R_3 = 9$ Iteration 1. III: count = 1, $R_1 = 1$ Iteration 5. I: count = 1, $R_3 = 1$ II. count = 1, $R_3 = 2$









Average Turn Around Time = ______ = 19.00

ILI QUESTION ANALYTICS