












**Nitish Kumar Gupta**  
 Course: GATE  
 Computer Science Engineering(CS)


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SINGLE SUBJECT : COMPUTER NETWORKS (GATE - 2019) - REPORTS

OVERALL ANALYSIS	COMPARISON REPORT	SOLUTION REPORT
ALL(33)	CORRECT(0)	INCORRECT(0)
		SKIPPED(33)

Q. 1

Which of the following is not a class A host address?

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- A

128.4.5.6

Correct Option
- Solution :**  
 (a)  
 Class A addresses are 0.0.0.0 to 127.255.255.255  
 So 128.4.5.6 is not a class A address.
- B

117.4.5.1
- C

117.0.0.0
- D


117.8.0.0


 QUESTION ANALYTICS

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Q. 2

Identify the correct sequence in which the following packets are transmitted on the network by a host when a browser requests a web page from a remote server, assuming that the host has just been restarted.

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- A

HTTP GET request, DNS query, TCP SYN
- B

DNS query, HTTP GET request, TCP SYN
- C

DNS query, TCP SYN, HTTP GET request

Correct Option
- Solution :**  
 (c)  
 DNS query, TCP SYN, HTTP GET request.
- D


TCP SYN, DNS query, HTTP GET request


 QUESTION ANALYTICS

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Q. 3

The reference polynomial used in a CRC scheme is  $x^4 + x^3 + 1$ . A data sequence 1010101010 is to be sent. Determine the actual bit string that is transmitted.

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- A

10101010100011
- B

10101010101110
- C

10101010100110
- D

10101010100010


Correct Option
- Solution :**  
 (d)  
 Reference polynomial  $x^4 + x^3 + 1 = 11001$   
 Data sequence = 1010101010  
 On dividing 1010101010 by 11001,  
 We get CRC (i.e. remainder) = 0010  
 On appending this CRC on data sequence, we get the actual message transmitted i.e. 10101010100010.


 QUESTION ANALYTICS

+

Q. 4

If HLEN field of IP is 1111 then what is the length of header?

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- A

15
- B

30
- C

45

**Solution :**  
(d)  
 $HLEN \times 4 = \text{Header Size}$   
 $15 \times 4 = 60$

Q. 5

Consider Bob uses an e-mail client (mail reader) such as Outlook to send an e-mail to Alice who uses a Web-based e-mail account (Gmail). The IP address of Alice's mail server is initially unknown to Bob's mail server. Which of the following is the correct sequence of activity for transport and application layer protocols that are involved from the time when Bob sends the e-mail to the time when Alice reads it?

- $S_1$  : DNS over UDP is used by Bob's mail server to get the IP address for Alice's mail server.  
 $S_2$  : SMTP over TCP is used to transfer the mail from Bob's mail server to Alice's mail server.  
 $S_3$  : SMTP over TCP is used to transfer the mail from Bob's e-mail client to his SMTP server.  
 $S_4$  : HTTP over TCP is used by Alice to read this e-mail.

 Solution Video

Have any Doubt ?



A  $S_1, S_2, S_3$  and  $S_4$

B  $S_1, S_3, S_2$  and  $S_4$

C  $S_3, S_1, S_2$  and  $S_4$

Correct Option

**Solution :**  
(c)  
1. SMTP over TCP is used to transfer the mail from Bobs e-mail client to his SMTP server.  
2. DNS over UDP is used by Bob's mail server to get the IP address for Alice's mail server.  
3. SMTP over TCP is used to transfer the mail from Bob's mail server to Alice's mail server.  
4. Finally HTTP over TCP is used by Alice to read this e-mail.

D  $S_3, S_2, S_1$  and  $S_4$

Q. 6

Which of the following system call adds the local socket address to an already created socket?

 Solution Video

Have any Doubt ?



A listen

B create

C bind

Correct Option

**Solution :**  
(c)  
The bind( ) system call is used to associate a local address with a already created socket.

D connect

Q. 7

Which protocol uses a connection-oriented service to deliver files between end systems?

 Solution Video

Have any Doubt ?



A TFTP

B DNS

C FTP

Correct Option

**Solution :**  
(c)  
FTP depends on TCP, is connection oriented, and provides reliable control i.e. FTP has two connection one for sending data and one for Sending control signal.

D SNMP

Q. 8

Which of the following statements is true about Binary exponential backoff, a mechanism used in MAC protocols?

 Solution Video

Have any Doubt ?



A It ensures that two or more nodes that experience a collision in a time slot will experience a lower probability of colliding with each there when they each retry that packet.

Correct Option

**Solution :**  
(a)  
It ensures that two or more nodes that experience a collision in a time slot will experience a lower probability of colliding with each there when they each retry that packet.

B It ensures that two nodes that experience a collision in a time slot will never collide with each other when they each retry that packet

C

It can be used with slotted Aloha but not with carrier sense multiple access.

D

It improves the fairness of the throughput achieved by different nodes compared to not using the mechanism over a short period of time.

QUESTION ANALYTICS

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Q. 9

Which of the timer is responsible for keeping the window size information flowing even if the other end closes its receiver window?

[Solution Video](#) | [Have any Doubt ?](#)

A

Retransmission timer

B

Persistent timer

Correct Option

**Solution :**

(b)

(b)

1. A retransmission timer is used when expecting an acknowledgment from the other end.

2. A persistent timer keeps window size information flowing even if the other end closes its receiver window.

3. A keep alive timer detects when the other end on an otherwise idle connection crashes or reboots.

4. A 2MSL timer measures the time a connection has been in the TIME\_WAIT state.

C

Keep alive timer

D

Time-wait timer

QUESTION ANALYTICS

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Q. 10

Consider a 10 km long broadcast LAN has  $10^{10}$  bps bandwidth and uses CSMA/CD. The signal travels along the wire at  $2.5 \times 10^8$  m/s. What is the minimum packet size that can be used on this network?

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A

$10^6$  bits

B

$10^7$  bytes

C

$10^4$  bits

D

$10^5$  bytes

Correct Option

**Solution :**

(d)

$$\text{Minimum packet size} = \frac{(2 \times \text{Distance} \times \text{Bandwidth})}{\text{Velocity}}$$

$$\text{distance} = 10 \text{ km}$$

$$\text{Velocity} = 2.5 \times 10^8 \text{ m/s}$$

$$\text{Bandwidth} = 10^{10} \text{ bps}$$

$$\begin{aligned} \text{Minimum packet size} &= \frac{(2 \times 10 \times 10^3 \times 10^{10})}{(2.5 \times 10^8) \text{ bits}} \\ &= 8 \times 10^5 \text{ bits} = 10^5 \text{ bytes} \end{aligned}$$

QUESTION ANALYTICS

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Q. 11

Consider Noisy station that detects transmissions and disrupts them by beginning a competing transmission as soon as it hears the beginning of the transmitted frame, thereby causing a collision. Assume Detector machine, which is on this Ethernet with bandwidth 10 Mbps, detects collision during the transmission of its 12<sup>th</sup> bit on the wire (including any preamble). If the speed of the signal in the wire is  $2 \times 10^8$  meter/second, then the distance (in meters) of the Noisy station from Detector machine is \_\_\_\_\_.

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120

Correct Option

**Solution :**

120

If we see the collision after 12 bit times, then that means that the Noisy host sent its competing transmission after 6 bit times. (i.e. one bit from original data and second bit from noisy station, it will be continue till 12<sup>th</sup> bit)

This means that the distance is

$$= \frac{6 \text{ Bits}}{10 \text{ Mbps}} \times (2 \times 10^8 \text{ m/s})$$

$$= 120 \text{ meters}$$

QUESTION ANALYTICS

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Q. 12

Consider an instance of TCP's Additive Increase Multiplicative Decrease (AIMD) algorithm where the window size at the start of the slow start phase is 2 MSS and the threshold at the start of the first transmission is 12 MSS. Assume that a time out occurs during the 7<sup>th</sup> transmission. The congestion window size at the end of the 12<sup>th</sup> transmission is \_\_\_\_\_. (in MSS)

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10

Correct Option

**Solution :**

10

Threshold = 12 MSS

Window size for 1<sup>st</sup> transmission = 2 MSS

Window size for 2<sup>nd</sup> transmission = 4 MSS

Window size for 3<sup>rd</sup> transmission = 8 MSS

Window size for 4<sup>th</sup> transmission = 12 MSS

Threshold reached, increase linearly (according to AIMD)

Window size for 4<sup>th</sup> transmission = 13 MSS

Window size for 5<sup>th</sup> transmission = 14 MSS

Window size for 6<sup>th</sup> transmission = 15 MSS

Window size for 7<sup>th</sup> transmission = 16 MSS

Time out occurs, resend 8<sup>th</sup> transmission with window size starting from 2 MSS and new threshold is 8 MSS.

Window size for 8<sup>th</sup> transmission = 2 MSS

Window size for 9<sup>th</sup> transmission = 4 MSS

Window size for 10<sup>th</sup> transmission = 8 MSS

Threshold reached, now increase linearly (according to AIMD)

Window size for 11<sup>th</sup> transmission = 9 MSS


Window size for 12<sup>th</sup> transmission = 10 MSS


 QUESTION ANALYTICS



### Q. 13

A leaky bucket with the capacity of bucket of 200 MB is at the host network interface. The data rate on the network is 2 Mbyte/s. If the host has 450 Mbytes to send onto the network and it sends the data in a burst then the maximum data speed from the host in order that no data is lost is \_\_\_\_\_ in Mbps.  
(Upto 1 decimal place)

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 3.6 (3.5 - 3.7)

Correct Option

**Solution :**

3.6 (3.5 - 3.7)

Bucket size = 450 Mbyte - Actual data to send

200 Mbyte = 450 Mbyte - 250 Mbyte

2 Mbyte → 1 sec

250 Mbyte → ??

Time for computer to transmit all data =  $\frac{250}{2} = 125$  sec

450 Mbytes → 125 sec

?? → 1 sec

Host Max data rate = (450 Mbytes)/125 sec  
= 3.6 Mbyte/sec


 QUESTION ANALYTICS



### Q. 14

The distance between two microwave towers, with the link capacity 100 Mbps, is 24 Km and the speed of the signal is  $3 \times 10^8$  m/sec. If the frame size is 50 KB in the GoBack-N protocol, The approximate link utilization is \_\_\_\_\_ in %. (Assume that the ACK packets are negligible in the size and there are no errors during communication) (Upto 2 decimal places)

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 75.75 (75.74 - 75.76)

Correct Option

**Solution :**

75.75 (75.74 - 75.76)

Transmission Time (T.T.) = Data size/Bandwidth

$$= \frac{(50 \times 10^3)}{100 \times 10^6} = 500 \text{ microsec}$$

$$\text{Propagation Time (P.T.)} = \frac{\text{Length}}{\text{Velocity}}$$

$$= \frac{(24 \times 10^3)}{3 \times 10^8} = 80 \text{ microsec}$$

$$\% \text{ Link utilisation} = \left[ \frac{\text{T.T.}}{(\text{T.T.} + 2\text{P.T.})} \right] \times 100$$

$$= \left[ \frac{500}{(500 + 2 \times 80)} \right] \times 100 = 75.75\%$$

 QUESTION ANALYTICS



### Q. 15

In a RSA cryptosystem, a participant uses two prime numbers  $p$  and  $q$  is 17 and 11 respectively to generate his/her public key and private keys. If the public key of participant is 7 and cipher text(C) is 11, then the original message(M) is \_\_\_\_\_.

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 88

Correct Option

**Solution :**

88

$p = 17$  and  $q = 11$

$n = p \times q = 187$

$\phi(n) = (p - 1) \times (q - 1)$

$= 16 \times 10 = 160$

$e = 7$

$d = e^{-1} \bmod \phi(n)$




$$\begin{aligned}
 a &= e^{-1} \bmod \phi(n) \\
 &= 7^{-1} \bmod 160 = 23 \\
 M &= C^d \bmod n \\
 &= 11^{23} \bmod 187 = 88
 \end{aligned}$$

 QUESTION ANALYTICS



#### Q. 16

Suppose a switch is built using a computer workstation and that it can forward packets at a rate of 500000 packets per second, regardless of size. Assume the workstation uses Direct Memory Access (DMA) to move data in and out of its main memory, which has a bandwidth of 2 Gbps and that the I/O bus has a bandwidth of 1 Gbps. The packet size at which the bus bandwidth become the limiting factor \_\_\_\_\_ in Bytes.

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 125

Correct Option

**Solution :**  
125

The workstation can handle i.e. limited by the I/O bus =  $\frac{1000}{2} = 500$  Mbps

Let the packet size be  $x$  bits;

To support 500000 packets/second we need a total capacity of  $500000 \times x$  bps;

Equating  $5 \times 10^5 \times x = 500 \times 10^6$  bps,

We get  $x = 1000$  bits  
= 125 bytes


 QUESTION ANALYTICS




#### Q. 17

Consider 3 employees Ramesh, Suresh and Mahesh using three different computers working under Account-Network serviced by a single mail server. Assume Ramesh wish to send e-mail to Sonali, Suresh to Pooja and Mahesh to Vidhi, where Sonali, Pooja and Vidhi working under HRNetwork. What is the worst case and best case number of TCP connections that should be opened to enable this e-mail exchange if all the persons use a non-HTTP e-mail service?

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 Worst case = 9, Best case = 5

 Worst case = 9, Best case = 7

Correct Option


**Solution :**  
(b)


**Worst Case:** When they sent e-mails at different times.

1 each from Ramesh, Suresh, Mahesh to their mail server, one each from Accounts-mail server to HR-mail server for the 3 e-mails, 1 each from Sonali, Pooja ,Vidhi to the HR-mail server) =  $3 + 3 + 3 = 9$  TCP connections.

**Best Case:** When all 3 persons send e-mail roughly the same time.

(1 each from Ramesh, Suresh, Mahesh to their mail server, one from Accounts-mail server to HR-mail server, each from Sonali, Pooja ,Vidhi to the HR-mail server) =  $3 + 1 + 3 = 7$  TCP connections.

 Worst case = 7, Best case = 5

 Worst case = 9, Best case = 3


 QUESTION ANALYTICS




#### Q. 18

Consider Dijkstra's algorithm in the link-state routing protocol at node  $u$ . Professor Ram first sets the route for each directly connected node  $v$  to be the link connecting  $u$  to  $v$ . Ram then implements the rest of the algorithm correctly, aiming to produce minimum-cost routes, but does not change the routes to the directly connected nodes. In this network,  $u$  has at least two directly connected nodes and there is more than one path between any two nodes. Assume that all link costs are non-negative. Which of the following statements is False of  $u$ 's routing table?

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 There are topologies and link costs where the majority of the routes to other nodes will be incorrect.

 There are topologies and link costs where no routing table entry (other than from  $u$  to itself) will be correct.

Correct Option

**Solution :**


(b)

(a) Is true since for example, all the neighbors but one could have very high cost, and all the other links have low cost, so all the routes could in fact be just one link.

(b) Is false since The lowest-cost neighbor's route will be the direct link, of course!

(c) Is true since A trivial example is when all the links have equal cost.

 There are topologies and link costs where all routing table entry (other than from  $u$  to itself) will be correct.

 Both (a) and (b)


 QUESTION ANALYTICS



#### Q. 19

Consider the Select Repeat sliding window protocol is used at datalink layer to transmit frames between Sender and Receiver machine. The receiver sends "ACK  $k$ " when it receives a packet with sequence number  $k$ . Denote the window size by  $W$ . The sender's packets start with sequence number 1. Which of the following is true of a correct implementation of this protocol over a besteffort network?

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 Any new previously unsent packet with sequence number greater than  $W$  is sent by the sender iff a new previously unseen ACK arrives.

Correct Option

**Solution :**

- (a) Any new previously unsent packet with sequence number greater than W is sent by the sender iff a new previously unseen ACK arrives.  
 (b) The sender can send more than one packet between the receipt of one ACK and the next because the sender could time-out and retransmit.  
 (c) The receiver can never discard any new, out-of-order packet it receives after sending an ACK for it because the sender thinks the receiver has delivered this packet to the Receiver.  
 (d) False since Packets or ACKs could get reordered in the network.

**B** The sender will never send more than one packet between the receipt of one ACK and the next.

**C** The receiver can discard any new, out-of-order packet it receives after sending an ACK for it.

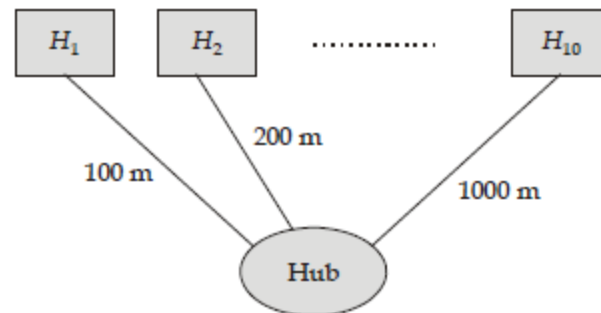
**D** Suppose that no packets or ACKs are lost and no packets are ever retransmitted. Then ACKs will arrive at the sender in non-decreasing order.

 QUESTION ANALYTICS



#### Q. 20

Consider the speed of propagation of a signal along a cable is  $2 \times 10^8$  m/s and network has 10 Mbps CSMA/CD network interconnecting ten computers shown in below image. Each computer is connected to the hub with a cable of different length. Computer  $H_1$  is connected via a 100 m cable, computer  $H_2$  via a 200 m cable, computer  $H_3$  via a 300 m cable and so on upto computer  $H_{10}$ , which is connected via a 1000 m cable (ignore the requirement of repeater due to signal degradation). Assume that the hub introduces a delay of 2.5 microseconds. What is the shortest packet length  $L_{\min}$  of this network in order to ensure that the CSMA/CD protocol functions properly?



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**A** 340 bits

**B** 240 bits

Correct Option

**Solution :**  
 (b)

$$\text{The maximum end-to-end propagation delay is} = \frac{(1000\text{m} + 900\text{m})}{(2 \times 10^8 \text{ m/s})} = 9.5 \mu\text{sec}$$

The hub introduces an additional 2.5 microseconds delay.

The time taken to transmit  $L_{\min} \geq 2 \times (9.5 \mu\text{sec} + 2.5 \mu\text{sec}) \geq 24 \mu\text{sec}$

$$\begin{aligned} \text{Hence, } L_{\min} &= 24 \mu\text{sec} \times 10 \text{ Mbps} \\ &= 240 \text{ bits} \end{aligned}$$

**C** 520 bits

**D** 120 bits

 QUESTION ANALYTICS



#### Q. 21

Consider the following statements regarding the slow start phase of the TCP congestion control algorithm. Note the cwnd stands for the TCP congestion window and MSS denotes the Maximum Segment Size.

- (i) The cwnd increases by 2 MSS on every successful acknowledgment.  
 (ii) The cwnd approximately doubles on every successful acknowledgment.  
 (iii) The cwnd increases by 1 MSS every round trip time.  
 (iv) The cwnd approximately doubles every round trip time.

Which one of the following is correct?

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**A** Only (ii) and (iii) are true

**B** Only (i) and (iii) are true

**C** Only (iv) is true

Correct Option

**Solution :**  
 (c)  
 cwnd approximately doubles every round trip time.

**D** Only (i) and (iv) are true

 QUESTION ANALYTICS



#### Q. 22

Consider the following statements regarding TCP's congestion control phases.

- (i) The size of the congestion window increases exponentially until it reaches a threshold (in slow start algorithm).  
 (ii) In multiplicative decrease procedure, threshold gets decreased to one-half of the previous window size.  
 (iii) The size of the congestion window increases exponentially after achieving threshold value till the timeouts.

Which of the following statements is true ?

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**A** (i) and (ii) only

Correct Option

**Solution :**

(a)


In the slow start algorithm, the size of the congestion window increases exponentially until it reaches a threshold, after this there is additive increases (one-one window) till the time outs. Statement (iii) is false.

- ☐ (i) and (iii) only
- ☐ (ii) and (iii) only
- ☐ (i), (ii) and (iii)


 QUESTION ANALYTICS
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Q. 23

ICMP has error reporting messages. Which of the following error reporting message describes the given statement ‘S’?  
 S : The packet is discarded due to the processing problem observing a change in the header format of the IP datagram.


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- ☐ Destination unreachable
- ☐ Source quench
- ☒ Parameter problem Correct Option
- Solution :**  
 (c)  
 • In redirection packet is not discarded but it is redirected to a n/w as the host doesn’t belong to this network.  
 • In source quench packet is discarded due to congestion in the n/w.  
 • Destination unreachable means host is not present in the n/w or the host is not responding to the request, then the packet is discarded.
- ☐ Redirection


 QUESTION ANALYTICS
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Q. 24

Suppose that  $x$  bits of user data are to be transmitted over K-hop path in a packet-switched network as a series of packets each containing p data bits and h header bits with  $x \gg (p + h)$ . The bit rate of lines is b bps and Propagation delay is negligible. What is the time taken by the source to transmit total bits?


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- ☐  $(p + h) x/b$  secs
- ☒  $(p + h) x/pb$  secs Correct Option
- Solution :**  
 (b)  
 The total number of packets needed is  $x/p$ .  
 Total data + header traffic is  $(p + h) x/p$  bits.  
 The source requires  $(p + h) x/pb$  seconds to transmit these bits.
- ☐  $p x/b$  secs
- ☐  $hx/pb$  secs


 QUESTION ANALYTICS
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Q. 25

PURE ALOHA and SLOTTED ALOHA are two multiple access protocols used to decrease collisions. The maximum efficiency obtained from PURE ALOHA and SLOTTED ALOHA can be written as  $a \times e^b$  and  $c \times e^d$ . What will be the value of  $8abcd$ ?

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- ☐ 2
- ☐ 3
- ☒ 4 Correct Option
- Solution :**  
 (c)  
 Maximum possible efficiency of PURE ALOHA =  $\frac{1}{2} \times e^{-1}$   
 Maximum possible efficiency of SLOTTED ALOHA =  $1 \times e^{-1}$   
 Therefore,  $a = \frac{1}{2}, b = -1, c = 1, d = -1$   

$$8 \times a \times b \times c \times d = 8 \times \frac{1}{2} \times (-1) \times 1 \times (-1)$$

$$= 4$$
- ☐ 5


 QUESTION ANALYTICS
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Q. 26

Consider the following:

Protocols	Facilities provided
A. TCP	(i) Flow control
B. UDP	(ii) Error control
C. Simple stop and wait	(iii) Congestion control

D. Stop and wait Automatic Repeat Request

Which of the following is a correct matching?

Codes:

(a) A-i, ii, iii, B-ii, iii, C-ii

(b) A-i, ii, B-ii, D-ii, iii

(c) A-i, iii, B-i, ii, iii, C-i, ii, iii

(d) A-i, ii, iii, B-ii, C-i, D-i, ii

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**A** a

**B** b

**C** c

**D** d

Correct Option

**Solution :**  
(d)

 QUESTION ANALYTICS

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**Q. 27**

Consider the following statements:

$S_1$  : Size of OPTIONS field in IPv4 header can be upto 40 bytes.

$S_2$  : "Version" field in IPv4 and IPv6 header contains the bit sequence 0100 and 0110 respectively.

Which of the option is correct?

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**A** Only  $S_1$  is true

**B** Only  $S_2$  is true

**C** Both  $S_1$  and  $S_2$  are true

Correct Option

**Solution :**  
(c)

$S_1$  : Size of OPTIONS field in IPv4 header can be upto 40 bytes and hence size of IPv4 header can be in between 20 bytes to 60 bytes.

$S_2$  : "Version" field in IPv4 and IPv6 header contains '0100' and '0110' respectively.

**D** Both  $S_1$  and  $S_2$  are false

 QUESTION ANALYTICS

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**Q. 28**

Consider sending a large file of F bits from Host A to Host B. There are two links (and one router) between A and B and the links are uncongested (that is, no queuing delays). Host A segments the file into segments of S bits each and adds 40 bits of header to each segment, forming packets of  $L = 40 + S$  bits. Each link has a transmission rate of R bps. What is the value of S that minimizes the delay of moving the file from Host A to Host B? (ignore the propagation delay)

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**A**  $(15F)^{-1/2}$

**B**  $(25F)^{-1/2}$

**C**  $(30F)^{-1/2}$

**D**  $(40F)^{-1/2}$

Correct Option

**Solution :**  
(d)

Number of packet required to send file of size F =  $\frac{\text{File size}}{\text{Size of segment}} = \frac{F}{S}$

Time at which the 1<sup>st</sup> packet is received at the destination =  $\frac{S+40}{R} \times 2$  sec. After this, one packet

is received by the destination every  $\frac{S+40}{R}$  sec because packets are transmitted back to back by Host A.

Thus delay in sending the whole file is,

$$\text{delay} = \frac{S+40}{R} \times 2 + \left( \frac{F}{S} - 1 \right) \times \left( \frac{S+40}{R} \right) = \frac{S+40}{R} \times \left( \frac{F}{S} + 1 \right)$$

To calculate the value of S which leads to the minimum delay, we take the derivative and equate it to zero,

$$\frac{\partial \text{delay}}{\partial S} = 0 \Rightarrow \frac{F}{R} \left( \frac{1}{S} - \frac{40+S}{S^2} \right) + \frac{1}{R} = 0$$

$$\Rightarrow S = \sqrt{40F}$$

 QUESTION ANALYTICS

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**Q. 29**

Consider TCP implements an extension that allows window sizes much larger than 64 KB. Suppose that you are using this extended TCP over a 1 Gbps link with a latency of 50 ms to transfer a 10 MB file and the TCP receive window is 1 MB. If TCP sends 1 KB packets and time to send the file is given by the number of required RTTs multiplied by the link latency, then the effective throughput for the transfer is \_\_\_\_\_ Mbps (assuming no congestion and no lost packets). [Upto 1 decimal place]



14.3 (14.2 - 14.4)

Correct Option

**Solution :**

14.3 (14.2 - 14.4)

In slow start, the size of the window doubles every RTT.

At the end of the  $i^{\text{th}}$  RTT, the window size is  $2^i$  KB.It will take 10 RTTs before the send window has reached  $2^{10}$  KB = 1 MB

After 10 RTTs, data transferred is = 1024 KB = 1 MB

The window size = 1 MB

It takes 4 more RTTs to transfer the remaining 9 MB.

Therefore, the file is transferred in 14 RTTs.

Transfer time = 50 ms

So for sending a full file it take =  $14 \times 0.05 = 0.7$  seconds to send the fileThe effective throughput is  $\left( \frac{10\text{MB}}{0.7\text{s}} \right) = 14.3 \text{ Mbps}$ 

QUESTION ANALYTICS

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**Q. 30**

Consider a TCP message that contains 1024 bytes of data and 20 bytes of TCP header is passed to IP for delivery across two networks interconnected by a router (i.e., it travels from the source host to a router to the destination host). The first network has an MTU of 1024 bytes and the second has an MTU of 576 bytes. If all packets are correctly delivered, the number of bytes, including headers, are delivered to the IP layer at the destination for TCP message, in the best

case is \_\_\_\_\_ bytes. (Assume all IP headers are 20 bytes)

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1104

Correct Option

**Solution :**

1104

**First network:** An MTU of 1024 means that is the largest IP datagram that can be carried, so a datagram has room for  $1024 - 20 = 1004$  bytes of IP-level data, because 1004 is not a multiple of 8,

each fragment can contain at most  $8 \times \left\lfloor \frac{1004}{8} \right\rfloor = 1000$  bytes.

We need to transfer  $1024 + 20 = 1044$  bytes of data (TCP header is included).

This would be fragmented into fragments of size 1000 bytes, and 44 bytes.

**Second network:** The 44 byte packet would be unfragmented but the 1000-data-byte packet would be fragmented as follows. The 576 byte MTU allows for up to  $576 - 20 = 556$  bytes of payload.

So rounding down to a multiple of 8 again allows = 552 bytes in the first fragment.

And remaining 448 bytes in the second fragment.

So total bytes =  $552 + 20 + 448 + 20 + 44 + 20$   
= 1104 bytes

QUESTION ANALYTICS

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**Q. 31**

Consider a router is blasting out IP packets whose total length is 1024 bytes. If packets live for 10 seconds, then the maximum line speed the router can operate at without danger of cycling through the IP datagram ID number space is \_\_\_\_\_ in Mbps. (Upto 1 decimal place)

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53.7 (53.6 - 53.8)

Correct Option

**Solution :**

53.7 (53.6 - 53.8)

The IP datagram ID number space is  $2^{16} = 65535$

Packet lifetime = 10 seconds

Therefore, a maximum of 65535 packets may be sent in 10 seconds. If any more packets were to be sent within the 10 seconds, the same id will be wrap around

$\text{floor} \left( \frac{65535}{10} \right) = 6553 \text{ packets/sec}$

Max line speed at 1024 bytes/packet is

=  $1024 \text{ bytes/packet} \times 6553 \text{ packets/sec} \times 8 \text{ bits/byte}$

= 53682176 bps

= 53.7 Mbps

QUESTION ANALYTICS

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**Q. 32**

Consider two Hosts A and B are each connected to a router R via 10 Mbps links. The propagation delay on each link is 20 microseconds. R is a store and forward device i.e transmission of the packet on the R-B link begin only if whole packet is received on the A-R link. Suppose R forwards a packet 35 microseconds (processing delay) after it has finished receiving it. The time saved when transmit 10000 bits from A to B as two 5000 bit packets sent one right after the other instead of as a single packet is \_\_\_\_\_ in  $\mu\text{s}$ .

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500

Correct Option

**Solution :**

500

Transmit delay of one link =  $10^4 \text{ bits}/10^7 \text{ (bits/sec)} = 1000 \mu\text{s}$

Transmission time for sending as single packet =  $2 \times 1000 + 2 \times 20 + 35 = 2075 \mu\text{s}$

When sending as 2 packets,

We have a total of one switch delay and two link delays;

Transmit delay =  $5000 \text{ bits}/10^7 \text{ (bits/sec)} = 500 \mu\text{s}$

Transmission time for sending as multiple packets =  $3 \times 500 + 2 \times 20 + 1 \times 35 = 1575 \mu\text{s}$


Time saved =  $(2075 - 1575) \mu\text{s} = 500 \mu\text{s}$

QUESTION ANALYTICS

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**Q. 33**

Consider a source using TCP-Reno (fast-retransmit and fast-recovery mechanisms) to send data to a destination. It is given the RTT of the link is 800 ms and the sender's window size is 8 segments. The sender sends segments at a regular rate of one every 100 ms, and the receiver sends ACKs back at the same rate without delay. A segment is lost, and the receiver sends 3 duplicate ACKs to trigger the fast-retransmit. If the sender waits for ACK of the retransmitted segment before advancing the window then the total time the sender lost (as compared to lossless transmission) is \_\_\_\_\_.

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 1100

Correct Option

**Solution :**

1100

Suppose packet number  $P$  is sent at  $t = 0$  by the sender's clock.

Suppose  $P$  is lost. So, 7 packets :  $P + 1$  to  $P + 7$  will be sent at  $t = 100, 200, \dots, 700$  respectively, since the window size is 8.

Sender stops sending at  $t = 800$ , since window size is 8 and ACK for  $P$  has not arrived.

ACK for  $P + 1$ ,  $P + 2$  and  $P + 3$  will be duplicate ACKs of  $P - 1$  (since 3 duplicate CK is used) and will arrive at  $t = 900, 1000$  and  $1100$ .

$P$  is retransmitted at  $t = 1100$  and ACK arrives at  $1100 + 800 = 1900$  ms

If no lost of segment then time = 800 ms

$$\text{Time lost} = 1900 - 800 = 1100$$

 QUESTION ANALYTICS

