CS 779 : Tutorial solutions

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§1. Tutorial 1

Exercise 1.1. Prove that the maximum number of subsets of $[n]$ with pairwise non-empty intersection is 2^{n-1}.

Solution

2^{n-1} is clearly attainable by taking $\{S \subseteq [n] : 1 \in S\}$. Furthermore, this is an upper bound since if $\mathcal{S}$ is a family of subsets with pairwise non-empty intersection, then $\mathcal{S}' = \{S^c : S \in \mathcal{S}\}$ has zero intersection with $\mathcal{S}$ and is of the same size, so $2|\mathcal{S}| = |\mathcal{S}'| + |\mathcal{S}| \leq 2^n$.

Exercise 1.2. Suppose you have a set system with m sets $(A_i)_{i=1}^m$ such that $|A_i|$ is odd for each i and $|A_i \cap A_j|$ is even for any $i \neq j$. Prove that $m \leq n$.

Solution

Consider the $m \times n$ matrix M where M_{ij} is 1 if $j \in A_i$ and is 0 otherwise. Then,

$$(MM^\top)_{ij} = \sum_{k \in [n]} M_{ik} M_{jk} = |A_i \cap A_j|.$$

In particular, all the diagonal entries of $MM^\top$ are odd and all off-diagonal entries are even. Using this, it is not too difficult to show that $\det(MM^\top) \neq 0$ (for an easy solution* of this, note that modulo 2, $MM^\top$ is congruent to the identity, which has nonzero determinant). Therefore, $m = \operatorname{rank}(MM^\top) = \operatorname{rank}(M)$, so $m \leq n$.

Exercise 1.3. Prove that for matrices A, B, $\operatorname{rank}(A + B) \leq \operatorname{rank}(A) + \operatorname{rank}(B)$.

Solution

It suffices to show that any column of $A + B$ is present in the space spanned by the column of A and B. This is straightforward since any column of $A + B$ is just the sum of the two corresponding columns in A and B.

Exercise 1.4. Suppose you have $A + A^\top = J - I$, where J is the all ones matrix. Prove that $\operatorname{rank}(A) \geq n/2$.

Solution

Using the previous exercise, we have $n = \operatorname{rank}(J - I) = \operatorname{rank}(A + A^\top) \leq \operatorname{rank}(A) + \operatorname{rank}(A^\top) = 2\operatorname{rank}(A)$.

Exercise 1.5. Suppose you have $A + A^\top = J - I$, where J is the all ones matrix. Show that if $\operatorname{rank}(A) < n - 1$, there is a vector x such that $Ax = 0$, $x \neq 0$, and $\mathbf{1}^\top x = 0$. Using this, prove that $\operatorname{rank}(A) \geq n - 1$.

Solution

Suppose $\operatorname{rank}(A) < n - 1$. Then, $\dim \ker A \geq 2$. We also have $\dim \mathbf{1}^\perp = n - 1$. Therefore, $\ker A$ and $\mathbf{1}^\perp$ have nonzero intersection, and say $x \neq 0$ is in both. x satisfies the conditions mentioned in the question. Now,

$$\begin{aligned} 0 &= x^\top (Ax) + (x^\top A^\top) x \\ &= x^\top (J - I) x \\ &= \left(\sum_i x_i \right)^2 - \left(\sum_i x_i^2 \right) = -\sum_i x_i^2, \end{aligned}$$

so $x = 0$, a contradiction. Therefore, $\operatorname{rank}(A) \geq n - 1$.

Exercise 1.6. Suppose $B_1, \ldots, B_m$ are complete bipartite graphs whose edge disjoint union yields the complete graph K_n. Show that $m \geq n - 1$.

Solution

Suppose that B_i corresponds to the complete bipartite graph between sets $X_i, Y_i \subseteq [n]$, where $X_i \cap Y_i = \varnothing$. As a graph on vertex set $[n]$, on setting $M_i = \mathbb{1}_X \mathbb{1}_Y^\top$, B_i has adjacency matrix $M_i + M_i^\top$. Note that $\operatorname{rank}(M_i) = 1$ for all i, since $\mathbb{1}_Y^\perp \subseteq \ker M_i$. Because the edge disjoint union of the B_i is K_n, we have $(\sum_i M_i) + (\sum_i M_i)^\top = J - I$. Using the previous exercise, $\operatorname{rank}(\sum_i M_i) \geq n - 1$. Using Exercise 1.3 and the observation that $\operatorname{rank}(M_i) = 1$ for all i, this implies that $m = \sum_{i=1}^m \operatorname{rank}(M_i) \geq n - 1$, completing the proof.

Exercise 1.7. Suppose you have a set system of m sets such that for every pair of sets, the intersection size is fixed as $\lambda \geq 1$. Prove that $m \leq n$.

Solution

Let the set system be $(A_i)_{i=1}^m$. The size of at most one set is equal to λ. Furthermore, if $|A_1| = \lambda$, then $A_i \setminus A_1$ are disjoint for distinct i, so $m - 1 \leq n - \lambda$. Thus, we may assume that the size of every set is greater than λ.
Define the matrix M exactly as in Exercise 1.2. We have that the off-diagonal entries of M are equal to λ. Now, $MM^\top = \lambda J + D$, for some diagonal matrix D with all positive diagonal entries. We wish to show that $\operatorname{rank}(\lambda J + D) = m$. Let $x \neq 0$ in $\mathbb{R}^n$, and let u, v be the components of x along and orthogonal to $\mathbf{1}$ respectively, such that $u = t\mathbf{1}$. Then,

$$\begin{aligned}(\lambda J + D)x &= (\lambda J + D)(u + v) \\ &= n\lambda u + D(u + v) \\ &= D(D^{-1}n\lambda u + u + v).\end{aligned}$$

When $t = 0$, this is clearly nonzero as $v \neq 0$. Otherwise, to conclude, note that

$$\sum_i (D^{-1}n\lambda u + u + v)_i = \sum_i (D_{ii}^{-1}n\lambda + 1)u_i + v_i = \sum_i t(D_{ii}^{-1}n\lambda + 1),$$

which is nonzero as $d_{ii}, \lambda > 0$ and $t \neq 0$.

§2. Tutorial 2

Exercise 2.1. Find the dimension of the space spanned by the following polynomials over the given field.

(a) $x_1, x_2, x_1x_2, x_1^2x_2, 1, (x_1 + x_2)^2, x_1^2 + x_2^2$ over $\mathbb{R}$ and over $\mathbb{F}_2$.

(b) $x_1^{i_1} x_2^{i_2} \cdots x_n^{i_n}$, where $i_1 + \cdots + i_n = m$ over $\mathbb{R}$ and over $\mathbb{F}_2$.

Solution

(a) Over $\mathbb{R}$, it is clear that $(x_1 + x_2)^2 = (x_1^2 + x_2^2) + 2(x_1x_2)$, and the collection formed by removing $(x_1 + x_2)^2$ is linearly independent, so the dimension of the space is 6.
We can again remove $(x_1 + x_2)^2$ in the $\mathbb{F}_2$ case. Note that $x^2 = x$ in $\mathbb{F}_2$. Therefore, over $\mathbb{F}_2$, the set of polynomials is $x_1, x_2, x_1x_2, x_1x_2, 1, x_1 + x_2$; it is clear a maximal linearly independent subset of these vectors is $x_1, x_2, x_1x_2, 1$, so the dimension is 4.

(b) Over $\mathbb{R}$, all these monomials are linearly independent, so the dimension is the number of ways of choosing n non-negative numbers that sum to m. This is a routine exercise in combinatorics, with the answer being $\binom{m+n-1}{m}$.
Over $\mathbb{F}_2$, since $x_1 = x_1^{i_1}$ for any $i_1 \geq 1$, we only care whether each i_r is zero or not. Let j_r be 1 if $i_r \geq 1$ and 0 otherwise. We then want to find the number of ways of choosing $(j_r)_{r=1}^n$ such that $0 < \sum_{r=1}^n j_r \leq m$. This is precisely $\binom{n}{1} + \binom{n}{2} + \cdots + \binom{n}{d}$, where $d = \min\{n, m\}$.

Exercise 2.2. Given m sets with sizes greater than d and pairwise intersection d, prove that $m \leq (n + 1)$.
Hint. Associate a polynomial to each set so that the polynomials are linearly independent. Give an upper bound on the space spanned by these polynomials.

Solution

Exercise 2.3.

(a) How do we define the distance between a pair of points in $\mathbb{R}^n$?

(b) Construct as many points as you can so that the distance between a pair is one of two distances, either d_1 or d_2. You may also choose d_1 and d_2 to maximize the number.

(c) Consider m points with exactly two pairwise distances. Associate polynomials $p_i(x)$ to each point such that the polynomials are linearly independent.

(d) Deduce an upper bound on the dimension of the span of your polynomials. What does this imply about the number of points with exactly two pairwise distances?

Solution

Exercise 2.4. A polynomial is called multilinear if the degree of each variable is at most one. What is the dimension of the space of multilinear polynomials of degree at most d over n variables?

Solution

The solution to this is near-identical to the second part of Exercise 2.1(b), with the answer being $\binom{n}{0} + \binom{n}{1} + \cdots + \binom{n}{d}$.

Exercise 2.5. Consider m sets $A_1, \ldots, A_m$ such that $|A_i| \equiv k \pmod{p}$ for some prime p. Assume that $|A_i \cap A_j| \in L \pmod{p}$ for some set L, such that $k \notin L$ and $|L| = \ell$. Show that $m \leq \binom{n}{0} + \cdots + \binom{n}{\ell}$.

Solution