BAGCHI'S CONJECTURE

Amit Rajaraman

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§1. Definitions

Definition 1.1 (Incidence System). An *incidence system* is a pair $(\mathcal{P}, \mathcal{L})$, where \mathcal{P} is a set and \mathcal{L} is a set of subsets of \mathcal{P} . Elements of \mathcal{P} are called *points* and elements of \mathcal{L} are called *lines*. A line ℓ is said to be *incident* on a point p if $p \in \ell$.

Definition 1.2 (Partial Linear Space). An incidence system $(\mathcal{P}, \mathcal{L})$ is said to be a partial linear space if

- 1. for each $\ell \in \mathcal{L}$, $|\ell| \geq 2$.
- 2. for distinct $x, y \in \mathcal{P}$, there is at most one $\ell \in \mathcal{L}$ such that $\{x, y\} \subseteq \mathcal{P}$.

An incidence system is said to be a *linear space* if in the second condition above, the "at most one" is replaced with "exactly one".

Definition 1.3. Given an incidence system $\mathcal{X} = (\mathcal{P}, \mathcal{L})$ and a field \mathbb{F} , we define the linear code $\mathcal{C}_{\mathbb{F}}(\mathcal{X})$ over $\mathbb{F}^{\mathcal{P}}$ as follows. Identify each line ℓ with the codeword in $\mathbb{F}^{\mathcal{P}}$ whose xth coordinate is 1 if $x \in \ell$ and 0 otherwise. $\mathcal{C}_{\mathbb{F}}(\mathcal{X})$ is then the space spanned by the codewords corresponding to the lines in \mathcal{L} . If $\mathbb{F} = \mathbb{F}_q$, we sometimes denote the above as $\mathcal{C}_q(\mathcal{X})$.

We also often denote this as $C_q(\mathcal{L})$ if the point set is clear from context. The incidence system \mathcal{X} is said to be *trivial* at q if $C_q(\mathcal{X})$ is all of $\mathbb{F}_q^{\mathcal{P}}$. **Definition 1.4** (Join). Given two partial linear spaces $(\mathcal{P}_1, \mathcal{L}_1)$ and $(\mathcal{P}_2, \mathcal{L}_2)$ with $\mathcal{P}_1 \cap \mathcal{P}_2 = \emptyset$, one can define the *join* of the two partial linear spaces by $(\mathcal{P}_1 \cup \mathcal{P}_2, \mathcal{L}_1 \cup \mathcal{L}_2 \cup \mathcal{L}_3)$, where

$$\mathcal{L}_3 = \{ \{x, y\} : x \in \mathcal{P}_1, y \in \mathcal{P}_2 \}.$$

§2. Inamdar's Result

Theorem 2.1. If a PLS $\mathcal{X} = (\mathcal{P}, \mathcal{L})$ is non-trivial at p and has at least p+1 lines through every point, then $|\mathcal{P}| \geq 2p$. Moreover, equality holds iff \mathcal{X} is the join of two p-lines.

2.1. Combinatorial Arguments

For the rest of this section, assume that $\mathcal{X} = (\mathcal{P}, \mathcal{L})$ is a partial linear space of the above prescribed format with $\mathcal{P} \leq 2p$. We wish to show that $\mathcal{P} = 2p$. Let $\mathcal{C} = \mathcal{C}_p(\mathcal{X})$.

It may be shown that it can be assumed that

- 1. Each point is incident on exactly p+1 lines (throw away extra lines).
- 2. \mathcal{C}^{\perp} is one-dimensional. Suppose that it is equal to $\langle w \rangle$ (restrict to the support of the minimum support word in \mathcal{C}^{\perp}).
- 3. For any line ℓ and any $\ell' \subsetneq \ell$, $\langle w, \ell' \rangle \neq 0$ (do the splitting procedure).

Also consider the colouring of \mathcal{P} wherein each point P is coloured w(P).

Proposition 2.2. If \mathcal{X} has a p-line, it is equal to a join of two p-lines.

Proof. Let $P_1 \cdots P_p$ be a p-line. Since each point has p lines remaining, there must be at least p points other than the P_i , say $(Q_i)_{i=1}^p$. Since $|\mathcal{P}| \leq 2p$, these constitute all the points. Further, since each P_i has p lines to the Q_j , there must be a 2-line P_iQ_j for each $1 \leq i, j \leq p$. Now, suppose that $w(P_1) = 1$. Because $w(P_1) + w(Q_j) = 0$ for all j (P_1Q_j forms a p-line), $w(Q_j) = -1$ for all j. Each of the Q_i now has one line not accounted for. This line must be contained within the (Q_i) . However, due to all of them having the same colour, the size of any such line must be p, completing the proof.

Proposition 2.3. If \mathcal{X} has no p-line, the largest line in \mathcal{X} is of size at most 2p/3.

Proof. Let ℓ be a line with $|\ell| > 2p/3$. Given a $P \in \ell$, let x_P be the number of 2-lines through P. We then have that

$$2p - |\ell| \ge \underbrace{x_P}_{\text{points in 2-lines}} + \underbrace{2(p - x_P)}_{\text{points in } \ge 3\text{-lines}},$$

so $x_P \ge \ell$. Observe that if $x_P > (2p - |\ell|)/2$ for all $P \in \ell$, it follows by a pigeonhole argument that any two points P, Q in ℓ have a "2-neighbour" (a point u such that uP and uQ are 2-lines) in common. This is indeed the case because $x_P \ge |\ell| > (2p/3) > (2p - |\ell|)/2$. This in turn implies that w(P) = w(Q), because w(P) = -w(u) = w(Q). Therefore, ℓ is monochromatic, so for any fixed $P \in \ell$,

$$0 = \sum_{P \in \ell} w(P) = |\ell| w(P).$$

As $w(P) \neq 0$, $|\ell| = p$, yielding a contradiction.

As a result, we are done if we manage to show that \mathcal{X} has a line of size greater than 2p/3.

Now, let us restrict ourselves to the case where \mathcal{X} is a linear space with exactly p+1 lines through each point.

2.2. Linear Algebra

Let C be the $m \times n$ incidence matrix of \mathcal{X} , where $m = |\mathcal{L}|$ and $n = |\mathcal{P}|$. Observe that $(C^tC)_{ij} = pI + J$, where J is the all 1s matrix. As a result, its characteristic equation is just $(X - (p+n))(X - p)^{n-1}$. Recall that the characteristic matrix of CC^t and C^tC only differ by a power of X. That is, the characteristic equation of CC^t is $X^{m-n}(X - (p+n))(X - p)^{n-1}$. Over \mathbb{F}_p , this becomes $X^{m-1}(X - n)$.

Suppose that y is an eigenvector with eigenvalue 0. That is, $CC^ty = 0$. A question to ask is: is C^ty always 0? If it is not, we can explicitly get a vector in C^{\perp} .

Proposition 2.4. There exists y in the kernel of CC^t such that $C^t y \neq 0$ (in \mathbb{F}_p).

Proof. Indeed, note that the kernel of CC^t is of dimension m-1, whereas the kernel of C^t is of dimension at most m-n. It follows that the kernel of C^t is a proper subset of CC^t .

Definition 2.1 (Irreducible Matrix). A matrix $A = (a_{ij})_{n \times n}$ is said to be *irreducible* if the (directed) graph G_A with vertex set [n] and edge from i to j iff $a_{ij} \neq 0$ is strongly connected.

Observe that CC^t is irreducible! Indeed, given any lines ℓ, ℓ' , there is a line ℓ'' through some (any) $x \in \ell$ and $y \in \ell'$.

Theorem 2.5 (Perron-Frobenius Theorem). Let $A = (a_{ij})_{n \times n}$ be a non-negative irreducible matrix. Then,

- 1. There exists a positive eigenvalue λ such that any other eigenvalue is of absolute value at most λ . Further, there is a number h called the *period* such that $\omega_h^k \lambda$ is an eigenvalue for every $0 \le k < h$. These are the only eigenvalues with absolute value λ .
- 2. λ is a simple root of the characteristic polynomial. As a result, its corresponding eigenspace is one-dimensional.
- 3. λ has an eigenvector v with all positive entries. Any eigenvector with all positive entries is in the span of v.

Since CC^t is irreducible and non-negative with maximum eigenvalue n + p, there is an eigenvector v of n + p with all positive entries.

For example, we now have that

$$m(n+p) = (n+p)\mathbf{1}^{\top}\mathbf{1}$$

$$\geq \mathbf{1}^{\top}CC^{\top}\mathbf{1}$$

$$= \sum_{\ell,\ell'\in\mathcal{L}} |\ell\cap\ell'|$$

$$= \sum_{\ell\in\mathcal{L}} (p+1)|\ell| \qquad \text{(each point on } \ell \text{ has } p+1 \text{ lines through it)}$$

$$= (p+1) \cdot n(p+1) = n(p+1)^2$$

That is,

$$m \ge \frac{n(p+1)^2}{n+p}.$$

The fact that $p + 2 < n \le 2p$ further gives that

$$m \ge \frac{(p+1)^2}{2}$$

The system has many lines! Can we use this to show that either there is a "large" line somewhere or the code is trivial?

Can we draw stronger conclusions by using another vector in place of 1? It seems that the ideal choice is to assign a larger weight to larger lines. What about the vector defined by $v_{\ell} = |\ell|$? Recall that

$$\sum_{\ell \in \mathcal{L}} |\ell| = n(p+1).$$

Proposition 2.6.

$$\sum_{\ell \in \mathcal{L}} |\ell|^2 = n(n+p).$$

Proof. Consider the vector v with $v_{\ell} = |\ell|$. Then,

$$\sum_{\ell \in \mathcal{L}} |\ell|^2 = v^t v$$

$$\geq \frac{1}{n+p} v^t C C^t v$$

$$= \frac{1}{n+p} \left(\sum_{\ell \in \mathcal{L}} \underbrace{v_\ell^2 |\ell|}_{|\ell|^3} + \sum_{\substack{\ell, \ell' \in \mathcal{L} \\ \ell \neq \ell' \\ \ell \cap \ell' \neq \varnothing}} v_\ell v_{\ell'} \right).$$

Let us compute the second quantity.

$$\begin{split} \sum_{\substack{\ell,\ell' \in \mathcal{L} \\ \ell \neq \ell' \\ \ell \cap \ell' \neq \varnothing}} v_{\ell} v_{\ell'} &= \sum_{P \in \mathcal{P}} \sum_{\substack{\ell,\ell' \ni P \\ \ell \neq \ell'}} |\ell| |\ell'| \\ &= \sum_{P \in \mathcal{P}} \left(\sum_{\substack{\ell,\ell' \ni P \\ \ell \ni P}} |\ell| |\ell'| - \sum_{\substack{\ell \ni P \\ \ell \ni P}} |\ell|^2 \right) \\ &= \sum_{P \in \mathcal{P}} \left(\sum_{\substack{\ell \ni P \\ \ell \ni P}} |\ell| \right)^2 - \sum_{P \in \mathcal{P}} \sum_{\substack{\ell \ni P \\ \ell \ni P}} |\ell|^2. \end{split}$$

The second summation is just

$$\sum_{P \in \mathcal{P}} \sum_{\ell \ni P} |\ell|^2 = \sum_{\ell \in \mathcal{L}} \sum_{P \in \ell} |\ell|^2 = \sum_{\ell \in \mathcal{L}} |\ell|^3.$$

The $\sum |\ell|^3$ terms cancel out! As a result,

$$\sum_{\ell \in \mathcal{L}} |\ell|^2 \ge \frac{1}{n+p} \sum_{P \in \mathcal{P}} \left(\sum_{\ell \ni P} |\ell| \right)^2.$$

On one hand, linearity and regularity imply that for any P, $\sum_{\ell \ni P} |\ell| = (n-1) + (p+1)$. Indeed, there is precisely one line through P and any point $Q \ne P$, and there are p+1 lines in all (in which P is present). Therefore,

$$\sum_{\ell \ni f} |\ell|^2 \ge \frac{1}{n+p} \sum_{P \in \mathcal{P}} (n+p)^2 = n(n+p).$$

Using the Cauchy-Schwarz inequality on the other hand,

$$\sum_{\ell \in \mathcal{L}} |\ell|^2 \ge \frac{1}{n+p} \sum_{P \in \mathcal{P}} \left(\sum_{\ell \ni P} |\ell| \right)^2$$

$$\ge \frac{1}{n(n+p)} \left(\sum_{P \in \mathcal{P}} \sum_{\ell \ni P} |\ell| \right)^2$$

$$= \frac{1}{n(n+p)} \left(\sum_{\ell \in \mathcal{L}} |\ell|^2 \right)^2$$

$$n(n+p) \ge \sum_{\ell \in \mathcal{L}} |\ell|^2,$$

completing the proof.

In fact, one similarly gets that for any m,

$$\sum_{\ell \in \mathcal{L}} |\ell|^{2m} \ge \frac{1}{n+p} \sum_{P \in \mathcal{P}} \left(\sum_{\ell \ni P} |\ell|^m \right)^2.$$

The right is easily computable for m=0 (because of regularity) and m=1 (because of linearity and regularity). Using the Cauchy-Schwarz inequality as we did in the previous proof, one further obtains that

$$\sum_{\ell \in \mathcal{L}} |\ell|^{2m} \ge \frac{1}{n(n+p)} \left(\sum_{\ell \in \mathcal{L}} |\ell|^{m+1} \right)^2.$$

Porism 2.7. The eigenvector (up to scalar multiplication) of CC^t corresponding to the eigenvalue n+p is the vector $v_{\mathcal{L}}: \mathcal{L} \to \mathbb{R}$ with $v_{\mathcal{L}}(\ell) = |\ell|$.

Thus, for any vector $x \in v_L^{\perp}$,

$$\left\| C^{\top} x \right\|_{2}^{2} \leq p \|x\|_{2}^{2}$$
.

2.3. Graphs

Consider the graph G with vertex set \mathcal{L} , and $u \leftrightarrow v$ iff $u \cap v \neq \emptyset$. What do our conditions on the linear space correspond to?

- 1. Each point has exactly p+1 lines through it: any maximal clique in G is of size p+1.
- 2. There is a line passing through any two points (linearity): any two maximal cliques have a single point in common.

We wish to count the number of (maximal) cliques in the graph.

Note that each vertex ℓ has degree $|\ell|$. As a result, the graph has $\frac{1}{2}n(p+1)$ edges. Consider the random variable X that chooses a vertex with probability proportional to its degree (the probability of choosing ℓ is $|\ell|/n(p+1)$). Note that this corresponds to the stationary distribution of a standard random walk on the graph! Further, we have that

$$\mathbb{E}d(v) = \frac{p+n}{p+1}.$$

Let D be the $m \times m$ diagonal matrix with $D_{ii} = |\ell_i|$. Let A_G and L_G be the adjacency matrix and ??? of G respectively. Then, observe that

$$L_G = pD - A_G$$
 and $CC^t = D + A_G$.