

# Searching

Linear Search

Binary Search

Ternary Search

Basic

non-sorted

Real no  $\Leftarrow$

Binary search on answers

mini-max

Binary search  
on graph

$$T(n) = \underbrace{n - n - 1 - n - 2}_{\rightarrow}$$

$$\begin{aligned} T(n) &= T(n-1) + O(1) \rightarrow \text{compariso} \\ T(n) &= \underline{O(n)} \end{aligned}$$

Binary Search  $\rightarrow$

Binary search says that if you have a search space

of size  $n$  then we can divide the space in  $n/2$  - by discarding

1 half

$$\underline{\left[ \begin{array}{c} p_1 \end{array} \right] \quad \left[ \begin{array}{c} p_2 \end{array} \right]}$$

$$T(n) = T\left(\frac{n}{2}\right) + O(1) \rightarrow \text{comparison}$$

$$T(n) = T(n/2) + O(1)$$

$$T(n/2) = T(n/4) + O(1)$$

$$T(n/4) = T(\cdot) + O(1)$$

$$\vdots$$

$$T(n) = T(1) + O(1)$$

$$n \rightarrow \frac{n}{2} \rightarrow \frac{n}{4} \rightarrow \frac{n}{8} \dots$$

$$\frac{n}{2^k} = 1 \quad n = 2^k$$

$$k = \log_2 n$$

Sorted

Ques → find the first index that is greater or equal target?

ex →

|    |    |    |    |    |    |    |
|----|----|----|----|----|----|----|
| 0  | 1  | 2  | 3  | 4  | 5  | 6  |
| 10 | 20 | 30 | 30 | 70 | 40 | 50 |

target = 30

ans = 2

→ lower bound

→ upper bound

target = 6

|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| 2 | 3 | 5 | 6 | 6 | 6 | 8 |
| F | F | F | T | T | T | T |

$n/2$

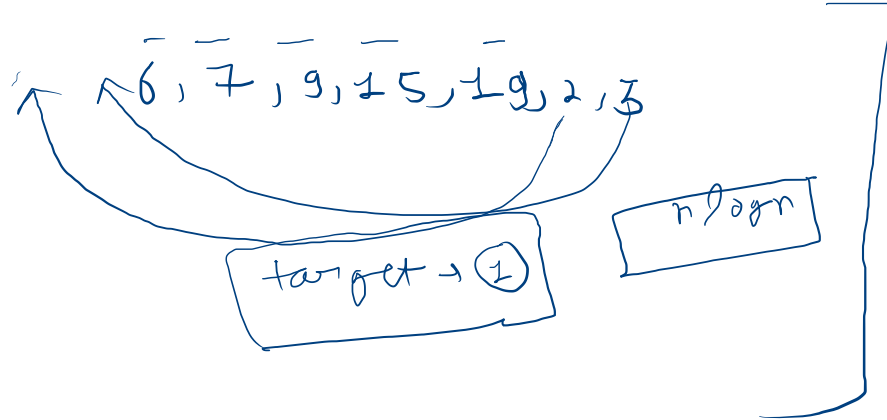
$$lo + \frac{(hi - lo)}{2}$$

$$\frac{lo + hi}{2} = \frac{lo + hi + lo - lo}{2}$$

$$= \frac{2lo}{2} + \frac{hi - lo}{2} =$$

$$lo + \frac{(hi - lo)}{2}$$

Ques → rotated sorted array find the index of target



Ques → Given an array, find any element in array such that it follows the following condition

$$\begin{array}{l} a[i] > a[i+1] \quad (\text{if } i+1 \text{ exist}) \\ \text{and} \\ a[i] < a[i-1] \quad (\text{if } i-1 \text{ exist}) \end{array}$$

Ques 2 .

## 278. First Bad Version

Easy

👍 5278

🗨 1963

♡ Add to List

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Ques Given a 2D array, find the Peak element in 2D array  
(no adjacent element is equal)

### Example 1:

|    |    |    |    |
|----|----|----|----|
| -1 | -1 | -1 | -1 |
| -1 | 1  | 4  | -1 |
| -1 | 3  | 2  | -1 |
| -1 | -1 | -1 | -1 |

Input: mat = [[1,4],[3,2]]

Output: [0,1]

Explanation: Both 3 and 4 are peak elements so [1,0]  
and [0,1] are both acceptable answers.

Ques → Given a number  $n$ , ( $n \leq 10^5$ ) find the square root of  $n$  (only int value)

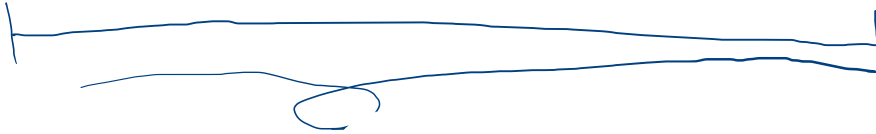
don't use inbuilt library

~~$O(n^2)$~~

$O(\log n)$

Binary search answer

$$\underline{n = 36 = 6}$$

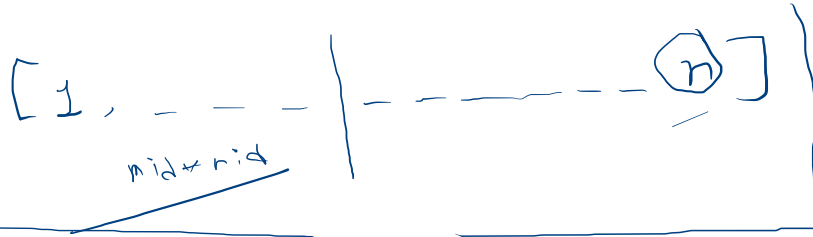


→ if the value of no.  $n$ , can i say the sort is  $\leq n$

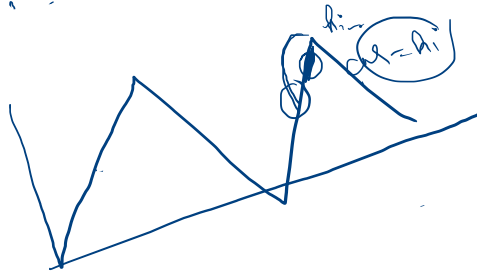
→ value of  $n$  is Positive

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space →



3/2



```
int lo=0;
int hi=n;
int ans=0;
while(lo<=hi){
    int mid=lo+(hi-lo)/2;
    if(mid*mid==n){
        return mid;
    }
    else if(mid*mid<n){
        //discard the left half
        lo=mid+1;
    }
    else{
        hi=mid-1;
        ans=hi;
    }
}
return ans;
```

| lo | hi | mid | ans |
|----|----|-----|-----|
| 0  |    | 1   | 6   |
|    |    |     | 7   |
|    |    |     | 3   |
|    |    |     |     |



There are n rectangle of same size ( $w \times h$ )

find a square of smallest size into which all of the  $n$  rectangle

can be placed (Rotation is not allowed)

$$(w, h, n) \leq 10^9$$

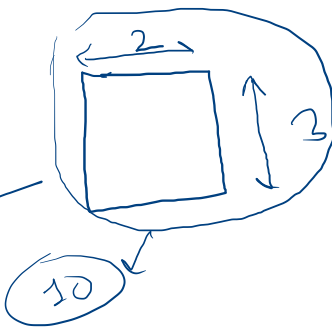
$$w=2 \quad h=3$$

$$n=10$$

$$ans = 9$$

$$w \neq h$$

$$n=10$$



$w, h \rightarrow$

$$3 \times 2$$

$$\begin{aligned} &10 \\ &= 6 \text{ area} \\ &\text{overall} = 60 \end{aligned}$$

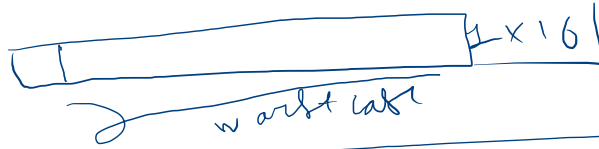
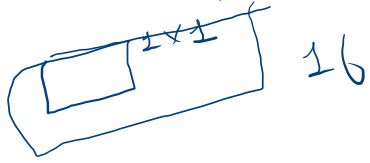
|   |    |   |   |  |
|---|----|---|---|--|
| 1 | 2  | 3 | 4 |  |
| 5 | 6  | 7 | 8 |  |
| 9 | 10 |   |   |  |

$$9 \times 9$$

|   |   |   |   |
|---|---|---|---|
| 1 | 2 | 3 | 4 |
|   |   |   |   |
|   |   |   |   |

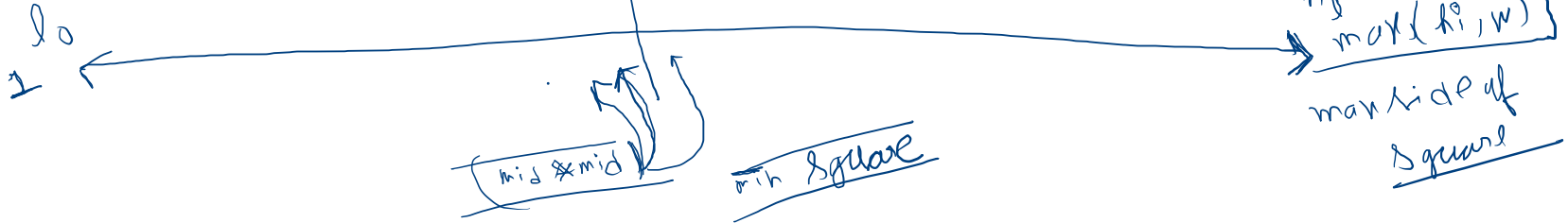
$$8 \times 8$$

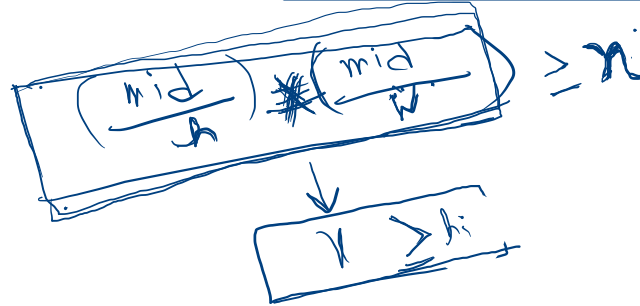
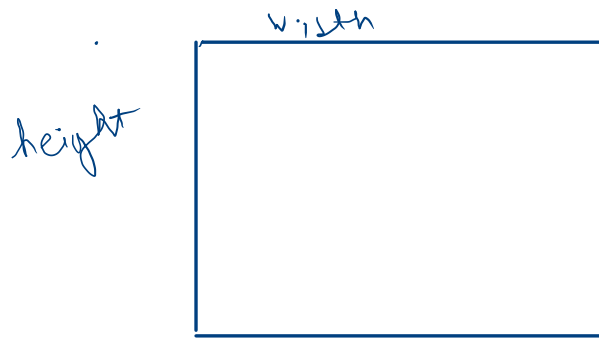
One of the sides of rectangle will cover a square side completely



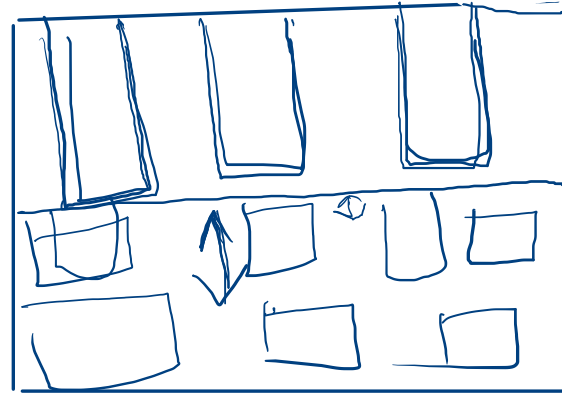
$\max(h, w)$

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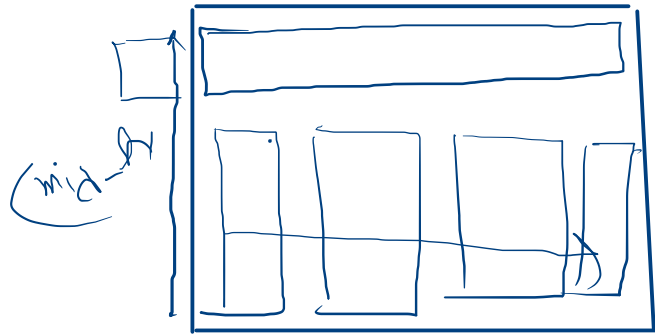
$\frac{(\text{mid} \times \text{mid})}{h \times w}$  rectangle can we fit  
 ① vertically



$\left(\frac{\text{mid}}{h}\right)$

horizontally  $\rightarrow \frac{(\text{mid})}{w}$

① horizontally



$$\frac{(mid)}{h}$$

$$\left( \frac{(mid)}{h} \times \frac{(mid)}{w} \right) \geq n$$

we can fit

n rectangles

② vertically

$$\frac{(mid)}{w}$$