

# Information & System Security

## Lecture 27



- >>Encryption
- >>Integrity
- >>Identification
- >>Authentication



**VIT-AP**  
**UNIVERSITY**

**Mathematics**  
**Related to**  
**Public Key**  
**Cryptography**

## 9-6 EXPONENTIATION AND LOGARITHM

**Exponentiation:**  $y = a^x \rightarrow$  **Logarithm:**  $x = \log_a y$

*Topics discussed in this section:*

**9.6.1 Exponentiation**

**9.6.2 Logarithm**

## 9.6.1 Exponentiation

$$y = a^x \bmod n$$

### Fast Exponentiation

$$y = a^{x_{n_b-1} \times 2^{n_b-1} + x_{n_b-2} \times 2^{n_b-2} + \dots + x_1 \times 2^1 + x_0 \times 2^0} \quad \text{in which } x_i \text{ is 0 or 1}$$

$$y = \boxed{a^{2^{n_b-1}} \text{ or } 1} \times \boxed{a^{2^{n_b-2}} \text{ or } 1} \times \dots \times \boxed{a^2 \text{ or } 1} \times \boxed{a \text{ or } 1}$$

Example:

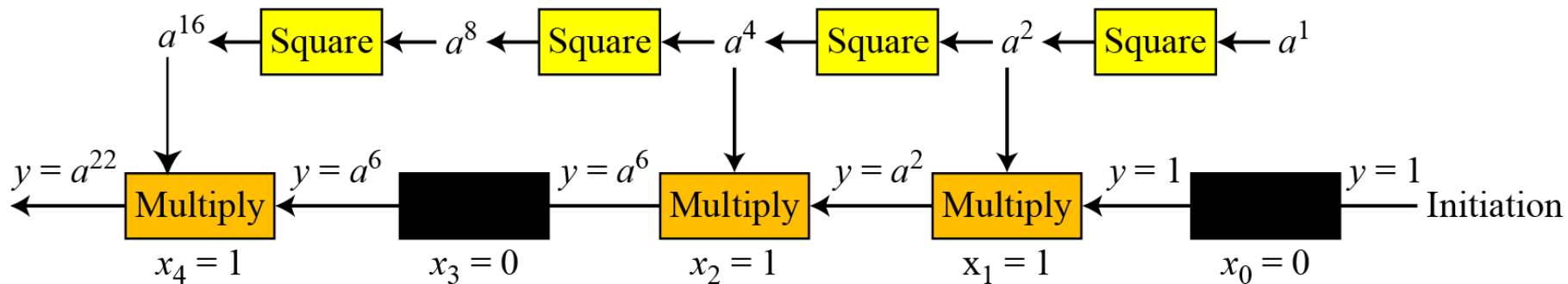
$$y = a^9 = a^{1001_2} = a^8 \times 1 \times 1 \times a$$

*The idea behind the **square-and-multiply** method*

## 9.6.1 Continued

### Example

Figure shows the process for calculating  $y = a^x$  using the Algorithm (for simplicity, the modulus is not shown). In this case,  $x = 22 = (10110)_2$  in binary. The exponent has five bits.



*Demonstration of calculation of  $a^{22}$  using square-and-multiply method*

## 9.6.1 Continued

Calculation of  $17^{22} \bmod 21$

$i$	$x_i$	Multiplication (Initialization: $y = 1$ )	Squaring (Initialization: $a = 17$ )
0	0	$\rightarrow$	$a = 17^2 \bmod 21 = 16$
1	1	$y = 1 \times 16 \bmod 21 = 16 \rightarrow$	$a = 16^2 \bmod 21 = 4$
2	1	$y = 16 \times 4 \bmod 21 = 1 \rightarrow$	$a = 4^2 \bmod 21 = 16$
3	0	$\rightarrow$	$a = 16^2 \bmod 21 = 4$
4	1	$y = 1 \times 4 \bmod 21 = 4 \rightarrow$	

**Note**

**The bit-operation complexity of the fast exponential algorithm is polynomial.**

## 9.6.2 Logarithm

*In cryptography, we also need to discuss modular logarithm.*

### *Exhaustive Search*

**Modular\_Logarithm** ( $a, y, n$ )

```
{  
    for ( $x = 1$  to  $n - 1$ ) //  $k$  is the number of bits in  $x$   
    {  
        if ( $y \equiv a^x \bmod n$ ) return  $x$   
    }  
    return failure  
}
```

## 9.6.2 Continued

### *Order of the Group*

#### Example

**What is the order of group  $G = \langle \mathbb{Z}_{21}^*, \times \rangle$ ?**

$$|G| = \phi(21) = \phi(3) \times \phi(7) = 2 \times 6 = 12.$$

**There are 12 elements in this group: 1, 2, 4, 5, 8, 10, 11, 13, 16, 17, 19, and 20.**

**All are relatively prime with 21.**



## 9.6.2 Continued

### *Order of an Element*

#### **Example**

**Find the order of all elements in  $G = \langle \mathbb{Z}_{10}^*, \times \rangle$ .**

#### **Solution**

**This group has only  $\phi(10) = 4$  elements: 1, 3, 7, 9. We can find the order of each element by trial and error.**

a.  $1^1 \equiv 1 \pmod{10} \rightarrow \text{ord}(1) = 1.$

b.  $3^4 \equiv 1 \pmod{10} \rightarrow \text{ord}(3) = 4.$

c.  $7^4 \equiv 1 \pmod{10} \rightarrow \text{ord}(7) = 4.$

d.  $9^2 \equiv 1 \pmod{10} \rightarrow \text{ord}(9) = 2.$

## 9.6.2 Continued

**Euler's Theorem:**

$$a^{\phi(n)} \equiv 1 \pmod{n}$$

### Example

The following Table shows the result of  $a^i \equiv x \pmod{8}$  for the group  $G = \langle \mathbb{Z}_8^*, \times \rangle$ .  $\phi(8)=4$ .

The elements are 1, 3, 5, and 7.

*Finding the orders of elements*

	$i = 1$	$i = 2$	$i = 3$	$i = 4$	$i = 5$	$i = 6$	$i = 7$
$a = 1$	<b>x: 1</b>	x: 1	x: 1	x: 1	x: 1	x: 1	x: 1
$a = 3$	x: 3	<b>x: 1</b>	x: 3	x: 1	x: 3	x: 1	x: 3
$a = 5$	x: 5	<b>x: 1</b>	x: 5	x: 1	x: 5	x: 1	x: 5
$a = 7$	x: 7	<b>x: 1</b>	x: 7	x: 1	x: 7	x: 1	x: 7

### *Primitive Roots*

*In the group  $G = \langle Z_n^*, \times \rangle$ , when the order of an element is the same as  $\phi(n)$ , that element is called the primitive root of the group.*

#### Example

The Table (previous) shows that there are no primitive roots in  $G = \langle Z_8^*, \times \rangle$  because **no element has the order equal to  $\phi(8) = 4$** . The order of elements are all smaller than 4.

## 9.6.2 Continued

### Example

The following Table shows the result of  $a^i \equiv x \pmod{7}$  for the group  $G = \langle \mathbb{Z}_7^*, \times \rangle$ . In this group,  $\phi(7) = 6$ .

		$i = 1$	$i = 2$	$i = 3$	$i = 4$	$i = 5$	$i = 6$
	$a = 1$	<b><math>x: 1</math></b>	$x: 1$	$x: 1$	$x: 1$	$x: 1$	$x: 1$
	$a = 2$	$x: 2$	$x: 4$	<b><math>x: 1</math></b>	$x: 2$	$x: 4$	$x: 1$
Primitive root $\rightarrow$	$a = 3$	$x: 3$	$x: 2$	$x: 6$	$x: 4$	$x: 5$	<b><math>x: 1</math></b>
	$a = 4$	$x: 4$	$x: 2$	<b><math>x: 1</math></b>	$x: 4$	$x: 2$	$x: 1$
Primitive root $\rightarrow$	$a = 5$	$x: 5$	$x: 4$	$x: 6$	$x: 2$	$x: 3$	<b><math>x: 1</math></b>
	$a = 6$	$x: 6$	<b><math>x: 1</math></b>	$x: 6$	$x: 1$	$x: 6$	$x: 1$

## 9.6.2 Continued

### Note

The group  $G = \langle Z_n^*, \times \rangle$  has primitive roots only if  $n$  is 2, 4,  $p^t$ , or  $2p^t$ .

### Example

For which value of  $n$ , does the group  $G = \langle Z_n^*, \times \rangle$  have primitive roots: 17, 20, 38, and 50?

### Solution

- a.  $G = \langle Z_{17}^*, \times \rangle$  has primitive roots, 17 is a prime.
- b.  $G = \langle Z_{20}^*, \times \rangle$  has no primitive roots.
- c.  $G = \langle Z_{38}^*, \times \rangle$  has primitive roots,  $38 = 2 \times 19$  prime.
- d.  $G = \langle Z_{50}^*, \times \rangle$  has primitive roots,  $50 = 2 \times 5^2$  and 5 is a prime.

## 9.6.2 Continued

*Note*

If the group  $G = \langle Z_n^*, \times \rangle$  has any primitive root, the number of primitive roots is  $\phi(\phi(n))$ .

### Example

Find number of primitive roots of 7 and 19.

### Solution

(a)  $\phi(\phi(7)) = \phi(7-1) = \phi(6) = \phi(2 \times 3) = \phi(2) \times \phi(3) = 2$ .

**7 has 2 primitive roots— 3 and 5.**

(b) **Try to find  $\phi(\phi(19))$ .**

## 9.6.2 Continued

### Powers of Integers, Modulo 19

$a$	$a^2$	$a^3$	$a^4$	$a^5$	$a^6$	$a^7$	$a^8$	$a^9$	$a^{10}$	$a^{11}$	$a^{12}$	$a^{13}$	$a^{14}$	$a^{15}$	$a^{16}$	$a^{17}$	$a^{18}$
1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
2	4	8	16	13	7	14	9	18	17	15	11	3	6	12	5	10	1
3	9	8	5	15	7	2	6	18	16	10	11	14	4	12	17	13	1
4	16	7	9	17	11	6	5	1	4	16	7	9	17	11	6	5	1
5	6	11	17	9	7	16	4	1	5	6	11	17	9	7	16	4	1
6	17	7	4	5	11	9	16	1	6	17	7	4	5	11	9	16	1
7	11	1	7	11	1	7	11	1	7	11	1	7	11	1	7	11	1
8	7	18	11	12	1	8	7	18	11	12	1	8	7	18	11	12	1
9	5	7	6	16	11	4	17	1	9	5	7	6	16	11	4	17	1
10	5	12	6	3	11	15	17	18	9	14	7	13	16	8	4	2	1
11	7	1	11	7	1	11	7	1	11	7	1	11	7	1	11	7	1
12	11	18	7	8	1	12	11	18	7	8	1	12	11	18	7	8	1
13	17	12	4	14	11	10	16	18	6	2	7	15	5	8	9	3	1
14	6	8	17	10	7	3	4	18	5	13	11	2	9	12	16	15	1
15	16	12	9	2	11	13	5	18	4	3	7	10	17	8	6	14	1
16	9	11	5	4	7	17	6	1	16	9	11	5	4	7	17	6	1
17	4	11	16	6	7	5	9	1	17	4	11	16	6	7	5	9	1
18	1	18	1	18	1	18	1	18	1	18	1	18	1	18	1	18	1

**Primitive roots of 19 are {2,3,10, 13,14,15}**

## 9.6.2 Continued

**Cyclic Group:** If  $g$  is a primitive root in the group, we can generate the set  $Z_n^*$  as  $Z_n^* = \{g^1, g^2, g^3, \dots, g^{\phi(n)}\}$

### Example

The group  $G = \langle Z_{10}^*, \times \rangle$  has two primitive roots because  $\phi(10) = 4$  and  $\phi(\phi(10)) = 2$ . It can be found that the primitive roots are 3 and 7. The following shows how we can create the whole set  $Z_{10}^*$  using each primitive root.

$$\begin{array}{llll} g = 3 & \rightarrow & g^1 \bmod 10 = 3 & g^2 \bmod 10 = 9 & g^3 \bmod 10 = 7 & g^4 \bmod 10 = 1 \\ g = 7 & \rightarrow & g^1 \bmod 10 = 7 & g^2 \bmod 10 = 9 & g^3 \bmod 10 = 3 & g^4 \bmod 10 = 1 \end{array}$$

The group  $G = \langle Z_n^*, \times \rangle$  is a cyclic group if it has primitive roots.  
The group  $G = \langle Z_p^*, \times \rangle$  is always cyclic.



## 9.6.2 Continued

### *The idea of Discrete Logarithm*

*Properties of  $G = \langle \mathbb{Z}_p^*, \times \rangle$ :*

- 1. Its elements include all integers from 1 to  $p - 1$ .*
- 2. It always has primitive roots.*
- 3. It is **cyclic**. The elements can be created using  $g^x$  where  $x$  is an integer from 1 to  $\phi(n) = p - 1$ .*
- 4. The primitive roots can be thought as the base of logarithm.*

If the group has  $k$  primitive roots, calculations can be done in  $k$  different bases. Given  $x = \log_g y$  for any element  $y$  in the set, there is an  $x$  that is the log of  $y$  in base  $g$ . This type of logarithm is called **discrete logarithm**.

## 9.6.2 Continued

### *Solution to Modular Logarithm using Discrete Logs*

Solve  $y = a^x \pmod n$  when  $y$  is given, and we need to find  $x$ .

### *Tabulation of Discrete Logarithms*

- One way to solve the above-mentioned problem is to use a table for each  $\mathbf{Z}_p^*$  and different bases.
- This type of table can be precalculated and saved.

*Discrete logarithm for  $\mathbf{G} = \langle \mathbf{Z}_7^*, \times \rangle$*

$y$	1	2	3	4	5	6
$x = L_3 y$	6	2	1	4	5	3
$x = L_5 y$	6	4	5	2	1	3

### Example

**Find  $x$  in each of the following cases:**

**1.**  $4 \equiv 3^x \pmod{7}.$

**2.**  $6 \equiv 5^x \pmod{7}.$

### *Solution*

**We can easily use the tabulation of the discrete logarithm in the previous Table.**

**1.**  $4 \equiv 3^x \pmod{7} \rightarrow x = L_3 4 \pmod{7} = 4 \pmod{7}$

**2.**  $6 \equiv 5^x \pmod{7} \rightarrow x = L_5 6 \pmod{7} = 3 \pmod{7}$

## 9.6.2 Continued

### *Using Properties of Discrete Logarithms*

*Comparison of traditional and discrete logarithms*

<i>Traditional Logarithm</i>	<i>Discrete Logarithms</i>
$\log_a 1 = 0$	$L_g 1 \equiv 0 \pmod{\phi(n)}$
$\log_a (x \times y) = \log_a x + \log_a y$	$L_g(x \times y) \equiv (L_g x + L_g y) \pmod{\phi(n)}$
$\log_a x^k = k \times \log_a x$	$L_g x^k \equiv k \times L_g x \pmod{\phi(n)}$

### *Using Algorithms Based on Discrete Logarithms*

- Cannot be used if  $n$  is very large.

**Note**

**The discrete logarithm problem has the same complexity as the factorization problem.**



## *References*

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- **Chapter 9** - Behrouz A Forouzan, Debdeep Mukhopadhyay, Cryptography and Network Security, Mc Graw Hill, 3rd Edition, 2015.
- **Chapter 8** - William Stallings, Cryptography and Network Security Principles and Practices, 7th Edition, Pearson Education, 2017.