

Problem:

Job-shop scheduling problem in a setting where **m** machines are separated by **(m-1)** bounded buffers. The objective is to schedule the tasks of **n** available jobs such that the total makespan of the jobs is minimum.

Every job goes through every machine.

The tasks in a job are constrained to be completed in a certain order.

Every task of a job takes a fixed amount of duration.

Every machine can carry out only one task at a time.

The buffers do not maintain the FIFO order of the jobs.

The machines and jobs are as follows –

	M₁	M₂	...	M_m
J₁	d ₁₁	d ₁₂	...	d _{1m}
J₂	d ₂₁	d ₂₂	...	d _{2m}
...
J_n	d _{n1}	d _{n2}	...	d _{nm}

And the buffers are bounded by the maximum capacities $b_1, b_2, b_3, \dots, b_{(m-1)}$.

Solution:

Variables –

1. $t_{11}, t_{12}, \dots, t_{nm}$

Start times of the tasks

$O(nm)$

2. $E_{11}, E_{12}, \dots, E_{1(m-1)}, \dots, E_{n(m-1)}$

Enqueue operation times for the tasks

$O(nm)$

3. $D_{12}, D_{13}, \dots, D_{1m}, \dots, D_{nm}$

De-queue operation times for the tasks

$O(nm)$

4. $B_{11}@E_{11}, B_{12}@E_{11}, \dots, B_{1(m-1)}@E_{11}, B_{21}@E_{11}, \dots, B_{n(m-1)}@E_{11},$
 $B_{11}@E_{12}, \dots, B_{1(m-1)}@E_{12}, B_{21}@E_{12}, \dots, B_{n(m-1)}@E_{12},$

...

$B_{11}@E_{n(m-1)}, \dots, B_{1(m-1)}@E_{n(m-1)}, B_{21}@E_{11}, \dots, B_{n(m-1)}@E_{n(m-1)}$

$O(n^2m^2)$

and

$B_{11}@D_{12}, B_{12}@D_{12}, \dots, B_{1(m-1)}@D_{12}, B_{21}@D_{12}, \dots, B_{n(m-1)}@D_{12},$

$B_{11}@D_{13}, \dots, B_{1(m-1)}@D_{13}, B_{21}@D_{13}, \dots, B_{n(m-1)}@D_{13},$

...

$B_{11}@D_{nm}, \dots, B_{1(m-1)}@D_{nm}, B_{21}@D_{nm}, \dots, B_{n(m-1)}@D_{nm}$

$O(n^2m^2)$

“Job-present-at-buffer” variables.

$B_{pq}@E_{rs}$ can be described as “whether job p is at buffer q, at the time of enqueue operation of task s of job r.

$B_{pq}@D_{rs}$ can be described as “whether job p is at buffer q, at the time of dequeue operation of task s of job r.

Constraints –

1. $(t_{11} > 0) \wedge (t_{12} \geq t_{11} + d_{11}) \wedge (t_{13} \geq t_{12} + d_{12}) \wedge \dots \wedge (t_{\text{horizon}} \geq t_{1m} + d_{1m})$
 $\wedge (t_{21} > 0) \wedge (t_{22} \geq t_{21} + d_{21}) \wedge (t_{23} \geq t_{22} + d_{22}) \wedge \dots \wedge (t_{\text{horizon}} \geq t_{2m} + d_{2m})$
 $\wedge \dots$
 $\wedge (t_{n1} > 0) \wedge (t_{n2} \geq t_{n1} + d_{n1}) \wedge (t_{n3} \geq t_{n2} + d_{n2}) \wedge \dots \wedge (t_{\text{horizon}} \geq t_{nm} + d_{nm})$

Constraints to ensure the order of the tasks of job don’t overlap.

$O(nm)$

2. $((t_{11} \geq t_{21} + d_{21}) \vee (t_{21} \geq t_{11} + d_{11})) \wedge \dots \wedge ((t_{n1} \geq t_{(n-1)1} + d_{(n-1)1}) \vee (t_{(n-1)1} \geq t_{n1} + d_{n1}))$
 $\wedge ((t_{12} \geq t_{22} + d_{22}) \vee (t_{22} \geq t_{12} + d_{12})) \wedge \dots \wedge ((t_{n2} \geq t_{(n-1)2} + d_{(n-1)2}) \vee (t_{(n-1)2} \geq t_{n2} + d_{n2}))$
 $\wedge \dots$
 $\wedge ((t_{1m} \geq t_{2m} + d_{2m}) \vee (t_{2m} \geq t_{1m} + d_{1m})) \wedge \dots \wedge ((t_{nm} \geq t_{(n-1)m} + d_{(n-1)m}) \vee (t_{(n-1)m} \geq t_{nm} + d_{nm}))$

Constraints to ensure that a machine does not get shared between tasks.

$O(nm)$

3. $(E_{11} == t_{11} + d_{11}) \wedge (E_{12} == t_{12} + d_{12}) \wedge \dots \wedge (E_{n(m-1)} == t_{n(m-1)} + d_{n(m-1)})$

Equality constraints on enqueue times.

$O(nm)$

4. $(D_{12} == t_{12}) \wedge (E_{13} == t_{13}) \wedge \dots \wedge (E_{nm} == t_{nm})$

Equality constraints on dequeue times.

$O(nm)$

5. $B_{11}@E_{11} = 1$ if $E_{11} \leq E_{11} < D_{12}$ (always true)
 0 otherwise
 $B_{12}@E_{11} = 1$ if $E_{12} \leq E_{11} < D_{13}$
 0 otherwise
...

ITE constraints on the “Job-present-at-buffer” variables at enqueue times.

$O(n^2m^2)$

6. $B_{11}@D_{12} = 1$ if $E_{11} \leq D_{12} < D_{12}$ (always false)
 0 otherwise
 $B_{12}@D_{12} = 1$ if $E_{12} \leq D_{12} < D_{13}$
 0 otherwise

...

ITE constraints on the “Job-present-at-buffer” variables at dequeue times.

$O(n^2m^2)$

7. $\Sigma (B_{11}@E_{11}, B_{21}@E_{11}, B_{31}@E_{11}, \dots, B_{n1}@E_{11}) \leq b_1$
 (i.e. Length of buffer 1 does not increase than its capacity at E_{11})
 $\Sigma (B_{11}@E_{21}, B_{21}@E_{21}, B_{31}@E_{21}, \dots, B_{n1}@E_{21}) \leq b_1$
 (i.e. Length of buffer 1 does not increase than its capacity at E_{11})
 ...
 $\Sigma (B_{11}@E_{n1}, B_{21}@E_{n1}, B_{31}@E_{n1}, \dots, B_{n1}@E_{n1}) \leq b_1$
 (i.e. Length of buffer 1 does not increase than its capacity at E_{n1})
 ...
 ...
 (i.e. Length of buffer m-1 does not increase than its capacity at $E_{1(m-1)}$ to $E_{n(m-1)}$)
 $O(mn)$
8. $\Sigma (B_{11}@D_{12}, B_{21}@D_{12}, B_{31}@D_{12}, \dots, B_{n1}@D_{12}) > 0$
 (i.e. Length of buffer 1 is not 0 at D_{11})
 ...
 ...
 $O(mn)$